

Assignment 1

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Group 407

1. **(0.5 points)** Give an example of a finite hypothesis class \mathcal{H} with $\text{VCdim}(\mathcal{H}) = 2022$. Justify your choice.

Solution

Let H_{rec}^d be the class of axis aligned rectangles in \mathbb{R}^d .

In *Lecture 6* we proved that $\text{VCdim}(H_{rec}^2) = 4$.

In *Seminar 3* we proved generally that $\text{VCdim}(H_{rec}^d) = 2 * d$.

Then, we can choose $d = 1011$, for which $\text{VCdim}(H_{rec}^{1011}) = 2022$ as the class of aligned rectangles in \mathbb{R}^{1011} .

Proof

In order to show that $\text{VCdim}(H_{rec}^d) = 2 * d$, we need to show that

- 1) there exists a set C of $2d$ points that is shattered by (H_{rec}^d) .
- 2) every set C of $2d + 1$ points is not shattered by (H_{rec}^d) .

$$\mathcal{H}_{rec}^d = \{h_{(a_1, b_1, a_2, b_2, \dots, a_d, b_d)} \mid a_i \leq b_i, i = \overline{1, d}\}$$

$$h_{(a_1, b_1, a_2, b_2, \dots, a_d, b_d)}(\underline{x}) = \begin{cases} 1, & a_i \leq x^i \leq b_i \quad \forall i = \overline{1, d} \\ 0, & \text{otherwise} \end{cases}$$

$\underline{x} = (x^1, x^2, \dots, x^d)$

In order to show that $\text{VCdim}(\mathcal{H}_{rec}^d) = 2d$, we need to show that:

- 1) There exists a set C of $2d$ points that is shattered by \mathcal{H}_{rec}^d (this will mean that $\text{VCdim}(\mathcal{H}_{rec}^d) \geq 2d$)
- 2) Every set C of $2d + 1$ points is not shattered by \mathcal{H}_{rec}^d (this will mean that $\text{VCdim}(\mathcal{H}_{rec}^d) < 2d + 1$)

Proof of 1) Consider $C = \{c_1, c_2, c_3, \dots, c_{2d-1}, c_{2d}\}$ where

$$\begin{aligned} c_1 &= (1, 0, 0, \dots, 0) &= e_1 \\ c_2 &= (0, 1, 0, \dots, 0) &= e_2 \\ &\vdots \\ c_d &= (0, 0, 0, \dots, 1) &= e_d \\ c_{d+1} &= (-1, 0, 0, \dots, 0) &= -e_1 \\ c_{d+2} &= (0, -1, 0, \dots, 0) &= -e_2 \\ &\vdots \\ c_{2d} &= (0, 0, 0, \dots, -1) &= -e_d \end{aligned} \quad \begin{aligned} c_i &= e_i = -c_{i+d} \\ \forall i &= \overline{1, d} \end{aligned}$$

For $d = 2$, we will have in 2 dimensions:

$$c_1 = (1, 0) \quad c_2 = (0, 1) \quad c_3 = (-1, 0) \quad c_4 = (0, -1)$$

We want to show that, for each labeling $(y_1, y_2, \dots, y_{2d})$ of the points $(c_1, c_2, \dots, c_{2d})$ (there are 2^{2d} possible labelings), there exists a function h in \mathcal{H}_{rec}^d such that $h(c_i) = y_i \forall i = \overline{1, 2d}$.

Consider a labeling $(y_1, y_2, \dots, y_{2d}) \in \{0, 1\}^{2d}$.

Each point c_i has all components $= 0$, apart from component i if $i \in \{1, \dots, d\}$ or $i - d$ if $i \in \{d+1, \dots, 2d\}$.

$$\begin{array}{c|c|c|c} c_1 = (1, 0, 0, \dots, 0) & c_2 = (0, 1, 0, \dots, 0) & \cdots & c_d = (0, 0, 0, \dots, 1) \\ c_{d+1} = (-1, 0, 0, \dots, 0) & c_{d+2} = (0, -1, 0, \dots, 0) & \cdots & c_{2d} = (0, 0, 0, \dots, -1) \end{array}$$

We want to find $h = h_{(a_1, b_1, a_2, b_2, \dots, a_d, b_d)}$ such that $h_{(a_1, b_1, a_2, b_2, \dots, a_d, b_d)}(c_i) = y_i$.

The choice of the interval $[a_i, b_i]$ is influenced by the labels y_i and y_{i+d} of the points c_i and c_{i+d} . As all other points $c_1, c_2, \dots, c_{i-1}, c_{i+1}, \dots, c_{i+d-1}, c_{i+d+1}, \dots, c_{2d}$ have 0 on the i -th component, we have that $[a_i, b_i]$ should contain 0, otherwise each point will be labelled with 0.

So $[a_i, b_i]$ depends on y_i and y_{i+d} , and $[a_i, b_i]$ decides basically the label of points c_i and c_{i+d} :

$$c_i = (0, \dots, 0, 1, 0, \dots, 0) \quad c_{i+d} = (0, \dots, 0, -1, 0, \dots, 0)$$

Possible cases:

1) $y_i = 0, y_{i+d} = 0$, then $[a_i, b_i] \cap \{-1, 1\} = \emptyset$

$[a_i, b_i]$ should not contain points -1 and 1.

In this case, take $a_i = -0.5, b_i = 0.5$ (many other choices are possible)

2) $y_i = 0, y_{i+d} = 1$, then $[a_i, b_i] \cap \{-1, 1\} = \{-1\}$

$[a_i, b_i]$ should contain only point -1 such that c_{i+d} will get label 1.

In this case, take $a_i = -2, b_i = 0.5$ (many other choices are possible)

3) $y_i = 1, y_{i+d} = 0$, then $[a_i, b_i] \cap \{-1, 1\} = \{1\}$

$[a_i, b_i]$ should contain only point +1 such that c_i will get label 1.

In this case, take $a_i = -0.5, b_i = 2$ (many other choices are possible)

4) $y_i = 1, y_{i+d} = 1$, then $[a_i, b_i] \cap \{-1, 1\} = \{-1, 1\}$

$[a_i, b_i]$ should contain both points $\{-1, 1\}$ such that c_i and c_{i+d} will get label 1.

In this case, take $a_i = -2, b_i = 2$ (many other choices are possible)

In all cases, we have that $h_{(a_1, b_1, a_2, b_2, \dots, a_d, b_d)}(c_i) = y_i, \forall i = \overline{1, 2d}$, where each interval $[a_i, b_i]$ was determined based on y_i and y_{i+d} , $i = \overline{1, d}$.

So, $VC \dim(\mathcal{H}_{rec}^d) \geq 2d$.

□

2. (0.5 points) What is the maximum value of the natural even number n , $n = 2m$, such that there exists a hypothesis class \mathcal{H} with n elements that shatters a set C of $m = \frac{n}{2}$ points? Give an example of such an \mathcal{H} and C . Justify your answer.

Solution

If class \mathcal{H} with n elements shatters a set C of $m = \frac{n}{2}$ points, then $VCdim(\mathcal{H}) \geq \frac{n}{2}$. (1)

Also, because \mathcal{H} is a class with n elements, then \mathcal{H} is a finite hypothesis class, which means that $VCdim(\mathcal{H}) \leq \log_2 n$. (2)

From (1) and (2), we can say that $\frac{n}{2} \leq \log_2 n$, where $n \in \mathbb{N}$, which means that $\frac{n}{2} - \log_2 n \leq 0$.

If we take into consideration that $n = 2 * m$, where $m \in \mathbb{N}$, the equation is transformed into $m - \log_2 2 * m \leq 0, \Rightarrow m - \log_2 m - 1 \leq 0$.

If $m = 0, n = 2 * m = 0$, the hypothesis class (\mathcal{H}) is \emptyset .

Let $f(m) = m - \log_2 m - 1, m \in \mathbb{N}$. We have to find the maximum m where $f(m) \leq 0, m \geq 1, m \in \mathbb{N}$.

For $m = 1, f(m) = 1 - 0 - 1 = 0 \leq 0$

For $m = 2, f(m) = 2 - 1 - 1 = 0 \leq 0$

For $m = 3, f(m) = 3 - \log_2^3 - 1 = 2 - \log_2^3 > 0$.

Also, for any $m \geq 3$, we can show that $f'(m) > 0$, because $f'(m) = 1 - \frac{1}{m * \ln m}$ and $1 - \frac{1}{m * \ln m} < 1$.

1. As both m and $\ln(m) \geq 1$ ($m \geq 3$), then $m * \ln m > 1$.

So we have to find an example for $m = 2, n = 4$, so a hypothesis class with n elements, which shatters a set C of 2 points.

3. (0.75 points) Let $\mathcal{X} = \mathbb{R}^2$ and consider \mathcal{H} the set of axis aligned rectangles with the center in origin $O(0, 0)$. Compute the $VCdim(\mathcal{H})$.

Solution

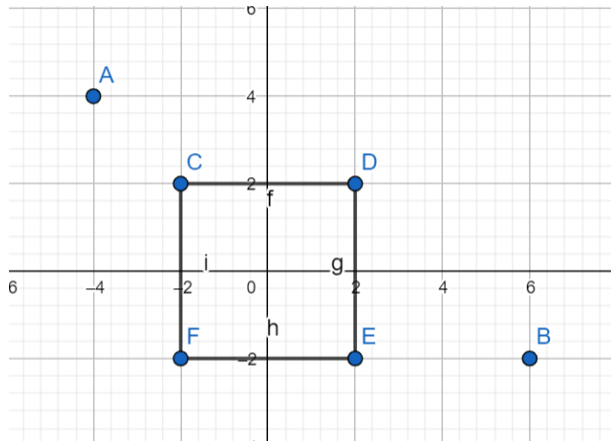
$VCdim(\mathcal{H}) = 2$

Proof

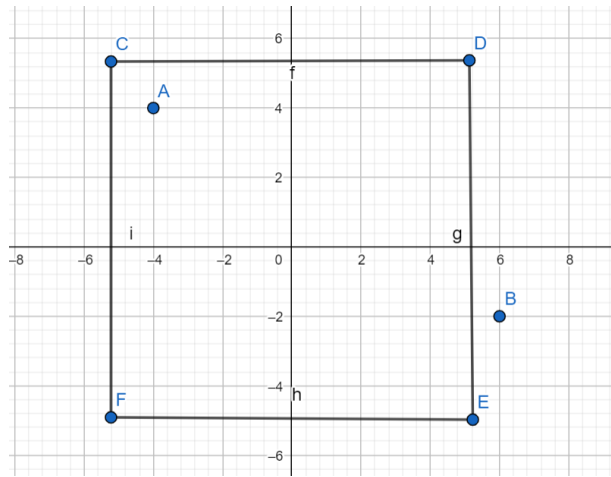
We will note $(H_{rec_{co}}^2)$ as the class of axis aligned rectangles with the center in origin $O(0, 0)$.

1) There exists a set C of 2 points that is shattered by $H_{rec_{co}}^2$ (this will mean that $H_{rec_{co}}^2 \geq 2d$).

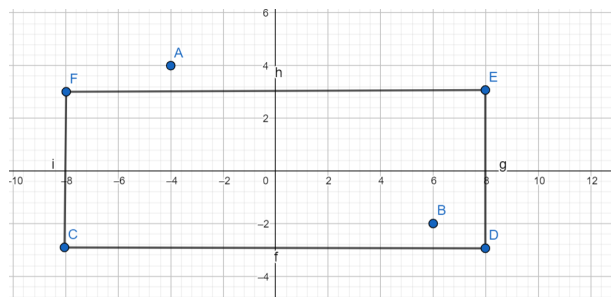
We will choose $A(-4, 4)$ and $B(-2, 6)$ from \mathbb{R}^2 . We will show that any labeling of these two points is valid by an axis aligned rectangle with the center in $O(0, 0)$.



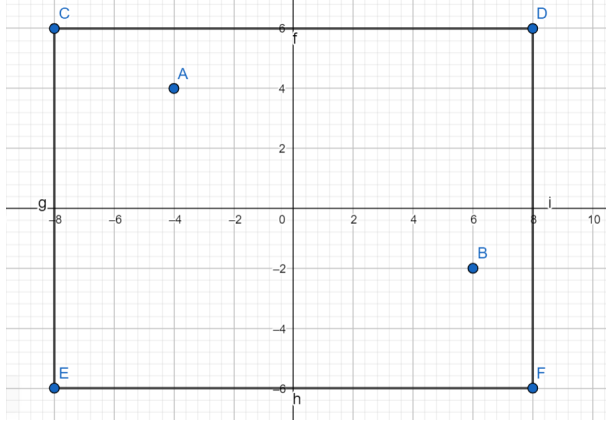
Label (A-0, B-0)



Label (A-1, B-0)



Label (A-0, B-1)



Label (A-1, B-1)

So, the given set of points is labelled by the class of axis aligned rectangles with the center in the origin.

2) Every set C of 3 points is not shattered by \mathcal{H}_{rec}^d (this will mean that $VC \dim(\mathcal{H}_{rec}^d) < 2d + 1$).

A rectangle with the center in the origin $O(0, 0)$ has the four points : $A1(-x, y)$, $A2(-x, -y)$, $A3(x, -y)$, $A4(x, y)$. So, it marks with label 1 any point which has coordinates (a, b) , where $|a| \leq |x|$ and $|b| \leq |y|$

Let $A = (x1, y1)$, $B = (x2, y2)$, $C = (x3, y3)$, where $A, B, C \in \mathcal{X} = \mathbb{R}^2$, considering $|x1| \leq |x2| \leq |x3|$.

Let's considering all the possibilities for the ordering of $y1, y2, y3$.

i) $|y1| \leq |y2| \leq |y3|$, in which the label $(0, 1, 0)$ can't be assigned by any rectangle (A2 can't be labelled as 1 without marking the A1 as well, as they have smaller x and y coordinates).

ii) $|y1| \leq |y3| \leq |y2|$, in which the label $(0, 1, 0)$ can't be assigned by any rectangle (A2 can't be labelled as 1 without marking A1 as well).

iii) $|y2| \leq |y1| \leq |y3|$, in which the label $(0, 0, 1)$ can't be assigned by any rectangle (A3 can't be labelled as 1 without marking A1 and A2).

iv) $|y2| \leq |y3| \leq |y1|$, in which the label $(0, 0, 1)$ can't be assigned by any rectangle (A3 can't be labelled as 1 without marking A2 as well).

v) $|y3| \leq |y1| \leq |y2|$, in which the label $(0, 1, 0)$ can't be assigned by any rectangle (A2 can't be labelled as 1 without marking A1 as well).

vi) $|y3| \leq |y2| \leq |y1|$, in which the label $(1, 0, 1)$ can't be assigned by any rectangle (in order to assign A1 and A3, the width and height of the rectangle must be bigger than $|x3|$ and $|y1|$ and then it labels A2 as well).

So, there can't be any set of 3 points shattered by \mathcal{H}_{rec}^2 , so $VCdim(\mathcal{H}) < 3$.

Form 1) and 2), $VCdim(\mathcal{H}) = 2$

4. **(1 point)** Let $\mathcal{X} = \mathbb{R}^2$ and consider \mathcal{H}_α the set of concepts defined by the area inside a right triangle ABC with two catheti AB and AC parallel to the axes (Ox and Oy), and with the ratio $AB/AC = \alpha$ (fixed constant > 0). Consider the realizability assumption. Show that the class \mathcal{H}_α is (ϵ, δ) -PAC learnable by giving an algorithm A and determining an upper bound on the sample complexity $m_H(\epsilon, \delta)$ such that the definition of PAC-learnability is satisfied.

Solution

Informally, \mathcal{H} is the set of similar triangles obtained from the base triangle with the points $A = (0, 0), B = (\alpha, 0), C = (0, 1)$ by doing the following operations: rescaling and translation.

Algorithm A: For each point $P \in S$ that is labelled as 1, generate the parallel lines through it to the sides of triangle ABC. Iterating through all the points, we obtain 3 sets of lines:

LAB = lines for which $\exists P \in S$ with $\text{label}(P) = 1, P \in d, d \parallel AB$

LAC = lines for which $\exists P \in S$ with $\text{label}(P) = 1, P \in d, d \parallel AC$

LBC = lines for which $\exists P \in S$ with $\text{label}(P) = 1, P \in d, d \parallel BC$

From LAB choose $d1$ that has smallest x. From LAC choose $d2$ that has the smallest y. From LBC choose $d3$ that has the biggest x.

Our algorithm will return the triangle that is at the intersection of the $d1, d2, d3$, which is similar to the ΔXYZ (because $d1, d2, d3$ are parallel with the lines of the ΔXYZ , it follows that this triangle is indeed in \mathcal{H}).

If there are no points labelled as 1, return any triangle that does not contain any points from S . Obviously, $L_S(h_S) = 0$.

We need to prove now that \mathcal{H} is PAC-learnable, doing the following:

For the following steps, we fix the $\epsilon, \delta > 0$.

In our case, all the positive points must be inside ΔABC , while the others outside...

5. **(1.25 points)** Consider $\mathcal{H} = \mathcal{H}_1 \cup \mathcal{H}_2 \cup \mathcal{H}_3$, where:

$$\mathcal{H}_1 = \{h_{\theta_1} : \mathbb{R} \rightarrow \{0, 1\} \mid h_{\theta_1}(x) = \mathbb{1}_{[x \geq \theta_1]}(x) = \mathbb{1}_{[\theta_1, +\infty)}(x), \theta_1 \in \mathbb{R}\},$$

$$\mathcal{H}_2 = \{h_{\theta_2} : \mathbb{R} \rightarrow \{0, 1\} \mid h_{\theta_2}(x) = \mathbb{1}_{[x < \theta_2]}(x) = \mathbb{1}_{(-\infty, \theta_2)}(x), \theta_2 \in \mathbb{R}\},$$

$$\mathcal{H}_3 = \{h_{\theta_1, \theta_2} : \mathbb{R} \rightarrow \{0, 1\} \mid h_{\theta_1, \theta_2}(x) = \mathbb{1}_{[\theta_1 \leq x \leq \theta_2]}(x) = \mathbb{1}_{[\theta_1, \theta_2]}(x), \theta_1, \theta_2 \in \mathbb{R}\}.$$

Consider the realizability assumption.

a) Compute $\text{VCdim}(\mathcal{H})$.

Solution

Initial consideration. If we pick $\theta_1 \leq \theta_2$, then every element is labelled as 1.

We first prove that $\text{VCdim}(\mathcal{H}) \geq 2$.

Let $\{x_1, x_2\} \in \mathbb{R}^2$, where $x_1 < x_2$

$H(x_1) = H(x_2) = 0$, pick $\theta_1 = x_2 + 1, \theta_2 = x_1 - 1$.

$H(x_1) = H(x_2) = 1$, pick $\theta_1 = \theta_2 = x_1$.

$H(x_1) = 0, H(x_2) = 1$, pick $\theta_1 = x_2, \theta_2 = x_1 - 1$.

$H(x_1) = 1, H(x_2) = 0$, pick $\theta_1 = x_2 + 1, \theta_2 = x_1$
 \mathcal{H} shatters any set $\{x_1, x_2\} \in \mathbb{R}^2$. So, $VCdim(\mathcal{H}) \geq 2$.

Now, we will prove that $VCdim(\mathcal{H}) < 3$

Let's consider $x_1, x_2, x_3 \in \mathcal{R}$, where $x_1 < x_2 < x_3$ and the following labelling $\mathcal{H}(x_1) = 0$, $\mathcal{H}(x_2) = 1, \mathcal{H}(x_3) = 0$

if $x_2 \in \mathcal{H}_1$ and $x_3 > x_2$, then x_3 will be labelled by \mathcal{H}_1 as 1.

if $x_2 \in \mathcal{H}_2$ and $x_1 < x_2$, then x_1 will be labelled by \mathcal{H}_2 as 1.

if $x_2 \in \mathcal{H}_3$, then $\theta_1 \leq \theta_2$, then any $x \in \mathcal{R}$ will be labelled as 1, so both x_1 and x_3 will be mislabelled.

This proves that $VCdim(\mathcal{H}) < 3$. Having proved previously that $VCdim(\mathcal{H}) \geq 2$

Then, prove that $VCdim(\mathcal{H}) = 2$.

[b)] Show that \mathcal{H} is PAC-learnable.

Solution

As is it shown in *Lecture 9*, any function which has a finite $VCdim$ is also *PAC-learnable*, within The Fundamental Theorem of Statistical Learning Theory.

[c)] Give an algorithm A and determine an upper bound on the sample complexity $m_{\mathcal{H}}(\epsilon, \delta)$ such that the definition of PAC-learnability is satisfied.

6. **(1 point)** A decision list may be thought of as an ordered sequence of if-then-else statements. The sequence of conditions in the decision list is tested in order, and the answer associated with the first satisfied condition is output.

More formally, a k -decision list over the boolean variables x_1, x_2, \dots, x_n is an ordered sequence $L = \{(c_1, b_1), (c_2, b_2), \dots, (c_l, b_l)\}$ and a bit b , in which each c_i is a conjunction of at most k literals over x_1, x_2, \dots, x_n and each $b_i \in \{0, 1\}$. For any input $a \in \{0, 1\}^n$, the value $L(a)$ is defined to be b_j where j is the smallest index satisfying $c_j(a) = 1$; if no such index exists, then $L(a) = b$. Thus, b is the "default" value in case a falls off the end of the list. We call b_i the bit associated with the condition c_i .

Show that the VC dimension of 1-decision lists over $\{0, 1\}^n$ is lower and upper bounded by linear functions, by showing that there exists $\alpha, \beta, \gamma, \delta \in \mathbb{R}$ such that:

$$\alpha \cdot n + \beta \leq VCdim(\mathcal{H}_{1\text{-decision list}}) \leq \gamma \cdot n + \delta$$

Hint: Show that 1-decision lists over $\{0, 1\}^n$ compute linearly separable functions (halfspaces).

Ex-officio: 0.5 points