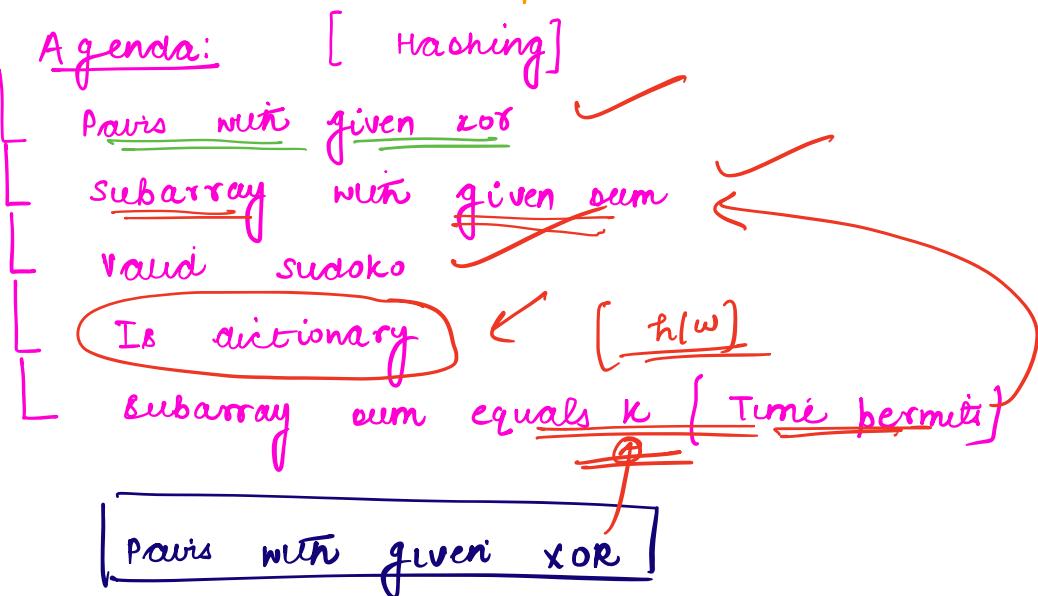


class starts at 7:37 a.m



Given: Given an array of distinct integers ↗ N

To find: NO of pairs with given XOR

Example: A = [ 3, 6, 8, 10, 15, 50 ]

$$B = \boxed{5}$$

3 ^ 6

$$\begin{array}{r} 0 \ 1 \ 0 \\ 1 \ 1 \ 0 \\ \hline 1 \ 0 \ 1 \end{array} \rightarrow 5.$$

$$10^{\wedge} 15$$

L ans=2

XOR: 1 → diff values  
0 → same values

## Brute force:

Go through all one pairs       $\text{for } (i=0; i < n; i+1) \{$   
                                          $\text{for } (j=i; j < n; j+1) \{$   
                                         find xor of that pair }  
                                         if xor == B  
                                         increase my ans

$$\text{T.C} = O(n^2)$$

## Optimized approach

observation:

$A[i] \wedge A[j] = B$

2 indexes:  $i$        $j$

$(n^2)$

1 index —  $O(n)$

## mathematical concept

$$[a[i] \wedge a[j] = B]$$

If and only if

$$a[i] \wedge B = a[j]$$

$$\underline{a[i]} \ ^n \ a[j] = B$$

$$a = a[i]$$

$$b = a[j]$$

$$x = B$$

$$a^{\wedge} \left( a^{\wedge} b \right) = a^{\wedge} x$$

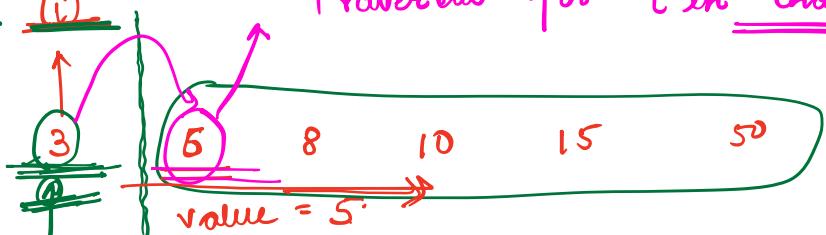
$$0^{\wedge} b = a^{\wedge} x$$

$$b = a^{\wedge} x$$

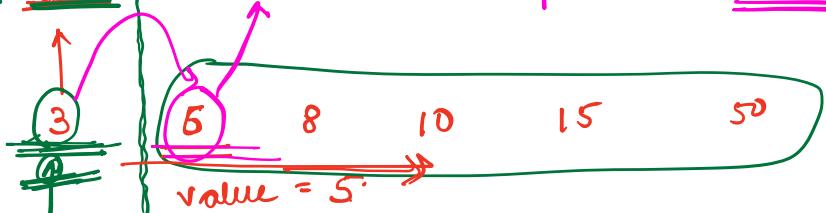
$$a[j] = a[i]^{\wedge} x$$

Algorithm

for [ init  $i=0$  ;  $i < \text{arr.length}$  ;  $i++$  ]

Dry run  $i =$  

Traversal for  $i$ th index



$$a[i] = \frac{3}{0} \ ^n \ \frac{5}{1} = \boxed{6}$$

$$6^{\wedge} 5 = 3$$

ans++

$$\begin{array}{ccc} 1 & 0 & 1 \\ \hline 1 & 1 & 0 \end{array} \text{ hash: } \left[ \begin{array}{c} \cancel{\frac{6}{3}} \\ \hline \end{array} \right]$$

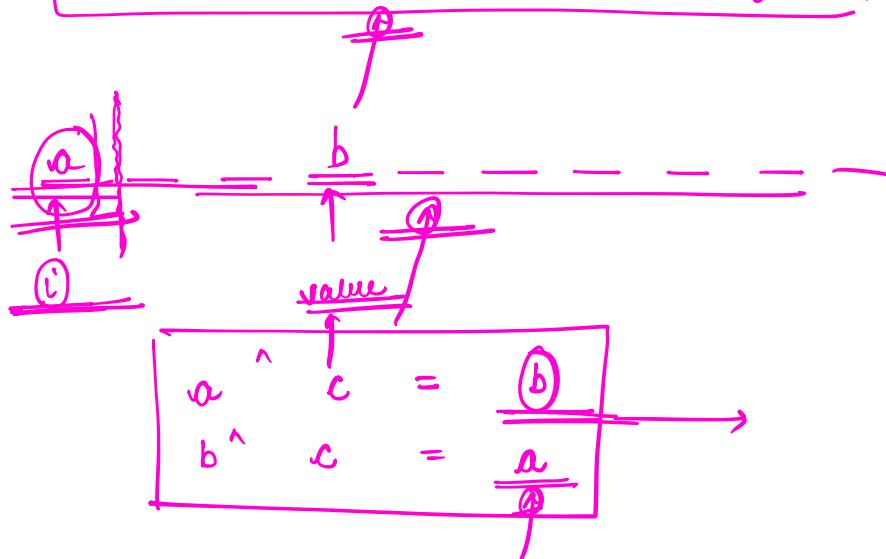
hashmap  
(or)  
hashset

count?

Distinct integers = hashset

Duplicates : hashmap

↳ frequency

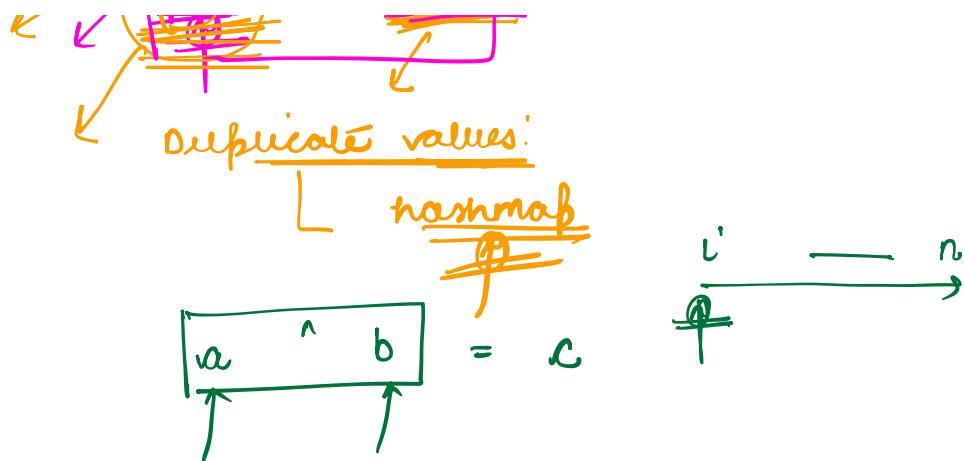


If insert every thing in map  
↳ duplicates pairs

[ 3, 6, 8, 10, 15, 30 ]

hashset: [ 3, 6, 8, 10, 15, 30 ]

$$x = \boxed{b}$$



Try to store that element which you are 100 %. sure, that is part of array

Pseudo code:

```

    solve ( int[] A , int x ) {
        Set < Integer > set = new HashSet();
        for ( i = 0 ; i < n ; i ++ ) {
            if ( set . contains ( A [ i ] ^ x ) )
                count++;
            set . put ( A [ i ] );
        }
        return count;
    }
}

```

Subarray with given sum

Given: A [] → positive integers

B

/ ~~if~~

To find find and return the {first}   
continuous subarray with sum B.

Example:  $A = [1, 2, 3, 4, 5]$   
 $B = 5$

↙  
subarray 1 =  $\underline{[2, 3]}$  → 5  
subarray 2:  $[5] = 5$

Brute force:

↳ go through all subarrays  
    | find sum  
    | if sum == B  
    | return the subarray

T.C:-

$O(n^3)$

$O(n^2)$

$n^3$

{  
    for( $i=0$ ;  $i < n$ ;  $i++$ ) |  
        for( $int j=i$ ;  $j < n$ ;  $j++$ ) |  
            int sum=0; |  
            for( $int k=i$ ;  $k < j$ ;  $k++$ ) |  
                sum += arr(k); |

$n^2$

for( $i=0$ ;  $i < n$ ;  $i++$ ) |  
    int sum=0; |  
    for( $int j=i$ ;  $j < n$ ;  $j++$ ) |  
        sum = sum + A[j]; |

checky

Bouté force app | Bouté approach)

Bouté force app | Bouté approach)

$$\begin{array}{r}
 \underline{\textcircled{1}} \\
 \hline
 \underline{1 + 2} \\
 \hline
 \underline{1 + 2 + 3} \\
 \uparrow \\
 \underline{o(n)} \\
 \cancel{f}
 \end{array}$$

A hand-drawn diagram illustrating a mathematical concept. At the top, the numbers 1 through 8 are written in red. Below this, a horizontal line has segments labeled 1, 2, and 3. An oval encloses the sum of these three segments. A pink arrow points from this oval to a second horizontal line below it, which also has segments labeled 1, 2, and 3. The number 4 is circled at the end of the third segment on this second line.

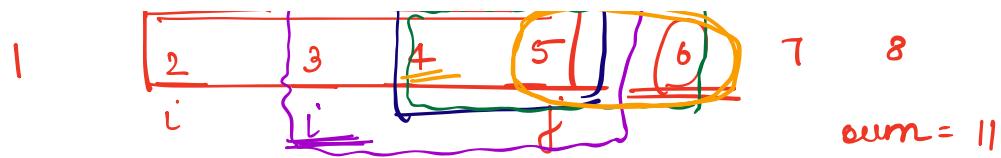
## optimization

start index  $i \rightarrow$  2

end index  $j \rightarrow$  5

calculated sum = 9

$9 < 11$



calculated sum = 14

$14 > 11$  [ j++ X ]

cal-sum(3, 4, 5) = 12

12 > 11

calculated -sum(4, 5) = 9.

9 < 11

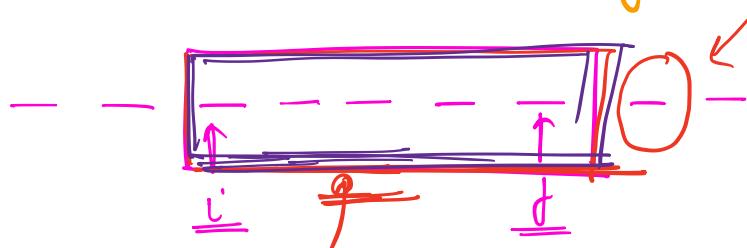
cal-sum(4, 5, 6) = 15

15 > 11

cal-sum(5, 6) = 11

11 == 11

L return my ans



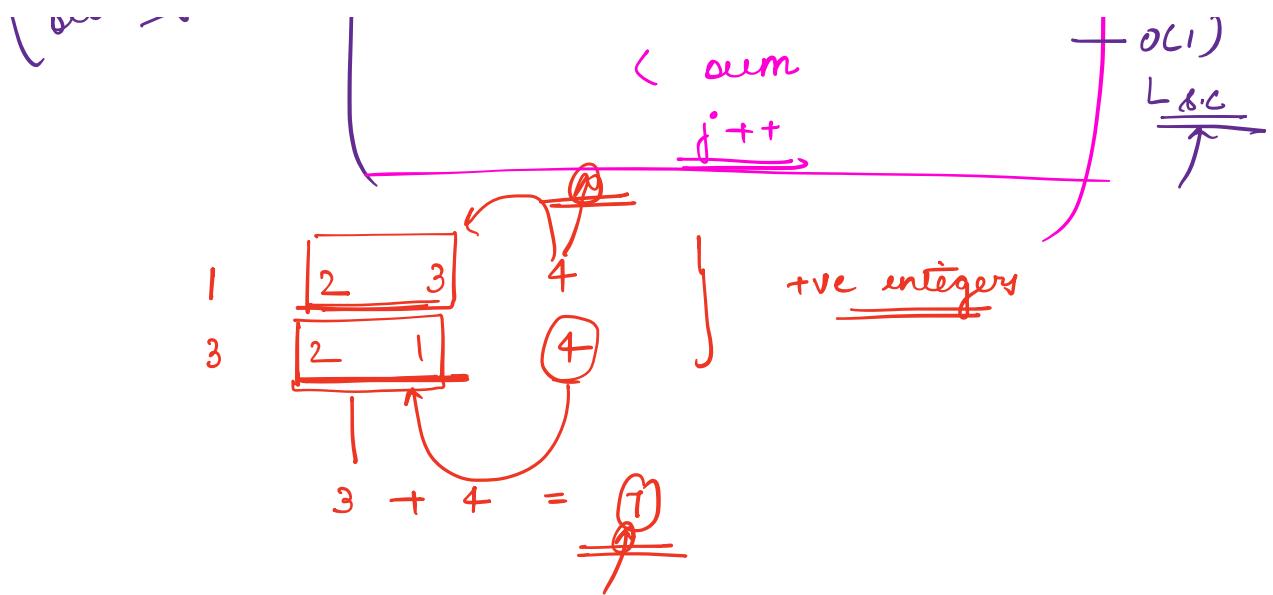
ans[i, j] = sum

( return my ans )

> sum  
i++

$O(n)$

~~algo~~  
~~sliding window~~



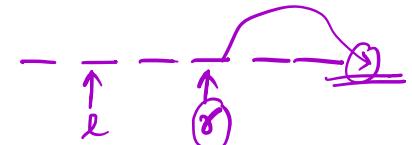
Pseudo code:

```
fun (int[] A, int x)
```

```
int  $l$  = 0;
```

1 2 3 4

```
int  $r$  = 0;
```



```
int sum = A[l];
```

```
while ( $r < n$ ) {
```

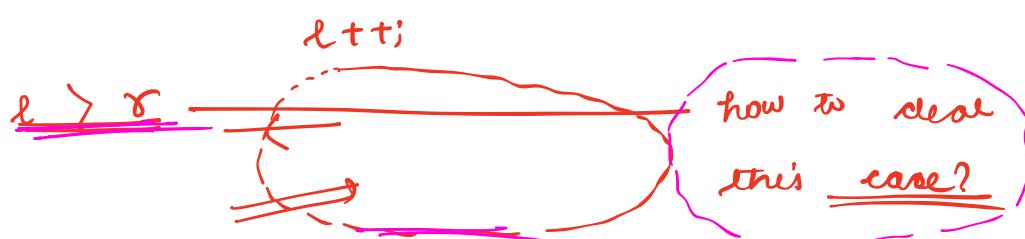
```
if (sum == x) {
```

return the subarray from  
 $l$  to  $r$ ;

```
}
```

```
else if ( $sum > x$ )
```

```
sum -= A[l]
```



else if ( $\text{sum} < x$ )  
~~=~~  $\text{sum} += A[r]$

Diagram illustrating a linked list with 6 nodes. The first node contains the value 1. A brace on the left groups the first two nodes. An arrow labeled "sum = 102" points to the third node, which contains the value 3. The fourth node contains 4, the fifth 5, and the sixth 6.

Break till: 9:00

Valid Sudoku

5	3	.	.	7	.	.	.	.
6	1	9	5	.	.	.	.	.
9	8	.	.	.	.	.	6	.
8	.	.	.	6	.	.	.	3
4	.	.	8	.	.	.	.	1
7	.	.	.	2	.	.	.	6
6	.	.	.	1	9	2	8	5
.	.	.	.	8	.	7	9	.

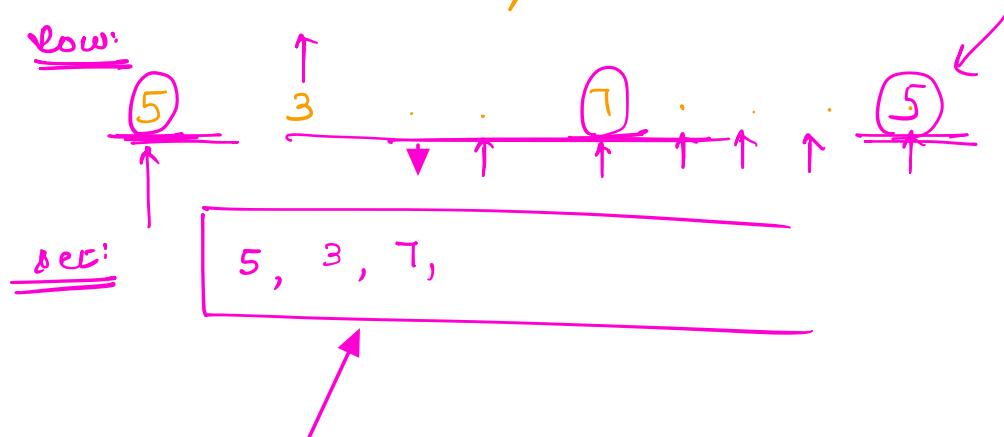
### Rules:

Every row should have distinct values  
(invalid) ① 2 3 ① 4 5 8 9 7

every col should have distinct values  
 Every box should have a distinct value

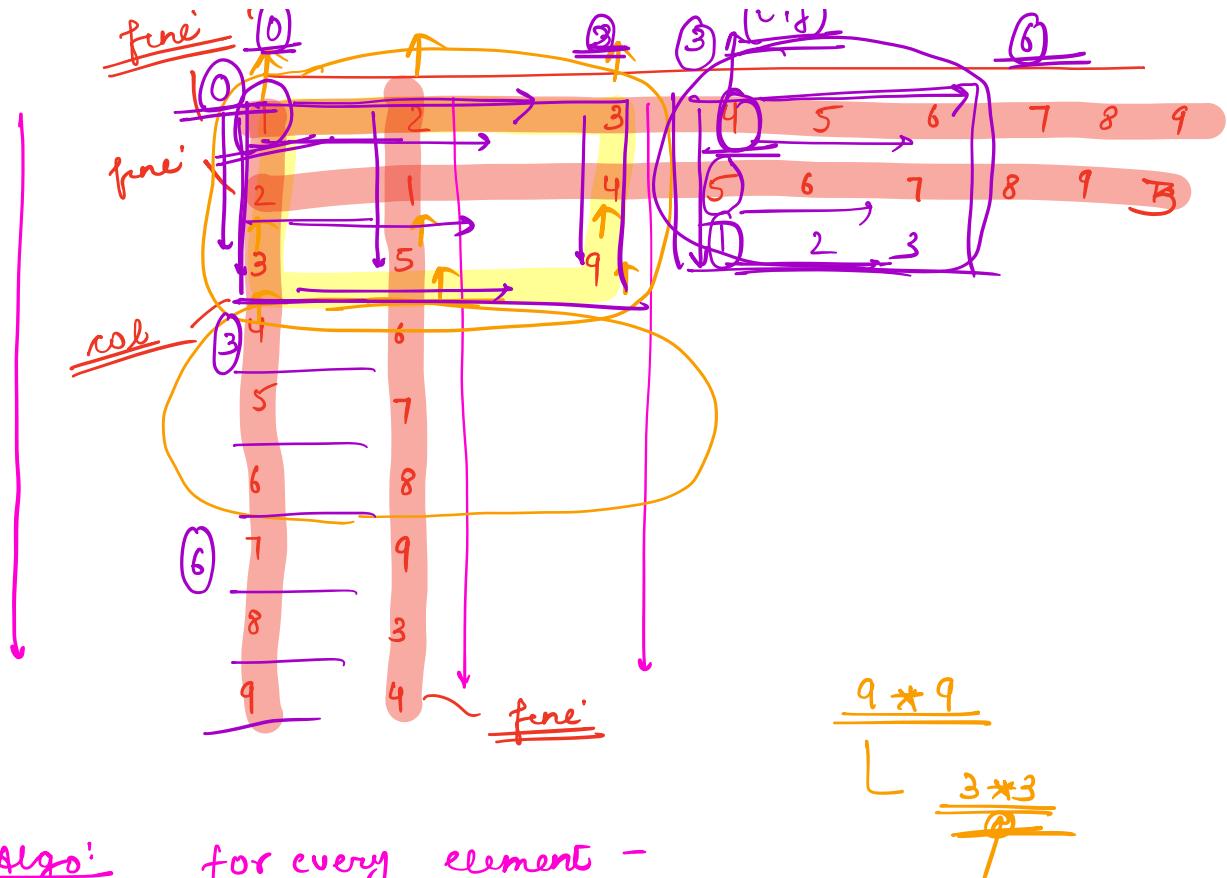
Please note, we need to check whether the row is valid for existing value only

DS :- Set | Hashset  
 ↗



This set stores the already visited values

T.C:  $O(n^2)$



Algo: for every element -  
check whether it exists in my  
row, col and box

$$\underline{\underline{TC:}} \quad O(n^2)$$

Pseudo code:      mat[][]      Hashset [set]

```

    // check for row
    for( i=0; i < n; i++ ) {
        set::clear();
        for(j=0; j < q; j++) {
            if ( mat[i][j] == '.' ) {
                continue;
            }
        }
    }
    
```

```

    if ( set.contains( mat[i][j] ) )
        return false;
    set.add( mat[i][j] );
}

```

// check for cols [Assignment]

// check for box

mat 16 \* 16 [ Box = 4 \* 4 ]

```

for ( int i = x; i < x + n; i++ ) {
    for ( int j = y; j < y + n; j++ ) {
        if ( c == mat[i][j] ) {
            set.add( mat[i][j] );
        }
    }
}

if ( set.contains(c) ) {
    return false;
}

```

*for one box*

$$\boxed{9 * 9} \rightarrow n = \underline{\underline{9}}$$

$\begin{array}{c} 0 \\ y \\ y+3 \\ \hline (y+3)+3 \end{array}$

$x+3$   
 $\underline{x+6}$        $\underline{\underline{6}} \quad (x+3) + 3$

$$\underline{k=0} \longrightarrow \underline{k=0} \quad x = 0 / 3 = \underline{\underline{0}} \quad \checkmark$$

$$y = 0 \% 3 = \underline{\underline{0}}$$

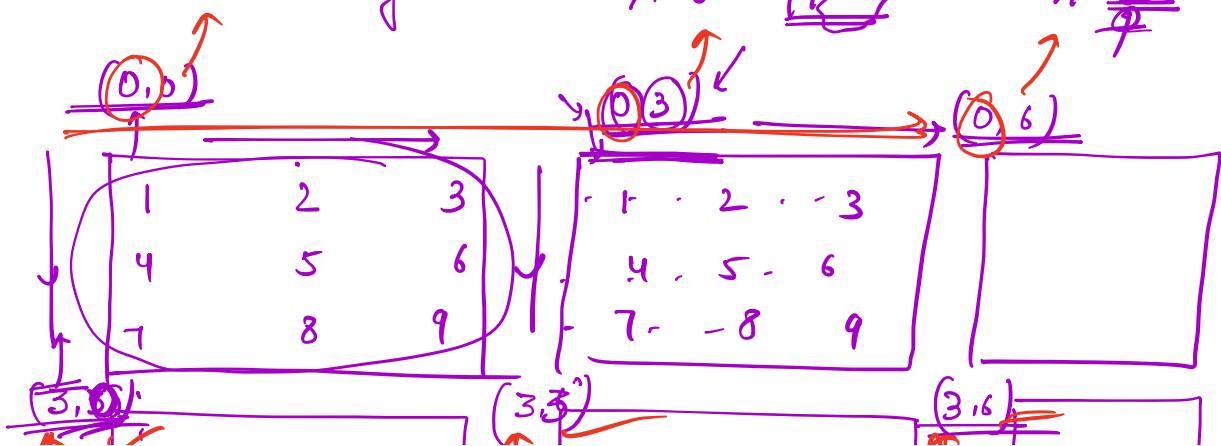
$$i = x \quad i < x + 3$$

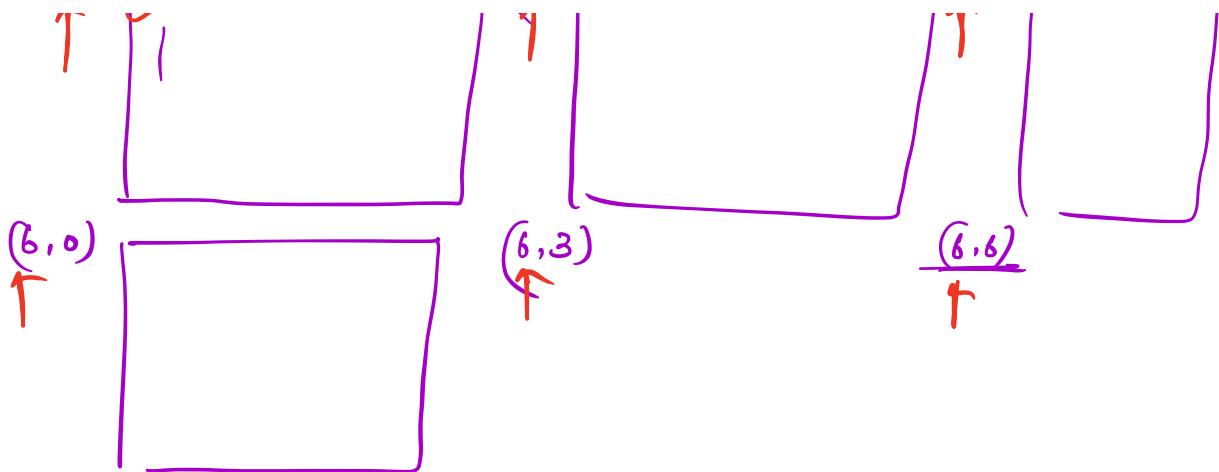
$$j = y : \quad j < y + 3$$

$$9 * 9$$

$\lfloor$  no of boxes =  $\underline{\underline{9}}$

for ( until  $box = \underline{\underline{0}}$ ;  $box < 9$ ;  $box++$ ) {  
 $x = \frac{box}{n}; \quad // 0$   
 $y = box \% n; \quad // \underline{\underline{0}}$



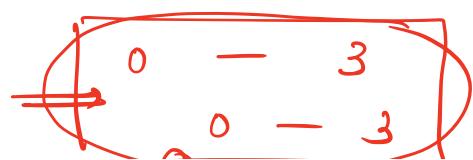


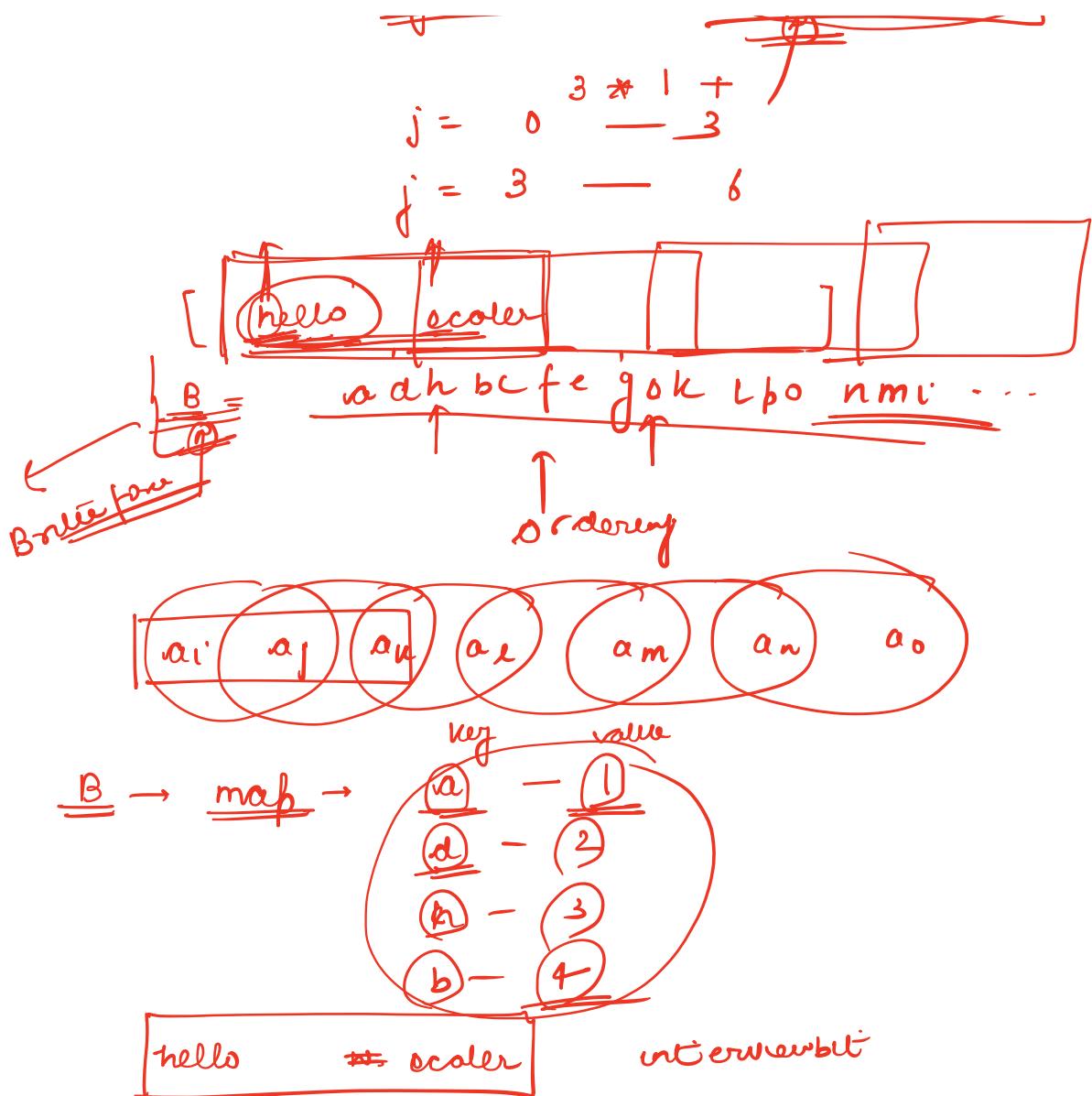
request: Go through the box with implementation

~~for( int box = 0; box < n\*n; box++ )~~  
 int ~~x = box / n; // ①~~  
 int ~~y = box % n; // ②~~  
 for( int i = ~~x \* n + 0~~; i < ~~n \* x + n~~; i++ )  
 for( int j = ~~y \* n + 0~~; j < ~~n \* y + n~~; j++ )  
 // all my checks

}

~~box = 0; x = 0~~  
~~y = 0~~





~~$O(n^2)$~~

for (int a = w[i-1], b = v[i]; smallest, smallest.length)

for (i = smallest.length; map.get(a.charAt(i)) != map.get(b.charAt(i)); i++)

for (int i=0; i < n.length; i++)  
 {  
 string a = word[i-1];  
 string b = word[i];  
 if (map.contains(a))  
 map.get(b)++;  
 else  
 map.put(b, 1);  
 if (map.get(b) == n)  
 return true;  
 }

ayesh

ayush

af