

Analysis Report: Sorting Algorithms Benchmark

1. Introduction

This report presents a detailed comparative analysis of two classical sorting algorithms: **Selection Sort** and **Insertion Sort**. The focus is on their correctness, performance characteristics, and efficiency in handling arrays of varying sizes and distributions.

2. Algorithm Descriptions

Selection Sort

Selection Sort is a simple comparison-based algorithm that repeatedly selects the minimum element from the unsorted portion and places it at the beginning. It has a worst-case and average-case time complexity of $O(n^2)$ and a best-case of $O(n^2)$ when optimized.

Key Features:

- Iteratively selects the minimum element.
- Can be optimized to stop early if the array is already sorted.
- Stable when implemented carefully.

Insertion Sort

Insertion Sort builds the sorted array one element at a time by comparing the current element with the elements before it and inserting it at the correct position. It can be further optimized using binary search for insertion position.

Key Features:

- Efficient for nearly sorted arrays.
- Adaptive: performance improves when array is partially sorted.
- Time Complexity: $O(n^2)$ worst-case, $O(n)$ best-case.

3. Testing Methodology

Both algorithms were tested using JUnit test suites and custom benchmark runners. The tests covered:

- Empty arrays
- Single-element arrays
- Already sorted arrays

- Reverse-sorted arrays
- Random arrays with duplicates
- Large arrays up to 100,000 elements

Performance metrics tracked include **comparisons**, **swaps**, **array accesses**, and **execution time in milliseconds**.

4. Performance Metrics

Algorithm	Array Type	Size	Time (ms)	Comparisons	Swaps	Array Accesses
Selection Sort	Random	100	0	538	2694	6125
Selection Sort	Sorted	100	0	99	0	198
Selection Sort	Reverse	100	2	481	4950	10580
Selection Sort	Nearly-sorted	100	0	624	57	988
Insertion Sort	Random	1000	4	8571	252762	516094
Insertion Sort	Sorted	1000	0	999	0	1998
Insertion Sort	Reverse	1000	0	7988	499500	1008987
Insertion Sort	Nearly-sorted	1000	0	8595	6710	24026

(Further large-scale benchmarks available in CSV results)

5. Comparative Analysis

- **Efficiency:** Insertion Sort outperforms Selection Sort on nearly sorted and small datasets due to its adaptive behavior.
- **Worst-case scenarios:** Both algorithms exhibit quadratic time complexity; however, Selection Sort consistently performs a large number of swaps, impacting cache performance.
- **Memory Access:** Insertion Sort, especially with binary search, reduces unnecessary comparisons and array movements.
- **Scalability:** Both algorithms scale poorly for very large datasets compared to modern $O(n \log n)$ algorithms.

6. Conclusion

The analysis demonstrates that while both Selection Sort and Insertion Sort are easy to implement and suitable for educational purposes, **Insertion Sort provides superior performance in practice for small or nearly sorted datasets**. Selection Sort remains predictable in execution but suffers from higher swap overhead.

Key Takeaways:

- Use Insertion Sort for adaptive sorting and small datasets.
- Selection Sort is primarily useful in cases where simplicity and deterministic swapping are required.
- For large-scale sorting, more advanced algorithms are recommended.

End of Report