Discussion Worksheet 4

Recall that to build the LR(1) parsing DFA:

- Every state is a closed set of LR(1) parsing items, formed with the Closure operation
- If S is the start production, The start state is the closure of $[S' \to \bullet S, \$]$, where S' is some dummy nonterminal if S has multiple production rules.
- If a state "State" contains the item $[X \to a \bullet yb, b]$, we add a transition labeled y to a state that contains the items Transition(State, y)

The algorithms are reproduced here for your reference, but you should learn these.

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\begin{array}{c} \textbf{procedure Closure}(\textbf{Items}) \\ \textbf{repeat} \\ \textbf{for each } [X \rightarrow a \bullet Yb, a] \text{ in Items do} \\ \textbf{for each production } Y \rightarrow g \textbf{ do} \\ \textbf{for each } b \in \textbf{First}(ba) \textbf{ do} \\ \textbf{add } [Y \rightarrow \bullet g, b] \text{ to Items} \\ \textbf{until Items is unchanged} \\ \textbf{return Items} \end{array} \\ \begin{array}{c} \textbf{procedure Transition}(\textbf{State}, y) \\ \textbf{Items} \leftarrow \emptyset \\ \textbf{for } [X \rightarrow a \bullet yb, b] \in \textbf{State do} \\ \textbf{add } [X \rightarrow ay \bullet b, b] \text{ to Items} \\ \textbf{return Closure}(\textbf{Items}) \end{array}
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1 LR(1) Conflicts

We are looking at strings of lowercase and uppercase letters. We want to separate the string into individual sequences of lowercase and uppercase letters. Here is a basic CFG to accomplish this task. lower and upper are tokens for lowercase and uppercase letters individually.

$$\begin{split} S &\to \varepsilon \mid S \; L \mid S \; U \\ L &\to \varepsilon \mid L \; \text{lower} \\ U &\to \varepsilon \mid U \; \text{upper} \end{split}$$

Exercise 1 Recall that a reduce/reduce conflict occurs when a parsing state contains at least two items:

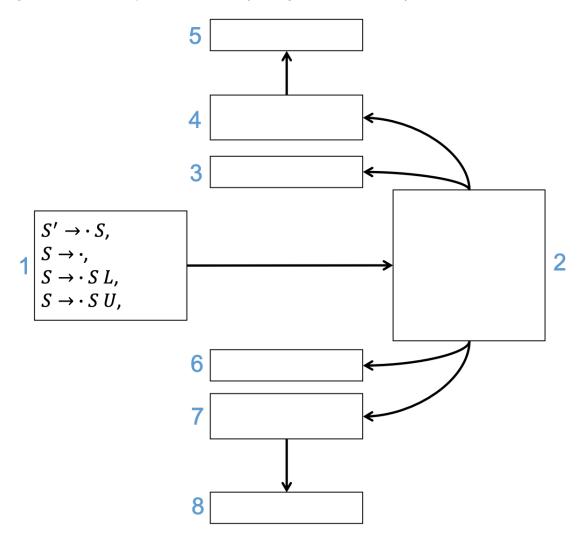
$$[X \to \alpha \bullet, a] \ [Y \to \beta \bullet, a]$$

1.1 There is a reduce/reduce conflict in this grammar. Identify the smallest string with reduce/reduce conflicts, the rules that can be used to parse it, and the ways the parser can reduce the string.

1.2 Generally, in order to remove reduce/reduce conflicts, we need to remove parsing rules or make them stricter. Change two of the parsing rules in the grammar to remove the reduce/reduce conflict.

Exercise 2 Recall that a shift/reduce conflict occurs when a state contains $[X \to \alpha \bullet a\gamma, b], [Y \to \gamma \bullet, a].$

2.1 There is still a shift/reduce conflict in the grammar from 1.2. In fact, there are two. Draw the DFA for the grammar and identify all the conflicts. (The right box has 7 items.)



2	For each	conflict	evnlain v	what	precedence v	blunds av	implement	and why
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2 LR Parsing

Recall that from the DFA we can create a parse table where:

- ullet a state s with an item $[X \to \alpha ullet, b]$ reduces with $X \to \alpha$ on b
- a state with transtion $s \to^b s'$ shifts on b

Exercise 3 Create a parse table from the DFA from 2.1, keeping in mind the precedences we chose in 2.2.

	lower	upper	\$ S	L	U
s1					
s2					
s3					
s4					
s5					
s6					
s7					
s8					

Exercise 4 Is our grammar (with precedence modifications from 2.2) LR(0)? Explain.