Project (Part 1)

1 Team Members

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2 The changes

Add a new feature: FOR loop

- Syntax: Lines: (9) and (17)
- Static Semantics: Line (9)
- Dynamic Semantics: densem(<for>)

3 The Language

3.1 Syntax

(1) `<prog>` → `<stmt_list>`
(2) `<stmt_list>` → ϵ
(3) | `<stmt>` ";" `<stmt_list>`
(4) `<stmt>` → `<print>`
(5) | `<input>`
(6) | `<assign>`
(7) | `<if>`
(8) | `<while>`
(9) | `<for>`
(10) `<print>` → "print" `<p-arg>`
(11) `<p-arg>` → **STRING**
(12) | `<expr>`
(13) `<input>` → "get" **ID**
(14) `<assign>` → **ID** "=" `<expr>`
(15) `<if>` → "if" `<expr>` "then" `<stmt_list>` "else" `<stmt_list>` "end"
(16) `<while>` → "while" `<expr>` "do" `<stmt_list>` "end"
(17) `<for>` → "for" `<expr>` `<assign>` "," `<expr>` "," `<expr>` `<stmt_list>`
(18) `<expr>` → `<n_expr>` `<b_expr>`
(19) `<b_expr>` → ϵ

(20) | "and" `<n_expr>`
(21) | "or" `<n_expr>`
(22) `<n_expr>` → `<term> <t_expr>`
(23) `<t_expr>` → ϵ
(24) | "+" `<n_expr>`
(25) | "-" `<n_expr>`
(26) `<term>` → `<factor> <f_expr>`
(27) `<f_expr>` → ϵ
(28) | "*" `<term>`
(29) | "/" `<term>`
(30) | "%" `<term>`
(31) `<factor>` → `<value> <v_expr>`
(32) `<v_expr>` → ϵ
(33) | ">" `<value>`
(34) | ">=" `<value>`
(35) | "<" `<value>`
(36) | "<=" `<value>`
(37) | "==" `<value>`
(38) | "!=" `<value>`
(39) `<value>` → "(" `<expr>` ")"
(40) | "not" `<value>`
(41) | "-" `<value>`
(42) | **ID**
(43) | **INT**

3.1.1 Tokens

This subsection describes the token used in the above grammar. Provided for each token is a regex and a description. The regex is for those that know regular expressions and prefer it as a description. The description says the same thing in English. Preprocessing describes how the lexeme is transformed before passing it to the parser.

STRING

As a regex: `"([^"]|\")*"`

Description: A quotation mark followed by zero or more characters, where quotation marks must be preceded by a backslash, followed by another quotation mark.

Preprocessing: The first and last quotation marks are removed. Scanning from left to right, "\\" is replaced with "\", "\t" is replaced with a tab, "\n" is replaced with a newline, "\"" is replaced with """, and any "\" that is followed by anything else is removed.

ID

As regex: [_a-zA-Z][_a-zA-Z0-9]*

Description: A letter or underscore followed by a combination of zero or more letters, underscores or digits.

INT

As Regex: (+|-)?[0-9]+

Description: an optional "+" or "-" followed by one or more digits.

3.2 Static Semantics

1 `<stmt_list>` $.ids = \{\}$

3 `<stmt_list>` $[1].ids = \{$ `<stmt>` $.id\} \cup$ `<stmt_list>` $[0].ids$
`<stmt>` $.ids =$ `<stmt_list>` $[0].ids$

4 `<print>` $.ids =$ `<stmt>` $.ids$

5 `<stmt>` $.id =$ `<input>` $.id$

6 `<stmt>` $.id =$ `<assign>` $.id$
`<assign>` $.ids =$ `<stmt>` $.ids$

7 `<if>` $.ids =$ `<stmt>` $.ids$

8 `<while>` $.ids =$ `<stmt>` $.ids$

9 `<for>` $.ids =$ `<stmt>` $.ids$

10 `<p-arg>` $.ids =$ `<print>` $.ids$

11 `<expr>` $.ids =$ `<p-arg>` $.ids$

12 `<input>` $.id =$ **ID** $.id$

13 `<assign>` $.id =$ **ID** $.id$

16 `<n_expr>` $.ids =$ `<expr>`
`<b_expr>` $.ids =$ `<expr>`

18 `<n_expr>` $.ids =$ `<b_expr>`

19 `<n_expr>` $.ids =$ `<b_expr>`

20 `<term>` $.ids =$ `<b_expr>`
`<t_expr>` $.ids =$ `<b_expr>`

22 `<n_expr>` $.ids =$ `<t_expr>` $.ids$

23 `<n_expr>` $.ids =$ `<t_expr>` $.ids$

24 <factor> $.ids =$ <term> $.ids$
<f_expr> $.ids =$ <term> $.ids$

26 <term> $.ids =$ <f_expr> $.ids$

27 <term> $.ids =$ <f_expr> $.ids$

28 <term> $.ids =$ <f_expr> $.ids$

29 <value> $.ids =$ <factor> $.ids$
<v_expr> $.ids =$ <factor> $.ids$

31 <value> $.ids =$ <v_expr> $.ids$

32 <value> $.ids =$ <v_expr> $.ids$

33 <value> $.ids =$ <v_expr> $.ids$

34 <value> $.ids =$ <v_expr> $.ids$

35 <value> $.ids =$ <v_expr> $.ids$

36 <value> $.ids =$ <v_expr> $.ids$

37 <expr> $.ids =$ <value> $.ids$

38 <value> $[1].ids =$ <value> $[0].ids$

39 <value> $[1].ids =$ <value> $[0].ids$

40 Predicate: **ID** $.id \in$ <value> $.ids$

3.3 Dynamic Semantics

This section gives the dynamic semantics of the language using denotational semantics. Consider the *demsem* function the denotational semantics for this language. We will use a mapping from variable name to value to represent the symbol table of the program during execution, and in code can be represented as a HashMap or similar datatype in your language of choice. We will use a sequence of characters to represent the output of a program, with ϵ representing the empty sequence. I will also assume that all strings will be represented as sequences of characters. Assume there is a function *append* that, when given two sequences, appends the second sequence to the first. Also assume, there is a function *seq* that takes an integer and gives a sequence of characters representing that integer as text. Assume there are the functions *head*, which maps a sequence to its first element, *tail*, which maps a sequence to a new one created by removing the first element, *clean*, which maps a sequence of input characters to a new sequence by removing any non-digits from the front of the sequence, and *int* that maps a sequence of digits to the corresponding integer. If the sequence is empty, *int* will give zero. A state, as well as the meaning of a program, will be a 3-tuple consisting of a variable name mapping function, a sequence of input characters and an output sequence. The initial state for any program is $(\{\}, i, \epsilon)$, where i is some sequence of characters the user will input. If a token (represented by all

caps and bold font) appears as a value on the right hand side of a function definition, then replace it with its lexeme. So if a **ID** was generated by the lexer from an x, then replace **ID** with x.

$$
\begin{aligned}
densem(\epsilon, (\theta, i, p)) &= (\theta, i, p) \\
densem(\texttt{<stmt>}\ \text{“;”}\ \texttt{<stmt\_list>}, (\theta, i, p)) &= densem(\texttt{<stmt\_list>}, densem(\texttt{<stmt>}, (\theta, i, p))) \\
densem(\text{“print”}\ \textbf{STRING}, (\theta, i, p)) &= (\theta, i, append(p, \textbf{STRING})) \\
densem(\text{“print”}\ \texttt{<expr>}, (\theta, i, p) &= (\theta, i, append(p, seq(out))) \\
&\quad \text{where } out = exprsem(\texttt{<expr>}) \\
densem(\text{“get”}\ \textbf{ID}, (\theta, i, p)) &= (\theta', i', p) \\
&\quad \text{where} \\
&\quad (x, i') = getInt(clean(i)) \\
&\quad \theta'(n) = \text{if } n = \textbf{ID} \text{ then } x \text{ else } \theta(n) \\
densem(\textbf{ID}\ \text{“=”}\ \texttt{<expr>}, (\theta, i, p)) &= (\theta', i, p) \\
&\quad \text{where} \\
&\quad \theta'(n) = \text{if } n = \textbf{ID} \text{ then } exprsem(\texttt{<expr>}, \theta) \text{ else } \theta(n) \\
densem(\texttt{<if>}, (\theta, i, p)) &= \text{if } exprsem(\texttt{<if>.<expr>}, \theta) \neq 0 \\
&\quad \text{then } densem(\texttt{<if>.<stmt\_list>}[0], (\theta, i, p)) \\
&\quad \text{else } densem(\texttt{<if>.<stmt\_list>}[1], (\theta, i, p)) \\
densem(\texttt{<while>}, (\theta, i, p)) &= \text{if } exprsem(\texttt{<while>.<expr>}, \theta) = 0 \\
&\quad \text{then } (\theta, i, p) \\
&\quad \text{else } densem(\texttt{<while>}, \\
&\qquad densem(\texttt{<while>.<stmt\_list>}, (\theta, i, p))) \\
densem(\texttt{<for>}, (\theta, i, p)) &= \text{if } exprsem(\texttt{<for>.<expr>.<assign>.“,”.} \\
&\qquad \texttt{<expr>.“,”.<expr>}, \theta) = 0 \\
&\quad \text{then } (\theta, i, p) \\
&\quad \text{else } densem(\texttt{<for>}, \\
&\qquad densem(\texttt{<for>.<stmt\_list>}, (\theta, i, p))) \\
exprsem(\texttt{<expr>}, \theta) &= \text{if } \texttt{<expr>.<b\_expr>} = \epsilon \\
&\quad \text{then } exprsem(\texttt{<expr>.<n\_expr>}, \theta) \\
&\quad \text{else } bexprsem(\texttt{<expr>.<b\_expr>}, \\
&\qquad exprsem(\texttt{<expr>.<n\_expr>}), \theta) \\
exprsem(\texttt{<n\_expr>}, \theta) &= \text{if } \texttt{<n\_expr>.<t\_expr>} = \epsilon \\
&\quad \text{then } exprsem(\texttt{<n\_expr>.<term>}, \theta) \\
&\quad \text{else } texprsem(\texttt{<n\_expr>.<t\_expr>}, \\
&\qquad exprsem(\texttt{<n\_expr>.<term>}), \theta) \\
exprsem(\texttt{<term>}, \theta) &= \text{if } \texttt{<term>.<f\_expr>} = \epsilon \\
&\quad \text{then } exprsem(\texttt{<term>.<factor>}, \theta) \\
&\quad \text{else } fexprsem(\texttt{<term>.<f\_expr>},
\end{aligned}
$$

$$
\begin{aligned}
& & & exprsem(\texttt{<term>}\,.\,\texttt{<factor>}),\theta) \\
exprsem(\texttt{<factor>},\theta) &= \text{if } \texttt{<factor>}\,.\,\texttt{<v\_expr>} = \epsilon \\
& & & \text{then } exprsem(\texttt{<factor>}\,.\,\texttt{<value>},\theta) \\
& & & \text{else } vexprsem(\texttt{<factor>}\,.\,\texttt{<v\_expr>}, \\
& & & \qquad exprsem(\texttt{<factor>}\,.\,\texttt{<value>}),\theta) \\
exprsem(\text{“(”}\ \texttt{<expr>}\ \text{“)”},\theta) &= exprsem(\texttt{<expr>},\theta) \\
exprsem(\text{“not”}\ \texttt{<value>},\theta) &= \text{if } exprsem(\texttt{<value>},\theta) = 0 \text{ then 1 else 0} \\
exprsem(\text{“-”}\ \texttt{<value>},\theta) &= -exprsem(\texttt{<value>},\theta) \\
exprsem(\textbf{ID},\theta) &= \theta(\textbf{ID}) \\
exprsem(\textbf{INT},\theta) &= \textbf{INT} \\
bexprsem(\text{“and”}\ \texttt{<n\_expr>},v,\theta) &= \text{if } v \neq 0 \text{ and } exprsem(\texttt{<n\_expr>},\theta) \neq 0 \text{ then 1 else 0} \\
bexprsem(\text{“or”}\ \texttt{<n\_expr>},v,\theta) &= \text{if } v \neq 0 \text{ or } exprsem(\texttt{<n\_expr>},\theta) \neq 0 \text{ then 1 else 0} \\
texprsem(\text{“+”}\ \texttt{<n\_expr>},v,\theta) &= v + exprsem(\texttt{<n\_expr>},\theta) \\
texprsem(\text{“-”}\ \texttt{<n\_expr>},v,\theta) &= v - exprsem(\texttt{<n\_expr>},\theta) \\
fexprsem(\text{“*”}\ \texttt{<term>},v,\theta) &= v \times exprsem(\texttt{<term>},\theta) \\
fexprsem(\text{“/”}\ \texttt{<term>},v,\theta) &= \frac{v}{exprsem(\texttt{<term>},\theta)} \\
fexprsem(\text{“\%”}\ \texttt{<term>},v,\theta) &= v \mod exprsem(\texttt{<term>},\theta) \\
vexprsem(\text{“>”}\ \texttt{<value>},v,\theta) &= \text{if } v > exprsem(\texttt{<value>},\theta) \text{ then 1 else 0} \\
vexprsem(\text{“>=”}\ \texttt{<value>},v,\theta) &= \text{if } v \geqslant exprsem(\texttt{<value>},\theta) \text{ then 1 else 0} \\
vexprsem(\text{“<”}\ \texttt{<value>},v,\theta) &= \text{if } v < exprsem(\texttt{<value>},\theta) \text{ then 1 else 0} \\
vexprsem(\text{“<=”}\ \texttt{<value>},v,\theta) &= \text{if } v \leqslant exprsem(\texttt{<value>},\theta) \text{ then 1 else 0} \\
vexprsem(\text{“==”}\ \texttt{<value>},v,\theta) &= \text{if } v = exprsem(\texttt{<value>},\theta) \text{ then 1 else 0} \\
vexprsem(\text{“!=”}\ \texttt{<value>},v,\theta) &= \text{if } v \neq exprsem(\texttt{<value>},\theta) \text{ then 1 else 0} \\
getInt(i) &= (int(x), i') \\
& & & \text{where } (x, i') = getIntSeq(\epsilon, i) \\
getIntSeq(i_1, i_2) &= \text{if } digit(head(i_2)) \\
& & & \text{then } getIntSeq(append(i_1, head(i_2)), tail(i_2)) \\
& & & \text{else } (i_1, i_2)
\end{aligned}
$$

3.4 State Transition Diagram

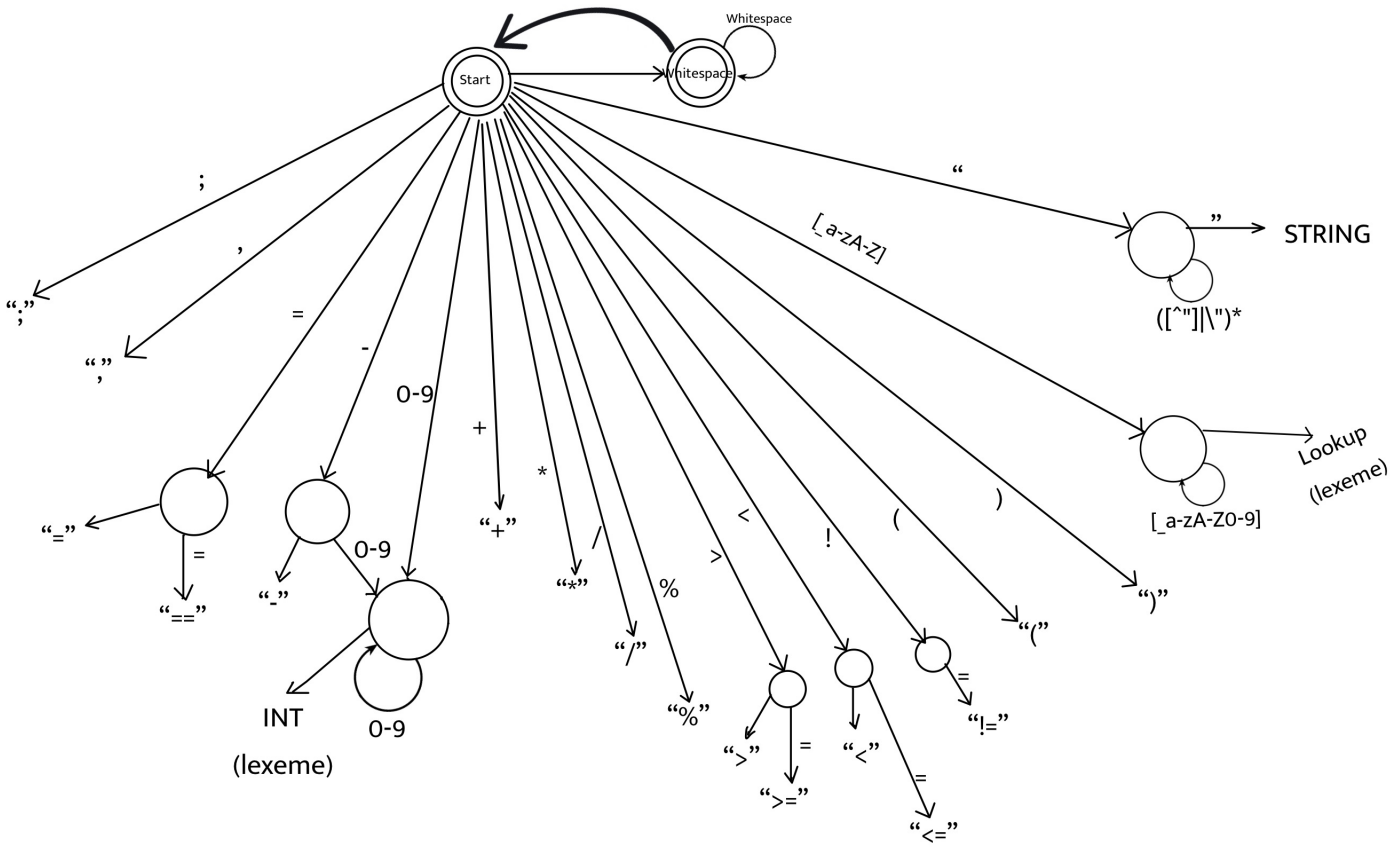

Figure 1: State Transition Diagram.

3.5 Test Programs: Output

This section contains a few test programs with output.

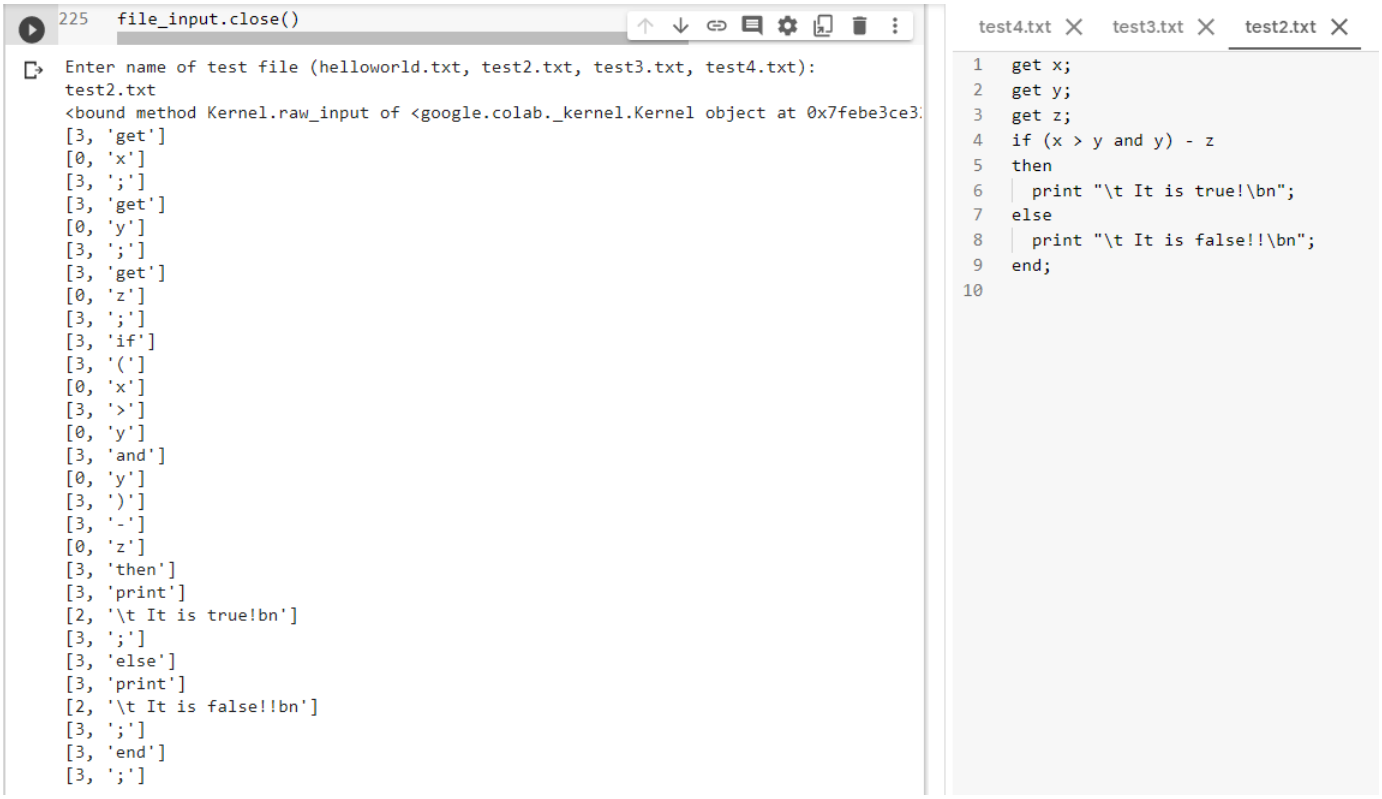

Figure 2: test2.txt's output.

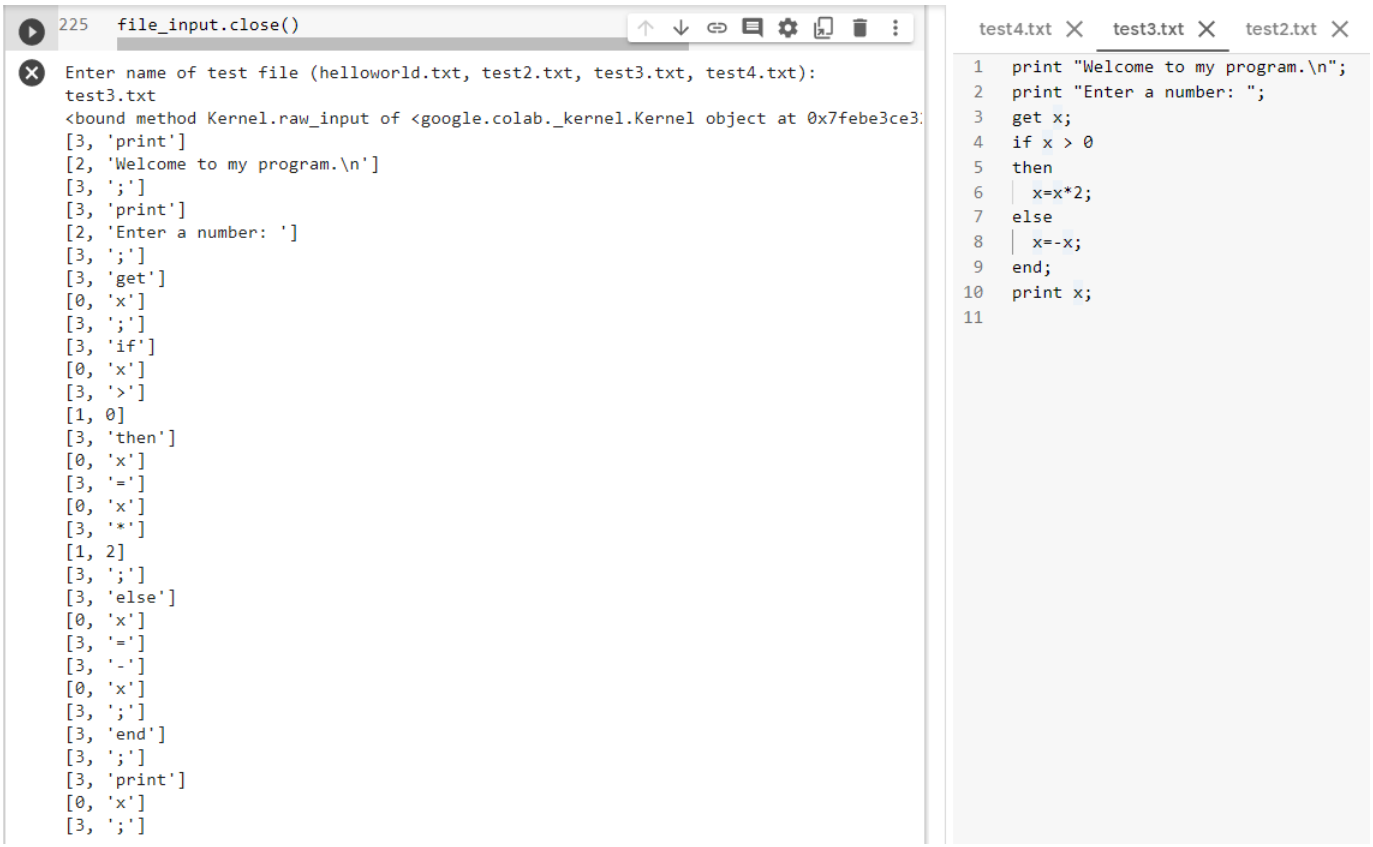

Figure 3: test3.txt's output.