The Relational Data Model and Relational Database Constraints

Part 1

Relational Model Concepts

- When a relation is thought of as a **table** of values, each row in the table represents a collection of related data values.
- A row represents a fact that typically corresponds to a real-world entity or relationship.
- The table name and column names are used to help to interpret the meaning of the values in each row.
- All values in a column are of the same data type.

- In the formal relational model terminology, a row is called a *tuple*, a column header is called an *attribute*, and the table is called a *relation*.
- The data type describing the types of values that can appear in each column is represented by a *domain* of possible values.

Domains, Attributes, Tuples, and Relations

- A **domain** D is a set of atomic values.
- By **atomic** we mean that each value in the domain is indivisible as far as the formal relational model is concerned.
- A common method of specifying a domain is to specify a data type from which the data values forming the domain are drawn.
- It is also useful to specify a name for the domain, to help in interpreting its values.
- Some examples of domains follow:
 - Usa_phone_numbers
 - Local phone numbers
 - Social security numbers
 - Names
 - Grade point averages
 - Employee ages
 - Academic department names
 - Academic department codes

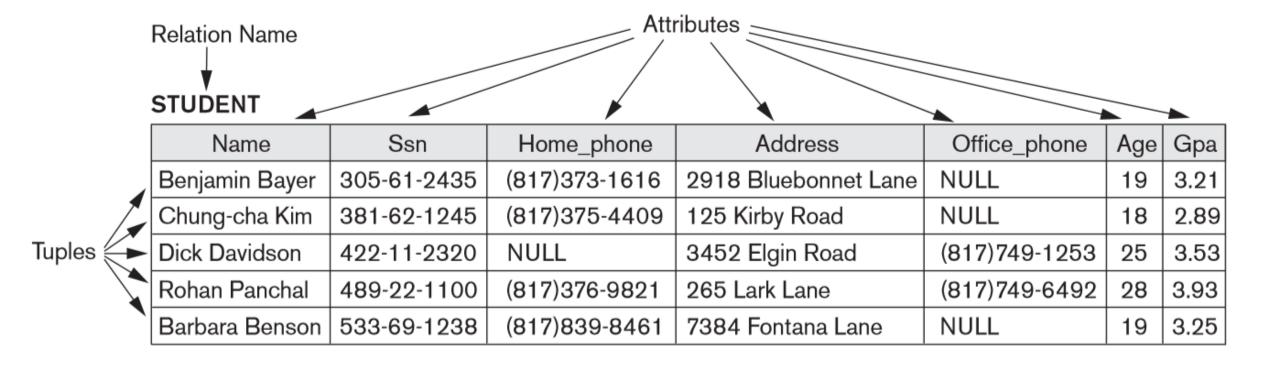
- A **relation schema** R, denoted by $R(A_1, A_2, ..., A_n)$, is made up of a relation name R and a list of attributes, $A_1, A_2, ..., A_n$.
- Each attribute A_i is the name of a role played by some domain D in the relation schema R.
- D is called the domain of A_i and is denoted by $dom(A_i)$.
- A relation schema is used to describe a relation; R is called the name of this relation.
- The **degree** (or **arity**) of a relation is the number of attributes *n* of its relation schema.

- A relation of degree seven, which stores information about university students, would contain seven attributes describing each student, as follows:
 - STUDENT(Name, Ssn, Home_phone, Address, Office phone, Age, Gpa)
- Using the data type of each attribute, the definition is sometimes written as:
 - STUDENT(Name: string, Ssn: string, Home_phone: string, Address: string, Office_phone: string, Age: integer, Gpa: real)

- More precisely, we can specify the following previously defined domains for some of the attributes of the STUDENT relation:
 - dom(Name) = Names
 - dom(Ssn) = Social_security_numbers
 - dom(HomePhone) = USA_phone_numbers
 - dom(Office phone) = USA phone numbers
 - dom(Gpa) = Grade point averages

• It is also possible to refer to attributes of a relation schema by their position within the relation; thus, the second attribute of the STUDENT relation is Ssn, whereas the fourth attribute is Address.

- A **relation** (or **relation state**) r of the relation schema $R(A_1, A_2, ..., A_n)$, also denoted by r(R), is a set of n-tuples $r = \{t_1, t_2, ..., t_m\}$.
- Each n-tuple t is an ordered list of n values $t = \langle v_1, v_2, ..., v_n \rangle$, where each value v_i , $1 \le i \le n$, is an element of dom (A_i) or is a special NULL value.
- The i^{th} value in tuple t, which corresponds to the attribute A_i , is referred to as $t[A_i]$ or $t.A_i$ (or t[i] if we use the positional notation).
- The terms **relation intension** for the schema R and **relation extension** for a relation state r(R) are also commonly used.



- The earlier definition of a relation can be restated more formally using set theory concepts as follows.
- A relation (or relation state) r(R) is a **mathematical relation** of degree n on the domains $dom(A_1)$, $dom(A_2)$, ..., $dom(A_n)$, which is a subset of the **Cartesian product** (denoted by \times) of the domains that define R:

$$r(R) \subseteq (\text{dom}(A_1) \times \text{dom}(A_2) \times ... \times \text{dom}(A_n))$$

Characteristics of Relations

Ordering of Tuples in a Relation

- A relation is defined as a *set* of tuples.
- Mathematically, elements of a set have *no order* among them; hence, tuples in a relation do not have any particular order.
- In other words, a relation is not sensitive to the ordering of tuples.

Ordering of Values within a Tuple and an Alternative Definition of a Relation

- According to the preceding definition of a relation, an n-tuple is an ordered list of n values, so the ordering of values in a tuple—and hence of attributes in a relation schema—is important.
- However, at a more abstract level, the order of attributes and their values is not that important as long as the correspondence between attributes and values is maintained.
- A **tuple** can be considered as a **set** of (<attribute>, <value>) pairs, where each pair gives the value of the mapping from an attribute A_i to a value v_i from dom(A_i).

Two identical tuples when the order of attributes and values is not part of relation definition.

t = < (Name, Dick Davidson),(Ssn, 422-11-2320),(Home_phone, NULL),(Address, 3452 Elgin Road), (Office_phone, (817)749-1253),(Age, 25),(Gpa, 3.53)>

t = < (Address, 3452 Elgin Road),(Name, Dick Davidson),(Ssn, 422-11-2320),(Age, 25), (Office_phone, (817)749-1253),(Gpa, 3.53),(Home_phone, NULL)>

Values and NULLs in the Tuples

- Each value in a tuple is an **atomic** value; that is, it is not divisible into components within the framework of the basic relational model.
- Hence, composite and multivalued attributes are not allowed.
- In general, we can have several meanings for NULL values, such as **value unknown**, **value** exists but is **not available**, or **attribute does not apply** to this tuple(also known as **value undefined**).
- The exact meaning of a NULL value governs how it fares during arithmetic aggregations or comparisons with other values.
- During database design, it is best to avoid NULL values as much as possible.

• Interpretation (Meaning) of a Relation

- The relation schema can be interpreted as a declaration or a type of assertion.
- An alternative interpretation of a relation schema is as a **predicate**.