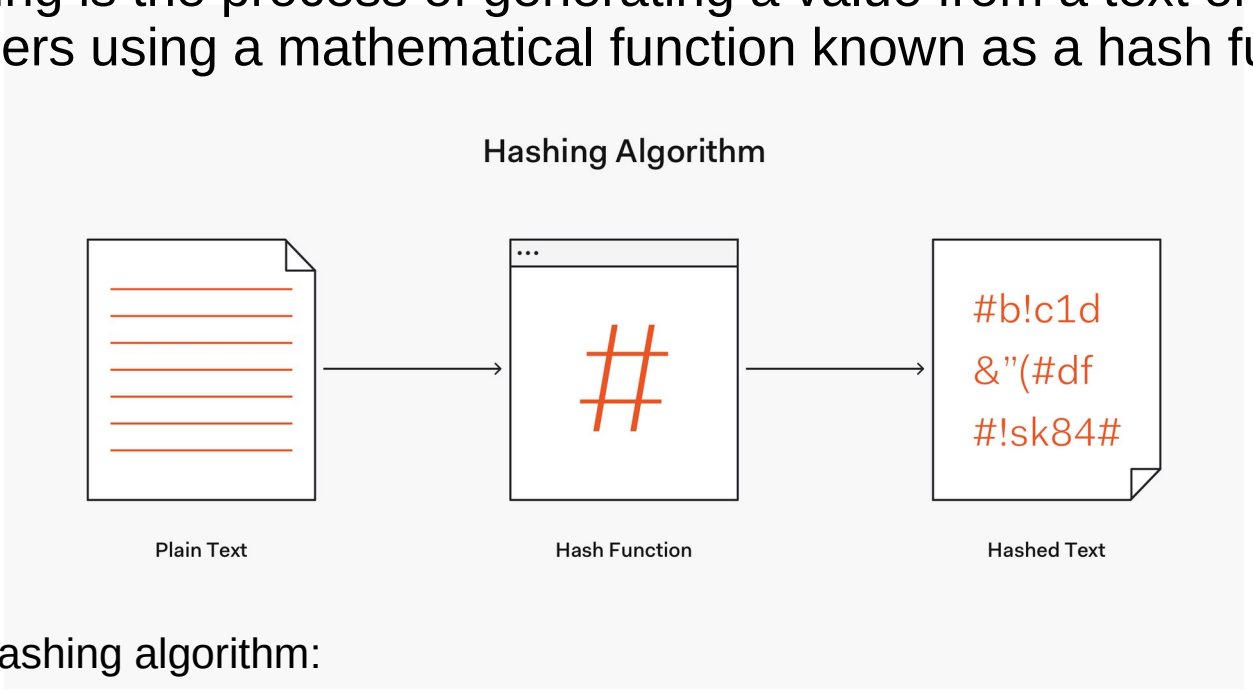
Hashing Algorithms

Hashing is the process of generating a value from a text or a list of numbers using a mathematical function known as a hash function.



A good hashing algorithm:

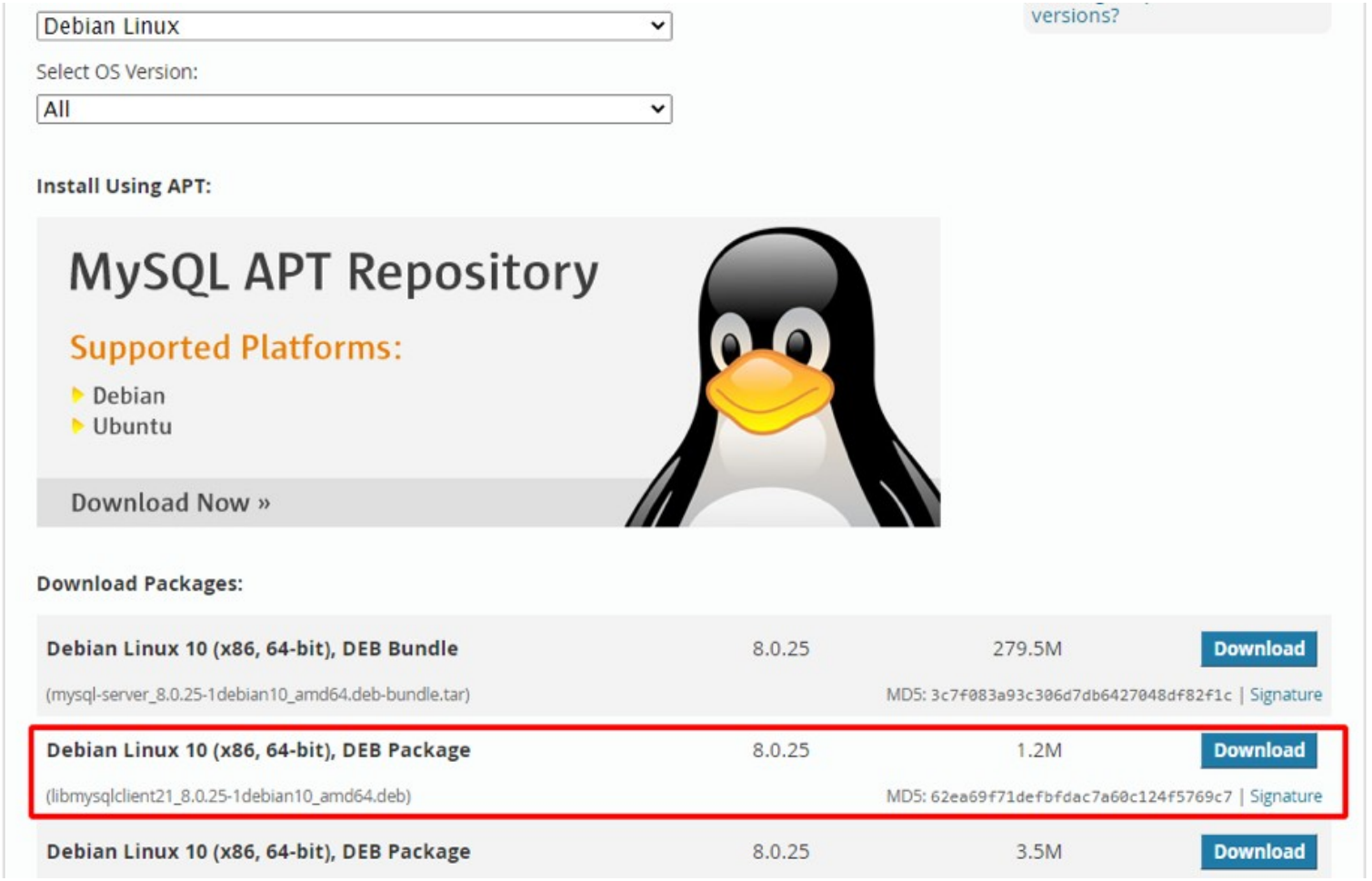
- Should be a one-way algorithm
- Should be efficiently computable
- Should uniformly distribute the data

Some Popular Hashing Algorithms

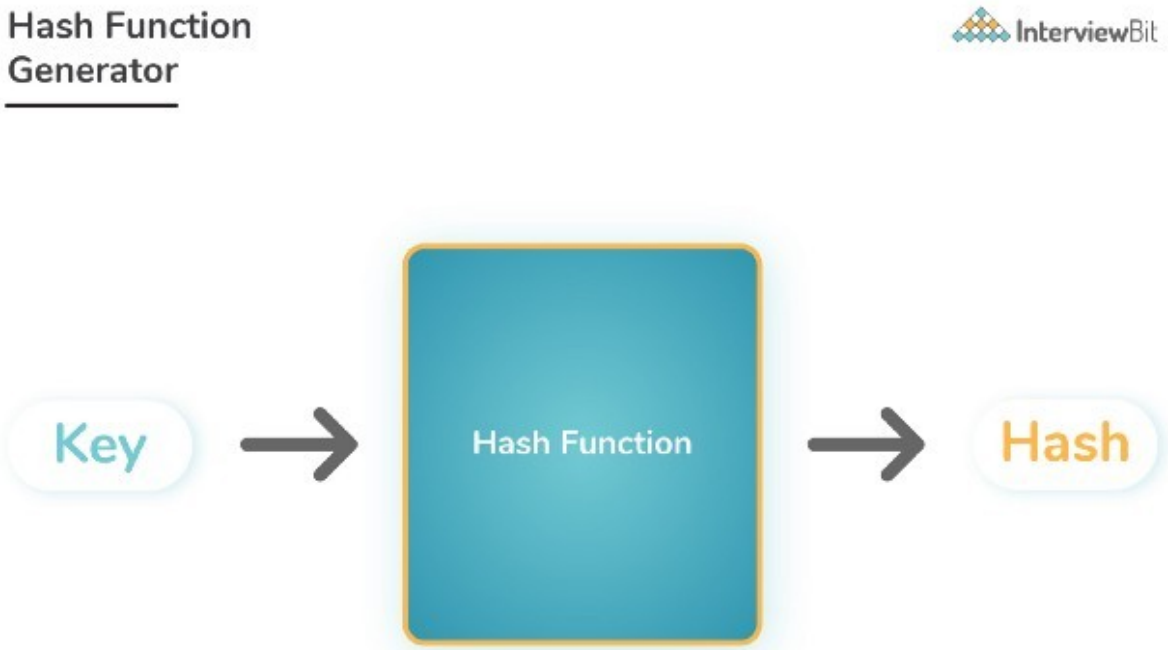
- **MD5** (message digest version 5): Designed in 1991, this hashing algorithm produces a 128-bit hash value. It's still one of the most commonly used despite being one of the most insecure algorithms. (It's susceptible to brute force attacks. Stay away from it!)
- **SHA** (secure hashing algorithm) family:
  - **SHA**-1: This hashing algorithm generates a 160-bit hash value. Vulnerable to brute force attacks, it's no longer considered a secure hashing algorithm. As a result, Microsoft, Google and Mozilla no longer accept SHA-1 SSL certificates (since 2017).
  - **SHA**-256: This hashing algorithm is a variant of the SHA2 hashing algorithm, recommended and approved by the National Institute of Standards and Technology (NIST). It generates a 256-bit hash value. Even if it's 30% slower than the previous algorithms, it's more complicated, thus, it's more secure.
  - **SHA**-384: This hashing algorithm is the latest member of the SHA family, it's much faster than the SHA-256 and it's based on a totally different approach (sponge construction).
- **Whirlpool**: This hashing algorithm is based on the advanced encryption standard (AES) and produces a 512-bit hash digest.

Hashing Algorithms

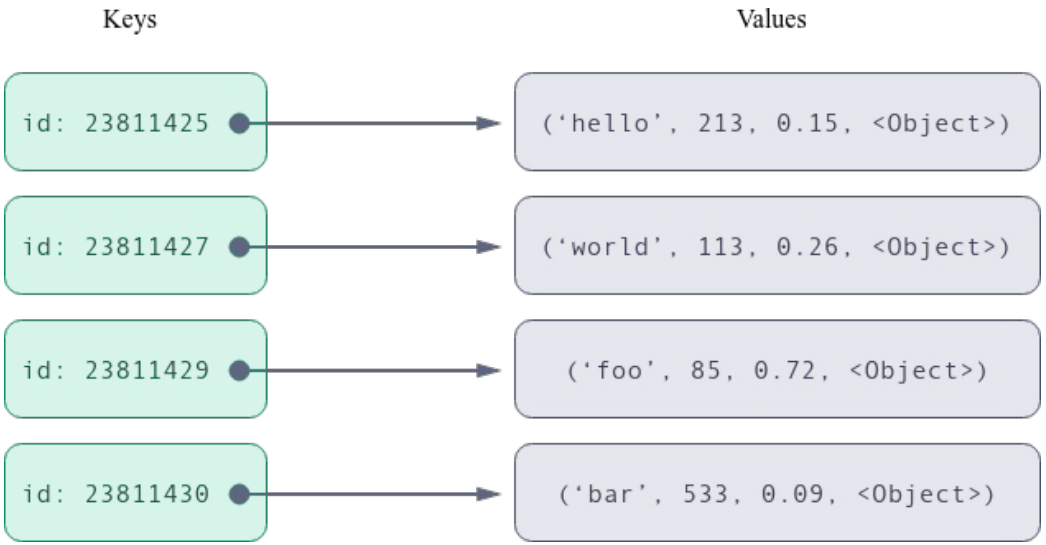
Hashing Algorithms



Key Hashing

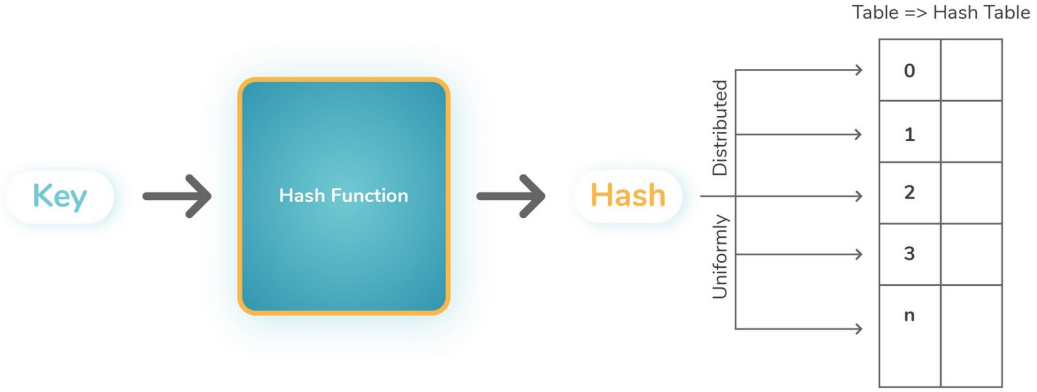

Hash Function
Generator
InterviewBit
Key
Hash Function
Hash

Hash Map/Table



InterviewBit

Hash Function Generator



Some basic hash functions

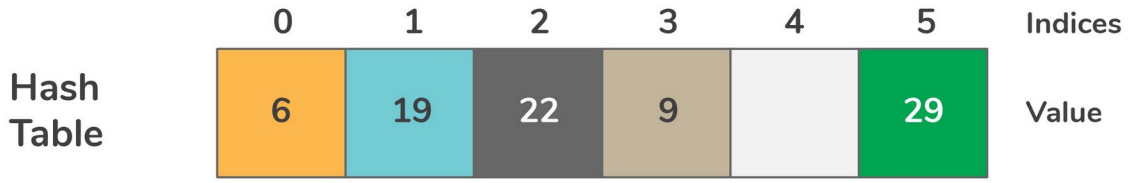
Hash Function:
Division Method

InterviewBit

Table_Length= 6
Values to be hashed = 6, 9, 19, 22, 29

hash(Key) = key % Table_Length

_> hash(6) = 6 % 6 = 0 => Place key **6** in **0th** position
_> hash(9) = 9 % 6 = 3 => Place key **9** in **3rd** position
_> hash(18) = 19 % 6 = 1 => Place key **19** in **1st** position
_> hash(21) = 22 % 6 = 2 => Place key **22** in **2nd** position
_> hash(29) = 29 % 6 = 5 => Place key **29** in **5th** position



Some basic hash functions

Hash Function:
Mid Square Method

InterviewBit

Key = 3101
Table_Size = 2000

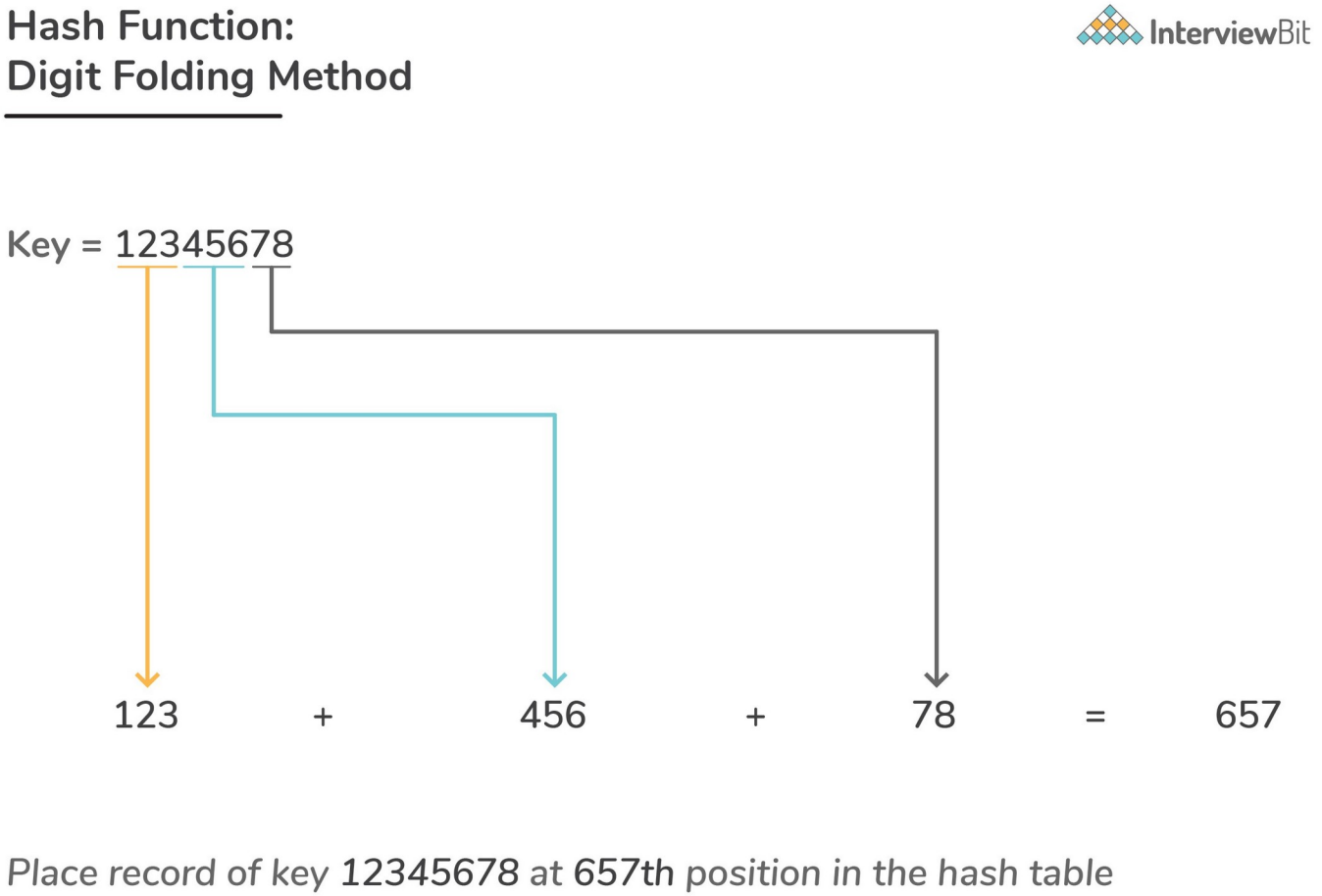
hash(Key) = middle numbers from key * key value

=> hash(3101):
3101 * 3101 = 9 6 1 6 2 0 1

Middle Number = 162

Place record of key 3101 at 162nd position in hash table

Some basic hash functions



Some basic hash functions

Hash Function:
Using Multiplication Method

InterviewBit

Key = 50
Assume c = 0.81, where 0 < c < 1
Assume Table_Size = 1000

hash(Key) = floor(Table_Size * fractional(k * c))

hash(50) = floor(1000 * fractional(50 * 0.81))

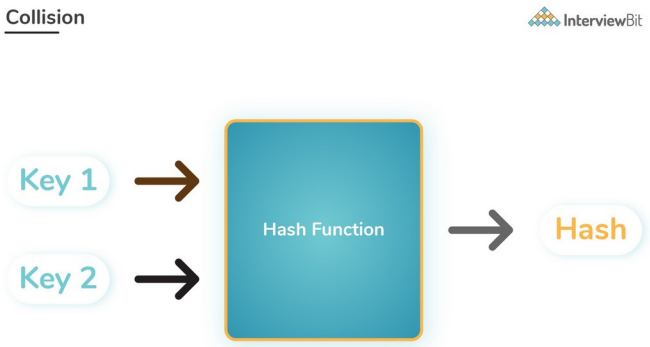
↓

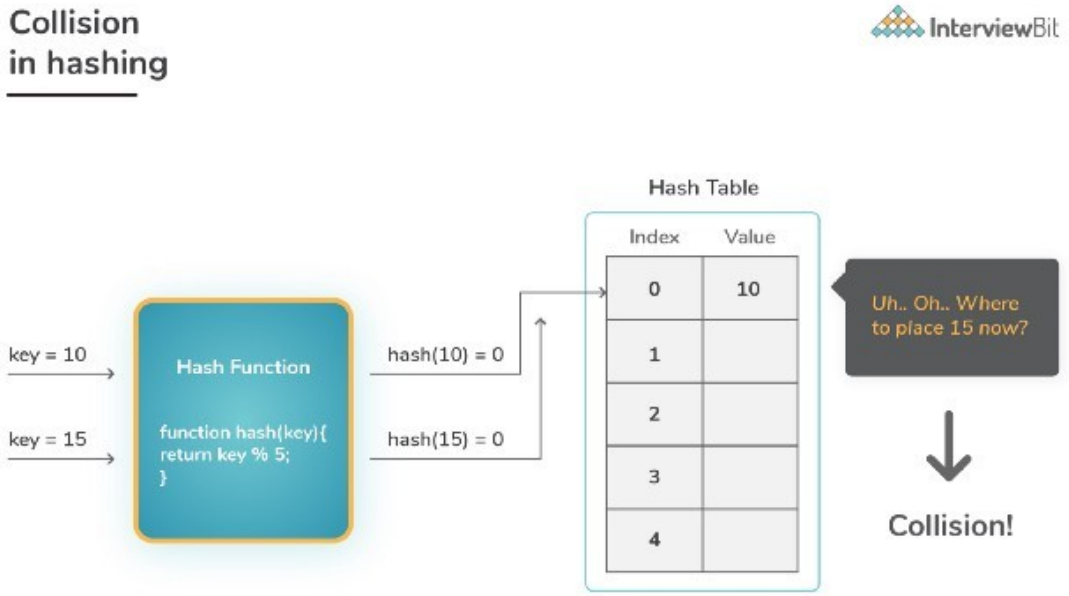
50 * 0.81 = 40.5 => Fractional Part = x - floor(x)
= 40.5 - floor(40.5)
= 40.5 - 40
= 0.5

hash(50) = floor(1000 * 0.5) = floor(500) = 500

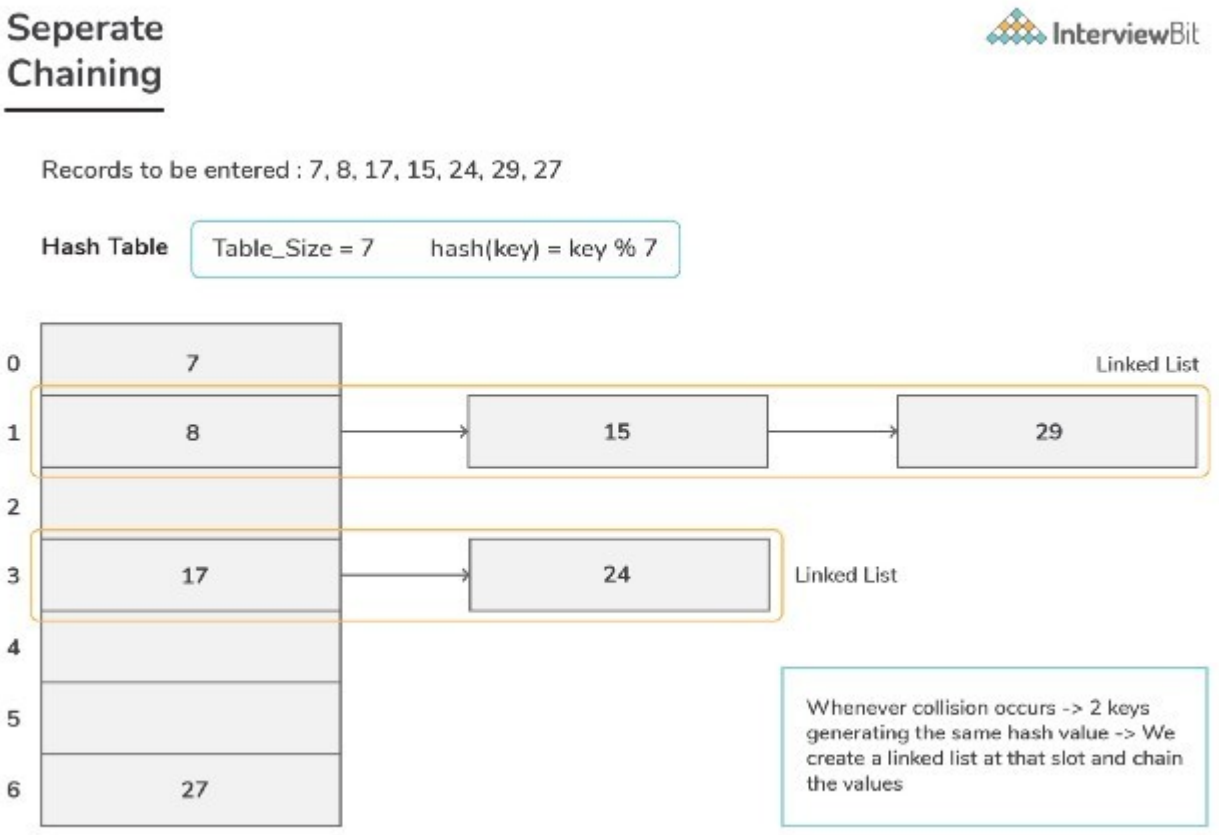
Place record of key 50 at 500th position in hash table

Collision


Collision
InterviewBit
Key 1
Key 2
Hash Function
Hash


Collision
in hashing
InterviewBit
Hash Table
Index
Value
0
10
1
2
3
4
key = 10
key = 15
Hash Function
function hash(key){
return key % 5;
}
hash(10) = 0
hash(15) = 0
Uh.. Oh.. Where
to place 15 now?
Collision!

Collision


Seperate Chaining
InterviewBit
Records to be entered : 7, 8, 17, 15, 24, 29, 27
Hash Table
Table_Size = 7
hash(key) = key % 7
0
1
2
3
4
5
6
7
8
17
27
15
29
24
Linked List
Linked List
Whenever collision occurs -> 2 keys generating the same hash value -> We create a linked list at that slot and chain the values

Collision

Open Addressing : Linear Probing

InterviewBit

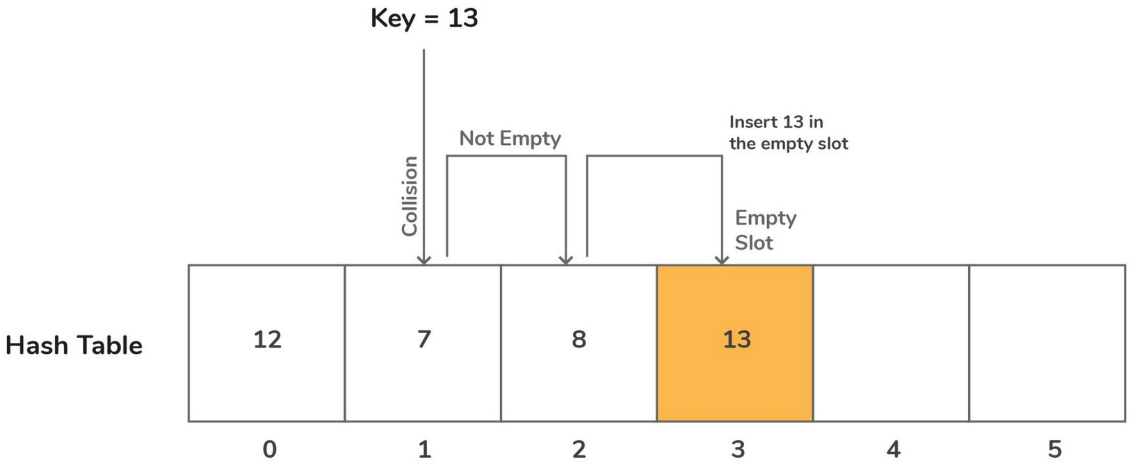
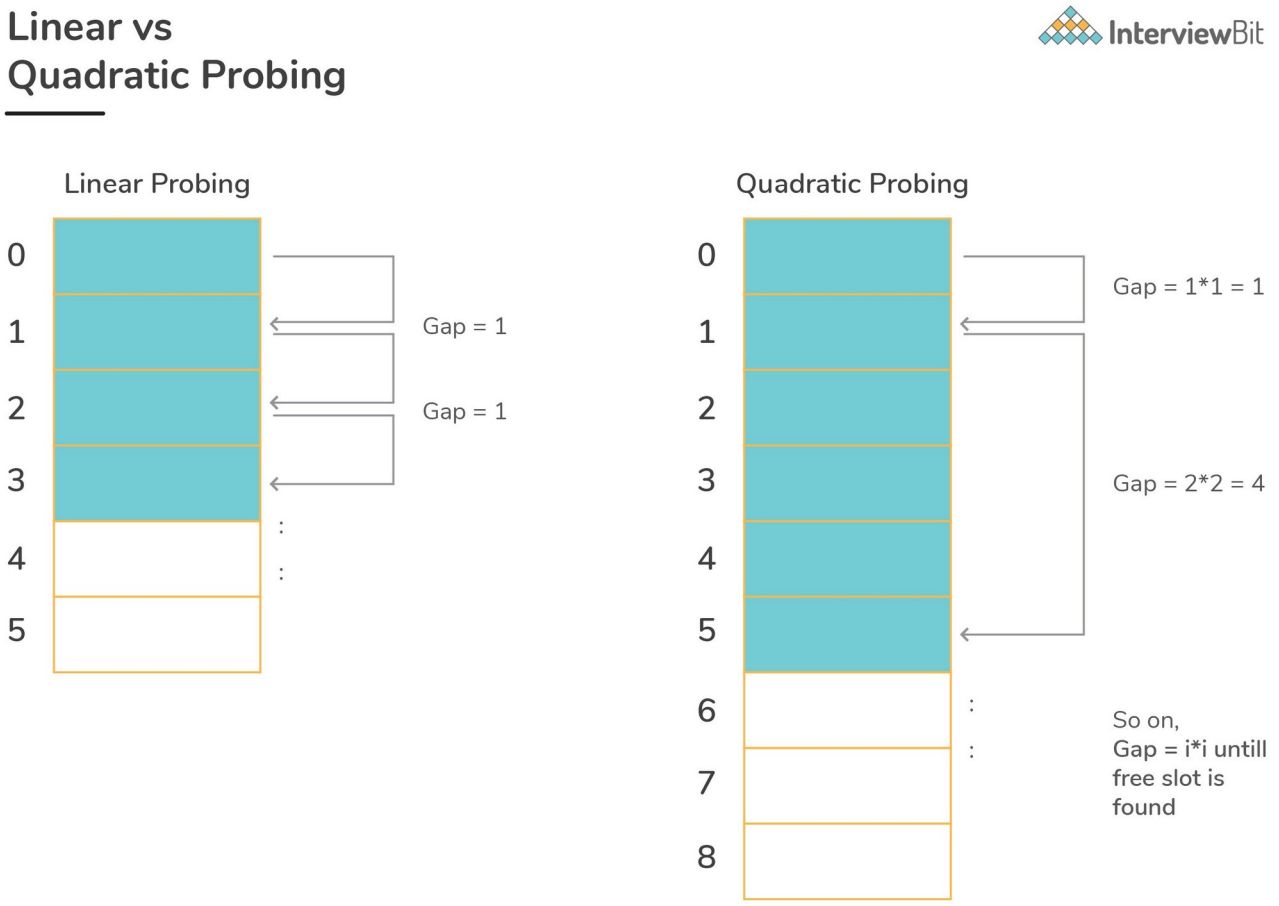


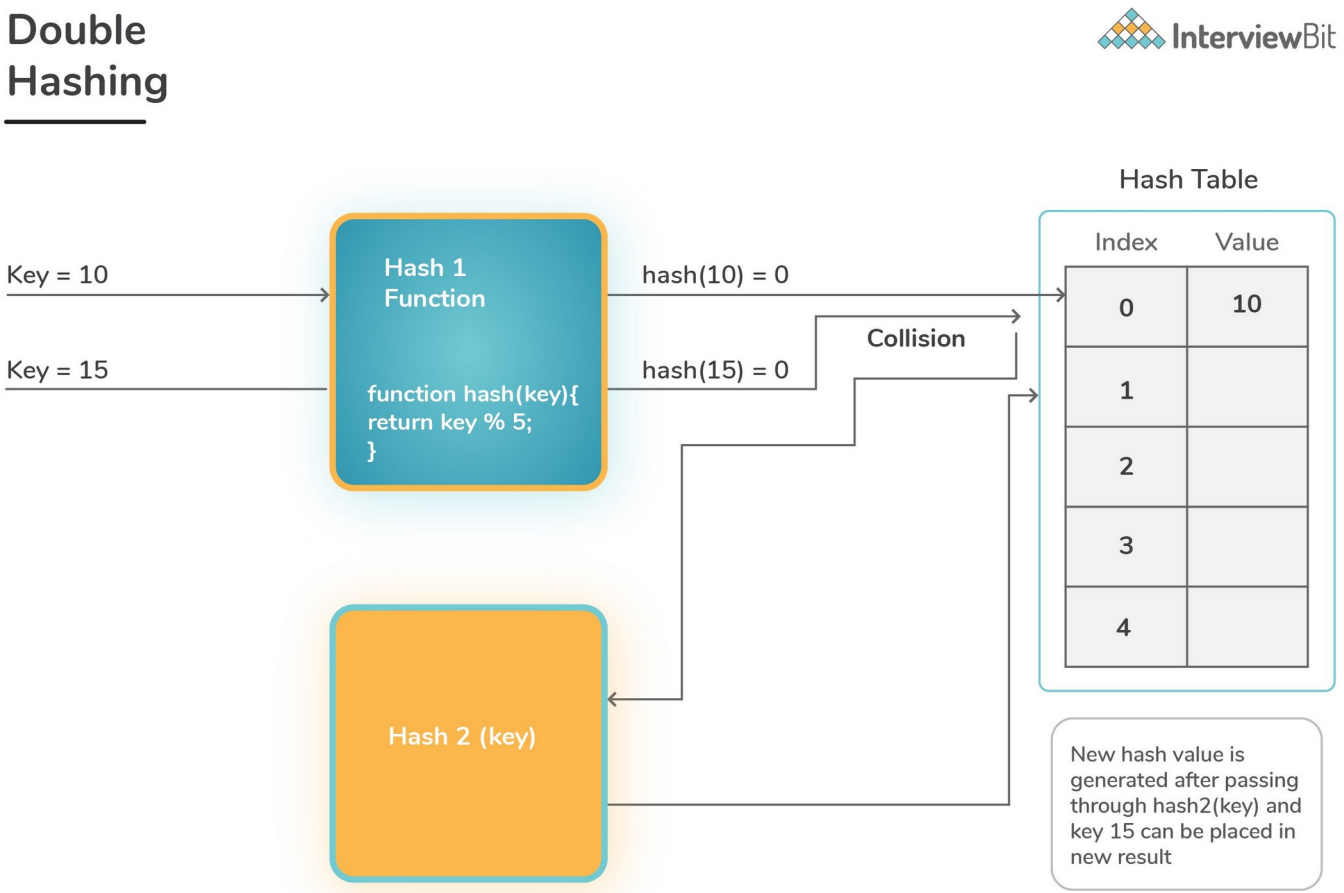
Table Size = 6

hash(key) = key % Table_Size

Collision


Linear vs
Quadratic Probing
InterviewBit
Linear Probing
0
1
2
3
4
5
Gap = 1
Gap = 1
Quadratic Probing
0
1
2
3
4
5
6
7
8
Gap = 1*1 = 1
Gap = 2*2 = 4
So on,
Gap = i*i untill
free slot is
found

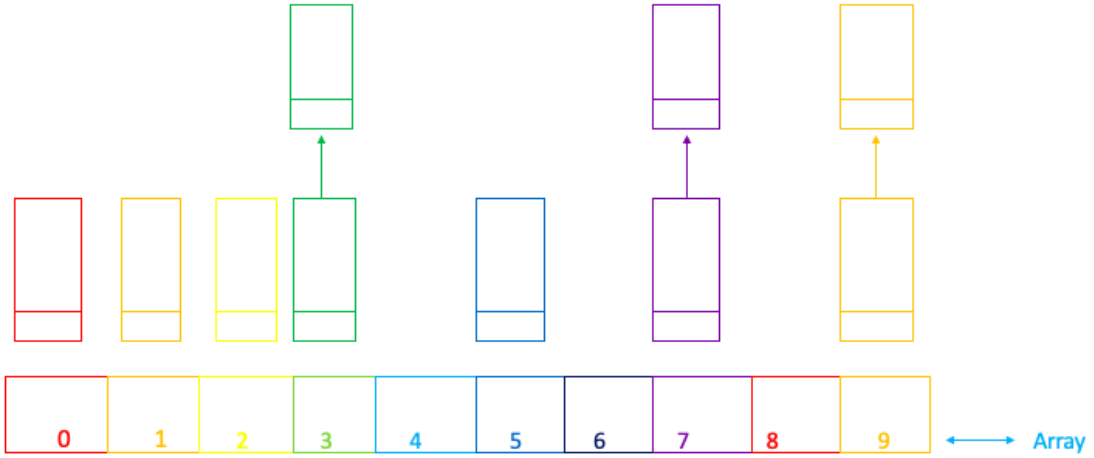
Collision

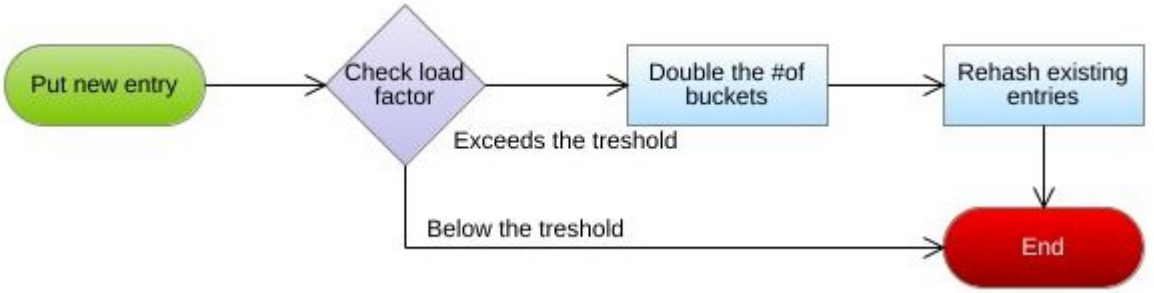

Double Hashing
InterviewBit
Hash Table
Index
Value
Key = 10
Key = 15
Hash 1 Function
function hash(key){
return key % 5;
}
hash(10) = 0
hash(15) = 0
Collision
0
10
1
2
3
4
Hash 2 (key)
New hash value is generated after passing through hash2(key) and key 15 can be placed in new result

Collision

Hash table properties:

- ***#of buckets(indexes)***
- **Load factor = *#of entries* / *#of buckets***





Java hash table implementations

```
Map<String, String> ht = new Hashtable<String, String>();
ht.put("ahmet", "312 1233212");
ht.put("mehmet", "212 1233212");
ht.put("ali", "412 1233212");
ht.put("veli", "512 1233212");
System.out.println(ht.get("ali"));
```

→ Thread safe

Allows nulls ←

```
Map<String, String> ht = new HashMap<String, String>();
ht.put("ahmet", "312 1233212");
ht.put("mehmet", "212 1233212");
ht.put("ali", "412 1233212");
ht.put("veli", "512 1233212");
System.out.println(ht.get("ali"));
```