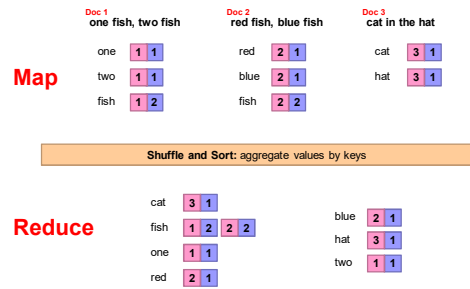


Illustration of the baseline algorithm



7

Inverted Indexing: Pseudo-Code

```

1: class MAPPER
2:   procedure MAP(docid n, doc d)
3:     H ← new ASSOCIATIVEARRAY
4:     for all term t ∈ doc d do
5:       H{t} ← H{t} + 1
6:     for all term t ∈ H do
7:       EMIT(term t, posting ⟨n, H{t}⟩)
1: class REDUCER
2:   procedure REDUCE(term t, postings [(a1, f1), (a2, f2)...])
3:     P ← new LIST
4:     for all posting (a, f) ∈ postings [(a1, f1), (a2, f2)...] do
5:       APPEND(P, (a, f))
6:       SORT(P)
7:       EMIT(term t, postings P)

```

What's the problem?

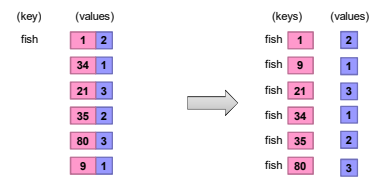
8

Scalability issue of the baseline implementation

- Initial implementation: terms as keys, postings as values
 - Reducers must buffer all postings associated with key (to sort)
 - What if we run out of memory to buffer postings?

9

Another try



- Value-to-key conversion

10

Revised implementation

```

1: class MAPPER
2:   method MAP(docid n, doc d)
3:     H ← new ASSOCIATIVEARRAY
4:     for all term t ∈ doc d do
5:       H{t} ← H{t} + 1
6:     for all term t ∈ H do
7:       EMIT(tuple (t, n), tf H{t})
1: class REDUCER
2:   method INITIALIZE
3:     tprev ← ∅
4:     P ← new POSTINGSLIST
5:   method REDUCE(tuple (t, n), tf [f])
6:     if t ≠ tprev ∧ tprev ≠ ∅ then
7:       EMIT(term t, postings P)
8:       P.RESET()
9:       P.ADD((n, f))
10:    tprev ← t
11:   method CLOSE
12:     EMIT(term t, postings P)

```

11