

Outline

1 More Operators for Binary Coded GA

- Selection Operators ✓
- Crossover Operators ✓
- Mutation Operators ✓

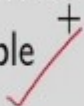
2 Simulation and Application ✓

- Rosenbrock Function
- Himmelblau Function
- Rastrigin Function
- Ackley Function

3 Closure⁺ ✓

Recap

- Binary-coded genetic algorithm
 - ▶ Generalized framework
 - ▶ Solution representation
 - ▶ Working principles through an example
 - ★ selection operator,
 - ★ crossover and mutation operators,
 - ★ survivor operator
 - ▶ Graphical example⁺



Selection Operators

- The purpose is to identify good (usually above-average) solutions in the population.
- In this process, we eliminate bad solutions in the population.
- We make multiple copies of good solutions.
- Selection can be used either before or after search/variation operators.
 - ▶ When selection is used before search/variation operators, it is called as reproduction or selection.
 - ▶ After search/variation operators: The process of choosing the next generation population from parent and off-spring populations is referred to as survivor or elimination. It is also being referred to as environmental selection.
- Common selection methods are
 - ▶ Fitness proportionate selection ✓
 - ▶ Tournament selection ✓
 - ▶ Ranking selection ✓

Fitness Proportionate Selection Operator


- **Probability** of selecting individual i from population P of size N is:

$$p_i = \frac{f_i}{\sum_{j=1}^N f_j}, \quad (1)$$

where f_i is the fitness of individual i .

- It is suitable for maximization problem. ✓
- It is also known as roulette wheel selection method. +

Fitness Proportionate Selection Operator

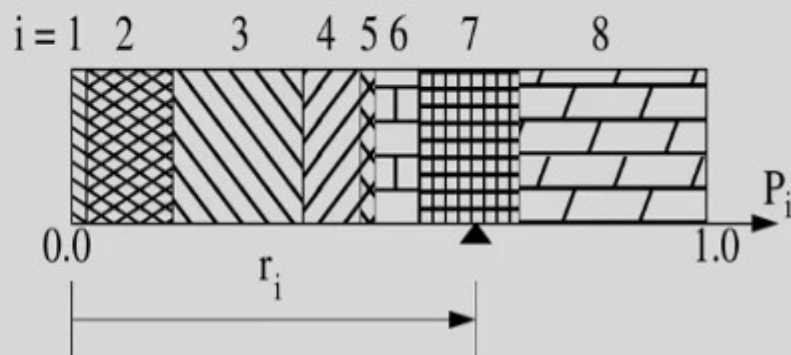


Index	$f(x_1, x_2)$	p_i	P_i
1	210.445	0.0041	0.0041
2	7536.407	0.1471	0.1512
3	14041.882	0.2740	0.4252
4	1546.603	0.0302	0.4554
5	189.732	0.0037	0.4591
6	671.924	0.0131	0.4122
7	7252.209	0.1415	0.6137
8	19793.183	0.3863	1.0000
sum	51242.386		

- Let us calculate probability of solution '1'.
 - ▶ Sum of all fitness values
 $= \sum_{j=1}^N f_j = 51242.386$
 - ▶ Probability, $p_1 = \frac{210.445}{51242.386} = 0.0041$
- Similarly, we can calculate probability of solution '2'.
 - ▶ Probability, $p_2 = \frac{7536.407}{51242.386} = 0.1471$
- We then calculate probability of other solutions.
- P_i is the cumulative probability.

Fitness Proportionate Selection Operator

- The figure is drawn using the cumulative probability.
- r_i is the random number between 0 to 1.



Index	$f(x_1, x_2)$	p_i	P_i
1	210.445	0.0041	0.0041
2	7536.407	0.1471	0.1512
3	14041.882	0.2740	0.4252
4	1546.603	0.0302	0.4554
5	189.732	0.0037	0.4591
6	671.924	0.0131	0.4722
7	7252.209	0.1415	0.6137
8	19793.183	0.3863	1.0000

- Let the first random number is $r_1 = 0.07218$. As per the cumulative probability, solution 2 is selected. ✓ (1)
- The second number random number is $r_2 = 0.68799$. Solution 8 is selected. + (2)

Fitness Proportionate Selection Operator

Index	$f(x_1, x_2)$	p_i	P_i
1	210.445	0.0041	0.0041
2	7536.407	0.1471	0.1512
3	14041.882	0.2740	0.4252
4	1546.603	0.0302	0.4554
5	189.732	0.0037	0.4591
6	671.924	0.0131	0.4722
7	7252.209	0.1415	0.6137
8	19793.183	0.3863	1.0

- The rest of the random numbers are

$\triangleright r_3 = 0.49976, r_4 = 0.31172, r_5 = 0.51961, r_6 = 0.48610, r_7 = 0.87648, r_8 = 0.99177.$ + N

- Selected solutions are '7', '3', '7', '7', '8', '8'.

Fitness Proportionate Selection Operator

Index	$f(x_1, x_2)$	p_i	P_i
1	210.445	0.0041	0.0041
2	7536.407	0.1471	0.1512
3	14041.882	0.2740	0.4252
4	1546.603	0.0302	0.4554
5	189.732	0.0037	0.4591
6	671.924	0.0131	0.4722
7	7252.209	0.1415	0.6137
8	19793.183	0.3863	1.0

- The rest of the random numbers are
 - $r_3 = 0.49976, r_4 = 0.31172,$
 $r_5 = 0.51961, r_6 = 0.48610,$
 $r_7 = 0.87648, r_8 = 0.99177.$
- Selected solutions are '7', '3', '7', '7', '8', '8'.

- Solution 2 gets 1 copy, solution 3 gets 1 copy, solution 7 gets 3 copies and solution 8 gets 3 copies.

Index	$f(x_1, x_2)$
2	7536.407
3	14041.882
7	7252.209
7	7252.209
7	7252.209
8	19793.183
8	19793.183
8	19793.183

- Two solutions become 'super-solution' in the population.

Stochastic Remainder Roulette-Wheel Selection Operator

- We calculate probability of each solution using the formula used with fitness proportionate selection operator.
- We then multiply the population size (N) with the probability of each solution.

Index	$f(x_1, x_2)$	p_i	Np_i
1	210.445	0.0041	0.0329
2	7536.407	0.1471	1.1766
3	14041.882	0.2740	2.1922
4	1546.603	0.0302	0.2415
5	189.732	0.0037	0.0296
6	671.924	0.0131	0.1049
7	7252.209	0.1415	1.1322
8	19793.183	0.3863	3.0901

- We assign a number of copies to the solution based on the integer value of product Np_i .
- For example, $Np_1 = 0.0329$ for solution '1', we do not assign any copy to it.
- For solution '2', $Np_2 = 1.17$. Therefore, we assign 1 copy of it.
- Similarly, solution '3' gets 2 copies, solution '7' gets 1 copy and solution '8' gets 3 copies.
- The total number of solutions selected is 7.
- We have to choose 1 more solution to keep the population size fixed, that is, $N = 8$.

Stochastic Remainder Roulette-Wheel Selection Operator

- We then use the fitness proportionate selection operator using only the decimal values.

Index	Np_i	Decimals	P_i
1	0.0329	0.0329	0.00329
2	1.1766	0.1766	0.2094
3	2.1922	0.1922	0.4017
4	0.2415	0.2415	0.6431
5	0.0296	0.0296	0.6728
6	0.1049	0.1049	0.7777
7	1.1322	0.1322	0.9099
8	3.0901	0.0901	1.0000
+	sum	1	

- Let the random number is $r_i = 0.76335$. We now select solution 6.
- We can see that the first part of stochastic remainder roulette-wheel selection operator is deterministic because the number of copies is decided using the integer value.
- It is considered less noisy than fitness proportionate selection operator in the sense of introducing less variance in its realization.

Stochastic Universal Sampling (SUS)

Index	$f(x_1, x_2)$	p_i	P_i
1	210.445	0.0041	0.0041
2	7536.407	0.1471	0.1512
3	14041.882	0.2740	0.4252
4	1546.603	0.0302	0.4554
5	189.732	0.0037	0.4591
6	671.924	0.0131	0.4722
7	7252.209	0.1415	0.6137
8	19793.183	0.3863	1.0000

- Only one random number r is required for selecting solutions.
- Since N different solutions have to be chosen, a set of N equi-spaced number is created.

- $R = \{r, r + 1/N, r + 2/N, \dots, r + (N - 1)/N\} \bmod 1$.
- Let $r = 0.40416$ and $1/N = 0.125$. Here, the population size is $N = 8$.
- Solution '3' is selected first using r .
- The second solution is selected using $r + 1/N = 0.40416 + 0.125 = 0.52916$, that is, solution '7'.
- The series of numbers for selecting solutions is
0.65416, 0.77916, 0.90416, 0.02916, 0.15416,
0.27916.
- The following solutions are then selected:
'8' '8', '8', '2', '3' '3'

Fitness Proportionate Selection

Limitations

- Domination of “super solution” in early generations
 - ▶ Suppose, the fitness (maximization) of five solutions in the population is given as $f_1 = 10, f_2 = 5, f_3 = 70, f_4 = 7, f_5 = 8$.
 - ▶ Because of the fitness proportional selection, solution ‘3’ will get more copies and will become a “super-solution”.
- Slower convergence in later generations
 - ▶ Suppose, the fitness of five solutions in the population is given as $f_1 = 19, f_2 = 21, f_3 = 22, f_4 = 18, f_5 = 20$.
 - ▶ Every solution will get a copy. It means that selection operator has no role in selecting solutions and thus, it becomes dummy.

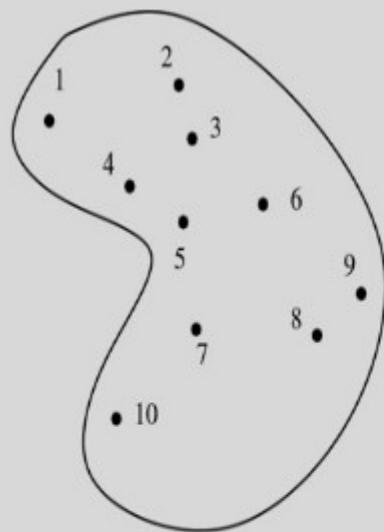
Possible remedies:

- Fitness scaling ✓
- Use ranking selection to avoid fitness scaling ✓
- Tournament selection: It does not have any scaling problem. ✓ +

Tournament Selection Operator

Tournament selection with tournament size ' k '

Randomly, we sample a subset \hat{P} of k solutions from the population P . Select solution in \hat{P} with the best fitness.



- Suppose $k = 4$, we choose four solutions randomly, say, $\{2, 6, 3, 9\}$ and their fitness values are $f_2 = 10, f_6 = 5.9, f_3 = 16.7, f_9 = 9.4$.
- For maximization problem, who will win the tournament?
Solution 3 ✓
- For minimization problem, who will win the tournament?
Solution 6 ✓
- Often, tournament size $k = 2$ is used, which is known as binary tournament selection operator.

Operators for Survivor or Elimination

$(\mu + \lambda)$ selection scheme

- Let the size of parent population is μ .
- Let λ number of offspring solutions are generated after variation operator.
- The next population is filled by choosing the μ best number of solutions from the combined populations of parent and offspring solutions.

(μ, λ) selection scheme where $(\lambda > \mu)$

- Let the size of parent population is μ .
- Let λ number of offspring solutions are generated after variation operator. Note that $(\lambda > \mu)$.
- The next population is filled by choosing the μ best number of solutions from the offspring population only.

Crossover Operators for Binary Variables

- Crossover operator creates offspring/new solutions from **more than one parent**.
- Crossover probability (p_c)
 - ▶ $100p_c\%$ strings are used in the crossover operator.
 - ▶ $100(1 - p_c)\%$ of the population are simply copied to the new population.
- Single-point crossover operator
 - ▶ Choose a random site on the two parents
 - ▶ Split parents at this crossover site
 - ▶ Create children by exchanging tails

parent 1:	1 0 1 0 1 1		1 0 1 0	
parent 2:	0 1 0 1 0 0		1 1 1 0	
offspring 1:	1 0 1 0 1 1		1 1 1 0	✓
offspring 2:	0 1 0 1 0 0		1 0 1 0	✓

Crossover Operators for Binary Variables

- n -point crossover:
 - ▶ Choose n random crossover sites
 - ▶ Split along those sites
 - ▶ Glue parts, alternating between parents
- 2-point crossover operator

parent 1:	1	0	1		0	1	1		1	0	1	0
parent 2:	0	1	0		1	0	0		1	1	1	0
offspring 1:	1	0	1		1	0	0		1	0	1	0
offspring 2:	0	1	0		0	1	1		1	1	1	0

$$\begin{array}{l} x_1 \quad 01101 \\ x_2 \quad 11100 \end{array} \quad \left. \vphantom{\begin{array}{l} x_1 \\ x_2 \end{array}} \right\} +$$

Crossover Operators for Binary Variables

- Uniform crossover

- ▶ Select a bit-string z of length n uniformly at random.
- ▶ for all i in 1 to n
 - ★ if $z_i = 1$ then bit i in offspring-1 is x_i else y_i .
 - ★ if $z_i = 1$ then bit i in offspring-2 is y_i else x_i .

parent 1:	1	0	1	0	1	1	1	0	1	0
parent 2:	0	1	0	1	0	0	1	1	1	0
$z =$	1	0	1	0	0	0	1	1	1	0
offspring 1:	1	1	1	1	0	0	1	0	1	0
offspring 2:	0	0	0	0	1	1	1	1	1	0

Mutation Operator for Binary String

The mutation operator introduces small, random changes to a solution's chromosome.

Local Mutation

- One randomly chosen bit is flipped.

• 1 0 1 0 1 1 1 0 1 0

• 1 0 1 **1** 1 1 1 0 1 0

Global Mutation

- Each bit flipped independently with a given probability p_m .
- It is also called the per bit mutation rate, which is often $= 1/n$, where n is the chromosome length.

• 1 0 1 0 1 1 1 0 1 0

• 1 0 **0** 0 1 1 1 **1** 1 0

• $\Pr[k \text{ bits flipped}] = \binom{n}{k} \cdot p_m^k \cdot (1 - p_m)^{n-k}$

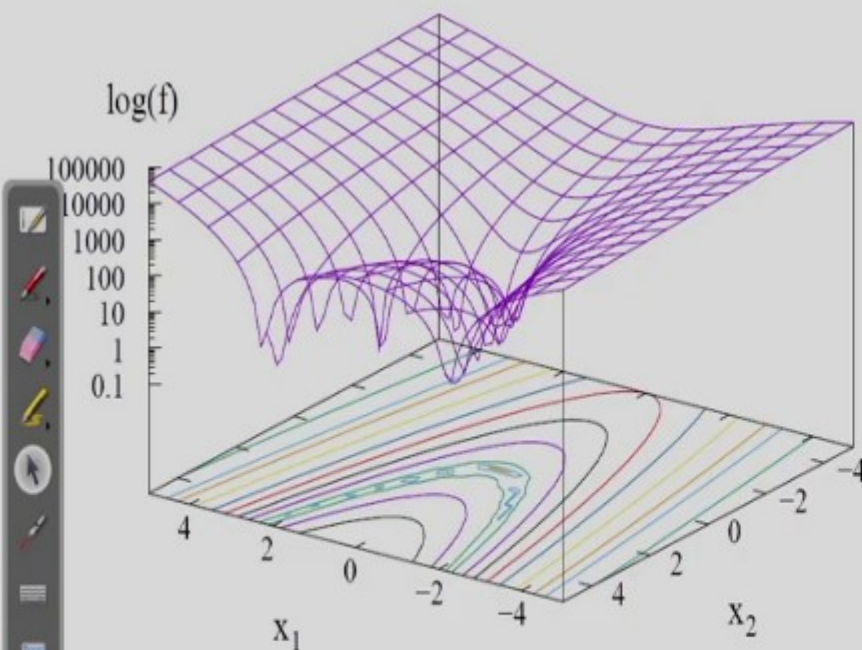
$$p_m = \frac{1}{n}$$

Simulation and Application

Rosenbrock Function

Rosenbrock Function

Minimize $f(x_1, \dots, x_N) = \sum_{i=1}^{N-1} (100(x_{i+1} - x_i^2)^2 + (1 - x_i)^2)$,
bounds $-5 \leq x_i \leq 5$ and $i = 1, \dots, N$.



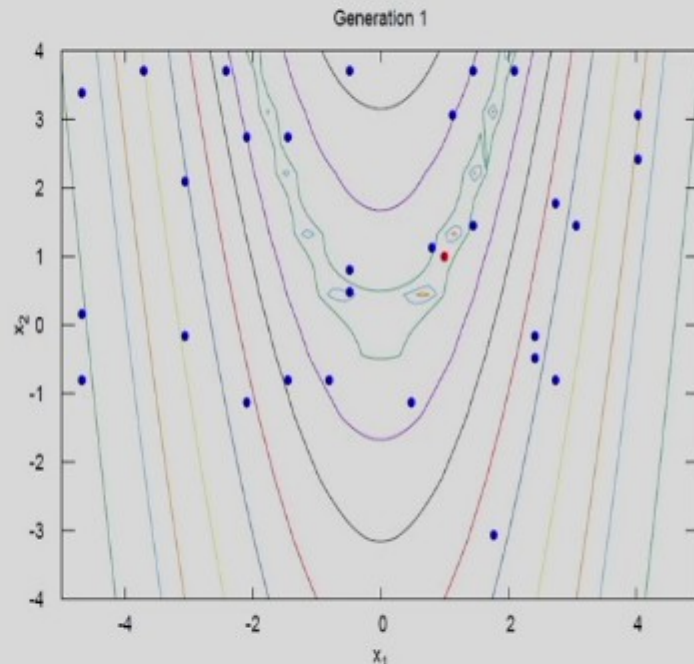
- Optimal solution is $x^* = (1, \dots, 1)^T$ and $f(x) = 0$

Rosenbrock Function

BGA Parameters

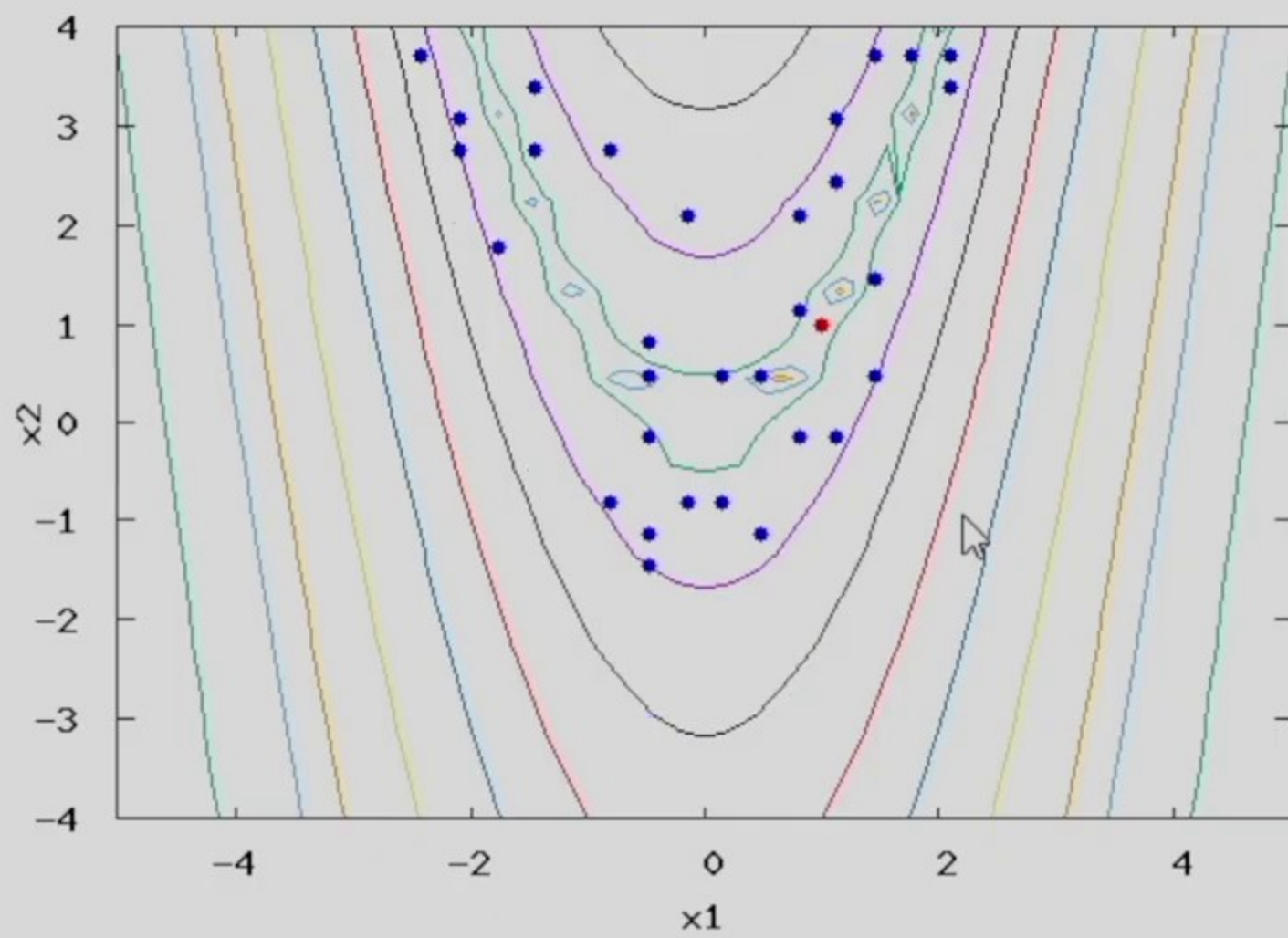
- Number of variables: $N = 2$
- Population size: 40
- No. of generations: 200
- Binary string length of x_1 is $l_1 = 5$
- Binary string length of x_2 is $l_2 = 5$
- Total length of binary string is $l = l_1 + l_2 = 10$
- Probability of crossover: $p_c = 1.0$
- Probability of mutation; $p_m = 1/N = 0.5$
- Binary tournament selection operator
- Single-point crossover operator
- Bit-wise mutation operator

- $(\mu + \lambda)$ -strategy

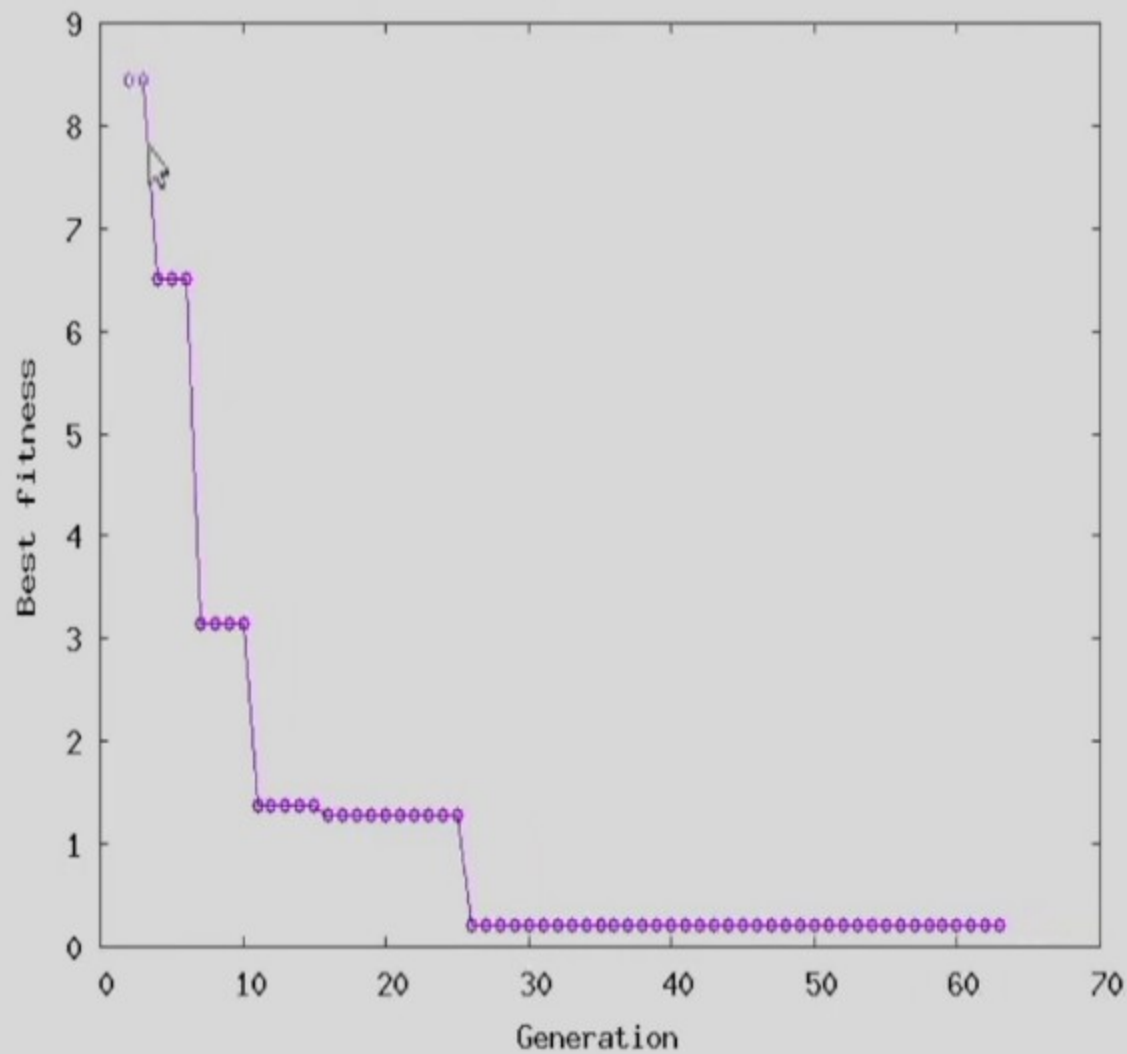


- Simulation [▶ Link](#)
- Progress [▶ Link](#)

Generation 5



Generation 67

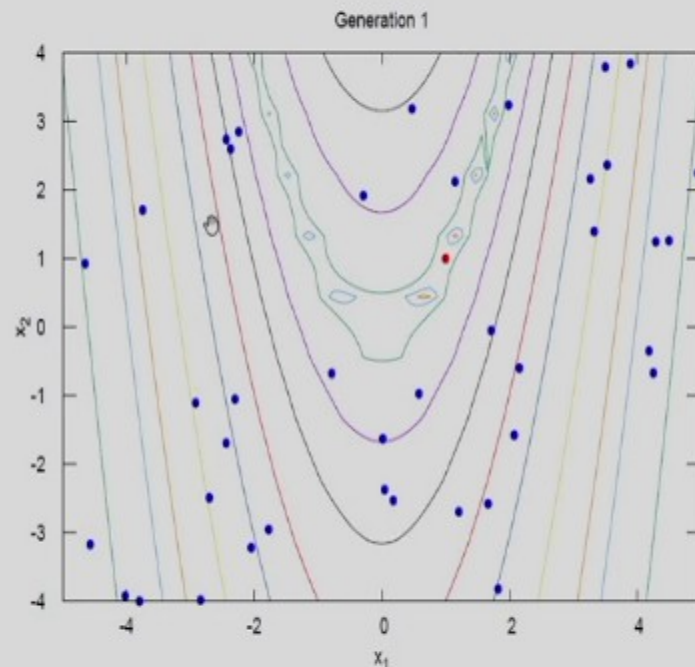


Rosenbrock Function

BGA Parameters

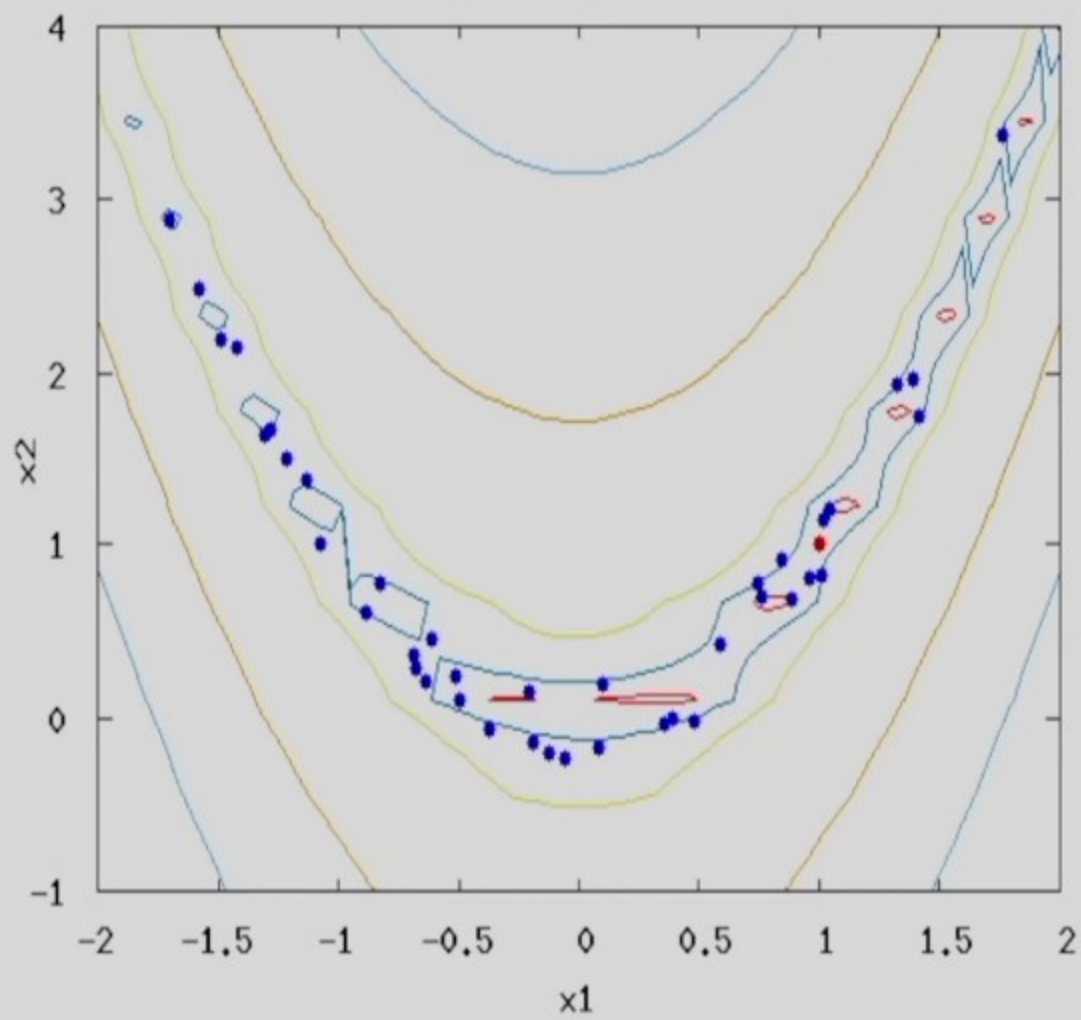
- Number of variables: $N = 2$
- Population size: 40
- No. of generations: 200
- Binary string length of x_1 is $l_1 = 10$
- Binary string length of x_2 is $l_2 = 10$
- Total length of binary string is $l = l_1 + l_2 = 20$
- Probability of crossover: $p_c = 1.0$
- Probability of mutation; $p_m = 1/N = 0.5$
- Binary tournament selection operator
- Single-point crossover operator
- Bit-wise mutation operator

- $(\mu + \lambda)$ -strategy

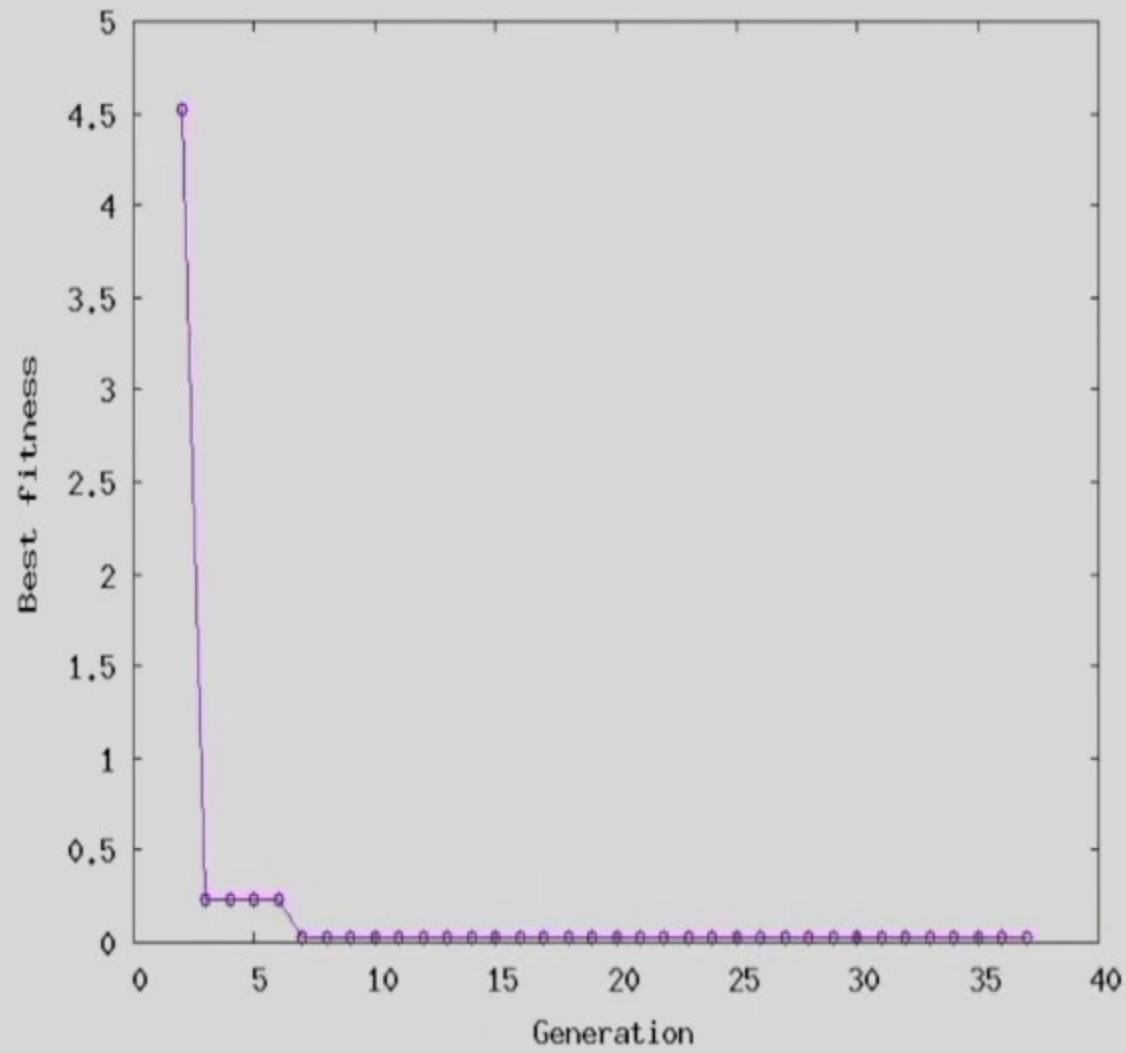


- Simulation [▶ Link](#)
- Progress [▶ Link](#)

Generation 40



Generation 36

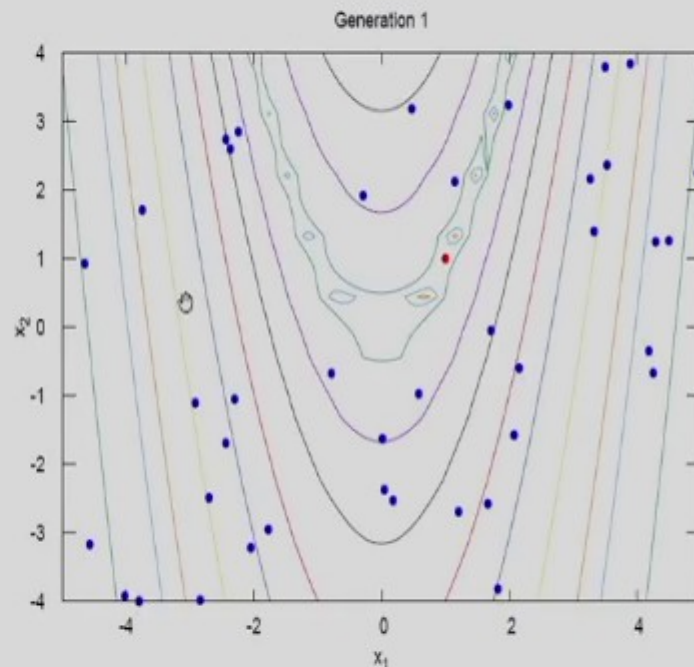


Rosenbrock Function

BGA Parameters

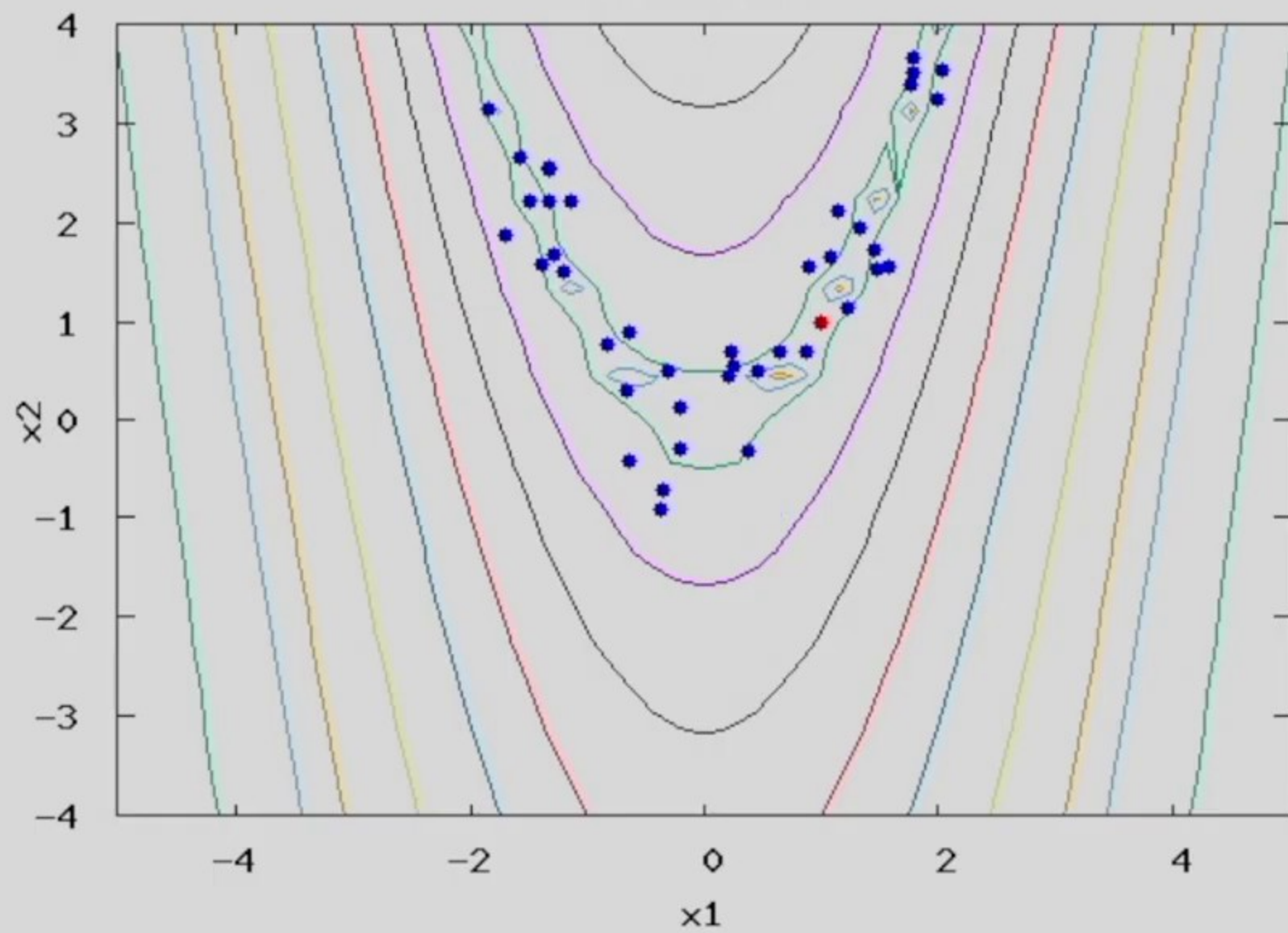
- Number of variables: $N = 2$
- Population size: 40
- No. of generations: 200
- Binary string length of x_1 is $l_1 = 20$
- Binary string length of x_2 is $l_2 = 20$
- Total length of binary string is $l = l_1 + l_2 = 40$
- Probability of crossover: $p_c = 1.0$
- Probability of mutation; $p_m = 1/N = 0.5$
- Binary tournament selection operator
- Single-point crossover operator
- Bit-wise mutation operator

- $(\mu + \lambda)$ -strategy

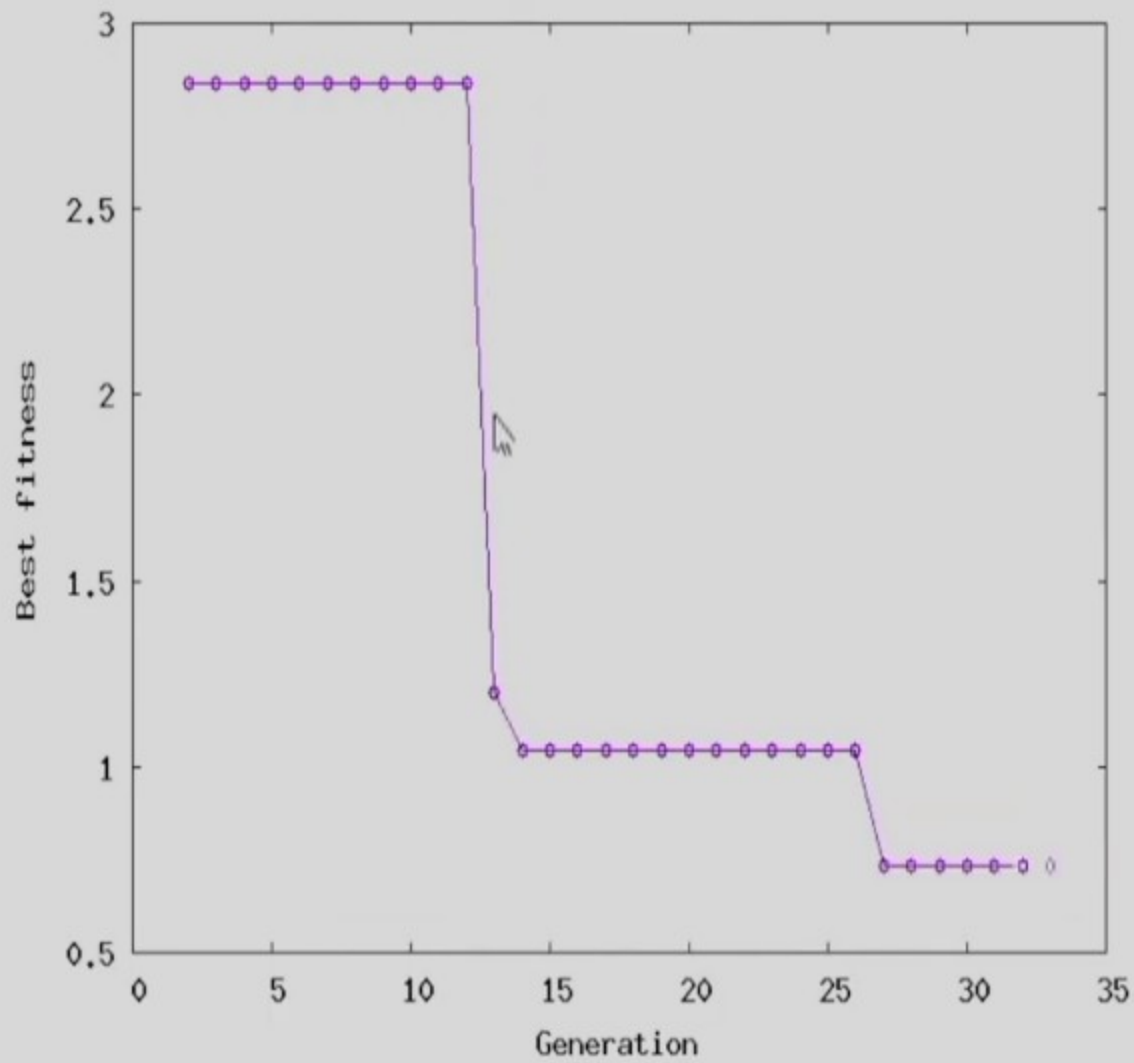


- Simulation [▶ Link](#)
- Progress [▶ Link](#)

Generation 8



Generation 37

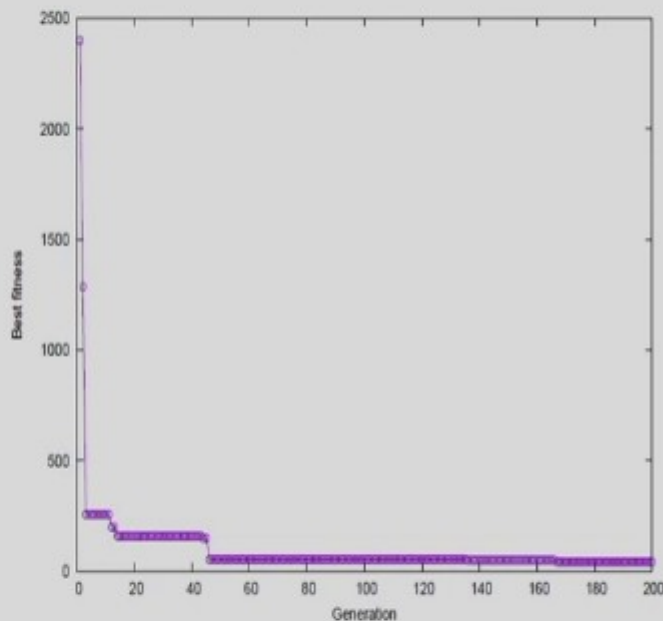


Rosenbrock Function

BGA Parameters

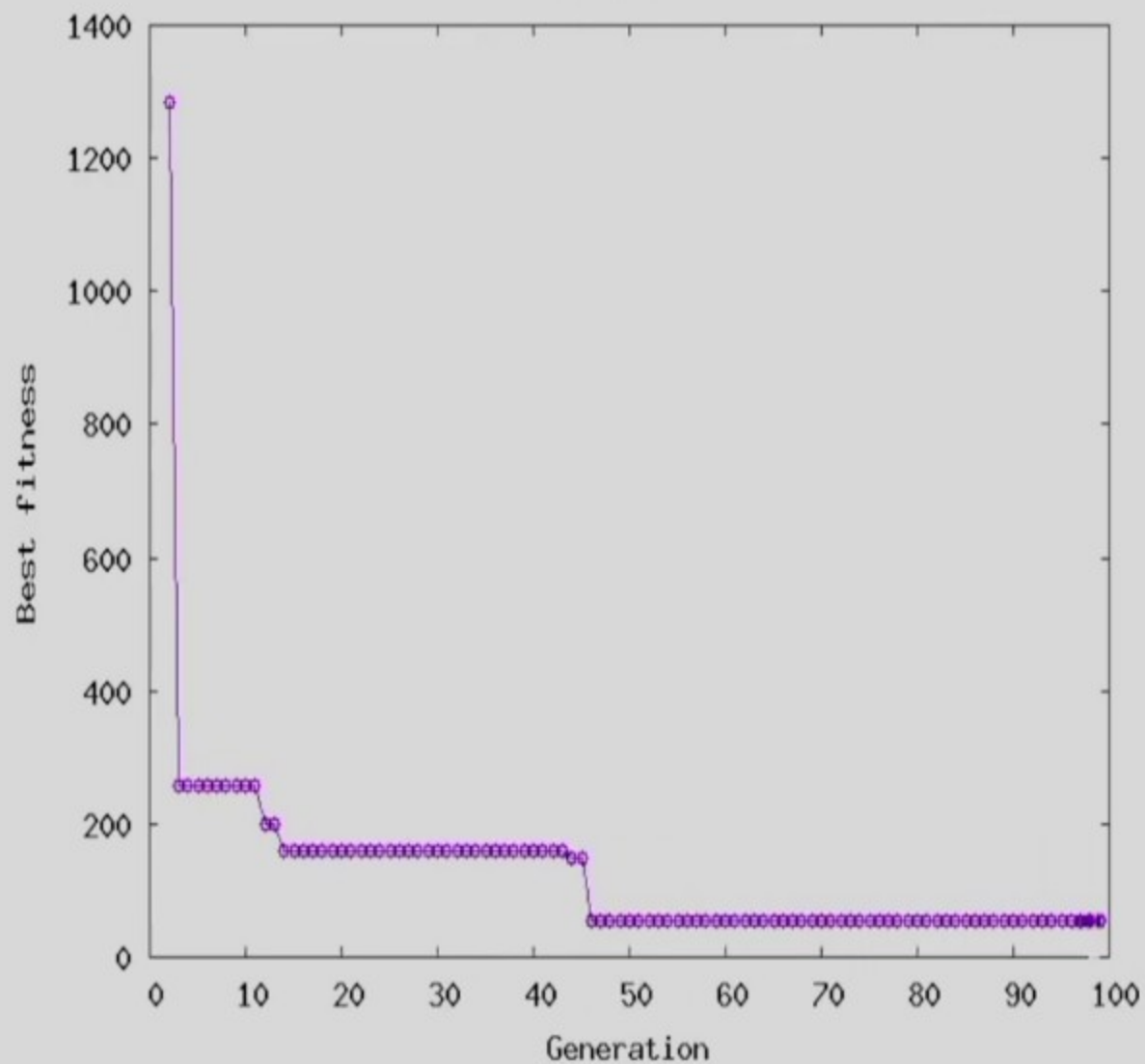
- Number of variables: $N = 4$
- Population size: 40
- No. of generations: 200
- Binary string length of x_1 is $l_1 = 20$
- Binary string length of x_2 is $l_2 = 20$
- Total length of binary string is $l = l_1 + l_2 = 40$
- Probability of crossover: $p_c = 1.0$
- Probability of mutation; $p_m = 1/N = 0.5$
- Binary tournament selection operator
- Single-point crossover operator
- Bit-wise mutation operator

- $(\mu + \lambda)$ -strategy



- Progress [▶ Link](#)

Generation 98

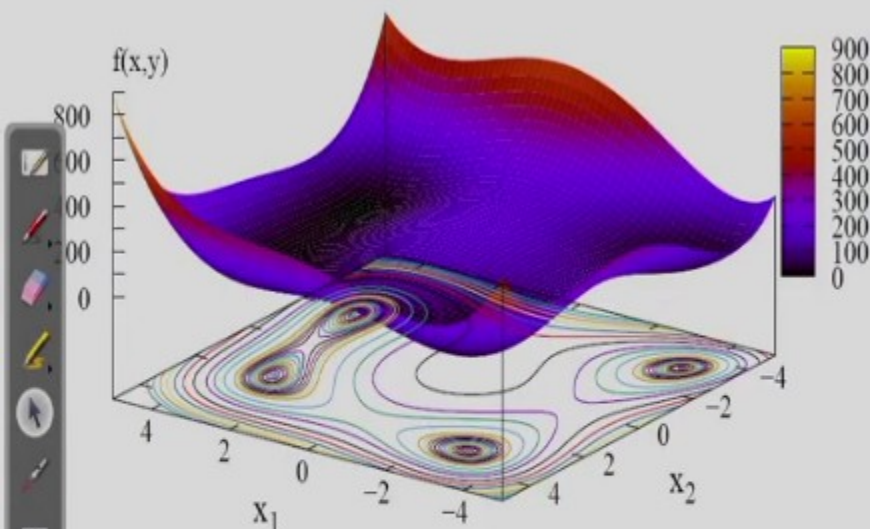


Himmelblau Function

Himmelblau Function

Minimize $f(x_1, x_2) = (x_1^2 + x_2 - 11)^2 + (x_1 + x_2^2 - 7)^2$,

bounds $-5 \leq x_1 \leq 5$ and $-5 \leq x_2 \leq 5$.

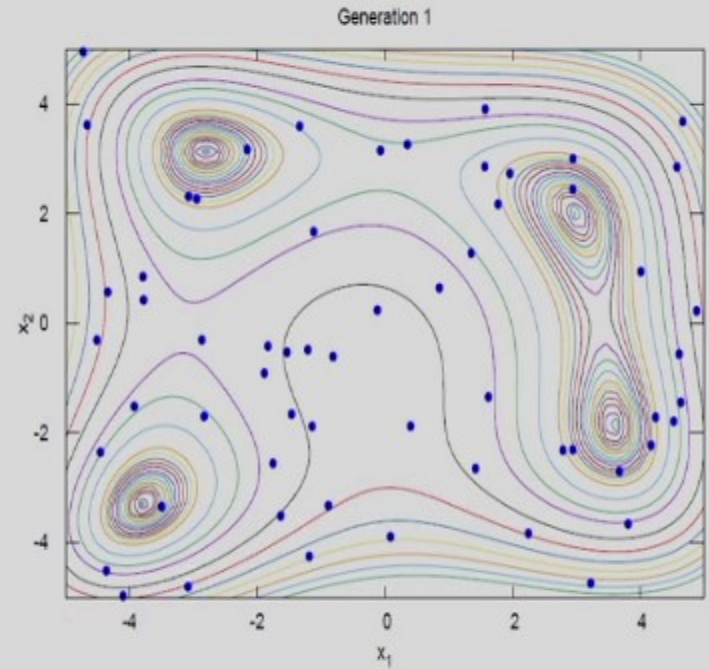


- Multi-modal function: it has 4 minimum points
- First optimal solution is $x^* = (3, 2)^T$ and $f(x) = 0$
- Second optimal solution is $x^* = (-2.805, 3.131)^T$ and $f(x) = 0$
- Third optimal solution is $x^* = (-3.779, -3.283)^T$ and $f(x) = 0$
- Fourth optimal solution is $x^* = (3.584, -1.848)^T$ and $f(x) = 0$

Himmelblau Function

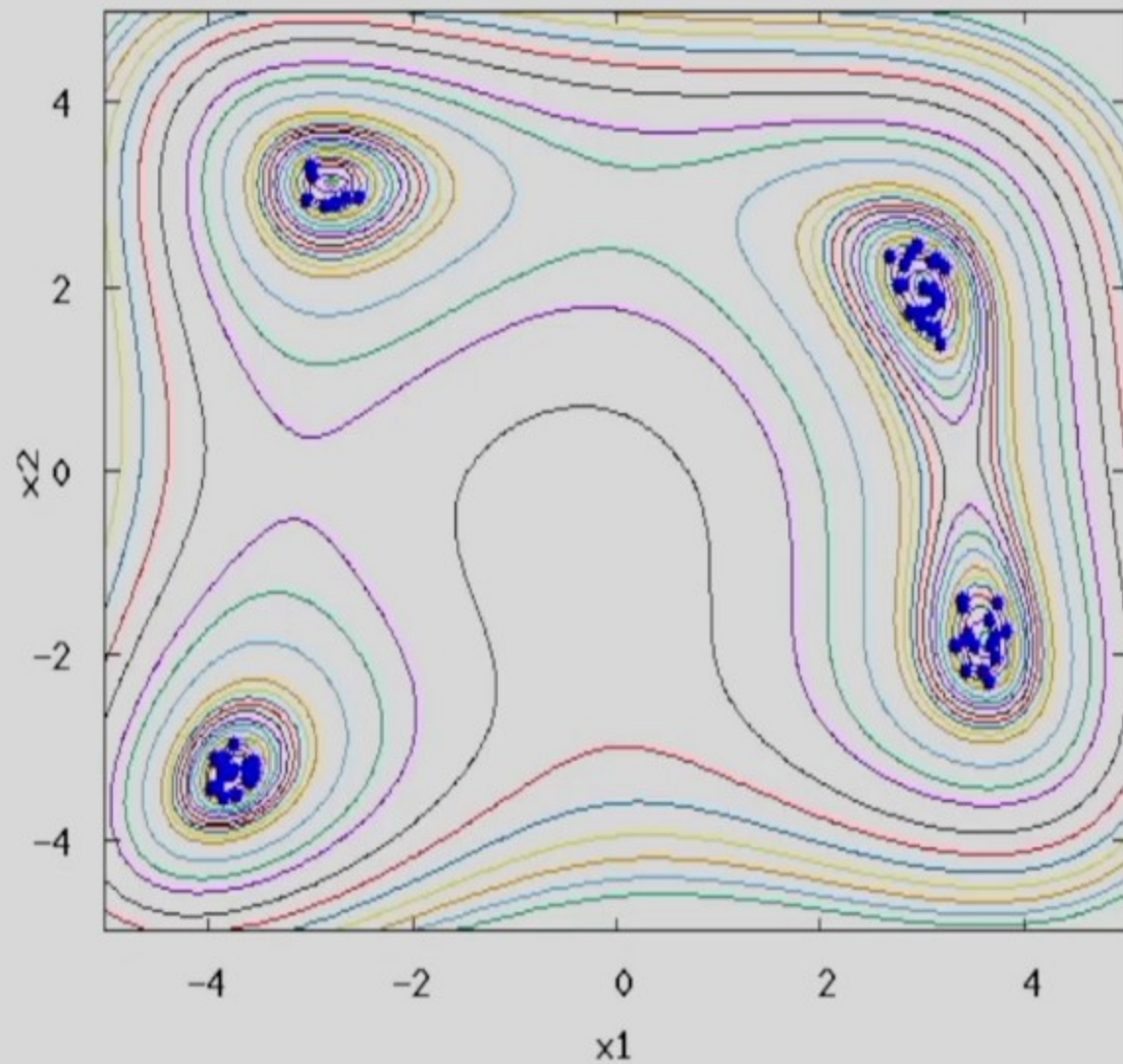
BGA Parameters

- Population size: 60
- No. of generations: 200
- Binary string length of x_1 is $l_1 = 20$
- Binary string length of x_2 is $l_2 = 20$
- Total length of binary string is $l = l_1 + l_2 = 40$
- Probability of crossover: $p_c = 1.0$
- Probability of mutation; $p_m = 1/N = 0.5$
- Binary tournament selection operator
- Single-point crossover operator
- Bit-wise mutation operator
- $(\mu + \lambda)$ -strategy

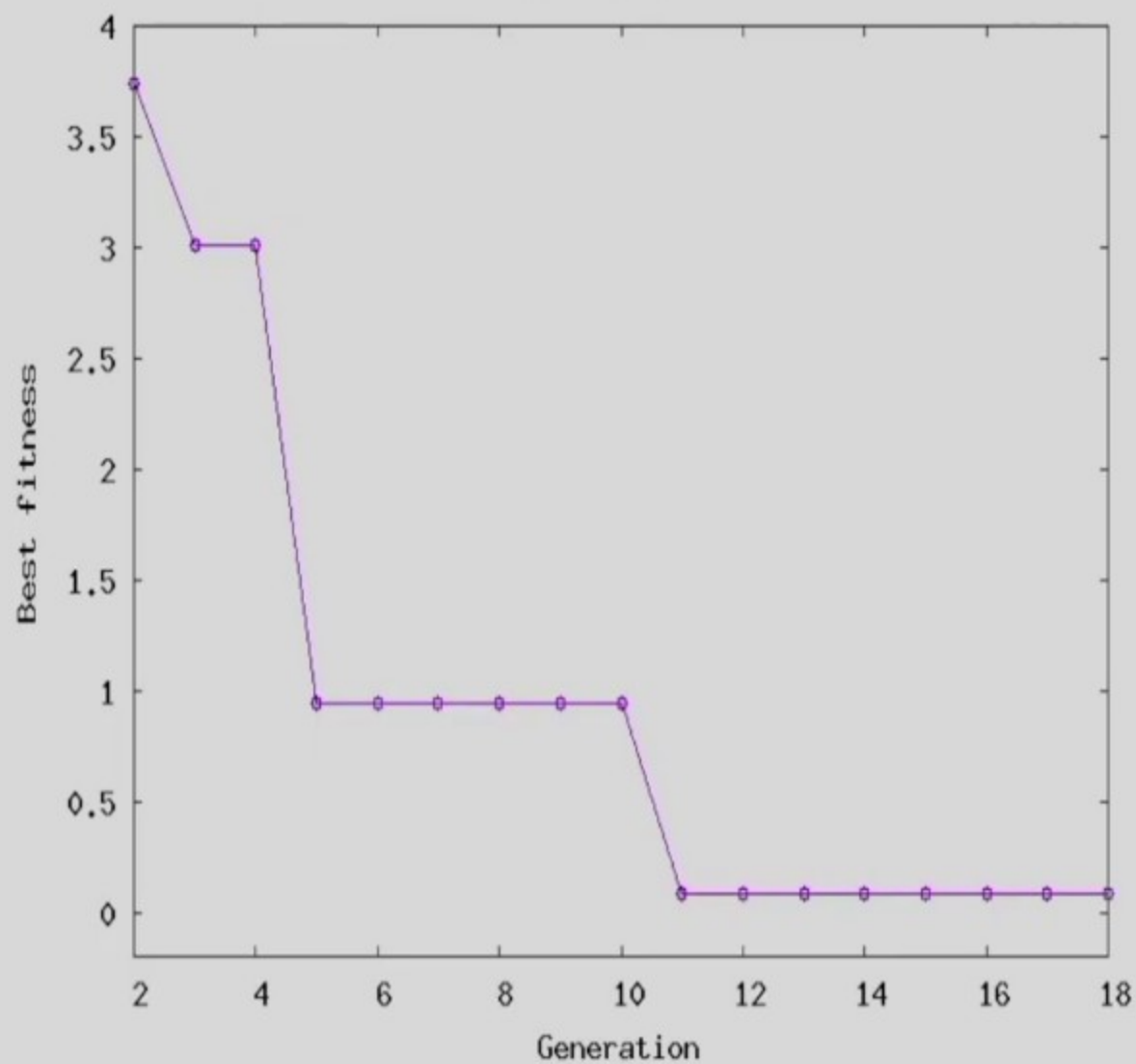


- Simulation [▶ Link](#)
- Progress [▶ Link](#)

Generation 55



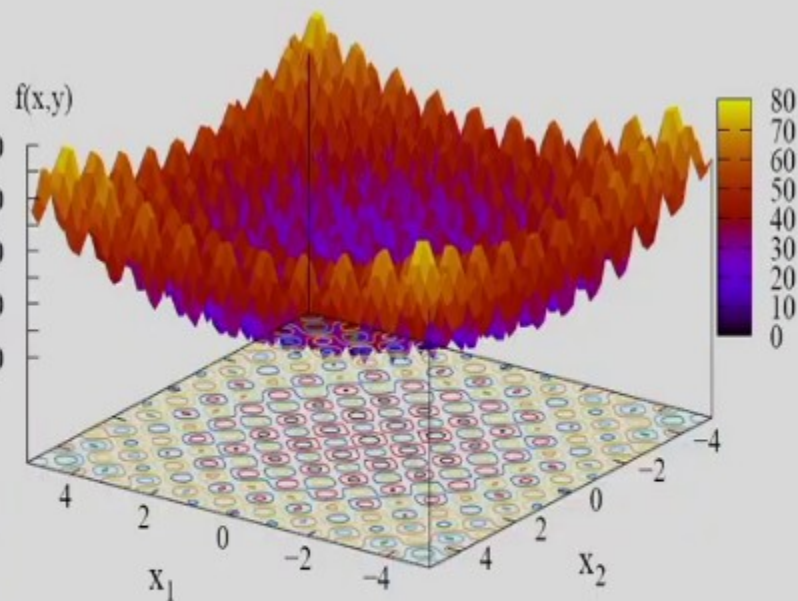
Generation 17



Rastrigin Function

Rastrigin Function

Minimize $f(x_1, \dots, x_N) = 10N + \sum_{i=1}^N (x_i^2 - 10 \cos(2 * \pi x_i))$,
bounds $-5.12 \leq x_i \leq 5.12$



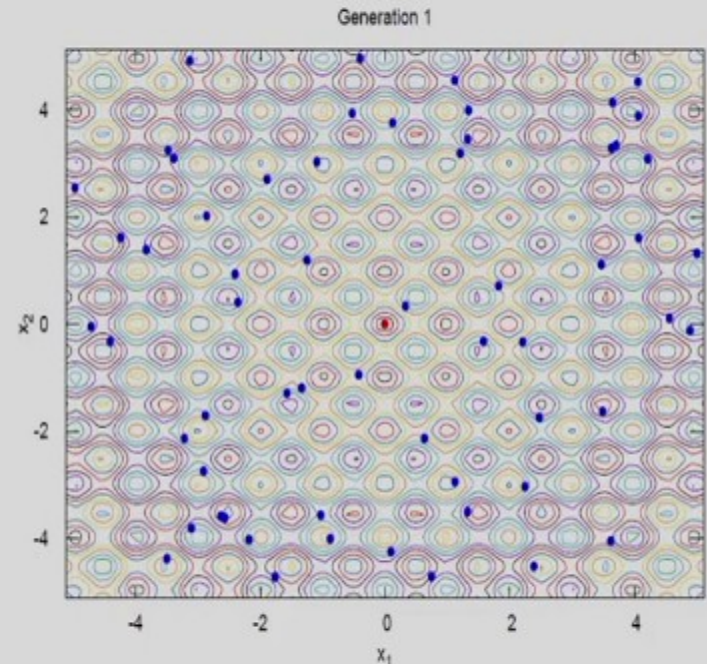
- Optimal solution is $x^* = (0, \dots, 0)^T$ and $f(x) = 0$

Rastrigin Function

BGA Parameters

- Number of variable: $N = 2$
- Population size: 60
- No. of generations: 200
- Binary string length of x_1 is $l_1 = 20$
- Binary string length of x_2 is $l_2 = 20$
- Total length of binary string is $l = l_1 + l_2 = 40$
- Probability of crossover: $p_c = 1.0$
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- Bit-wise mutation operator

- $(\mu + \lambda)$ -strategy

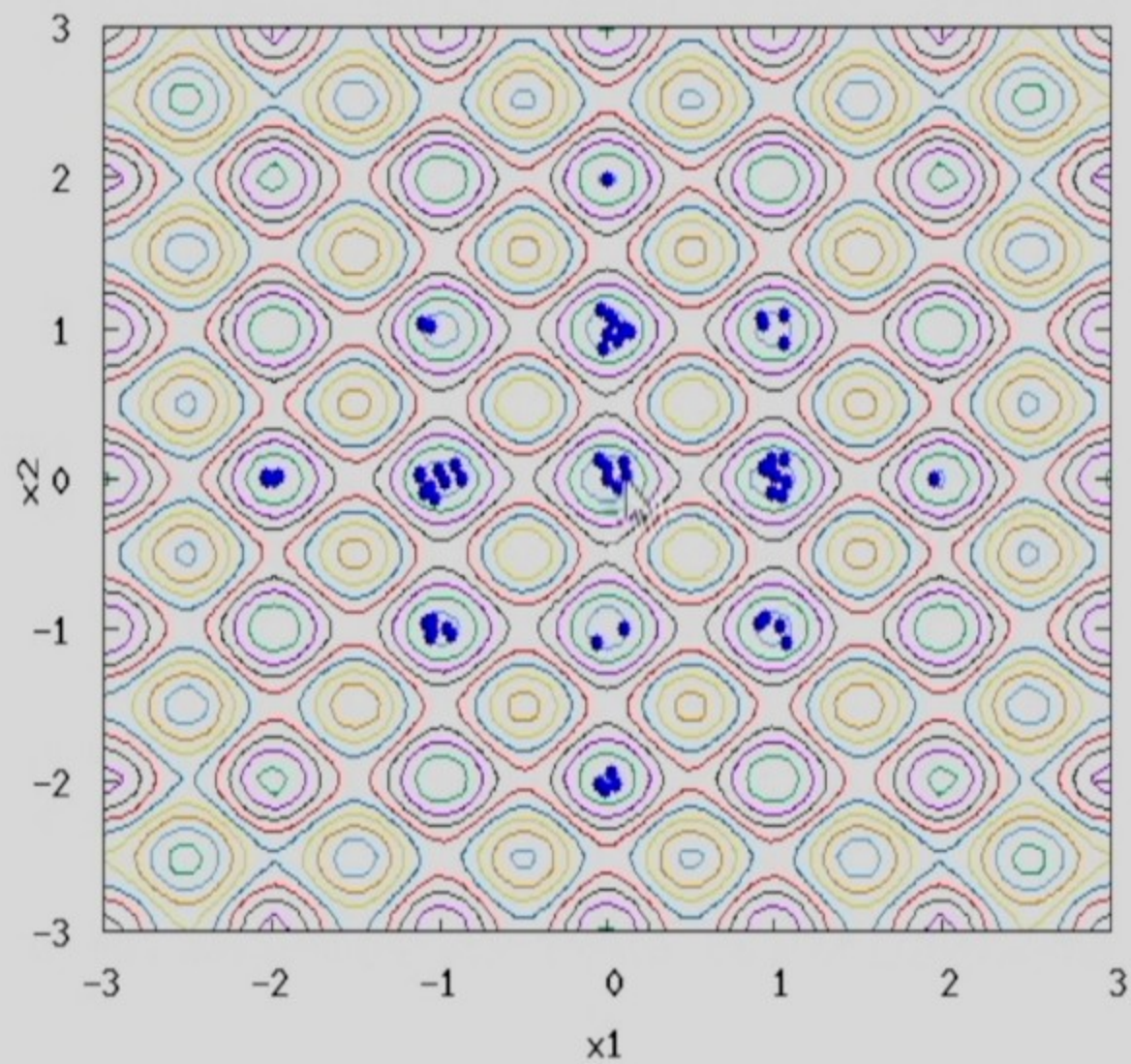


- Simulation [▶ Link](#)

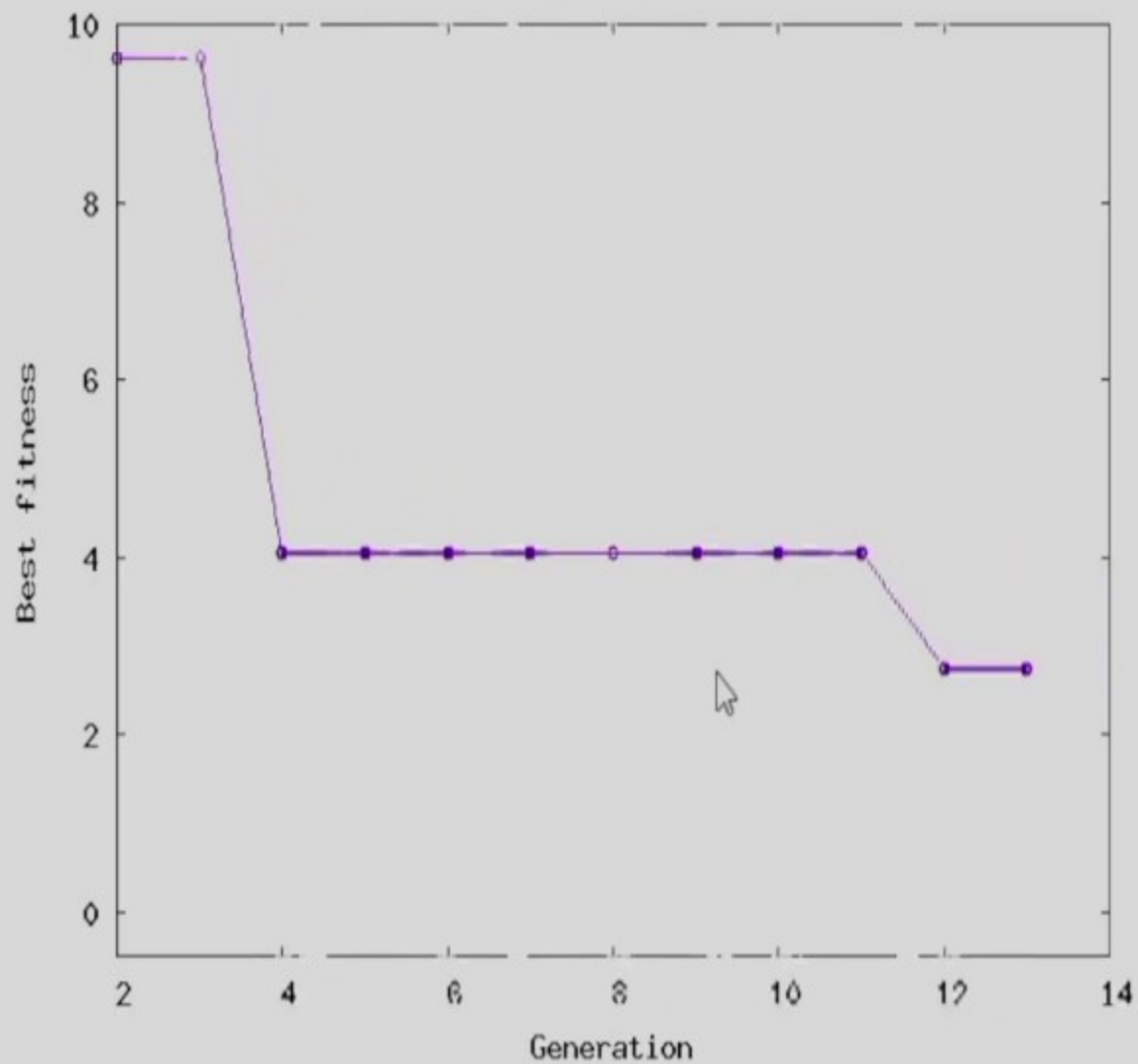
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Progress

Generation 195



Generation 17

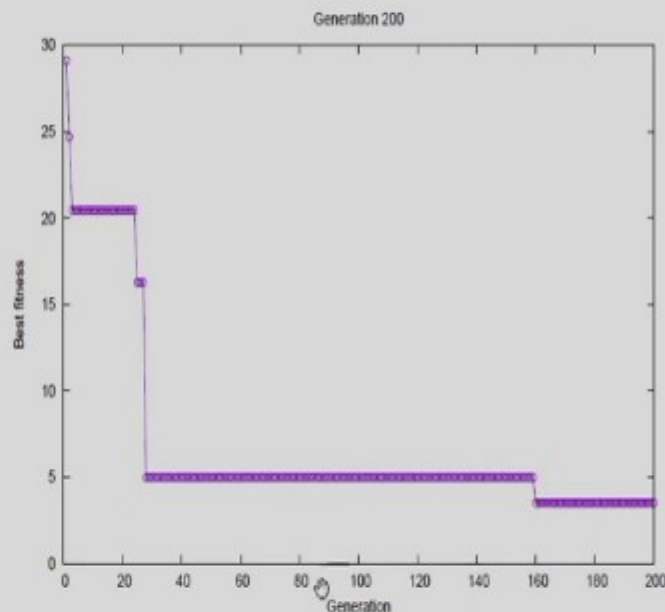


Rastrigin Function

BGA Parameters

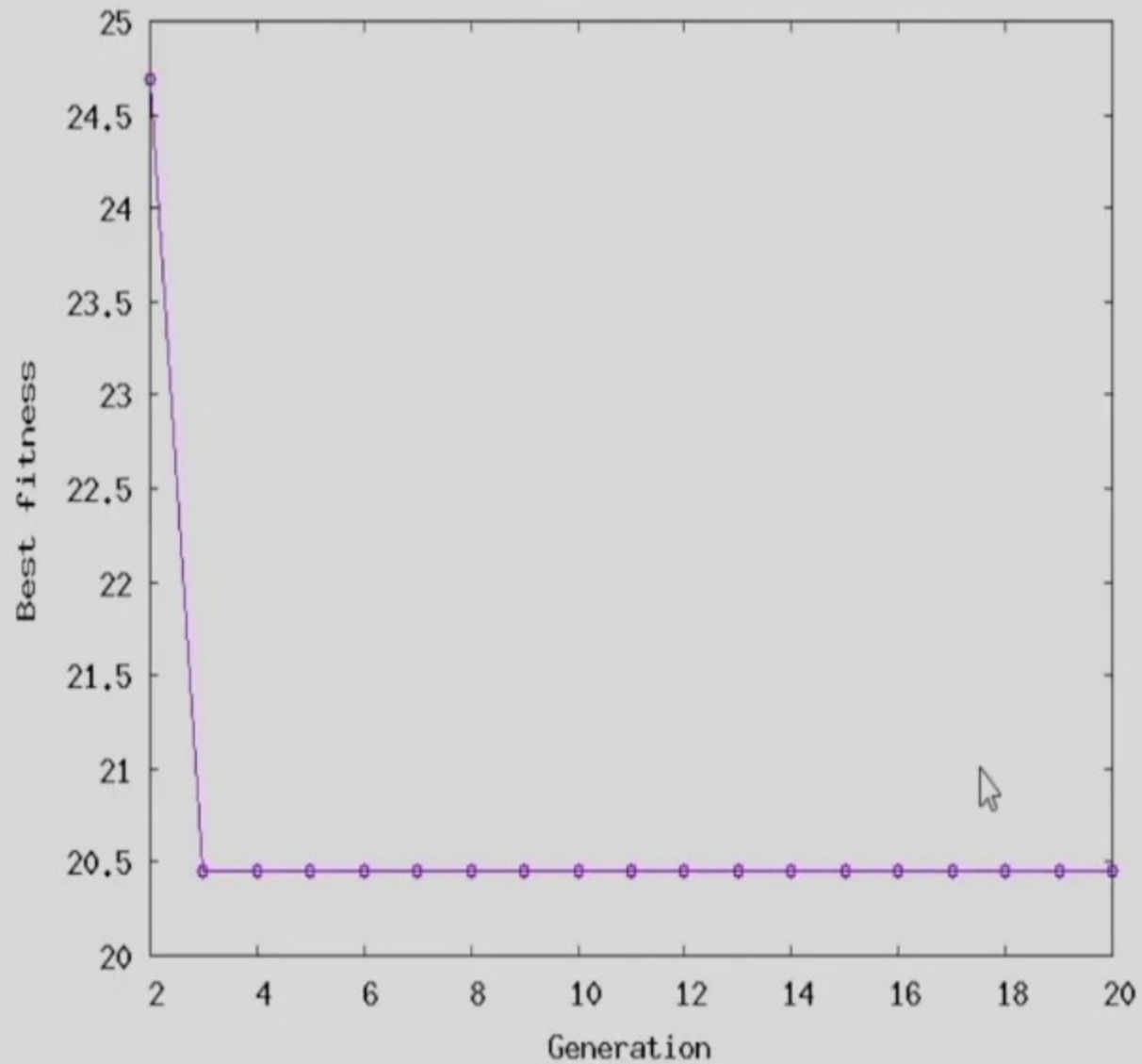
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- Progress [▶ Link](#)

Generation 19

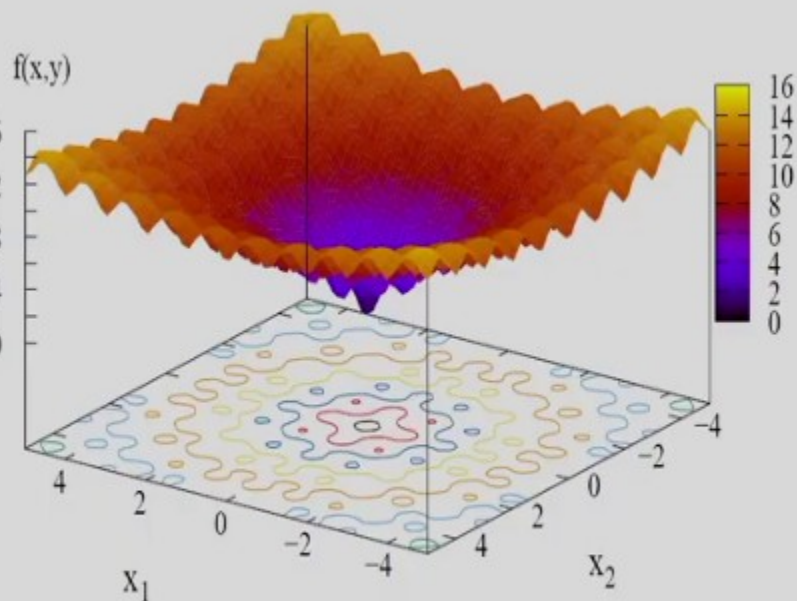


Ackley Function

Ackley Function

Minimize $f(x_1, x_2) = -20 \exp \left(-0.2 \sqrt{0.5(x_1^2 + x_2^2)} \right) - \exp(0.5(\cos(2\pi x_1) + \cos(2\pi x_2))) + \exp(1) + 20,$
bounds $-5 \leq x_1 \leq 5$ and $-5 \leq x_2 \leq 5$.

0

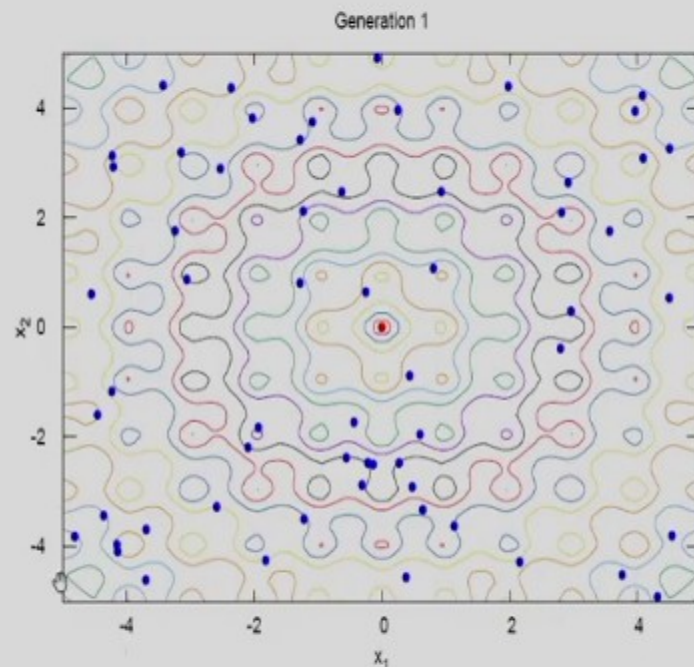


- Optimal solution is $x^* = (0,0)^T$ and $f(x) = 0$

Ackley Function

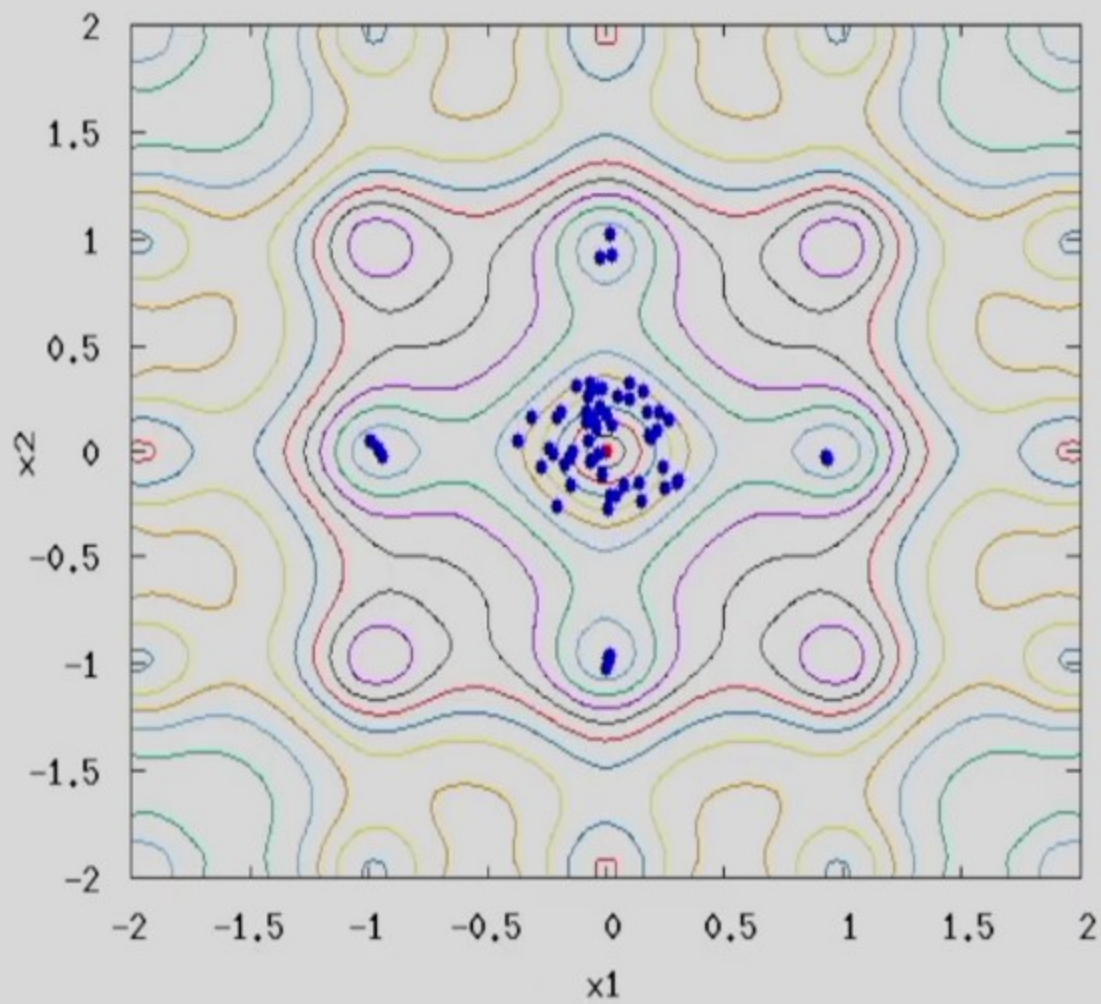
BGA Parameters

- Population size: 60
- No. of generations: 200
- Binary string length of x_1 is $l_1 = 20$
- Binary string length of x_2 is $l_2 = 20$
- Total length of binary string is $l = l_1 + l_2 = 40$
- Probability of crossover: $p_c = 1.0$
- Probability of mutation; $p_m = 1/N = 0.5$
- Binary tournament selection operator
- Single-point crossover operator
- Bit-wise mutation operator
- $(\mu + \lambda)$ -strategy

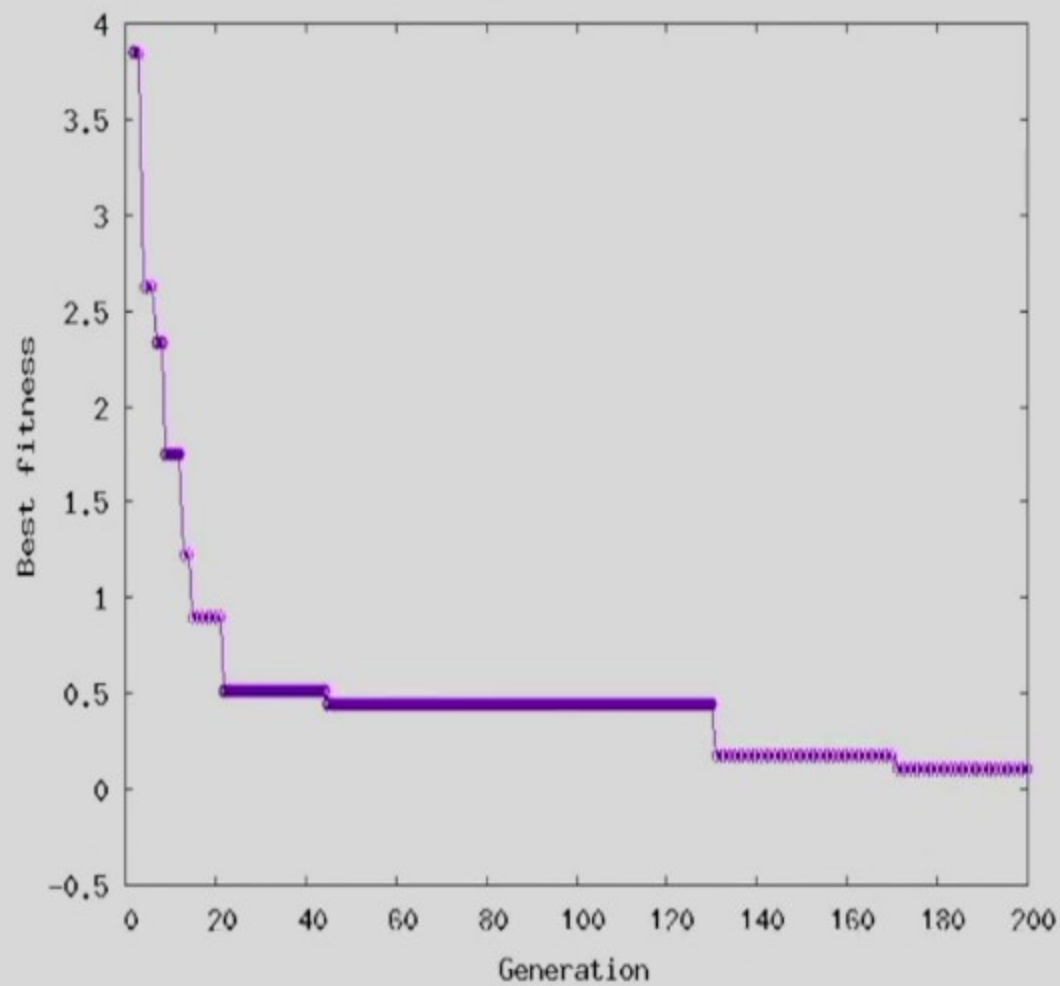


- Simulation [▶ Link](#)
- Progress [▶ Link](#)

Generation 170



Generation 200



- Binary-coded genetic algorithm
 - ▶ Solution representation
 - ▶ Working principles through an example: selection operator, crossover and mutation operators, survivor operator
 - ▶ Graphical example
- More operators
 - ▶ selection operator, crossover and mutation operators, survivor operator
- Simulation and application of BGA on four mathematical problems

