

Denoising, Decomposition and Demixing of Moment Sequences by Convex Optimization

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Preliminary Examination

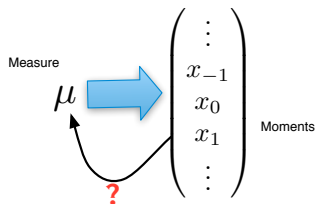
July 26, 2012

Moments

The moments of a measure μ are given by

$$x_m = \int z^m d\mu(z)$$

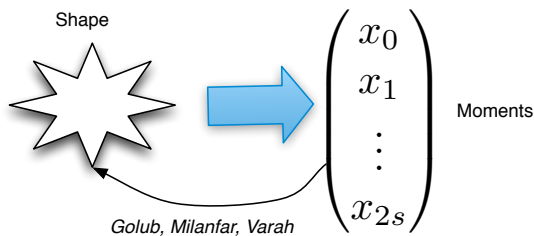
Moment Problem



Trigonometric Moments μ on \mathbb{T}

Power Moments μ on \mathbb{R}

Shape from Moments

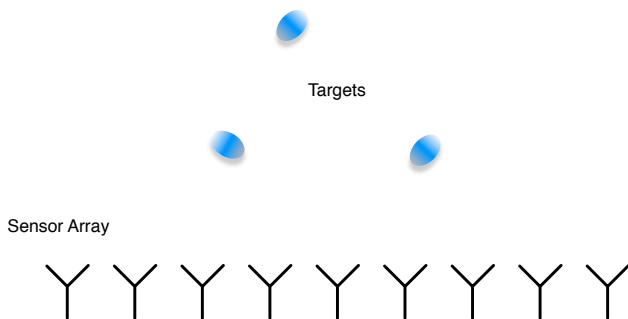


Theorem (Motzkin-Schoenberg)

Suppose \mathcal{P} has vertices z_1, \dots, z_s .

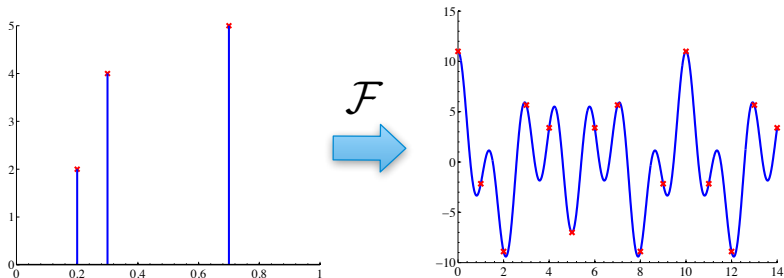
$$k(k-1) \int_{\mathcal{P}} z^{k-2} dx dy = \sum_{j=1}^k a_j z_j^k = \int z^k \left(\sum_{j=1}^k a_j \delta(z - z_k) \right) dx dy$$

Array Signal Processing



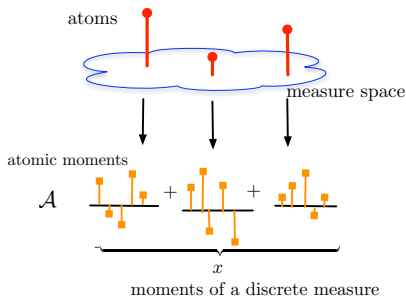
- ▶ Estimate location and the direction of arrival
- ▶ Deploy as few sensors as possible

Line Spectrum Estimation



- Discrete spectrum = Mixture of sinusoids

Structure and Simplicity



- ▶ Nonlinear estimation problem.
- ▶ Ill-posed – need structure.
- ▶ Simple objects: few atoms.
- ▶ Continuously many atoms.

$$x = \sum_{a \in T} c_a a$$

“weights” c_a “features” a

simple $\Rightarrow T \subset \mathcal{A}$ is small

Examples of Simple Objects

Object

Vectors

Matrices

Bandlimited Signals

Linear Systems

Notion of Parsimony

Sparsity

Rank

No. of Frequencies

McMillan Degree

Atoms (\mathcal{A})

1-sparse vectors

Rank-1 matrices

Complex sinusoids

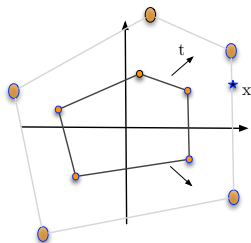
Single pole systems

Atomic Norm

(Chandrasekaran, Recht, Parillo, Willsky, 2010)

- The atomic norm at x is given by:

$$\begin{aligned}\|x\|_{\mathcal{A}} &= \inf \{t > 0 \mid x \in t \operatorname{conv}(\mathcal{A})\} \\ &= \inf \left\{ \sum_a c_a \mid x = \sum_{a \in \mathcal{A}} c_a a, c_a > 0 \right\}\end{aligned}$$



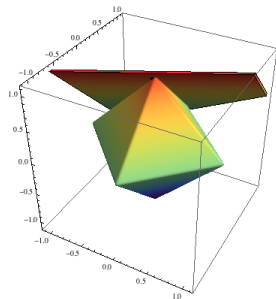
- The dual norm is the support function of \mathcal{A} :

$$\|z\|_{\mathcal{A}}^* = \sup_{x \in \mathcal{A}} \langle z, x \rangle.$$

Examples

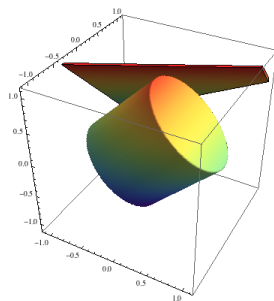
- ▶ x^* is sparse. $y = \Phi x^*$.
- ▶ $\text{minimize}_x \|x\|_1$
subject to $y = \Phi x$

Figure: ℓ_1 norm ball $\left\| \begin{bmatrix} x & y & z \end{bmatrix}^T \right\|_1 \leq 1$

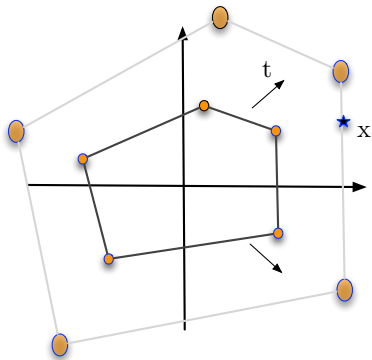


- ▶ X^* is low rank. $Y = \mathcal{A}(X^*)$.
- ▶ $\text{minimize}_X \|X\|_*$
subject to $Y = \mathcal{A}(X)$.

Figure: Nuclear Norm ball $\left\| \begin{bmatrix} x & y \\ y & z \end{bmatrix} \right\|_* \leq 1$



Denoising with Atomic Norms



Atomic Soft Thresholding

Suppose x^\star is simple and we observe $y = x^\star + w$.

Primal Problem

$$\underset{x}{\text{minimize}} \quad \frac{1}{2} \|y - x\|_2^2 + \tau \|x\|_{\mathcal{A}}$$

Atomic Soft Thresholding

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$$\begin{aligned} & \underset{z}{\text{minimize}} \quad \|y - z\|_2 \\ & \text{subject to} \quad \|z\|_{\mathcal{A}}^* \leq \tau \end{aligned}$$

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$$\begin{aligned} & \underset{z}{\text{minimize}} \quad \|y - z\|_2 \\ & \text{subject to} \quad \|z\|_{\mathcal{A}}^* \leq \tau \end{aligned}$$

- ▶ The primal-dual solution pair (\hat{x}, \hat{z}) satisfy
(i) $y = \hat{x} + \hat{z}$, (ii) $\langle \hat{x}, \hat{z} \rangle = \tau \|\hat{x}\|_{\mathcal{A}}$, (iii) $\|\hat{z}\|_{\mathcal{A}}^* \leq \tau$.

Atomic Soft Thresholding

Suppose x^\star is simple and we observe $y = x^\star + w$.

Primal Problem

$$\underset{x}{\text{minimize}} \quad \frac{1}{2} \|y - x\|_2^2 + \tau \|x\|_{\mathcal{A}}$$

Dual Problem

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- ▶ The primal-dual solution pair (\hat{x}, \hat{z}) satisfy
(i) $y = \hat{x} + \hat{z}$, (ii) $\langle \hat{x}, \hat{z} \rangle = \tau \|\hat{x}\|_{\mathcal{A}}$, (iii) $\|\hat{z}\|_{\mathcal{A}}^* \leq \tau$.
- ▶ Correct choice of regularization parameter. $\tau = \mathbb{E} \|w\|_{\mathcal{A}}^*$

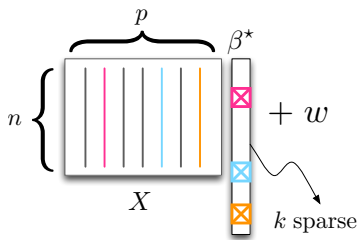
Proposition (MSE)

If $\tau \geq \tau_0 \geq \mathbb{E}(\|w\|_{\mathcal{A}}^*)$, then $\mathbb{E}(\|x^\star - \hat{x}\|_2^2) \leq \tau \|x^\star\|_{\mathcal{A}}$.

Proposition (MSE)

If $\tau \geq \tau_0 \geq \mathbb{E}(\|w\|_{\mathcal{A}}^*)$, then $\mathbb{E}(\|x^* - \hat{x}\|_2^2) \leq \tau \|x^*\|_{\mathcal{A}}$.

Application to sparse combination of columns



$$\blacktriangleright \frac{1}{n} \mathbb{E} \left(\|X\hat{\beta} - X\beta^*\|_2^2 \right) \leq \sigma \sqrt{\frac{k \log(p)}{n}} \|\beta^*\|_2.$$

\blacktriangleright No assumptions on $X \implies$ Slow rate.

Fast Rates

Conditions on approximate descent cone of \mathcal{A} :

- ▶ $T_\gamma(x^\star, \mathcal{A}) = \text{cone}(\{z \mid \|x^\star + z\|_{\mathcal{A}} \leq \|x^\star\|_{\mathcal{A}} + \gamma\|z\|_{\mathcal{A}}\})$.
- ▶ $\phi_\gamma(x^\star, \mathcal{A}) = \inf \left\{ \frac{\|z\|_2}{\|z\|_{\mathcal{A}}} \mid z \in T_\gamma(x^\star, \gamma) \right\} > 0$

Theorem

If $\tau \geq 2\tau_0$ and $\phi_{1/2}(x^\star, \mathcal{A}) > 0$, then with high probability,

$$\|x^\star - \hat{x}\|_2^2 = O\left(\frac{\tau^2}{\phi_{1/2}^2}\right).$$

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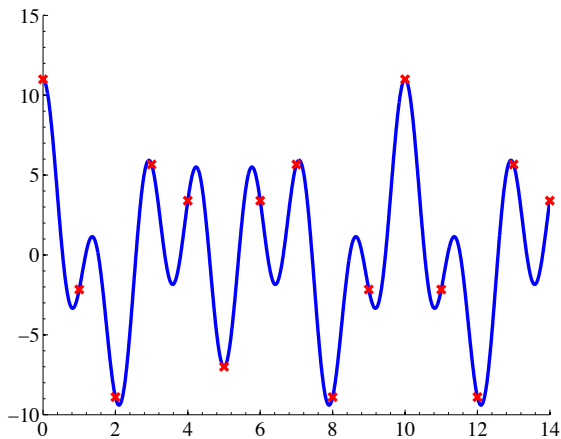
$$\|x^\star - \hat{x}\|_2^2 = O\left(\frac{\tau^2}{\phi_{1/2}^2}\right).$$

For the sparse combination of columns,

- ▶ With no assumptions on X , $\text{MSE} = O\left(\sqrt{\frac{k \log(p)}{n}}\right)$.
- ▶ If $\phi_{1/2}(X\beta^\star) > 0$, $\text{MSE} = O\left(\frac{k \log(p)}{n}\right)$.

Atomic norm denoising with applications to line spectrum estimation

B, Gongguo Tang, Ben Recht



Trigonometric Moments and Line Spectrum

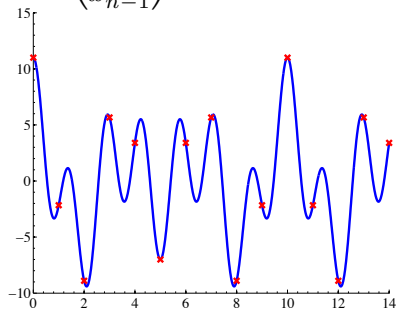
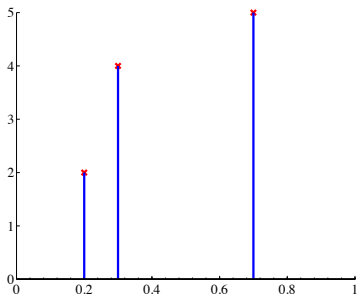
Measure

$$\mu = \sum_{j=1}^s c_j^* \delta_{u_j^*}$$



Trigonometric Moments

$$\begin{pmatrix} x_0^* \\ \vdots \\ x_k^* \\ \vdots \\ x_{n-1}^* \end{pmatrix} = \sum_{j=1}^s c_j^* e^{i2\pi u_j^* k}$$



$$\underbrace{\begin{pmatrix} x_0^* \\ x_1^* \\ \vdots \\ x_{n-1}^* \end{pmatrix}}_{x^*} = \sum_{j=1}^s c_j^* \underbrace{\begin{pmatrix} 1 \\ \exp(i2\pi u_j^*) \\ \vdots \\ \exp(i2\pi(n-1)u_j^*) \end{pmatrix}}_{a(u_j^*)}$$

$$x^* = \sum_{j=1}^s c_j^* a(u_j^*)$$

After a little algebra,

$$Tx^* = \begin{pmatrix} x_0^* & x_1^* & \cdots & x_{n-1}^* \\ x_{-1}^* & x_0^* & \ddots & x_n^* \\ \vdots & \ddots & \ddots & \vdots \\ x_{-n+1}^* & \cdots & \cdots & x_0^* \end{pmatrix} = \sum_{j=1}^s c_j^* a(u_j^*) a(u_j^*)^*.$$

► $Tx^* \in \mathbb{C}^{n \times n}$ is always rank s , as long as $n \geq s$.

Classical Line Spectrum Estimation

- ▶ Recall $Tx^* = \sum_{j=1}^s c_j^* a(u_j^*) a(u_j^*)^*$
- ▶ $n = s + 1 \Rightarrow$ null vector q satisfies $q^* T x^* q = 0$.

$$q^* T x^* q = \sum_{j=1}^s c_j^* |q^* a(u_j^*)|^2 = \sum_{j=1}^s c_j^* \left| \sum_{k=0}^n q_k^* e^{i2\pi u_j^* k} \right|^2$$

Roots of $q \Rightarrow$ frequencies (PRONY, 1795).

- ▶ Sensitive to noise.

Cadzow's Heuristic

Suppose $y = x^\star + w$. Recall that Tx^\star must have rank s .

Algorithm 1 Cadzow's Alternating Projections

$i \leftarrow 0$

$X^{(i)} \leftarrow Ty.$

while termination conditions aren't met **do**

$X^{(i+1)} \leftarrow \arg \inf_{\text{rank}(X)=k} \|X - X^{(i)}\|$

$X^{(i+1)} \leftarrow \text{TOEPLITZISE}(X^{(i+1)})$

$i \leftarrow i + 1$

end while

return $T^{-1}(X^{(i)})$

Denoising Samples of Line Spectral Signals

The diagram illustrates the signal model $x^* = \sum_{j=1}^s c_j^* a(u_j^*)$. Annotations include:

- A green arrow points from the label "amplitudes" to the coefficient c_j^* .
- A blue arrow points from the label "frequencies" to the frequency u_j^* .
- An orange arrow points from the atomic moment vector $a(u_j^*)$ to a red-bordered box.

Atomic Moment Vectors

$$a(u) = \begin{pmatrix} 1 \\ \exp(i2\pi u) \\ \vdots \\ \exp(i2\pi(n-1)u) \end{pmatrix}.$$

Positive Atoms $\mathcal{A}_+ = \{a(u) \mid u \in [0, 1]\}$.

Centrosymmetric $\mathcal{A} = \{e^{i\phi} a(u) \mid u \in [0, 1], \phi \in [0, 2\pi]\}$

Atomic Soft Thresholding

$$\text{minimize}_x \frac{1}{2} \|x - y\|_2^2 + \tau \|x\|_{\mathcal{A}}$$

Semidefinite characterization (known)

Figure: Moment Curve $(\cos(t), \cos(2t), \cos(3t))$

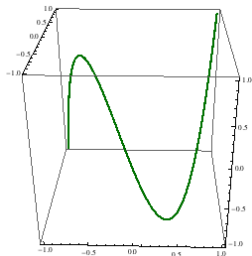
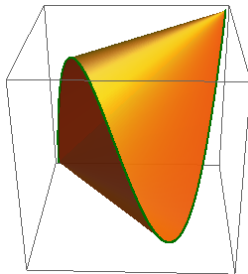


Figure: Convex Hull of the Positive Moment Curve



$$\|x\|_{\mathcal{A}} = \begin{cases} nx_0 & Tx \succeq 0 \\ \infty & \text{otherwise.} \end{cases}$$

Semidefinite characterization (our result)

Figure: Phase Symmetric Moment Curve

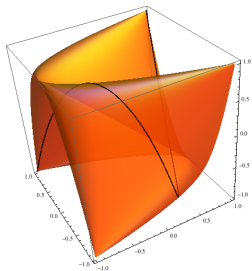
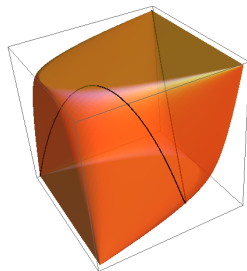


Figure: Convex Hull



$$\|x\|_{\mathcal{A}} = \inf \left\{ \frac{1}{2} \text{trace}(Tu + \frac{1}{2}t) \mid \begin{bmatrix} Tu & x \\ x^* & t \end{bmatrix} \succeq 0 \right\}.$$

Choice of Regularization Parameter

The “correct” choice of regularization parameter is given by the expected dual atomic norm of noise. We have,

$$\begin{aligned}\mathbb{E}(\|w\|_{\mathcal{A}}^*) &= \mathbb{E} \sup_{u, \phi} \left\langle w, e^{i\phi} a(u) \right\rangle \\ &= \mathbb{E} \sup_u \left| \sum_{k=0}^{n-1} w_k e^{i2\pi u k} \right|\end{aligned}$$

which is the expected maximum modulus of a random polynomial.

- ▶ We show

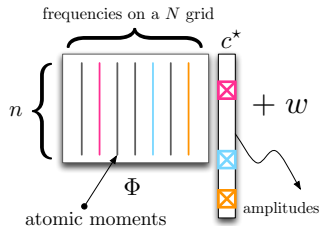
$$\mathbb{E}(\|w\|_{\mathcal{A}}^*) \leq \sigma \left(1 + \frac{1}{\log(n)} \right) \sqrt{n \log(n) + n \log(4\pi \log(n))}.$$

- ▶ We also show $\mathbb{E}(\|w\|_{\mathcal{A}}^*) \geq \sigma \sqrt{n \log(n) - \frac{n}{2} \log(4\pi \log(n))}.$
- ▶ Using MSE proposition (slow rate),

$$\mathbb{E} \|\hat{x} - x^*\|_2^2 \lesssim \sigma \sqrt{\frac{\log(n)}{n}} \sum_{l=1}^k |c_l^*|.$$

Discretized Atomic Soft Thresholding (DAST)

- ▶ For $N > 2\pi n$, put $\mathcal{A}_N = \{a_{m/N, \phi} \mid m = 0, \dots, N-1\}$.
- ▶ We show $(1 - \frac{2\pi n}{N}) \|x\|_{\mathcal{A}_N} \leq \|x\|_{\mathcal{A}} \leq \|x\|_{\mathcal{A}_N}$.
- ▶ Now a Lasso problem:

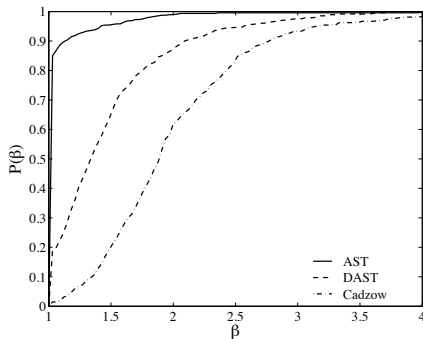


- ▶ *Remark:* DAST is $O(n \log N)$ since Φ is a DFT matrix
- ▶ Coherence doesn't matter for denoising.

Performance Profiles

$P(\beta)$ = Fraction of experiments with MSE less than $\beta \times$ minimum MSE.

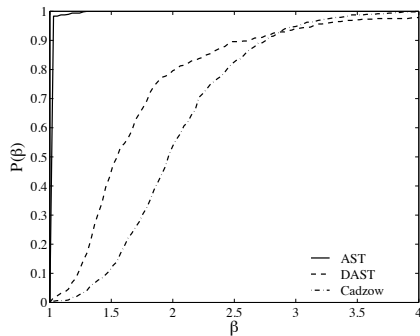
Figure: Random frequencies



Performance Profiles

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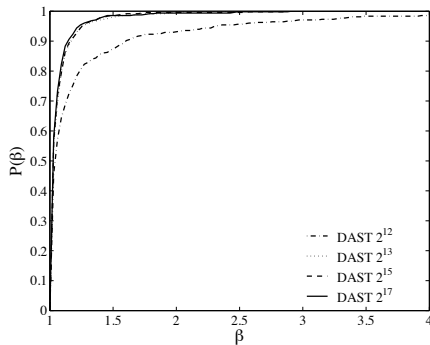
Figure: Equispaced frequencies



Performance Profiles

$P(\beta)$ = Fraction of experiments with MSE less than $\beta \times$ minimum MSE.

Figure: Comparing Discretization Levels



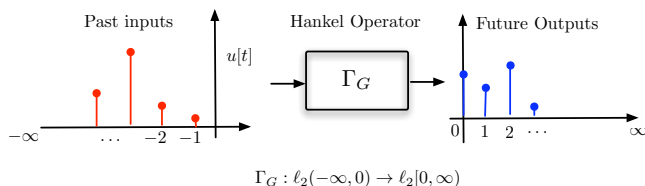
Linear System Identification via Atomic Norm Regularization

Parikshit Shah, B., Gongguo Tang, Ben Recht



Simple LTI systems

$$G(z) = \sum_{k=1}^{\infty} g_k z^{-k} = \sum_{j=1}^s \frac{c_j(1 - |w_j|^2)}{z - w_j}$$



McMillan degree = number of poles = $\text{rank}(\Gamma_G)$

Goal: Infer a system of low McMillan degree from linear measurements.

Hankel Nuclear Norm

- ▶ Minimize “nuclear norm” in place of rank!
- ▶ Spectral Series Representation:

$$\Gamma_G(x) = \sum_j \sigma_j \langle v_j, x \rangle u_j$$

Hankel Nuclear Norm

$$\|\Gamma_G\|_* = \sum_j \sigma_j$$

- ▶ LMI characterization of $\|\Gamma_G\|_*$?

Atomic Norm

Partial Fraction Expansion

$$G(z) = \sum_{j=1}^s \frac{c_j(1 - |w_j|^2)}{z - w_j}$$

Pick the atomic set

$$\mathcal{A} = \left\{ \frac{1 - |w|^2}{z - w} : w \in \mathbb{D} \right\}.$$

Atomic Norm

$$\|G(z)\|_{\mathcal{A}} = \inf \left\{ \sum_{w \in \mathbb{D}} |c_w| : G(z) = \sum_{w \in \mathbb{D}} \frac{c_w(1 - |w|^2)}{z - w} \right\},$$

Theorem (Peller, Coifman, Rochberg, Bonsall, Walsh)

$$\frac{\pi}{8} \|G\|_{\mathcal{A}} \leq \|\Gamma_G\|_1 \leq \|G\|_{\mathcal{A}}.$$

Atomic Norm Minimization

- ▶ Observe $y = \mathcal{L}(G) + w$.
- ▶ Examples
 - ▶ FIR coefficients
 - ▶ Frequency samples
 - ▶ Output to a known input

- ▶ Natural Minimization Problem to solve is

$$\text{minimize}_G \frac{1}{2} \|\mathcal{L}(G) - y\|_2^2 + \mu \|G\|_{\mathcal{A}}.$$

- ▶ We argue it is enough to solve the denoising problem

$$\text{minimize}_x \frac{1}{2} \|x - y\|_2^2 + \mu \|x\|_{\mathcal{L}(\mathcal{A})}.$$

where

$$\|x\|_{\mathcal{L}(\mathcal{A})} = \inf \left\{ \sum_{w \in \mathbb{D}_\rho} |c_w| \ : \ x_i = \sum_{w \in \mathbb{D}_\rho} c_w \mathcal{L}_i \left(\frac{1 - |w|^2}{z - w} \right) \right\}.$$

Our Result

Let us assume that \hat{G} is an approximate solution to the discretized version of the atomic soft thresholding problem.

Theorem

Suppose G^\star contains all poles within $\{z : |z| \leq \rho\}$ for some $\rho < 1$. Suppose we observe n samples

$$y_k = G^\star(e^{i2\pi k/n}) + w_k$$

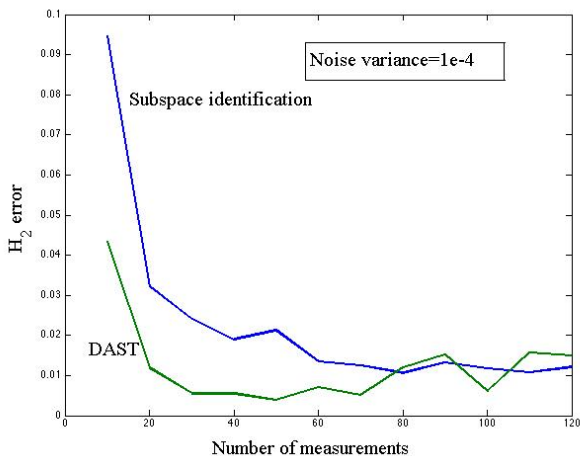
for $k = 0, \dots, (n-1)$ where w is i.i.d. $\mathcal{N}(0, \sigma^2)$. There is a quantity C depending on ρ and σ such that for sufficiently large n

$$\left\| \hat{G}(z) - G_\star(z) \right\|_{\mathcal{H}_2}^2 \leq C \|\Gamma_{G_\star}\|_1 n^{-\frac{1}{2}}$$

with probability exceeding $1 - e^{-o(n)}$.

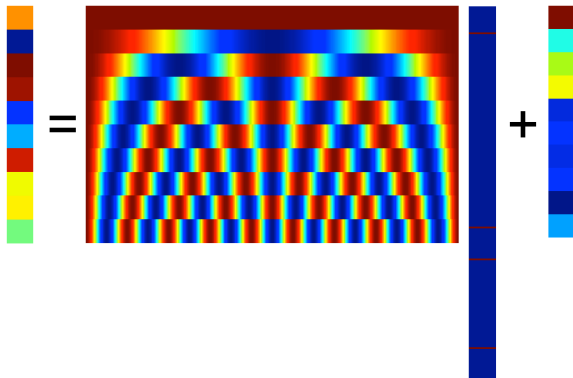
Experiments

- ▶ Compares favorably to Subspace ID for system identification.
- ▶ Does not need an estimate of model order.



Compressed Sensing off the Grid

Gongguo Tang, B., Parikshit Shah, Ben Recht



- ▶ Recover sparse vectors from a few incoherent measurements.
- ▶ CS Problem: s sparse DFT from $O(s \log(n))$ samples.
- ▶ *Continuous Case*: What if frequencies don't lie on a grid?

Theorem (Gongguo,B,Parikshit,Ben)

Let Δ be the minimum frequency separation in the signal. Exact recovery is possible with just $s \text{ polylog}(1/\Delta)$ random time samples by atomic norm minimization

- ▶ PRONY 1795: $2s$ samples sufficient. noise sensitive
- ▶ CANDÈS, FERNANDES-GRANDA 2012 : If $n > 2s$, convex optimization can recover spectra with $\Delta > \frac{4}{n}$.
- ▶ OUR RESULT: $O(s \log(s) \log(n))$ **random** samples are sufficient for $\Delta > \frac{4}{n}$.
 - ▶ Missing Data Problem
 - ▶ Sub-Nyquist sampling

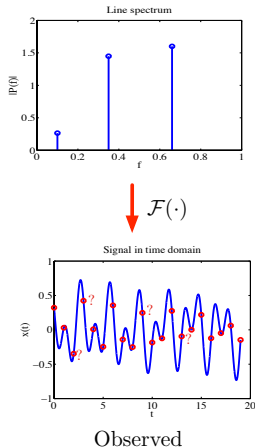


Figure: $\mathcal{A}_+ = (\cos(t), \cos(2t), \cos(3t))$

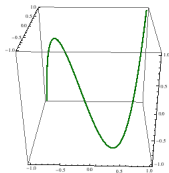


Figure: $\{\mathcal{A}_+, -\mathcal{A}_+\}$

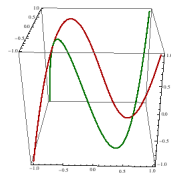


Figure: Convex Hull of \mathcal{A}_+

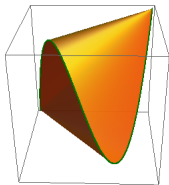
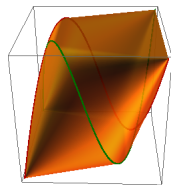


Figure: Convex Hull of $\{\mathcal{A}, -\mathcal{A}\}$



- ▶ Interpolate missing entries by

$$\begin{aligned} & \text{minimize}_x \|x\|_{\mathcal{A}} \\ & \text{subject to } x_j = x_j^*, \text{ for } j \in T. \end{aligned}$$

- ▶ Semidefinite formulation

$$\begin{aligned} & \text{minimize}_{u,t} \text{trace}(X) \\ & \text{subject to } X = \begin{pmatrix} Tu & x \\ x^* & t \end{pmatrix} \succeq 0 \\ & \quad x_j = x_j^*, \text{ for } j \in T. \end{aligned}$$

- ▶ Relaxation of

$$\begin{aligned} & \text{minimize}_{u,t} \text{rank}(X) \\ & \text{subject to } X = \begin{pmatrix} Tu & x \\ x^* & t \end{pmatrix} \succeq 0 \\ & \quad x_j = x_j^*, \text{ for } j \in T. \end{aligned}$$

Phase Transition Plots

Demixing

- ▶ Suppose x is sparse with respect to the atomic set \mathcal{A} and y is sparse with respect to the atomic set \mathcal{B}
- ▶ We observe $z = x + y$.
- ▶ When does

$$\begin{aligned} & \text{minimize}_{x,y} \|x\|_{\mathcal{A}} + \|y\|_{\mathcal{B}} \\ & \text{subject to } z = x + y. \end{aligned}$$

succeed in recovering x and y ?

Uncertainty Principles and Previous Work

- ▶ DONOHO AND STARK (1989) If a s sparse signal has a k sparse **DFT**, then

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- ▶ CANDÈS AND ROMBERG (2004) show that for most signals

$$k + s \sim \frac{n}{\sqrt{\log(n)}}$$

and in fact, when

$$k + s \leq C \frac{n}{\sqrt{\log(n)}}$$

ℓ^1 minimization succeeds with high probability.

Demixing Conjecture

Suppose we observe $z = x^\star + y^\star$, where $x^\star \in \mathbb{C}^n$ is a k -sparse combination of atoms from

$$\mathcal{A} = \left\{ [e^{i\phi} \ \dots \ e^{i(2\pi(n-1)t+\phi)}]^T, \ t \in [0, 1], \ \phi \in [0, 2\pi] \right\}.$$

with phases chosen uniformly randomly and $y^\star \in \mathbb{C}^n$ is a s -sparse vector. Then, (x^\star, y^\star) is the unique optimal solution to demixing atomic norm minimization if the minimum separation between frequencies in x^\star exceeds $\frac{1}{n}$ and

$$k + s \leq \frac{Cn}{\sqrt{\log(n)}}$$

where C is some numerical constant.

Accelerated Convergence Rates

- ▶ Experiments suggest much faster rates of convergence than predicted when frequencies are far apart.
- ▶ Our bound is pessimistic – Phase transition plots suggest the presence of a fast regime
- ▶ We can recover support at least “approximately” when frequencies are well separated.
- ▶ Difficult to connect separation with the condition on descent cones.

Fast Rate Conjecture

Suppose we observe noisy samples $y = x^\star + w$ where $w \in \mathcal{N}(0, \sigma^2 I_n)$ is Gaussian noise and for $j \in \{0, \dots, n-1\}$,

$$x_j^\star = \sum_{l=1}^k c_l e^{i2\pi j f_l}$$

where $c = [c_1 \dots c_k]^T \in \mathbb{C}^k$ are unknown amplitudes, $[f_1 \dots f_k] \in [0, 1]^k$ are normalized frequencies satisfying the separation condition

$$\Delta_{\min} = \min_{\substack{p \neq q \\ 1 \leq p, q \leq n}} (f_p - f_q) > \frac{1}{n},$$

the optimal solution \hat{x} to the atomic soft thresholding problem satisfies

$$\mathbb{E} \|\hat{x} - x^\star\|_2^2 \leq \frac{Ck\sigma^2 \log(n)}{n}.$$

where C is a numerical constant. The same rate holds even if we observe a fraction $s \log^2(n)$ of entries according to a Bernoulli model, instead of all n samples.

Minimax Rates

- ▶ Need to understand the fundamental limits of estimation in different regimes.
- ▶ Worst case error (or adversarial error) for the best estimator
- ▶ Can be information theoretically determined
- ▶ Coincides with the “fast rates” for other high dimensional inference problems.

Information Theoretic Determination

Basic Approach:

- ▶ Recast into a multiple hypothesis problem.
- ▶ Pack as many hypotheses at a critical separation so they are nearly indistinguishable.
- ▶ Find a lower bound on the probability of error by Fano's inequality.
- ▶ Fano's inequality requires an estimate of mutual information. Find a suitable upper bound for mutual information by a cover of the model space.
- ▶ This method provides optimal minimax rates for a number of problems.

Minimax Rate Conjecture for Line Spectrum Estimation

- ▶ Let Δ be the minimum separation between the normalized frequencies.
- ▶ Define the minimax rate in terms of the mean square error is given by

$$M(\Delta, k) = \min_{\hat{x}} \max_{x^* \in \mathcal{X}} \frac{1}{n} \mathbb{E} \|x^* - \hat{x}\|_2^2.$$

- ▶ *Conjecture:* If $\Delta > \frac{1}{n}$,

$$M(\Delta, k) \geq \frac{Ck\sigma^2 \log(n)}{n}$$

where C is independent of k, σ and n .

Scaling for Big Data

- ▶ Examples
 1. Heart rate and gas exchange data
 2. Time Series forecasting
 3. Neuronal spikes from Physiological Data
- ▶ Large scale problems - Challenges
- ▶ Parallel and Distributed Algorithms
- ▶ First order methods?
- ▶ Active set methods?

Trigonometric Moments

Theorem (Herglotz)

A sequence of numbers $\{x_k\}_{k=-\infty}^{\infty}$ form the trigonometric moments of a positive measure if and only if the sequence is positive definite. In other words, $Tx \succeq 0$ for every n and the sequence is Hermitian symmetric ($x_{-k} = x_k^*$).

Theorem (Caratheodory, Toeplitz, Fejer (1917))

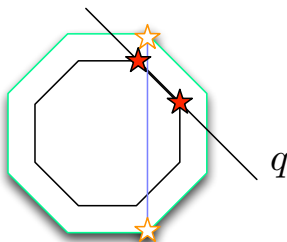
Suppose $x \in C^n$ such that $Tx \succeq 0$. Then, there exists r frequencies $f_1, \dots, f_r \in [0, 1]$ and positive numbers c_1, \dots, c_r such that

$$x_k = \sum_{j=1}^r c_j \exp(i2\pi f_j k) \text{ for each } k = 0, \dots, (n-1)$$

$$r = \text{rank}(Tx)$$

Exact Recovery

- ▶ Guarantee that a given decomposition has minimum atomic norm amongst all decompositions.
- ▶ Show it lies on a simplicial, exposed face of $\text{conv}(\mathcal{A})$



Dual Certificate for Line Spectral Estimation

Definition (Dual Certificate)

A dual certificate is a q satisfying

- ▶ $\langle q, a(f_j) \rangle = \text{sign}(c_j)$, for $j = 1, \dots, k$.
- ▶ $\langle q, a(f) \rangle < 1$ if $f \notin \{f_1, \dots, f_k\}$.

Dual Polynomial

Note that

$$\langle q, a(f) \rangle = \sum_{k=1}^n q_k e^{i2\pi k f} = Q(f)$$

is a trigonometric polynomial. So, we identify the “dual vector” with the “dual polynomial”

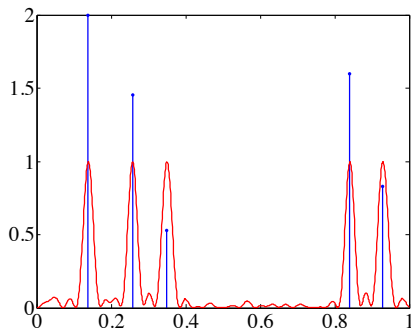
Dual Certificate for Line Spectrum Estimation

Necessary conditions for q to be a dual certificate:

$$Q(f_k) = \text{sign}(c_k) \quad \text{for each } k = 1, \dots, s$$

$$Q'(f_k) = 0 \quad \text{for each } k = 1, \dots, s$$

Figure: Dual Polynomial and Recovered Frequencies



Result of Candes and Fernandes-Granda

Candes and Fernandes-Granda (2012) propose writing

$$Q(f) = \sum_j \alpha_j K(f - f_j) + \beta_j K'(f - f_j)$$

where $K(f)$ is a trigonometric kernel and $\alpha \in \mathbb{C}^k$ and $\beta \in \mathbb{C}^k$ are undetermined coefficients.

Minimum separation:

$$\Delta = \min_{j \neq k} |f_j - f_k| \geq \Delta_{\min} = \frac{2}{n}.$$

If $\Delta \geq \frac{2}{n}$, convex optimization can exactly recover the correct decomposition from $n > 2s$ Nyquist samples.

- ▶ Prony's technique does not have a resolution limitation.
- ▶ Positive case doesn't either

Convex Hull of $\{\mathcal{A}, -\mathcal{A}\}$ – Insight

Figure: Moment Curve \mathcal{A} and $-\mathcal{A}$

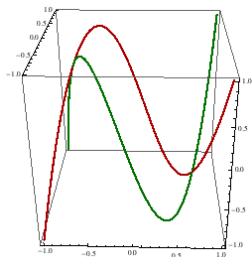
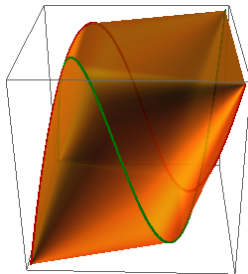


Figure: Convex Hull of $\{\mathcal{A}, -\mathcal{A}\}$



Compressed Sensing

Recall that the trigonometric moment sequence is given by

$$x_j^* = \sum_{k=1}^s c_k^* \exp(i2\pi u_k^* j)$$

- ▶ Candes and Fernandes-Granda (2012): Explicit dual construction for exact recovery from x_0, \dots, x_{n-1} if frequencies are at least $4/n$ apart.
- ▶ *Sampling Model*: With a probability $p = m/n$, observe each entry independently, so we observe m entries on an average out of n . Let $T \subset \{0, \dots, n-1\}$ index the observed entries.
- ▶ *Uniform Phases*: Assume that the phases of c_k are uniformly distributed in $[0, 2\pi]$.
- ▶ Interpolate the missing entries by

$$\begin{aligned} & \text{minimize}_x \|x\|_{\mathcal{A}} \\ & \text{subject to } x_j = x_j^*, \text{ for } j \in T. \end{aligned}$$

Dual Certificate for missing data

Definition (Dual Certificate)

A dual certificate q must satisfy

- ▶ $\langle q, a(f_j) \rangle = \text{sign}(c_j)$, for $j = 1, \dots, k$.
- ▶ $\langle q, a(f) \rangle < 1$ if $f \notin \{f_1, \dots, f_k\}$.
- ▶ $q_{T^c} = 0$.

- ▶ Define a random kernel

$$K(f) = \sum_{j=-(n-1)}^{n-1} u_j \delta_j \exp(i2\pi j f)$$

where the Bernoulli variable δ_j only sums observed entries in place of the deterministic kernel

$$\bar{K}(f) = \sum_{j=-(n-1)}^{n-1} u_j \exp(i2\pi j f)$$

- ▶ Proof relies on Matrix concentration, polynomial Bernstein inequality and proof techniques in Romberg et al.