



Trinity College Dublin
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CSU22012: Data Structures and Algorithms II

Extra Lecture 1: Recursion

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Complexity

```
function argmax(array):
```

Input: an array of size n

Output: the index of the maximum value

```
index  $\leftarrow$  0 #assignment, 1 op
```

```
foreach i in [1, n-1] do #1 op per Loop
```

```
    if array[i] > array[index] then #3 ops per Loop
```

```
        index  $\leftarrow$  i #1 op per Loop, sometimes
```

```
    endif
```

```
endfor
```

```
return index #1 op
```

How many operations if the list has 10 elements? 10,000 elements?

- *Varies proportional to the size of the input list:* $5n + 2$
- We'll be in the for loop longer and longer as the input list grows

Complexity

```
function argmax(array):  
  Input: an array  
  Output: the index of the maximum value  
  index ← 0 #assignment, 1 op  
  foreach i in [1,array.length[ do #1 op per loop  
    if array[i] > array[index] then #3 ops per loop  
      index ← i #1 op per loop, sometimes  
    endif  
  endfor  
  return index #1 op
```

$$T(n) = 5n + 2$$

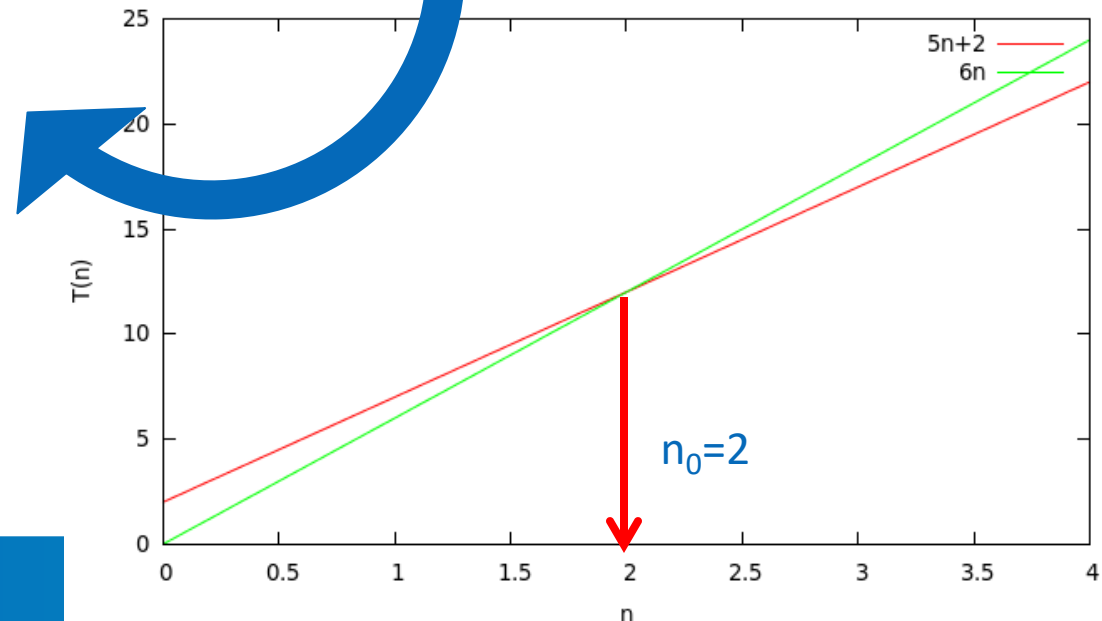
$$T(n) \leq 6n$$

$$g(n) = n$$

$$c = 6$$

$$n_0 = 2$$

$T(n)$ is $O(n)$



Outline

Recursion?

Base Case

Call Stack

Recursive vs. Iterative Functions

Tail Recursion

Turning a recursive algorithm into an iterative one

Complexity of recursive functions

Take home message:

Recursion is a method to divide a problem in similar sub-problems

— *Allows a function to call itself*

Recursion?

Recursion is a way of decomposing problems into smaller, simpler sub-tasks that are similar to the original.

Thus, each sub-task can be solved by applying a similar technique.

The whole problem is solved by combining the solutions to the smaller problems.

Requires a ***base case*** (a case simple enough to solve without recursion) to ***end recursion***.



Recursion Example

Factorial: $n! = 1 \times 2 \times \dots \times (n-1) \times n$

Or: **$n! = n \times (n-1)!, 1!=1$**

```
function factorial(n)
```

```
Input: n a natural number
```

```
Output: the n-th factorial number
```

```
if n = 1 then
```

```
    return 1
```


```
else
```

```
    return n * factorial(n-1)
```


```
endif
```

Recursion Simulation

```
factorial(3)
  if 3=1 then
    return 1
  else
    return 3 * factorial(2)
  endif
```

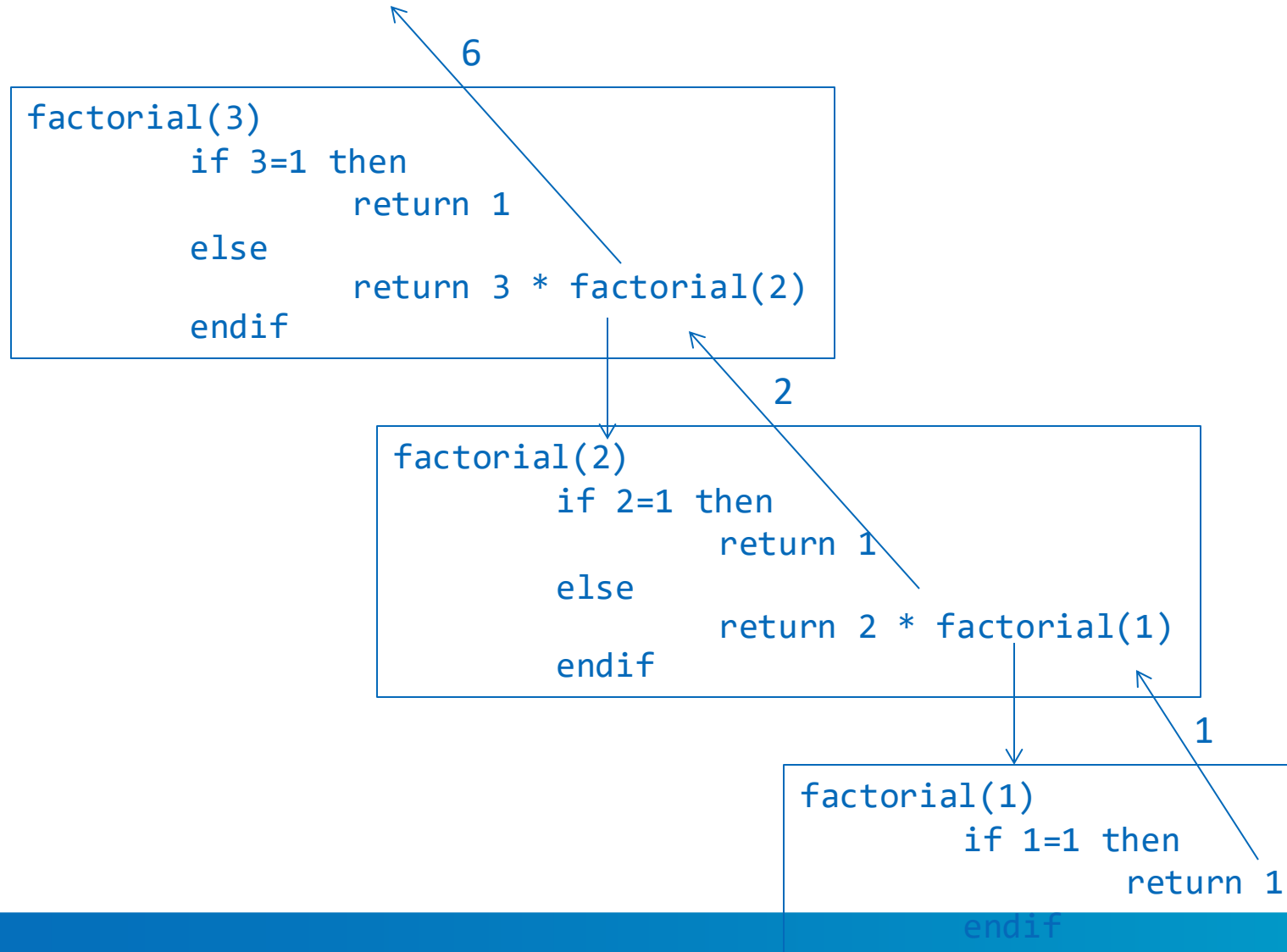


```
factorial(2)
  if 2=1 then
    return 1
  else
    return 2 * factorial(1)
  endif
```



```
factorial(1)
  if 1=1 then
    return 1
  endif
```

Recursion Simulation



Stopping Case/Base Case

As a recursive function calls itself, it is ***crucial to have a base case/stopping case*** – or the process will never stop!

The basic principle is to ***test first the stopping condition*** and then raise the recursive call if the condition is not met

```
function factorial(n)
```

Input: n a natural number

Output: the n-th factorial number

```
if n = 1 then  
    return 1
```

```
else
```

```
    return n * factorial(n-1)
```

```
endif
```

This is bad! Factorial without a Base Case

```
function factorial(n)
```

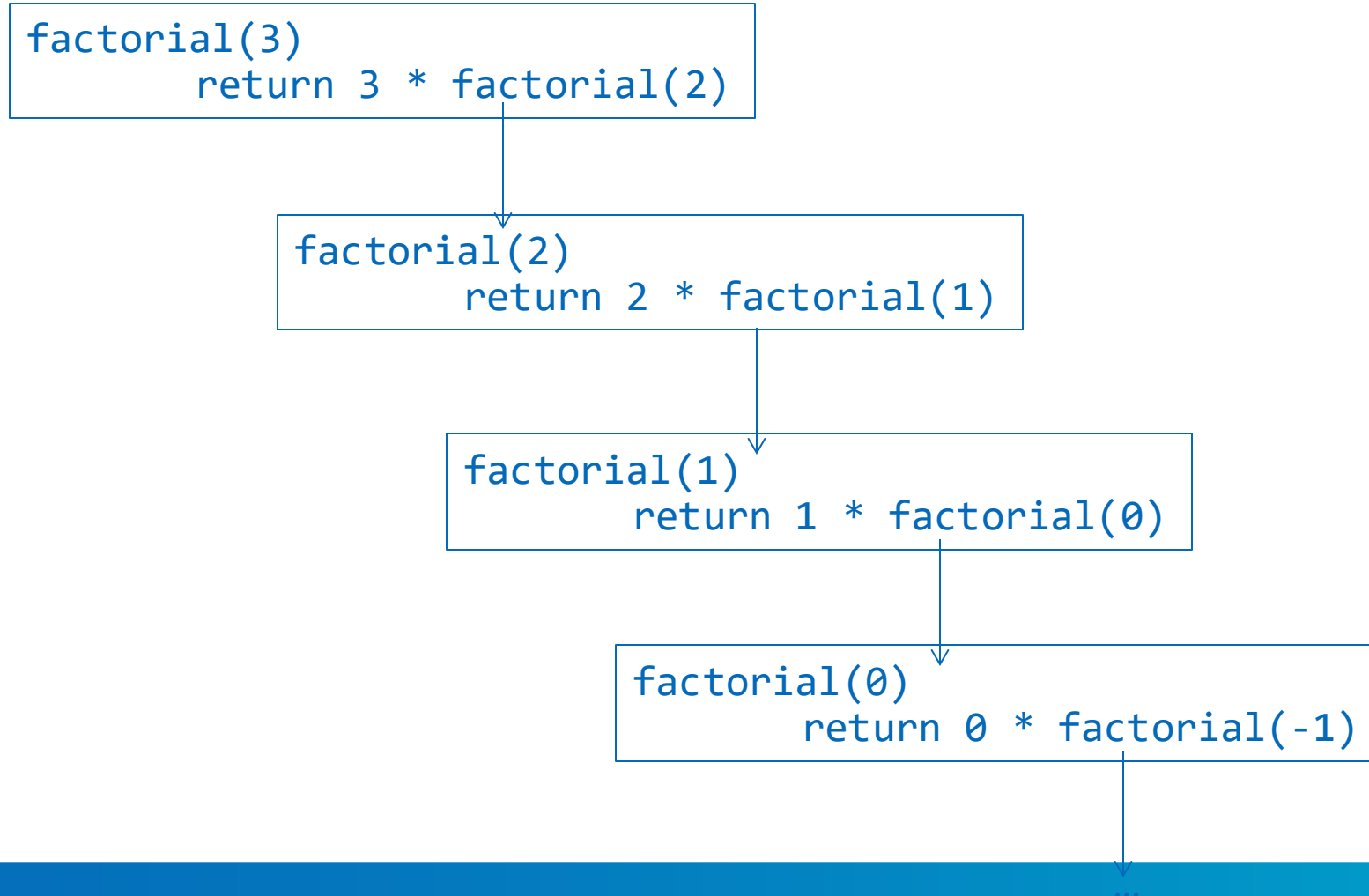
Input: n a natural number

Output: the n-th factorial number

```
    return n * factorial(n-1)
```

```
endif
```

This is bad! Recursion Simulation: Factorial without a Base Case



Call Stack

The basic idea behind recursion is that every call has a unique context (own memory address, own values for parameters and variables)

The ***call stack*** contains all this information and in the context of recursive function, this keeps track of the recursive calls

Recursion Simulation

Calculate 3 factorial

This is a call to factorial(3), so we put factorial(3) on the call stack

```
function factorial(n)
Input: n a natural number
Output: the n-th factorial number
if n = 1 then
    return 1
else
    return n * factorial(n-1)
endif
```



Call Stack

Recursion Simulation

Calculate 3 factorial

This is a call to factorial(3), so
we put factorial(3) on the
call stack

```
function factorial(n)
Input: n a natural number
Output: the n-th factorial number
if n = 1 then
    return 1
else
    return n * factorial(n-1)
endif
```



Call Stack

Recursion Simulation

$n \neq 1$, so we return
 $n * \text{factorial}(n-1)$, which includes a call to $\text{factorial}(2)$.
Remember, the call to $\text{factorial}(3)$
has not returned yet, so it is still
on the call stack!

```
function factorial(n)
Input: n a natural number
Output: the n-th factorial number
if n = 1 then
    return 1
else
    return n * factorial(n-1)
endif
```



Call Stack

Recursion Simulation

$n \neq 1$, so we return
 $n * \text{factorial}(n-1)$, which includes a call to
 $\text{factorial}(2)$.

Remember, the call to $\text{factorial}(3)$
has not returned yet, so it is still

on the call stack!

```
function factorial(n)
  Input: n a natural number
  Output: the n-th factorial number
  if n = 1 then
    return 1
  else
    return n * factorial(n-1)
  endif
```



Call Stack

Recursion Simulation

n still $\neq 1$, so we return
 $n * \text{factorial}(n-1)$, which includes a call to
 $\text{factorial}(1)$.

Neither $\text{factorial}(2)$ nor $\text{factorial}(3)$
has returned at this point!

```
function factorial(n)
Input: n a natural number
Output: the n-th factorial number
if n = 1 then
    return 1
else
    return n * factorial(n-1)
endif
```



Call Stack

Recursion Simulation

n still $\neq 1$, so we return
 $n * \text{factorial}(n-1)$, which includes a call to
 $\text{factorial}(1)$.

Neither $\text{factorial}(2)$ nor $\text{factorial}(3)$ has returned
at this point!

```
function factorial(n)
Input: n a natural number
Output: the n-th factorial number
if n = 1 then
    return 1
else
    return n * factorial(n-1)
endif
```



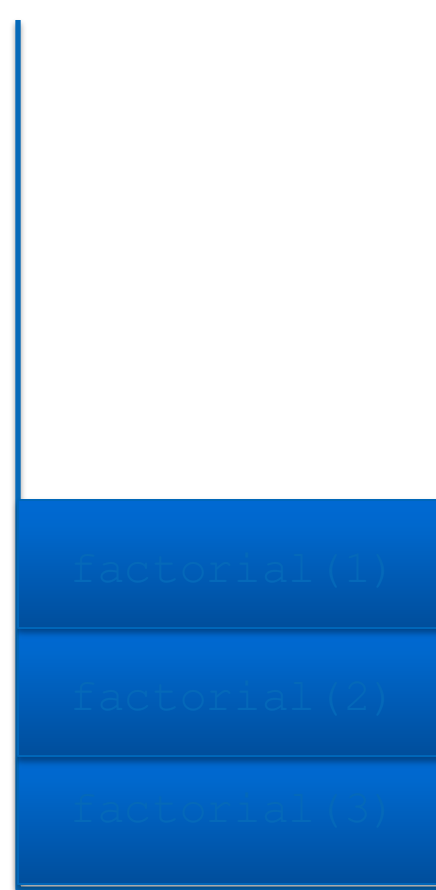
Call Stack

Recursion Simulation

Now $n = 1$, so we return 1!

This is not a recursive call, so factorial(1) returns, and we take it off of the call stack.

```
function factorial(n)
Input: n a natural number
Output: the n-th factorial number
if n = 1 then
    return 1
else
    return n * factorial(n-1)
endif
```



Call Stack

Recursion Simulation

Now factorial(2) is at the top of the call stack, so we return to where we were in factorial(2).

So we return $2 * \text{factorial}(1)$, which we now know is $2 * 1$, so factorial(2) returns 2 and is removed from the call stack!

```
function factorial(n)
Input: n a natural number
Output: the n-th factorial number
if n = 1 then
    return 1
else
    return n * factorial(n-1)
endif
```



Call Stack

Recursion Simulation

Now factorial(2) is at the top of the call stack, so we return to where we were in factorial(2).

So we return $2 * \text{factorial}(1)$, which we now know is $2 * 1$, so factorial(2) returns 2 and is removed from the call stack!

```
function factorial(n)
Input: n a natural number
Output: the n-th factorial number
if n = 1 then
    return 1
else
    return n * factorial(n-1)
endif
```



Call Stack

Recursion Simulation

Now factorial(3) is at the top of the call stack, so we're back in factorial(3). Return $3 * \text{factorial}(2)$, which we now know is $3 * 2$.

Factorial(3) returns 6 and removes itself from the call stack.

```
function factorial(n)
  Input: n a natural number
  Output: the n-th factorial number
  if n = 1 then
    return 1
  else
    return n * factorial(n-1)
  endif
```



Call Stack

Recursion Simulation

Now factorial(3) is at the top of the call stack, so we're back in factorial(3). Return $3 * \text{factorial}(2)$, which we now know is $3 * 2$.

Factorial(3) returns 6 and removes itself from the call stack.

```
function factorial(n)
  Input: n a natural number
  Output: the n-th factorial number
  if n = 1 then
    return 1
  else
    return n * factorial(n-1)
  endif
```



Call Stack

Recursion Simulation

Now our call stack is empty,
and we know that factorial(3)
is 6!

```
function factorial(n)
Input: n a natural number
Output: the n-th factorial number
if n = 1 then
    return 1
else
    return n * factorial(n-1)
endif
```



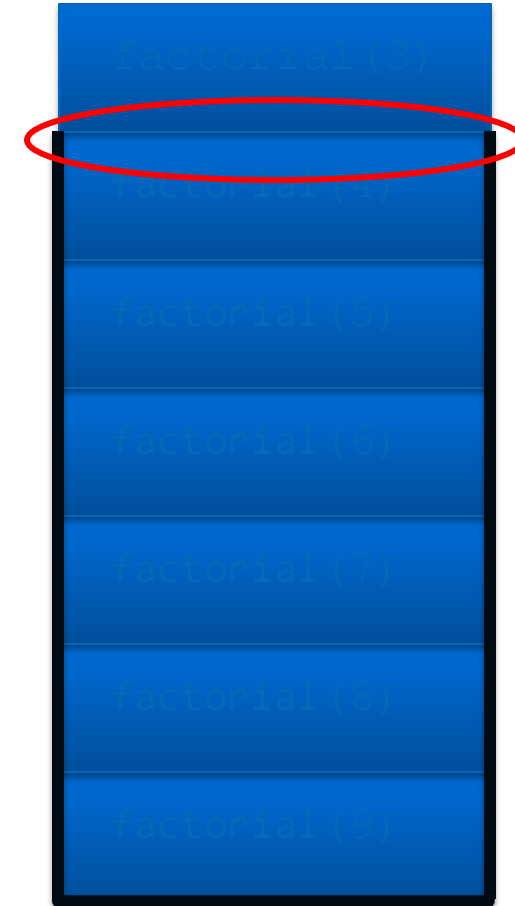
Call Stack

Call Stack

It is difficult to predict the number of calls – and the system needs to do dynamic allocation

Which can be a problem: the famous ***stack overflow*** problem being always around the corner in case there are too many calls)

Stack overflow happens when the call stack reaches the stack bound



Call Stack

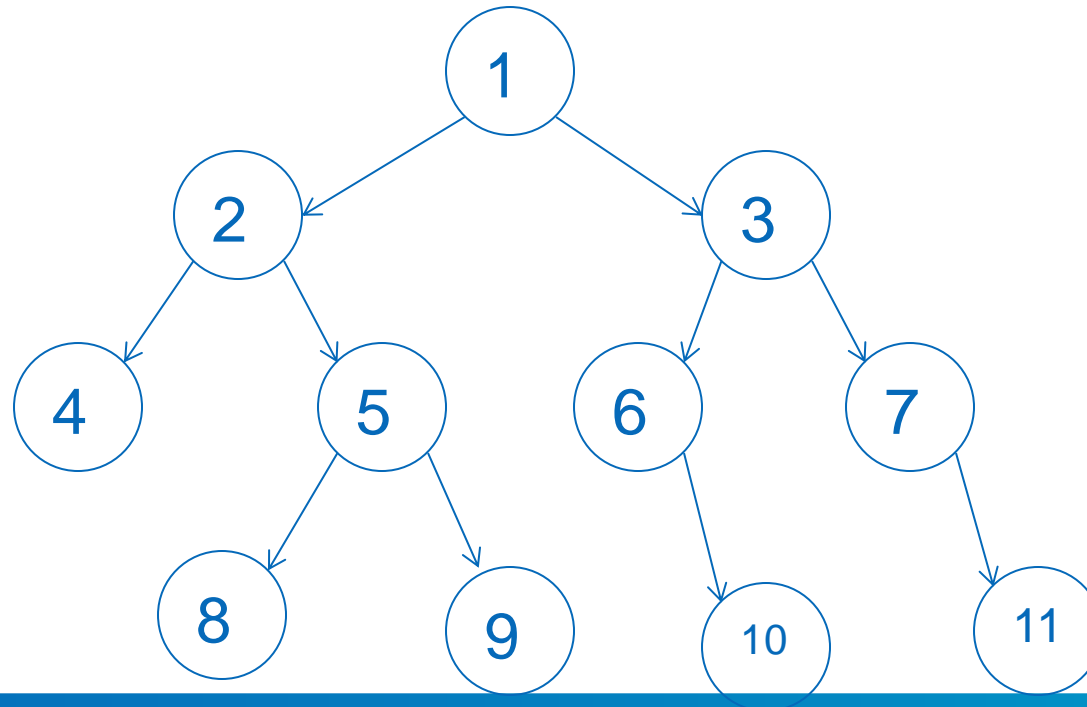
Recursive vs. Iterative Functions

It is often possible to write the same algorithm using recursive or iterative functions

```
function factorial_iterative(n)
Input: n a natural number
Output: the n-th factorial number
fact <- n
while n > 1 do
    n <- n-1
    fact <- fact * n
endwhile
return fact
```

Recursive vs. Iterative Functions

Some data structures are naturally recursive, i.e., it's much easier to write recursive functions for them than iterative ones



Recursive vs. Iterative Functions

Recursive functions are sometimes slower (calls are expensive in practice)
(Bad) recursive algorithms can generate a large number of calls

Tail Recursion

A function call is said to be ***tail recursive*** if there is nothing to do after the function returns except return its value.

A function is non tail recursive if there is some processing done after the function returns.

Example of non Tail Recursion

Factorial: $n! = 1 \times 2 \times \dots \times (n-1) \times n$

Or: **$n! = n \times (n-1)!, 1!=1$**

```
function factorial_non_tail(n)
```

Input: n a natural number

Output: the n-th factorial number

```
if n == 1 then
```

```
    return 1
```


```
else
```

```
    return n * factorial_non_tail(n-1)
```


```
endif
```

Recursion Simulation

```
factorial_non_tail(3)
  if 3=1 then
    return 1
  else
    return 3 * factorial_non_tail (2)
  endif
```

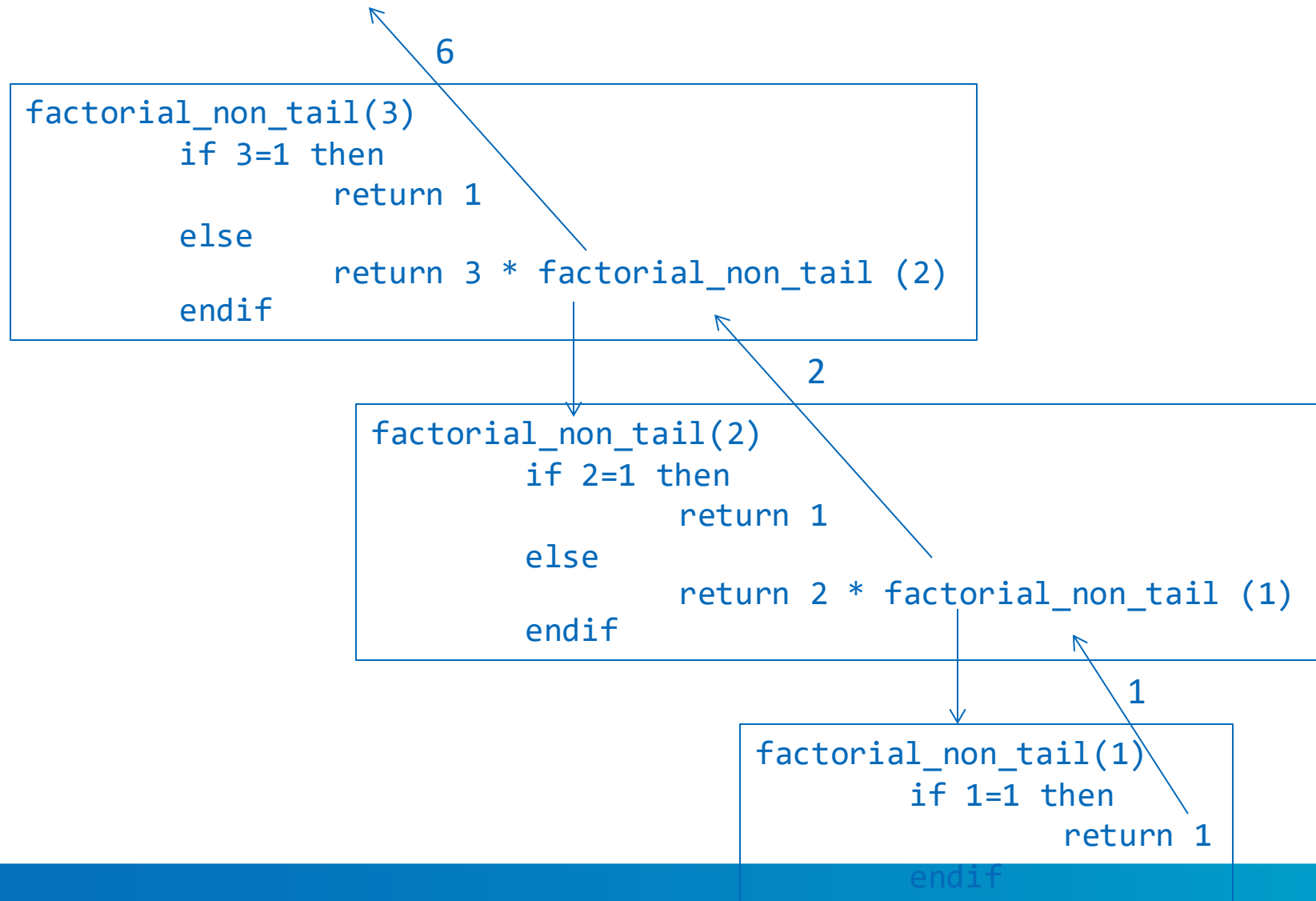


```
factorial_non_tail(2)
  if 2=1 then
    return 1
  else
    return 2 * factorial_non_tail (1)
  endif
```



```
factorial_non_tail(1)
  if 1=1 then
    return 1
  endif
```

Recursion Simulation



Example of Tail Recursion

Factorial: $n! = 1 \times 2 \times \dots \times (n-1) \times n$

Or: **$n! = n \times (n-1)!, 1!=1$**

```
function factorial_tail(n, accumulator)
```

Input: n and accumulator, two natural numbers

Output: the n-th factorial number

```
if n = 1 then
```

```
    return accumulator
```


```
else
```

```
    return factorial_tail(n-1, n*accumulator)
```


```
endif
```

Recursion Simulation

```
factorial_tail(3,1)
  if 3=1 then
    return 1
  else
    return factorial_tail(2,3*1)
  endif
```



```
factorial_non_tail(2,3)
  if 2=1 then
    return 3
  else
    return factorial_tail(1,2*3)
  endif
```



```
factorial_non_tail(1,6)
  if 1=1 then
    return 6
  endif
```

Why Tail Recursion?

Tail recursion is usually more efficient (although more difficult to write) than non tail recursion

The recursive calls do not need to be added to the call stack: there is only one, the current call, in the stack

It is possible to turn tail recursions into iterative algorithms

Recursion -> Iterative Algorithm

The general form of every *tail recursion* is:

- ret is the returned type
- Para is a list of parameters
- cond is the base case
- state0, state1 and state2 are statements
- f is a function transforming the parameters

```
function recursion(Para)
Input: set of parameters Para
Output: ret
state0
if cond then
    state1
else
    state2
    recursion(f(Para))
endif
```

Recursion -> Iterative Algorithm

The *iterative version of a tail recursion* is:

- ret is the returned type
- Para is a list of parameters
- cond is the base case
- state0, state1 and state2 are statements
- f is a function transforming the parameters

```
function iterative(Para)
Input: set of parameters Para
Output: ret
state0
while non cond do
    state2
    Para ← f(Para)
    state0
endwhile
state1
endif
```

Factorial Recursive

```
function factorial_tail(n, accumulator)
```

Input: n and accumulator two natural numbers

Output: the n-th factorial number

```
if n = 1 then
```

```
    return accumulator
```

```
else
```

```
    return factorial_tail(n-1, n*accumulator)
```

```
endif
```

cond

state1

f(Para)

Factorial Iterative

```
function factorial_iter(n, result)
```

Input: n and accumulator two natural numbers

Output: the n-th factorial number

```
while n > 1 do
```

```
    accumulator ← n * accumulator
```

```
    n ← n - 1
```

```
done
```

```
return accumulator
```

```
endif
```

non cond

f(Para)

state1

Recursion -> Iterative Algorithm

When you want to write iteratively a recursive function, the first technique is to come up with a tail recursion and then use the solution presented in previous slides

Otherwise you need to store context of the calls (in a way, re-doing the call stack in the program)

- use extra structures (e.g., arrays) to store the intermediary results.

This is an exemple of what's called dynamic programming

Recursion -> Iterative Algorithm

```
function factorial_dynamic(n)
```

Input: n a natural number

Output: the n-th factorial number

```
array <- array of size n
```

```
array[0] <- 1
```

```
for i from 2 till n do
```

```
    array[i-1] = array[i-2] * i
```

```
endfor
```

```
return array[n-1]
```

Complexity of Recursive Algorithms

We cannot apply the exact same mechanism we used for iterative algorithms: counting the number of basic operations and loops

Here we need to assess two things:

- the number of basic operations in each activation of the recursion (*this is easy*)
- the number of activations (*this is a little more difficult*)

Recursion Example

Factorial: $n! = 1 \times 2 \times \dots \times (n-1) \times n$

Or: **$n! = n \times (n-1)!, 1!=1$**

```
function factorial(n)
```

```
Input: n a natural number
```

```
Output: the n-th factorial number
```

```
if n = 1 then #1 operation
```

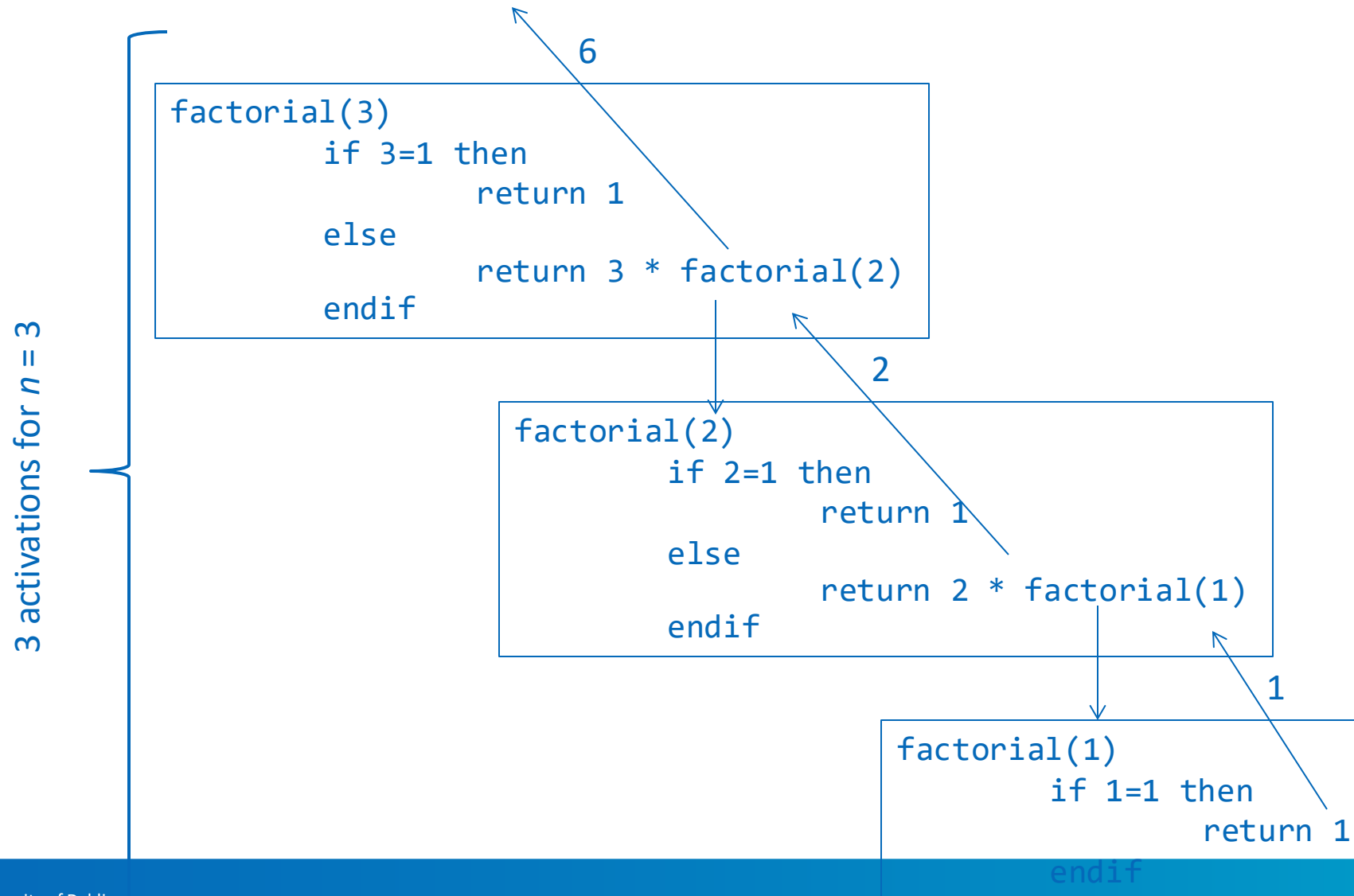
```
    return 1 #1 operation
```

```
else
```

```
    return n * factorial(n-1) #3 operations
```

```
endif
```

Recursion Simulation



Complexity of Recursive Algorithms

number of activations: **$O(n)$** (3 for $n=3$)

number of operations: 2 for base case, 4 otherwise (constant running time anyway) ->

$O(4)$ = $O(1)$

Total: **$O(n)$** (or $T(n) = 4n$, $O(4n) = O(n)$)