

CS127 Homework 6

Due: November 19, 2014 3:00PM

Warmup 1 (Textbook Problem 12.3)

Let relations $r_1(A,B,C)$ and $r_2(C,D,E)$ have the following properties: r_1 has 20,000 tuples, r_2 has 45,000 tuples, 25 tuples of r_1 fit on one block, and 30 tuples of r_2 fit on one block. Estimate the number of block transfers required, using each of the following join strategies for $r_1 \bowtie r_2$:

1. Nested-loop join
2. Block nested-loop join
3. Merge join
4. Hash join

Warmup 2 (Textbook Problem 12.6)

Consider the following bank database, where the primary keys are underlined:

branch(branch_name, branch_city, assets)
customer(customer_name, customer_street, customer_city)
loan(loan_number, branch_name, amount)
borrower(customer_name, loan_number)
account(account_number, branch_name, balance)
depositor(customer_name, account_number)

Suppose that a B+Tree index on *branch_city* is available on relation *branch*, and that no other index is available. List different ways to handle the following selections that involve negation:

1. $\sigma_{\neg(\text{branch_city} < \text{"Brooklyn"})}(\text{branch})$
2. $\sigma_{\neg(\text{branch_city} = \text{"Brooklyn"})}(\text{branch})$
3. $\sigma_{\neg(\text{branch_city} < \text{"Brooklyn"}) \vee \text{assets} < 5000}(\text{branch})$

Warmup 3 (Textbook Problem 13.4)

Consider the relations $r_1(A,B,C)$, $r_2(C,D,E)$, and $r_3(E,F)$, with primary keys A , C , and E , respectively. Assume that r_1 has 1000 tuples, r_2 has 1500 tuples, and r_3 has 750 tuples. Estimate the size of $r_1 \bowtie r_2 \bowtie r_3$ and give an efficient strategy for computing the join.

Warmup 4 (Textbook Problem 14.12)

List the ACID properties. Explain the usefulness of each.

Problem 5 (To Be Graded)

Consider again the simplified university registrar database from the previous homeworks:

| Student | | | Course | | | Enrollment | | | |
|---------|----------|------|--------|----------|------------|------------|-------|----------|-------|
| name | gradyear | gpa | title | semester | instructor | name | title | semester | grade |
| Amy | 2016 | 3.95 | CS33 | 2014F | Doeppner | Eliza | CS33 | 2014F | A |
| Ben | 2015 | 3.87 | CS127 | 2014F | Zdonik | Eliza | CS127 | 2014F | A |
| Carl | 2016 | 3.29 | CS195 | 2013F | Kraska | Ben | CS127 | 2012F | A |
| Dan | 2017 | 3.43 | CS127 | 2012F | Zdonik | Carl | CS195 | 2013F | C |
| Eliza | 2015 | 4.0 | CS136 | 2012S | Fonseca | Carl | CS127 | 2014F | B |

The database has the following statistics:

- $Student(\underline{name}, gradyear, gpa)$: $n_{Student} = 10,000$, $f_{Student} = 100$
- $Course(\underline{title}, semester, instructor)$: $n_{Course} = 200$, $f_{Course} = 50$
- $Enrollment(\underline{name}, \underline{title}, \underline{semester}, grade)$: $n_{Enrollment} = 250,000$, $f_{Enrollment} = 25$

Answer the questions about the following queries:

1. *select name, gpa from student where gpa > 3.75*
 - a. Draw a plan for this query.
 - b. Assume *Student* is sorted. Does this affect how the query should be executed? Explain.
 - c. Instead, we add an index to speed up the query. Should we use a B-tree or hash index? Explain.
2. *select * from student s, course c, enrollment e where s.name = e.name and c.title = e.title and c.semester = e.semester*
 - a. Draw all possible plans for this query.
 - b. Choose the best plan and give its cost with a block nested-loop join and a buffer size of 3 blocks.
 - c. Choose the best plan and give its cost with a block nested-loop join and a buffer size of 6 blocks.