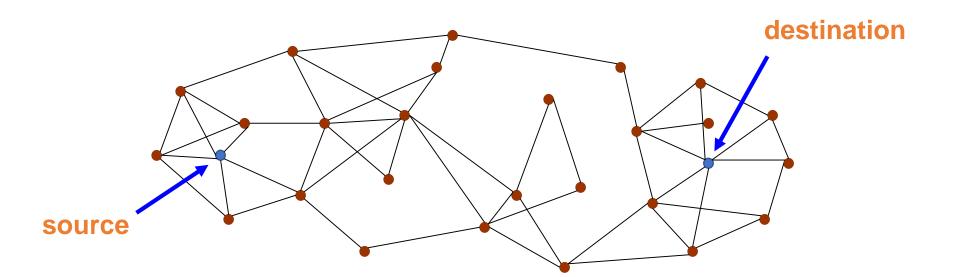
Data Structure Programming Project #1

郭建志

Background

- Large-scale geographic routing
- Every node only has the local information (i.e., its neighbors' information) and the destination's information



Background

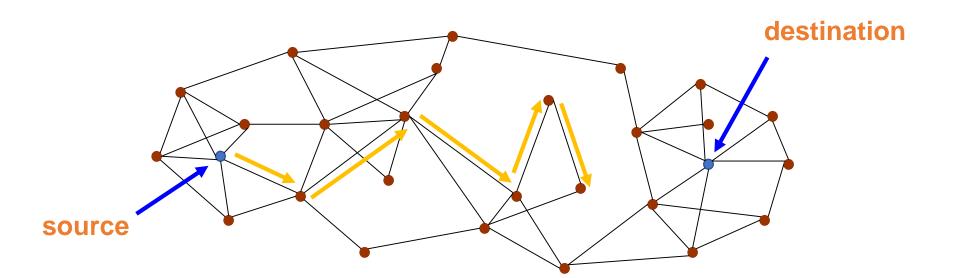
- Large-scale geographic routing
- Every node only has the local information (i.e., its neighbors' information) and the destination's information

 So, how should a source route a packet to a designated destination?



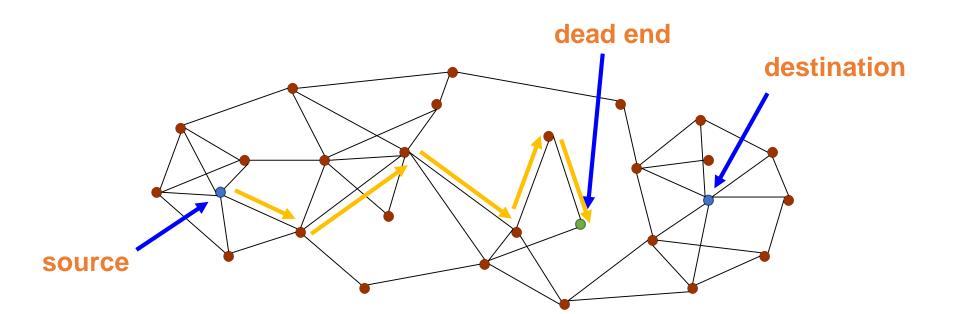
Naïve Routing Scheme: Greedy

• Forward the packet towards the neighboring node that is closer to the destination



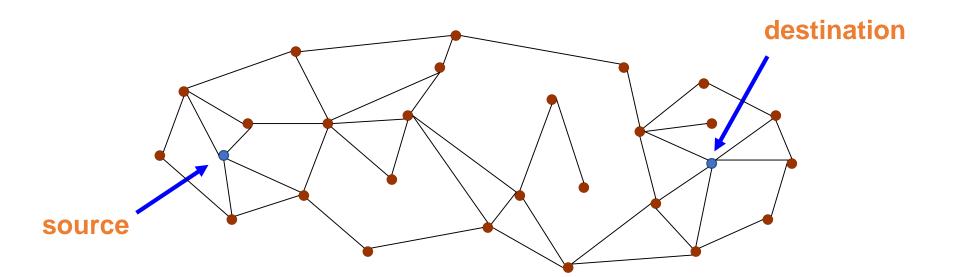
Naïve Routing Scheme: Greedy

- Forward the packet towards the neighboring node that is closer to the destination
- However, the packet may be stuck at the dead end



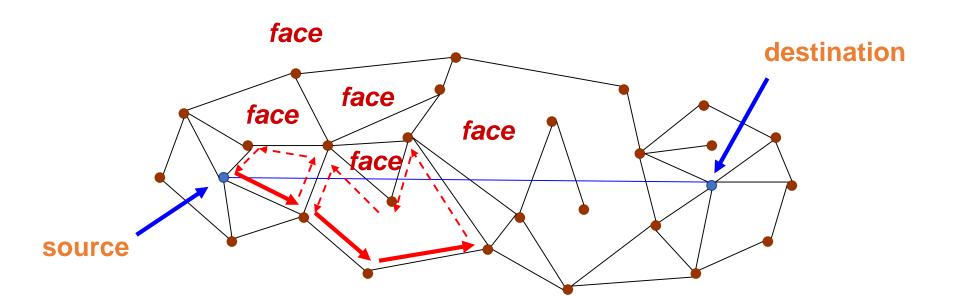
Face Routing

- Generate a planar subgraph: no crossing links
- Right-hand rule
- Change faces

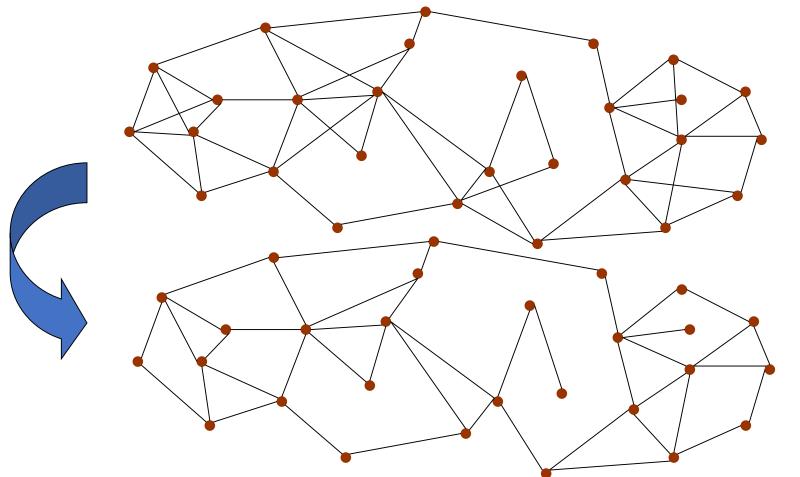


Face Routing

- Generate a planar subgraph: no crossing links
- Right-hand rule
- Change faces



Generate a planar subgraph: no crossing links



Programming Project #1:	Input file:		
Generate a Planar Subgraph			
	5		
• Input:	0	1.5	2.3
•	1	1.2	2.6
 Number of nodes 	2	1.3	1.2
 Nodes with non-negative coordinates (x, y) (the input graph is connected 	•••		
when we add links if dist(u,v) ≤ 1)			
• Procedure:	Out	put file:	
 Add a link between any two nodes u, v as dist(u,v) ≤ 1 	12		
 Remove some links to generate a planar 	0	0	1
graph	1	1	2
• Output:	2	1	3
• The edges before and after planarization	8		
 The grade is proportional to the 	0	0	1
number of remaining edges	2	1	3
namber of remaining eages	4	3	4

• Input:

- Number of nodes
- Nodes with non-negative coordinates (x, y) (the input graph is connected when we add links if $dist(u,v) \le 1$)

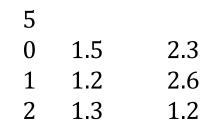
• Procedure:

- Add a link between any two nodes u, v as dist(u,v) ≤ 1
- Remove some links to generate a planar graph

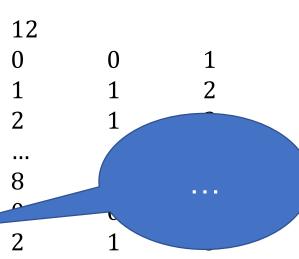
Output:

- The edges before and after planarization $\frac{1}{8}$
- The grade is proportional to the number of remaining edges

Input file:



Output file:



• Input:

- Number of nodes
- Nodes with non-negative coordinates (x, y) (the input graph is connected when we add links if $dist(u,v) \le 1$)

• Procedure:

- Add a link between any two nodes u, v as dist(u,v) ≤ 1
- Remove some links to generate a planar graph

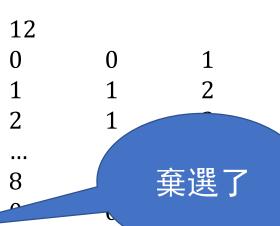
Output:

- The edges before and after planarization
- The grade is proportional to the number of remaining edges

Input file:

5		
0	1.5	2.3
1	1.2	2.6
2	1.3	1.2

Output file:



Programming Project #1: Generate a Planar Subgraph	Inp	ut file:	
cerrerate a riamar suspirapin	5		
Υ	0	1.5	2.3
• Input:	1	1.2	2.6
 Number of nodes 	2	1.3	1.2
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 Implement a designated algorithm 	0	0	1
to remove the links	2	1	3
	4	3	4

Input file: node.txt

• (O_1	ut	p'	u1	t:
			1		

0	1.5	2.3
1	1.2	2.6
2	13	1 2

• The edges before and after planarization

...

5

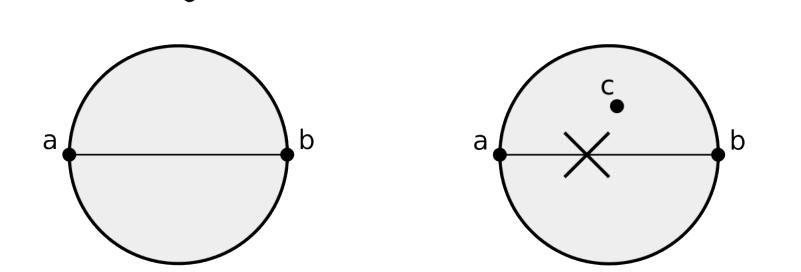
Notice that:

- The first node ID should be smaller than the second one for each link
- The links should be sequentially indexed in ascending order of the first node ID
- If there is a tie, the two links are indexed in ascending order of the second node ID
- The IDs of remaining links after planarization should be identical to the ones before the planarization

Output file: link.txt

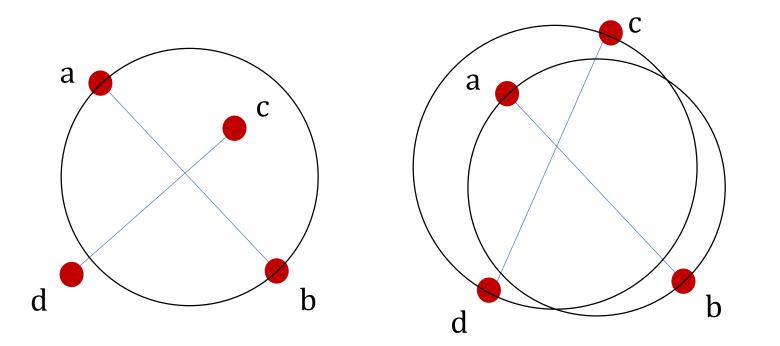
12		
0	0	1
1	1	2
2	1	3
8		
0	0	1
2	1	3
4	3	4

- Implement a designated algorithm to remove the links
- Cut a link between two neighboring nodes a and b, if there is other node c is within their the closed disc, where line segment \overline{ab} is a diameter



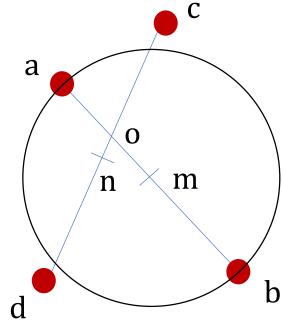
Why does the algorithm work?

- Prove by contradiction
- Assume that there is an intersection node on the segments \overline{ab} and \overline{cd}



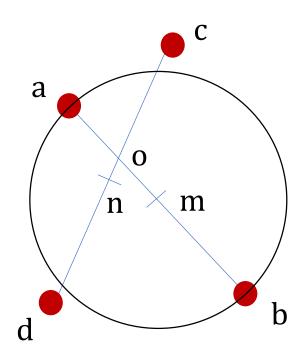
Why does the algorithm work?

- Let m and n denote the middle points of \overline{ab} and \overline{cd} , respectively
- Without loss of generality, we assume that \overline{am} and \overline{cn} intersect at node o
- Since node a is not in the circle of which \overline{cd} is the diameter, $\overline{cn} < \overline{an}$...(1)
- Similarly, since node c is not in the circle of which \overline{ab} is the diameter, $\overline{am} < \overline{cm}...(2)$



Why does the algorithm work?

- $\overline{cn} < \overline{an}...(1)$
- $\overline{am} < \overline{cm}...(2)$
- $(1)+(2) = \overline{cn} + \overline{am} < \overline{an} + \overline{cm}...(3)$
- Due to triangular inequality, $\overline{ao} + \overline{on} \ge \overline{an}...(4)$
- Similarly, $\overline{co} + \overline{om} \ge \overline{cm}...(5)$
- $(4)+(5) = \overline{ao} + \overline{om} + \overline{co} + \overline{on}$ = $\overline{am} + \overline{cn} \ge \overline{an} + \overline{cm}...(6)$
- \bullet (3) $\rightarrow \leftarrow$ (6)



Input Sample: node.txt

```
11
0
              1.1534
    1.4142
    1.97
              1.85
2
    0.8823
              1.3926
3
    1.9996
              1.9484
4
    0.6301
              1.9145
5
    0.3329
              1.1756
    1.952
              0.5866
6
7
    1.8181
              0.4043
8
    0.4949
              1.946
9
    0.2079
              1.4758
10
    0.1364
              1.4084
```

Output Sample: link.txt

22		
0	0	1
1	0	2
2	0	3
3	0	6
4	0	7
5	1	3
6	2	4
7	2	5
8	2	8
9	2	9
10	2	10
11	4	5
12	4	8
13	4	9
14	4	10
15	5	8
16	5	9
17	5	10
18	6	7
19	8	9
20	8	10
21	9	10

11		
0	0	1
1	0	2
3	0	6
5	1	3
6	2	4
7	2	5
12	4	8
17	5	10
18	6	7
19	8	9
21	9	10

Note

- Deadline: 10/11 Thu
- E-course
- C Source code
- Readme including:
 brief description and
 time complexity
 analysis of you
 implementation

