

Today's Quote →

**The man who asks a question
is a fool for a minute,
the man who does not ask
is a fool for life.**

{ Pen and }
Paper

Content

no. of iterations

- Time and Space Complexity
- Asymptotic Analysis
- Big.O Notation
- T.L.E → Time Limited Exceeded]

Tc-2.

① Sum of first N natural numbers = $\frac{N(N+1)}{2}$.

② How many no's are there in this range [3, 10].
↳ 8.

$$[a, b] = b - a + 1$$

$$[3, 10] \Rightarrow 10 - 3 + 1 = 8$$

$$[a, b) = b - a$$

{ [→ inclusive]
C → exclusive }

$$(a, b) = b - a - 1$$

③ How many times do we need to divide N by 2 to reduce it to 1?

$$N = 64 \rightarrow 32 \rightarrow 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1$$

$$N = 64 = 2^6 \quad \text{No. of steps} \rightarrow 6$$

$$N = 1024 = 2^{10} \quad \text{No. of steps} \rightarrow 10$$

$[\log_2 N \text{ times}]$

④ $\boxed{\log_a a^x = x}$

$$\log_2 2^{10} = 10$$

$$\log_3 3^5 = 5.$$

Arithmetic Progression

\Rightarrow difference b/w any two consecutive terms is fixed/constant

$$d = \underbrace{4, 7,}_{3} \underbrace{10, 13,}_{3} \underbrace{16, 19,}_{3} \dots$$

$$a \quad a+d \quad a+2d \quad a+3d \quad a+4d \quad \dots \quad a+(n-1)d$$

first term = a

common diff = d

$\left[\text{Sum of } N \text{ terms of an A.P} = \frac{n}{2} [2a + (n-1)d] \right]$

Geometric Progression (G.P)

ratio of 2 consecutive terms \rightarrow constant.

$$3, 6, 12, 24, 48, 96, \dots$$

$\curvearrowright \curvearrowright \curvearrowright \curvearrowright \curvearrowright$

Common ratio = r

$$a, ar, ar^2, ar^3, ar^4, \dots \quad a r^{n-1}$$

first term = a

common ratio = r

$$\left[\text{Sum of } N \text{ terms of G.P} = \frac{a [r^n - 1]}{r - 1} \right] \quad r > 1$$

$$= \frac{a [1 - r^n]}{(1 - r)} \quad r < 1$$

$\rightarrow \{ \text{class-10 NCERT} \}$

Q1 void fun(int N) {

 s = 0 ;

 for (i = 1 ; i <= N ; i++) { i : [1 , N] }

 s = s + i ;

 } $\Rightarrow N - 1 + 1 = N$ iterations.

 return s ;

}

1

Q2 int fun(int N) {

 s = 0 ;

 for (i = 1 ; i <= N ; i = i + 2) {

 s = s + i ;

}

iterations -

N = 12

i → 1, 3, 5, 7, 9, 11

$\frac{12}{2} = [6]$

$\left. \begin{array}{l} \\ \\ \\ \\ \end{array} \right\} \frac{(n+1)}{2}$

N = 13

i → 1, 3, 5, 7, 9, 11, 13

$7 = \frac{(13+1)}{2} = \frac{14}{2}$

no. of iterations = $\frac{(n+1)}{2}$.

Q3) void func (int N, int M) {

```
    for (i=1 ; i <= N ; i++) { : i → [1, N]
        if (i % 2 == 0) {
            print(i)
        }
    }

    for (i=1 ; i <= M ; i++) { : i → [1, M]
        if (i % 2 == 0) {
            print(i)
        }
    }

    } # total no. of iterations = N + M.
```

Q4) int fun (int N) {

```
    s = 0;
    for (i=0 ; i <= 100 ; i++) { : i → [0, 100]
        s = s + i + i^2;
    }

    return s;
}

# no. of iterations = 101
```

Q.s.)

```

void fun(N){
    for( i=1; i*i <= N; i++) {
        s = 1+i*i;
    }
    return s;
}

```

$i \geq 1 \leq N$
 $i^2 \leq N$
 $i \leq \sqrt{N}$
 $i : [1, \sqrt{N}]$

no. of iterations = \sqrt{n} .

Q.b)

```

void fun(int N){
    i = N;
    while (i > 1) {
        i = i / 2;
    }
}

```

iterations	value of i
1	$N/2^1$
2	$N/2^2$
3	$N/2^3$
4	$N/2^4$
5	$N/2^5$
...	
	1.

$$\frac{N}{2^0} \rightarrow \frac{N}{2^1} \rightarrow \frac{N}{2^2} \rightarrow \frac{N}{2^3} \rightarrow \frac{N}{2^4} \rightarrow \frac{N}{2^5} \rightarrow \dots 1$$

if $i = 1$, loop will break.

After k iterations, loop will break.

$$\Rightarrow i = 1 = \frac{N}{2^k}$$

$$\begin{aligned}
 2^k &= N \\
 \log_2 2^k &= \log_2 N \\
 \Rightarrow K &= \boxed{\log_2 N} \\
 \text{# no. of iterations} &= \log_2 N
 \end{aligned}$$

Q) void fun(int N){

```

    i = N;
    while (i > 1){
        i = i / 3;
    }
  
```

$$N \rightarrow \frac{N}{3} \rightarrow \frac{N}{9} \rightarrow \frac{N}{27} \rightarrow \frac{N}{81} \rightarrow \dots \quad \textcircled{1}$$

After K iterations, loop will break.

$$1 = \frac{N}{3^K}$$

$$3^K = N$$

$$\log_3 3^K = \log_3 N \Rightarrow \boxed{K = \log_3 N}$$

Amazon MCQ → Multiple choice question

Q7) void fun(int N){

```

    {
        s=0;
        for( i=0 ; i < N ; i=i*2){
            {
                s=s+i;
            }
        }
    }
  
```

$0 \rightarrow 0 \rightarrow 0 \rightarrow 0 \rightarrow 0$

If no. of iterations → infinite.

Q8) void fun(N){

```

    {
        s=0;
        for( i=1 ; i < N ; i=i*2){
            {
                s=s+i;
            }
        }
    }
  
```

<u>iterations</u>	<u>value of i</u>
1	2
2	2^2
3	2^3
4	2^4
5	2^5
⋮	⋮

$1 \rightarrow 2 \rightarrow 4 \rightarrow 8 \rightarrow 16 \rightarrow$

if $i == N$, loop will break.

After k iterations, loop will break.

$$N = 2^k$$

$$\log_2 N = \log_2 2^k \Rightarrow \boxed{\log_2 N = k}$$

no. of iterations = $\log_2 N$.

Nested loops

```
(Q9) void fun (int N){  
    for ( i=1 ; i <=10 ; i++) {  
        for ( j=1 ; j <=N ; j++) {  
            print(--)  
        }  
    }  
}
```

table:

value of i	j	iterations
1	[1, N]	$1 \times N$
2	[1, N]	$2 \times N$
3	[1, N]	$3 \times N$
...		\dots
10	[1, N]	$10 \times N$

total no. of iterations = $10N$

```
(Q10) void func ( int N){  
    for ( i=1 ; i <=N ; i++) {  
        for ( j=1 ; j <=N ; j++) {  
            print(i*j);  
        }  
    }  
}
```

table

i	j	iterations
1	[1, N]	$1 \times N$
2	[1, N]	$2 \times N$
3	[1, N]	$3 \times N$
...		\dots
N	[1, N]	$N \times N$

no. of iterations = $N \times N = N^2$.

Q11)

```

void fun( int N ) {
    for( i=1 ; i <= N ; i++ ) {
        for( j=1 ; j <= i ; j++ ) {
            print( i+j );
        }
    }
}

```

i	j	iterations
1	[1,1]	1
2	[1,2]	2
3	[1,3]	3
...		-
N	[1,N]	N

total no. of iterations = $\frac{N(N+1)}{2}$.

Q12)

```

void fun( int N ) {
    for( i=1 ; i <= N ; i++ ) {
        for( j=1 ; j <= N ; j=j*2 ) {
            print( i+j );
        }
    }
}

```

i	j	iterations
1	[1,N]	$\log_2 N$
2	[1,N]	$\log_2 N$
3	[1,N]	$\log_2 N$
...		-
N	[1,n]	$\log_2 N$

total no. of iterations = $\log_2 N * N$
 $= N \log_2 N$.

Q14)

```

void fun( int N ) {
    for ( i=1 ; i <= N ; i++ ) {
        for ( j=1 ; j <= 2i ; j++ ) {
            printf( "%d" );
        }
    }
}

```

i	j	iterations
1	[1, 2]	2
2	[1, 2 ²]	$4 = 2^2$
3	[1, 2 ³]	$8 = 2^3$
\vdots	\vdots	\vdots
N	[1, 2 ^N]	2^{2^N}

total no. of iterations = $2^1 + 2^2 + 2^3 + \dots + 2^N$

$$= \frac{2 [2^N - 1]}{2 - 1} = \underline{\underline{2 [2^N - 1]}}$$

$$\frac{2 [2^N - 1]}{2 - 1}$$

Q1

```

for( i = N ; i > 0 ; i = i / 2) {
    for( j = 1 ; j <= i ; j++) {
        print(i * j)
    }
}

```

\Leftarrow

		i	j	iterations
		N	[1, N]	$N/2^0$
		$N/2$	[1, $N/2$]	$N/2^1$
		$N/4$	[1, $N/4$]	$N/2^2$
		:		$N/2^3$
				1
				$\frac{N}{2^{\log_2 N}}$

total no. of iterations = $N + \frac{N}{2^1} + \frac{N}{2^2} + \frac{N}{2^3} + \dots + \frac{N}{2^{\log_2 N}}$

$$\begin{aligned}
&= N + N \left[\frac{1}{2^1} + \frac{1}{2^2} + \frac{1}{2^3} + \dots + \frac{1}{2^{\log_2 N}} \right] \\
&\quad \underbrace{\qquad\qquad\qquad}_{S = \frac{N-1}{2}} \\
&= N + \cancel{N} \left[\frac{N-1}{\cancel{N}} \right] = \boxed{2N-1}
\end{aligned}$$

$$S = \frac{1}{2} + \frac{1}{2^2} + \frac{1}{2^3} + \dots + \frac{1}{2^{\log_2 N}}$$

$$S = \frac{\cancel{1/2} \left[1 - \left(\frac{1}{2}\right)^{\log_2 N} \right]}{\cancel{\left(1 - \frac{1}{2}\right)}}$$

$$\frac{a[1 - r^n]}{[1 - r]}$$

$$\begin{cases} a = \frac{1}{2} \\ r = \frac{1}{2} \\ \text{term} = \log_2 N \end{cases}$$

$$S = 1 - \frac{1}{2^{\log_2 N}} = 1 - \frac{1}{2} = \frac{N-1}{N}.$$

$\Rightarrow \boxed{\log_a x = y}$ \Rightarrow Property of log

$$\Rightarrow \frac{\log_2 2^2}{2} = \frac{\log_2^2 2}{2} = \frac{2}{2} = 1$$

Comparison of terms { large value of N }

$$N [\log_2 N] < [\sqrt{N}] \cdot N$$

$$N [\sqrt{N}] < [N] \cdot N$$

$$\left\{ \log_2 N < \sqrt{N} < N < N \log N < N \sqrt{N} < N^2 < 2^N \right\}$$

$$\underline{\underline{N \rightarrow 2^{10}}}$$

$$\underline{\underline{\log_2 2^{10}}} \\ \Downarrow \\ 10$$

$$\sqrt{2^{10}} \\ \Downarrow \\ 2^5 = \underline{\underline{32.}}$$

$$\log N < \sqrt{N}.$$

N^2	2^N
$(10)^2$	(2^{10})
$\underline{100}$	$\underline{1024.}$

How to calculate Big-O? what? why?] → Next class.

- ① Calculate no. of iterations
- ② Ignore lower order terms.
- ③ Ignore constant / coefficients.

$$f(N) = 10N^2 + \underbrace{100N^1 + 10^3 N^0}_{\times} \\ = 10N^2 \\ \Rightarrow O(N^2)$$

(Q) $f(N) = \underbrace{4N^2}_{\times} + \underbrace{3N^1}_{\times} + \underbrace{10^5 N^0}_{\times} \Rightarrow O(N^2)$

(Q) $f(N) = \underbrace{4N}_{\times} + \underbrace{3N \log N}_{\times} + \underbrace{10^6}_{\times} \Rightarrow O(N \log N)$

(Q) $f(N) = \underbrace{4N \log N}_{\times} + \underbrace{3N \sqrt{N}}_{\times} + \underbrace{10^6}_{\times} \Rightarrow O(N \sqrt{N})$

→ {21 Question.}

Assignment / H.W question.

→ fable is the ultimate truth.

i=6

while ($i > 1$) {
 $i = i / 2$
}

$\log_2 6$ ↗

→ {Send me the code.}

for ($i = 1 ; i \leq n ; i *= 2$) {
 for ($j = 1 ; j \leq n ; j++$) {
 //
 }
}

i	j	iterations
1	[1, N]	2
2	[1, N]	2 + 2
4	[1, N]	2 + 2 + 2
8	[1, N]	2 + 2 + 2 + 2
16		-

$$\boxed{N} \quad | \quad (1, N) \quad | \quad \overset{+}{N}$$

1, 2, 4, 8, 16, ——
K-terms.

$$N * (\log_2 N + 1)$$

$$\Rightarrow O(N \log N)$$

$$2^{K-1} = N$$

$$\log_2 2^{K-1} = \log_2 N$$

$$\left\{ K = \log_2 N + 1 \right\}$$

\Rightarrow Quiz \rightarrow $\left\{ \begin{array}{l} \text{class performance} \\ \text{interactive} \end{array} \right\}$

\rightarrow All class.