LECTURE 9

Last time:

- Channel capacity
- Binary symmetric channels
- Erasure channels
- Maximizing capacity

Lecture outline

- Maximizing capacity: Arimoto-Blahut
- Convergence
- Examples

Lemma 1:

$$I(X;Y) = \max_{\widehat{P}_{X|Y}} \sum_{y \in \mathcal{Y}} \sum_{x \in \mathcal{X}} P_X(x) P_{Y|X}(y|x)$$
$$\log \left(\frac{\widehat{P}_{X|Y}(x|y)}{P_X(x)}\right)$$

Proof:

$$I(X;Y) = \sum_{x \in \mathcal{X}} \sum_{x \in \mathcal{Y}} P_{X|Y}(x|y) P_Y(y) \log \left(\frac{P_{X|Y}(x|y)}{P_X(x)}\right)$$

Recall:

$$P_{X|Y}(x|y) = \frac{P_X(x)P_{Y|X}(y|x)}{\sum_{x' \in \mathcal{X}P_X(x')P_{Y|X}(y|x')}}$$

and

$$P_Y(y) = \sum_{x' \in \mathcal{X}} P_X(x') P_{Y|X}(y|x')$$

$$\begin{split} &I(X;Y) - \sum_{y \in \mathcal{Y}} \sum_{x \in \mathcal{X}} P_X(x) P_{Y|X}(y|x) \\ &\log \left(\frac{\widehat{P}_{X|Y}(x|y)}{P_X(x)} \right) \\ &= &I(X;Y) - \sum_{y \in \mathcal{Y}} \sum_{x \in \mathcal{X}} P_Y(y) P_{X|Y}(x|y) \\ &\log \left(\frac{\widehat{P}_{X|Y}(x|y)}{P_X(x)} \right) \\ &= &\sum_{y \in \mathcal{Y}} \sum_{x \in \mathcal{X}} P_Y(y) P_{X|Y}(x|y) \log \left(\frac{P_{X|Y}(x|y)}{\widehat{P}_{X|Y}(x|y)} \right) \\ &\text{(using log}(x) \geq 1 - \frac{1}{x}) \\ &\geq &\sum_{y \in \mathcal{Y}} \sum_{x \in \mathcal{X}} P_Y(y) P_{X|Y}(x|y) \\ &- &\sum_{y \in \mathcal{Y}} \sum_{x \in \mathcal{X}} P_Y(y) \widehat{P}_{X|Y}(x|y) \\ &- &\sum_{y \in \mathcal{Y}} \sum_{x \in \mathcal{X}} P_Y(y) \widehat{P}_{X|Y}(x|y) \\ &- &\sum_{y \in \mathcal{Y}} \sum_{x \in \mathcal{X}} P_Y(y) \widehat{P}_{X|Y}(x|y) \end{split}$$

Capacity is

$$C = \max_{P_X} \max_{\widehat{P}_{X|Y}} \sum_{y \in \mathcal{Y}} \sum_{x \in \mathcal{X}} P_X(x) P_{Y|X}(y|x)$$
$$\log \left(\frac{\widehat{P}_{X|Y}(x|y)}{P_X(x)}\right)$$

For fixed P_X , RHS is maximized when

$$\hat{P}_{X|Y}(x|y) = \frac{P_X(x)P_{Y|X}(y|x)}{\sum_{x' \in \mathcal{X}} P_X(x')P_{Y|X}(y|x')}$$

For fixed $\widehat{P}_{X|Y}$, RHS is maximized when

$$P_X(x) = \frac{e^{\sum_{y \in \mathcal{Y}} P_{Y|X}(y|x) \log(\widehat{P}_{X|Y}(x|y))}}{\sum_{x' \in \mathcal{X}} \left(e^{\sum_{y \in \mathcal{Y}} P_{Y|X}(y|x') \log(\widehat{P}_{X|Y}(x'|y))}\right)}$$

Combining the two means maximization when

$$P_X(x) = \frac{e^{\sum_{y \in \mathcal{Y}} P_{Y|X}(y|x) \log(\widehat{P}_{X|Y}(x|y))}}{\sum_{x' \in \mathcal{X}} \left(e^{\sum_{y \in \mathcal{Y}} P_{Y|X}(y|x') \log(\widehat{P}_{X|Y}(x'|y))}\right)}$$

$$= \frac{P_X(x)e^{\sum_{y \in \mathcal{Y}} P_{Y|X}(y|x) \log\left(\frac{P_{Y|X}(y|x)}{\sum_{x' \in \mathcal{X}} P_{X}(x')P_{Y|X}(y|x')}\right)}}{\sum_{x' \in \mathcal{X}} P_X(x') \left(e^{\sum_{y \in \mathcal{Y}} P_{Y|X}(y|x') \log\left(\frac{P_{Y|X}(y|x)}{\sum_{x' \in \mathcal{X}} P_{X}(x')P_{Y|X}(y|x')}\right)}\right)$$

Note also that $\sum_{x \in \mathcal{X}} P_X(x) = 1$.

This may be very hard to solve.

Proof:

The first two statements follow immediately from our lemma

For any value of x where $P_{X|Y}(x|y) = 0$, $P_X(x)$ should be set to 0 to obtain the maximum.

To find the maximum over the PMF P_X , let us first ignore the constraint of positivity and use a Lagrange multiplier for the $\sum_x P_X(x) = 1$

Then

$$\frac{\partial}{\partial P_X(x)} \{ \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} P_X(x) P_{Y|X}(y|x) \log \left(\frac{\widehat{P}_{X|Y}(x|y)}{P_X(x)} \right) + \lambda \left(\sum_{x \in \mathcal{X}} P_X(x) - 1 \right) \} = 0$$

Equivalently

$$-\log(P_X(x)) - 1 + \sum_{y \in \mathcal{Y}} P_{Y|X}(y|x) \log \left(\widehat{P}_{X|Y}(x|y) \right) + \lambda = 0$$

SO

$$P_X(x) = \frac{e^{\sum_{x \in \mathcal{X}} P_{Y|X}(y|x) \log\left(\widehat{P}_{X|Y}(x|y)\right)}}{\sum_{x \in \mathcal{X}} e^{\sum_{x \in \mathcal{X}} P_{Y|X}(y|x) \log\left(\widehat{P}_{X|Y}(x|y)\right)}}$$

(this ensures that λ is such that the sum of the $P_X(x)$ s is 1)

What about the constraint we did not use for positivity?

The solution we found satisfies that.

Let ${\cal P}_X^0$ be a PMF and let

$$P_X^{r+1}(x) = P_X^r(x) \frac{c_x(P_X^r(x_1), \dots, P_X^r(x_{|\mathcal{X}|}))}{\sum_{x' \in \mathcal{X}} c_x(P_X^r(x_1), \dots, P_X^r(x_{|\mathcal{X}|})) P_X^r(x')}$$

where

$$c_x\left(P_X^r(x_1),\ldots,P_X^r(x_{|\mathcal{X}|})\right)$$

$$= e^{\sum_{y \in \mathcal{Y}} P_{Y|X}(y|x) \log \left(\frac{P_{Y|X}(y|x)}{\sum_{x' \in \mathcal{X}} P_{X}(x') P_{Y|X}(y|x')}\right)}$$

the sequence I^r of I(X;Y) for X taking the PMF P_X^R for I^r converges to C from below

Proof:

For any given P_X^r , we can increase mutual information by taking

$$P_{Y|X}^{r} = \frac{P_{X}^{r}(x)P_{Y|X}(y|x)}{\sum_{x' \in \mathcal{X}} P_{X}^{r}(x')P_{Y|X}(y|x')}$$

With $P_{Y|X}^r$ fixed, then choose P_X^{r+1} by

$$P_X^{r+1}(x) = \frac{e^{\sum_{y \in \mathcal{Y}} P_{Y|X}(y|x) \log(P_{X|Y}^r(x|y))}}{\sum_{x' \in \mathcal{X}} e^{\sum_{y \in \mathcal{Y}} P_{Y|X}(y|x') \log(P_{X|Y}^r(x'|y))}}$$

If we define

$$J^r = \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} P_X^{r+1}(x) P_{Y|X}(y|x) \log \left(\frac{P_{X|Y}^r(x|y)}{P_X^{r+1}(x)} \right)$$

Then
$$I^r \leq J^r \leq I^{r+1} \leq J^{r+1} \leq \dots$$

This an upper bounded non-decreasing sequence, therefore it reaches a limit

Why is the limit C?

Let P_X^* be a capacity achieving PMF

$$\begin{split} &\sum_{x \in \mathcal{X}} P_X^*(x) \log \left(\frac{P_X^{r+1}(x)}{P_X^r(x)} \right) \\ &= \sum_{x \in \mathcal{X}} P_X^*(x) \\ &\log \left(\frac{c_x \left(P_X^r(x_1), \dots, P_X^r(x_{|\mathcal{X}|}) \right)}{\sum_{x' \in \mathcal{X}} c_x \left(P_X^r(x_1), \dots, P_X^r(x_{|\mathcal{X}|}) \right) P_X^r(x')} \right) \\ &= \sum_{x \in \mathcal{X}} P_X^*(x) \sum_{y \in \mathcal{Y}} P_{Y|X}(y|x) \\ &\log \left(\frac{P_{Y|X}(y|x)}{\sum_{x' \in \mathcal{X}} P_X^r(x') P_{Y|X}(y|x')} \right) \\ &- \sum_{x \in \mathcal{X}} P_X^*(x) \sum_{y \in \mathcal{Y}} P_{Y|X}(y|x) \\ &\log \left(\sum_{x' \in \mathcal{X}} P_X^r(x') \\ &\sum_{x' \in \mathcal{X}} P_X^r(x') \log \left(\frac{P_{Y|X}(y'|x')}{\sum_{x'' \in \mathcal{X}} P_X^r(x'') P_{Y|X}(y'|x'')} \right) \right) \\ &e \end{split}$$

By considering the K-L distance, we have that

$$\sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} P_X^*(x) P_{Y|X}(y|x)$$

$$\log\left(\frac{\sum_{x'\in\mathcal{X}}P_X^*(x')P_{Y|X}(y|x')}{\sum_{x'\in\mathcal{X}}P_X^r(x')P_{Y|X}(y|x')}\right) \geq 0$$

SO

$$\sum_{x \in \mathcal{X}} P_X^*(x) \log \left(\frac{P_X^{r+1}(x)}{P_X^r(x)} \right)$$

$$\geq \sum_{x \in \mathcal{X}} P_X^*(x) \sum_{y \in \mathcal{Y}} P_{Y|X}(y|x)$$

$$\log \left(\frac{P_{Y|X}(y|x)}{\sum_{x' \in \mathcal{X}} P_X^*(x') P_{Y|X}(y|x')} \right)$$

$$- \sum_{x \in \mathcal{X}} P_X^*(x) \sum_{y \in \mathcal{Y}} P_{Y|X}(y|x)$$

$$\log \left(\sum_{x' \in \mathcal{X}} P_X^r(x') \right)$$

$$\sum_{x' \in \mathcal{X}} P_{Y|X}^r(x') \log \left(\frac{P_{Y|X}(y'|x')}{\sum_{x'' \in \mathcal{X}} P_X^r(x'') P_{Y|X}(y'|x'')} \right)$$

Hence

$$\sum_{x \in \mathcal{X}} P_X^*(x) \log \left(\frac{P_X^{r+1}(x)}{P_X^r(x)} \right)$$
> $C - J^r$

Sum over r

$$\sum_{r=0}^{m} (C - J^{r})$$

$$\leq \sum_{r=0}^{m} \sum_{x \in \mathcal{X}} P_{X}^{*}(x) \log \left(\frac{P_{X}^{r+1}(x)}{P_{X}^{r}(x)} \right)$$

$$= \sum_{x \in \mathcal{X}} P_{X}^{*}(x) \log \left(\frac{P_{X}^{m+1}(x)}{P_{X}^{0}(x)} \right)$$

$$\leq \sum_{x \in \mathcal{X}} P_{X}^{*}(x) \log \left(\frac{P_{X}^{*}(x)}{P_{X}^{0}(x)} \right)$$

 $C-J^r \geq$ 0 and non increasing, with bounded sum, so it goes to 0, hence J^r converges to C

In practice, **convergence can be very slow**

Example

Other types of maximization

Interior point methods

Cutting plane algorithms

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