

Discrete Optimization

1. Greedy Algorithms & MATROIDS

1.1. Basics

Def.: An Independent Set System is a pair
 (E, \mathcal{F}) , $\mathcal{F} \subseteq \mathcal{P}(E)$
where E is a finite set.

$$M1) \emptyset \in \mathcal{F}$$

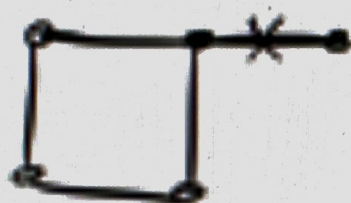
$$M2) \text{ If } X \subseteq Y, Y \in \mathcal{F} \\ \text{then } X \in \mathcal{F}.$$

$X \in \mathcal{F}$ is called independent
 $X \in \mathcal{P}(E) \setminus \mathcal{F}$ is called dependent

BASIS:

A maximally (inclusion-wise) independent set is called a BASIS.

CIRCLE: inclusion-wise minimal dependent sets are called CIRCLES.



← not a circle

Examples

B1) $E = \mathbb{R}^n$ or $E \subseteq \mathbb{R}^n$
 $E = K^n$, finite field K

$$\mathcal{F} = \{X \subseteq \mathbb{R}^n, X \text{ endl\u00fcz} :$$

$$\exists Y (\in \mathbb{R}^{|X|}), Y \neq 0 :$$

$$\left\{ \sum_{x \in X} \gamma_x x = 0 \right\}$$

↑
lin. independence

"representable over the field
 $K (= \mathbb{R})$."

$$B_2) \quad E = \{w_i\}_{[n]} \subseteq \mathbb{N}_0$$

↑ $:= \{1, \dots, n\}$

$$B \in \mathbb{N}$$

$$\mathcal{F} = \{x \subseteq E \mid \sum_{w_i \in x} w_i \leq B\}$$

"Knapsack"

$$B_3) \quad \bar{E}, \quad |\bar{E}| = n$$

$$m \in \mathbb{N}$$

$$\mathcal{F} = \{X \subseteq E \mid |X| \leq m\}$$

"Uniform Matroid"

B4) $E = E(G)$

$$\mathcal{F} = \{X \subseteq E \mid X \text{ is cycle-free}\}$$

"Graphic Matroid"

Def.: Rank & Closure
of an ISS

let (E, \mathcal{F}) be an ISS.

For $X \subseteq E$ we call

$$r(X) := \max\{|Y| : Y \subseteq X, Y \in \mathcal{F}\}$$

the RANK of X .

$$v \left(\begin{array}{c} \text{Diagram of a bipartite graph with 6 nodes and 9 edges. The left side has nodes 1, 2, 3 and the right side has nodes 4, 5, 6. Edges are (1,4), (1,5), (1,6), (2,4), (2,5), (2,6), (3,4), (3,5), (3,6). The edges (1,4), (2,5), and (3,6) are highlighted in red.} \end{array} \right) = 3$$

$$|Y| = 3$$

Def.: Closure of $X \subseteq E$

$$d(X) := \{y \in E : r(X \cup \{y\}) = v(X)\}$$

1.2. Optimization over ISS

Problem 1:

Input (E, \mathcal{F}) & $c: E \rightarrow \mathbb{R}^+$

Out : $X \in \mathcal{F}$:

$c(x) = \sum_{e \in x} c(e)$ is maximal

in \mathcal{F} .

Problem 2:

Input: (E, \mathcal{F}) & $c: E \rightarrow \mathbb{R}$

Out: Basis B with $c(B)$ minimal.

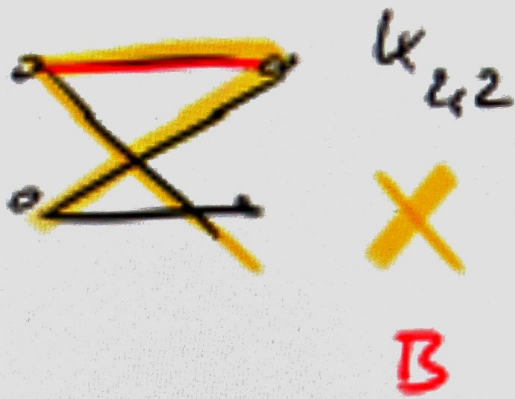
Remark:

There is no term BASIS for
an $(SS(E, \mathcal{F}))$.

For $X \subseteq E$ a BASIS B is

a set $B \subseteq X$ if $B \in \mathcal{F}$

and cannot be extended in X .
independently



1.3. Matroid

Def.: Matroid

Let (E, \mathcal{F}) be an ISS with

M3) IF: $\forall X, Y \in \mathcal{F}$ & $|X| > |Y|$

THEN: $\exists x \in X \setminus Y$ such that

$$Y \cup \{x\} \in \mathcal{F}.$$

(E, \mathcal{F}) fulfills M1-M3 is called a Matroid.

Examples:

E set of vectors in K^n