1. I approached the choulenge by employing the tecnique that in volved two utility tunctions, mamely find and union." which efficiently managed disjoint sets. After instializing a list to trepresent parcent cities, I organized the troads in ascending order based on their associated maintainance costs. While progressing through these shoreted roads, I incremented an accumulating total when the cities at the ends belonged to separate sets. Subsequently, I utilized the union function to merige these sets. The input details where extreacted from imput 1_2. txt, and the resulting output, signifying the minimized maintenance cost achieved by thoughtfully selecting a subset of moads, was neconded within output? _. 2 txt.

2. I used dynamic programming to calculate the number of unique ways friedly can achine's achieve this. The count-ways () function recursively determines the count while utilizing an array dp to store previously computed. results, minimizing tredundancy. The input is read from "input 2-4+xt", and the calculated tresult is written to output 2-4+xt. This approach

optimizes the solution by avoiding repetitive

3. I defined a list dp of length X+1 to storce the minimum number of coins required to make up each amount from O to X. I initialize dp to 0 and all other elements to a large number. INF. Then, for each coins denomination C[i], I Aterrated over all amounts from C[i] to X and. update dp [j] as the minimum of dp[j] and dp[j-C[i]] +1, since I com either use the con C[i] on not to make up the amount j. Finally, I returned dp[x] as the minimum number of coins trequired to make up the tanget amount, on-1 if dp[x] is still INF.