

Data Structures and Algorithms

Mohammad GANJTABESH

*Department of Computer Science,
School of Mathematics, Statistics and Computer Science,
University of Tehran,
Tehran, Iran.*

mgtabesh@ut.ac.ir



*Dept. of Computer Science
University of Tehran*

Perfectly - Balanced

BST



AVL - Tree

strictly Balance
 $(-1 \leq BF \leq 1)$



B - Tree



2 - 3 - 4 Tree

Red - Black Tree

(Self - balancing)
Almost Balance

Binary

Data Structures and Algorithms

Undergraduate course



Mohammad GANJTABESH
mgtabesh@ut.ac.ir



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University of Tehran

Motivations

2-3-4 Tree



Red-Black tree



search : $O(\log n)$

$2 \log n$ | ~~2.5n~~

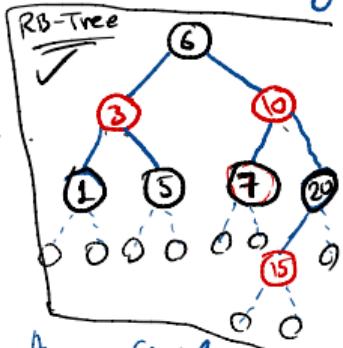
insertion : $O(\log n)$

deletion : $O(\log n)$

Basic Definitions/Rules

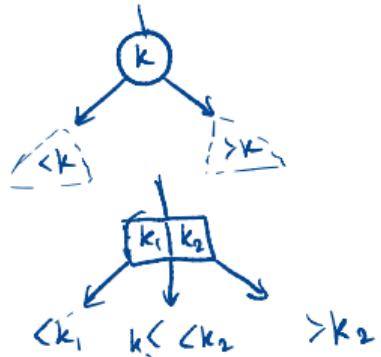
Red-Black tree is a binary search tree with the following rules:

- Every node is either RED or Black.
- Root is always Black.
- New insertions are always RED.
- Every path from the root to leaf has the same number of Black nodes.
- No path can have two consecutive RED nodes.
- Nulls are Black.



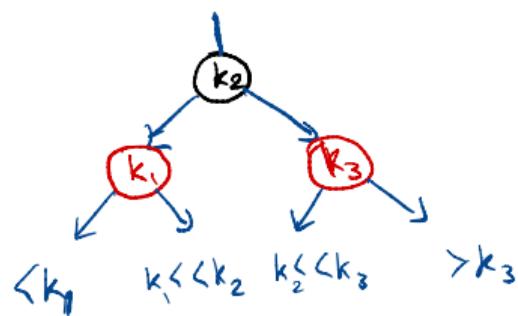
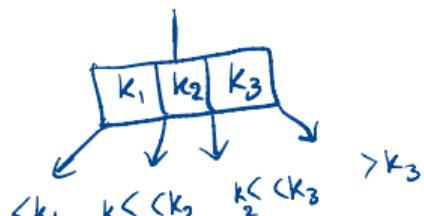
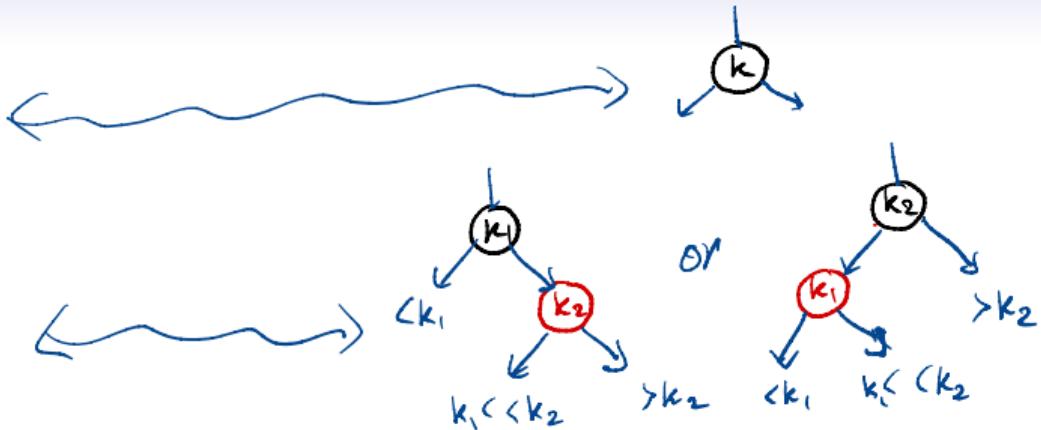
- ★ Fixing the tree if the above rules are violated.
- ★ Fixing is done by Rotation and Color flips.

2-3-4 Tree



vs

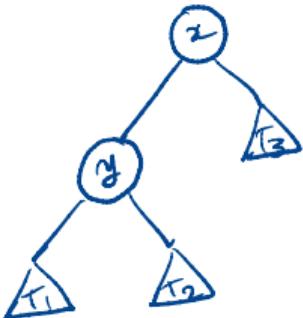
Red-Black Tree



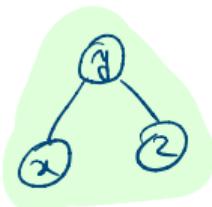
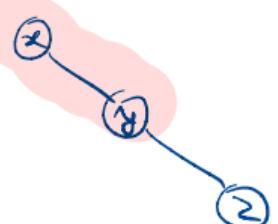
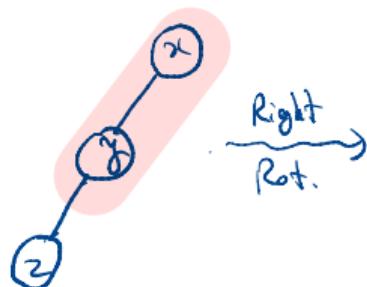
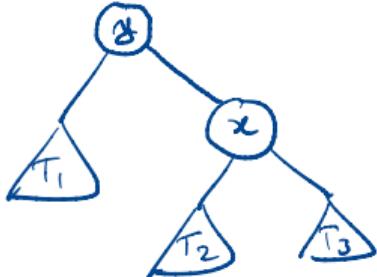
Rotations and Color Flippings

(exactly the same)
as AVL-tree

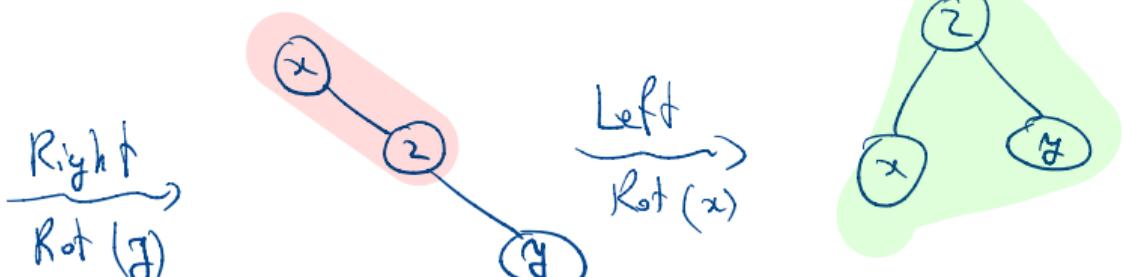
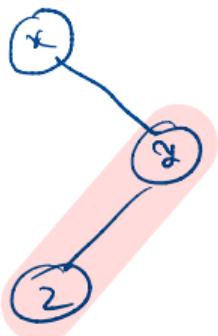
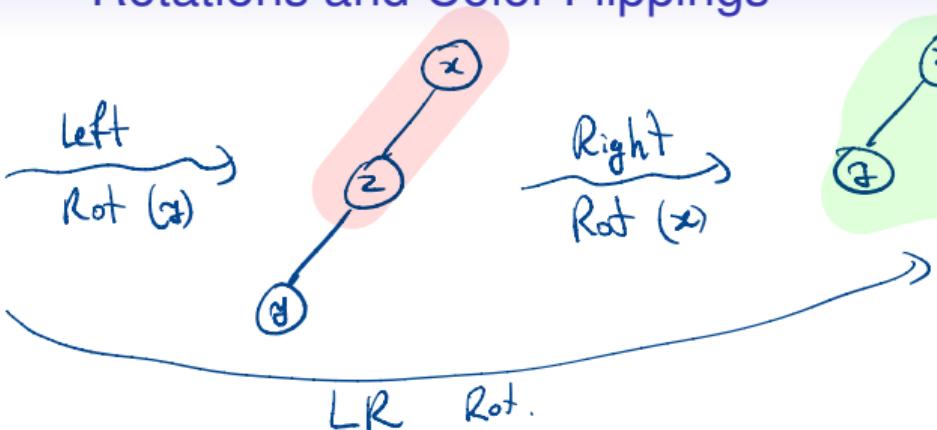
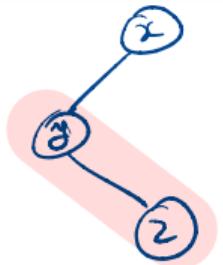
new operation



Left - Rotation
Right Rotation



Rotations and Color Flippings



Red-Black Tree Operations : Insert

① Insert a new node as we done in BST
(new node is always RED)

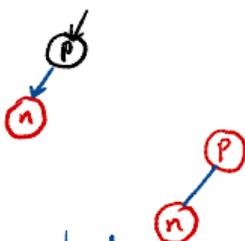
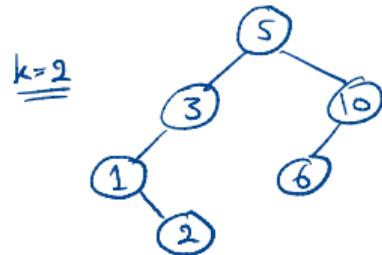
- if the node is root, flip its color.



- else, check the parent's color :

- if the color is Black , done.

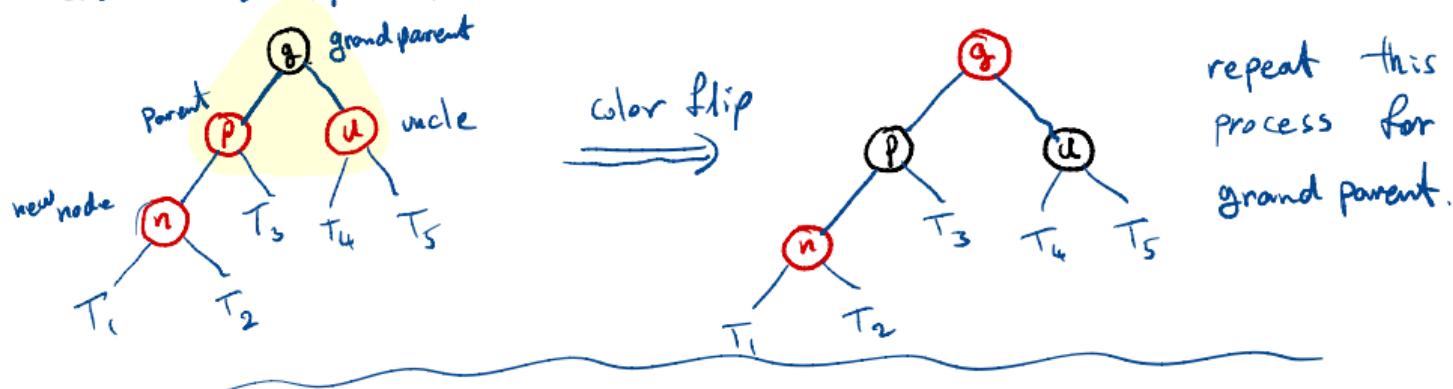
- else (if the parent's color is Red),
check the color of the node's uncle !



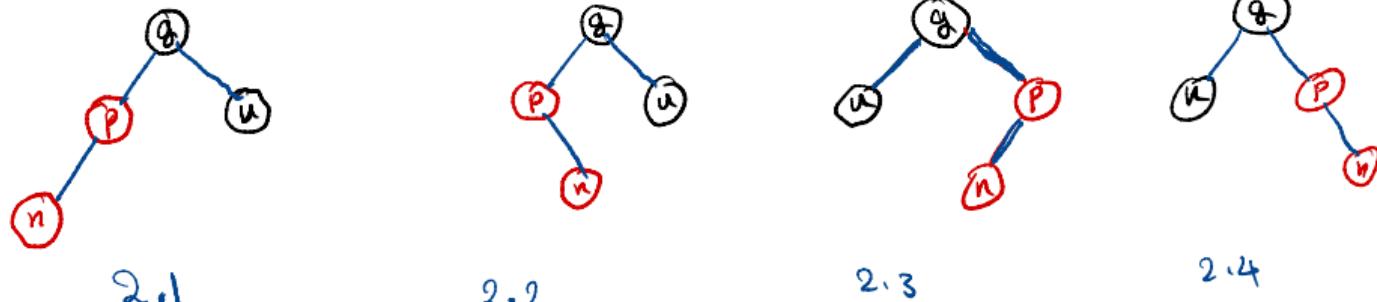
We have several cases

Red-Black Tree Operations : Insert

- Case 1 : if the node's uncle is Red.

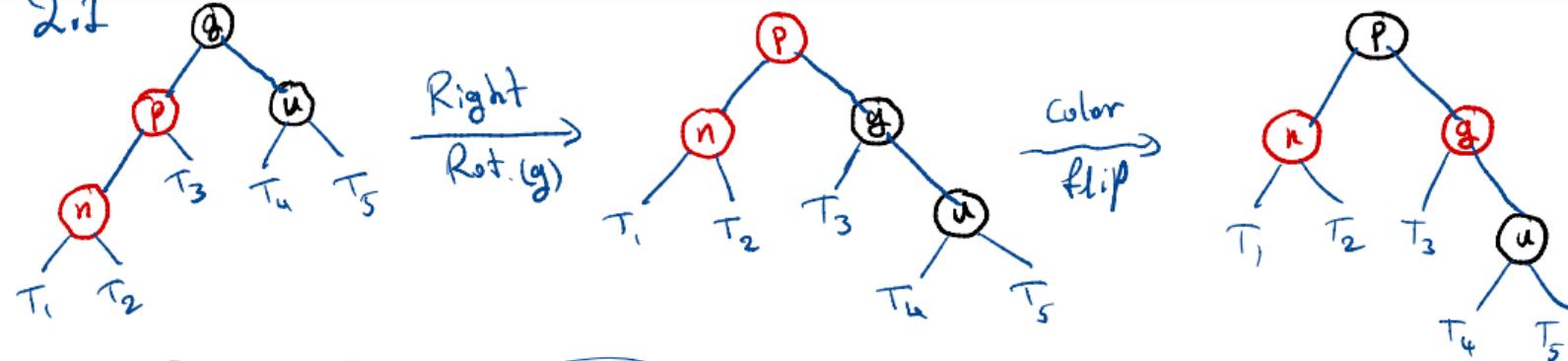


- Case 2 : if the node's uncle is Black :

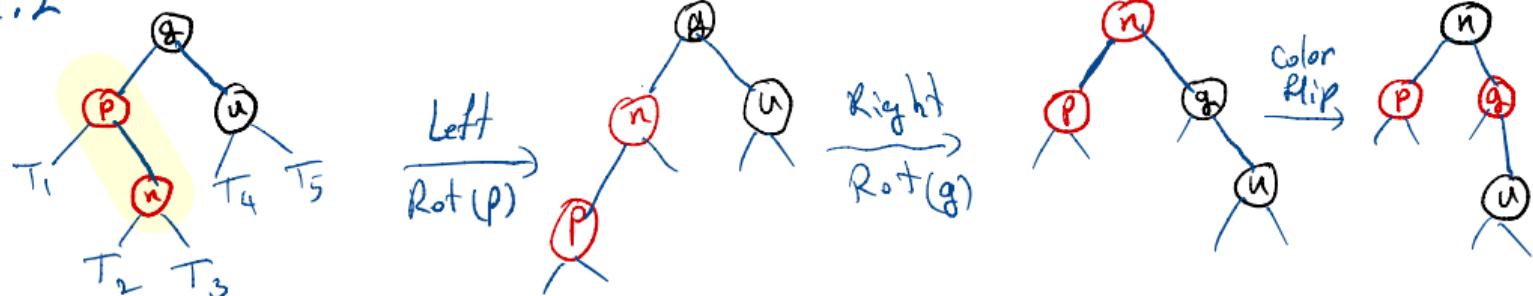


Red-Black Tree Operations : Insert

2.1

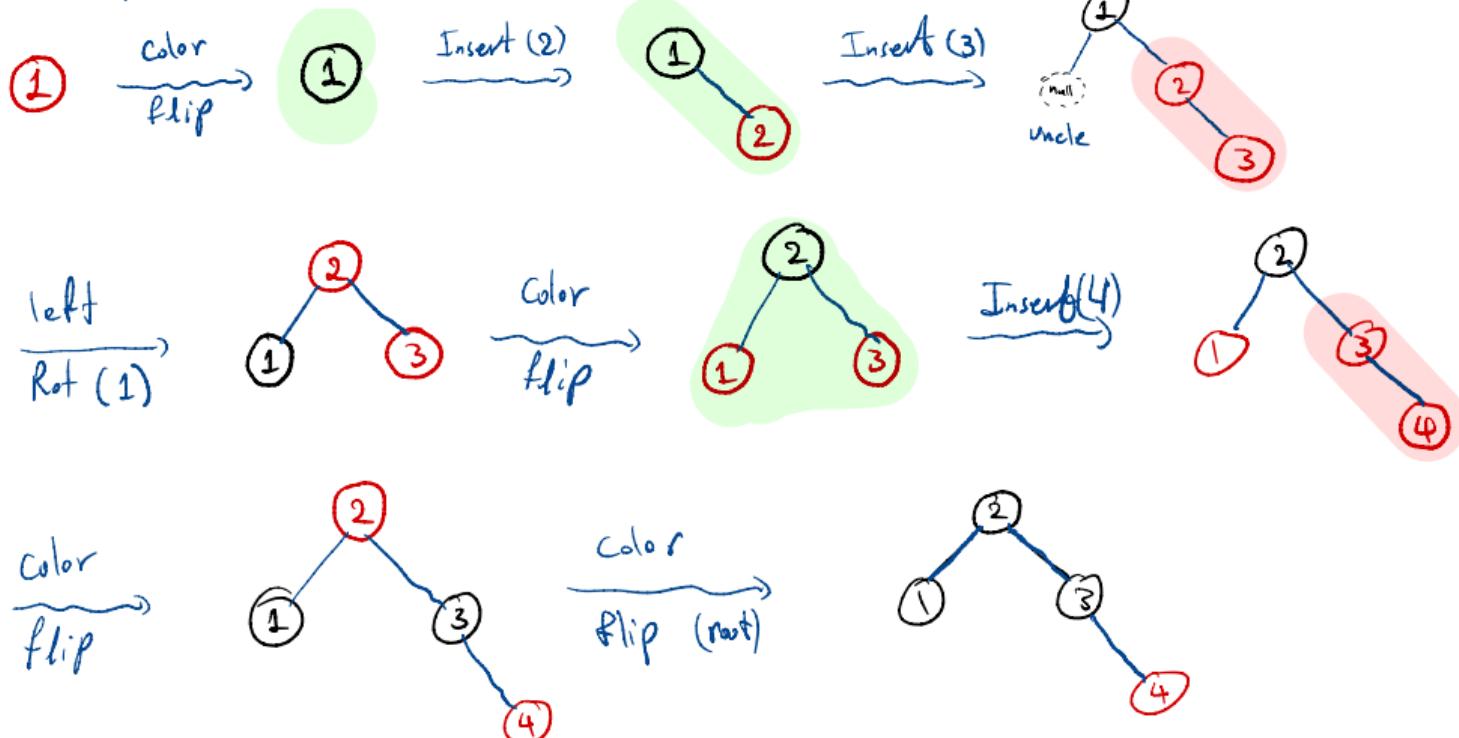


2.2



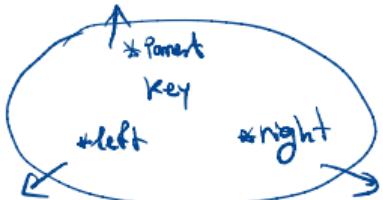
Red-Black Tree Operations : Insert

Example : 1, 2, 3, 4, 5, 6, ...



Red-Black Tree Implementation

```
struct Node {  
    int key;  
    Node *left, *right, *parent;  
    bool isBlack;  
    bool isLeftChild;  
    Node (int k) {  
        key ← k;  
        left ← right ← parent ← null;  
        isBlack ← false;  
        isLeftChild ← false;  
    }  
}
```



```
class Red-Black-Tree {  
    Node *root;  
    int size;  
    Red-Black-Tree () {  
        root ← null;  
        size ← 0;  
    }  
    // methods  
}
```



Red-Black Tree Implementation

```

public insert (int k) {
    Node *node ← new Node(k);
    if (root = null) {
        root ← node;
        root. isBlack ← true;
        size ← size + 1;
        return ;
    }
}

```

```
insert (root , node);
```

```
size ← size + 1;
```

```

private → insert (Node *parent, Node *node) {
    if (node.key > parent.key) {
        if (parent.right = null) {
            parent.right ← node;
            node.parent ← parent;
        } else {
            insert (parent.right , node);
        }
    } else { if (parent.left = null)
        parent.left ← node;
        node.parent ← parent;
    } node.isLeftChild ← true;
    else {
        insert (parent.left , node);
    }
}

```

Fix (node);

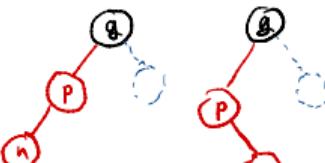
Red-Black Tree Implementation

```
Fix (Node *node){  
    if (node == root) return;  
    if (node.isBlack == false and node.parent.isBlack == false)  
        Fix-Tree (node);  
    Fix (node.parent);  
}
```



```
Fix-Tree (Node *node){
```

```
    if (node.parent.isLeftChild) {  $\Rightarrow$ 
```



```
        if (node.parent.parent.right == null or  
            node.parent.parent.right.isBlack)
```

```
            Rotation (node);
```

```
        if (node.parent.parent.right != null) { node.parent.parent.right.isBlack  $\leftarrow$  true; }
```

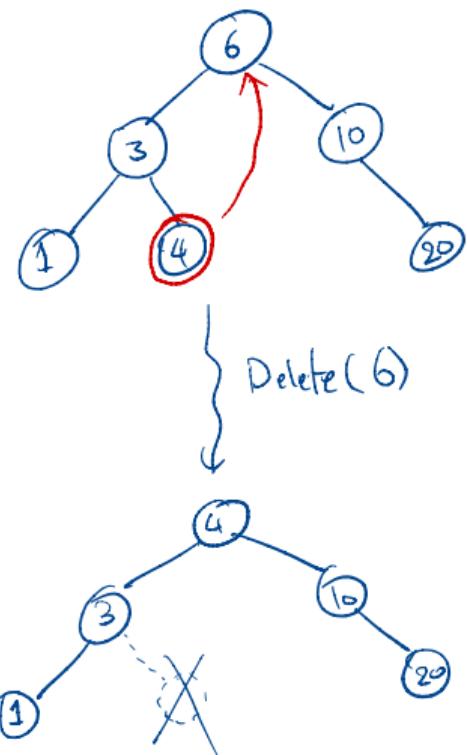
```
        node.parent.parent.isBlack  $\leftarrow$  false;
```

```
} node.parent.isBlack  $\leftarrow$  true;
```

```
} else {
```

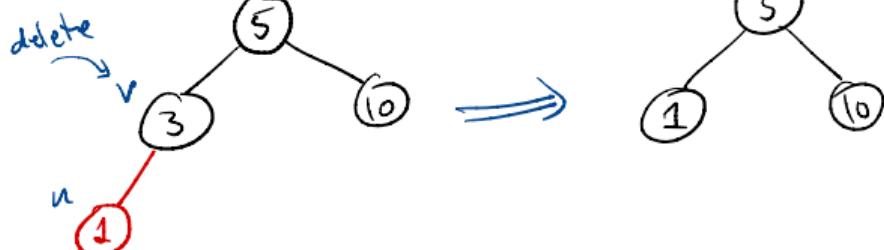


delete :



Let v be the node to be deleted
and u be the node to be replaced.

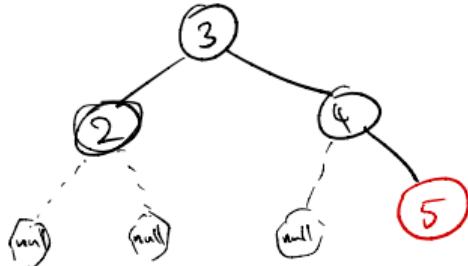
- ① if either v or u is Red:



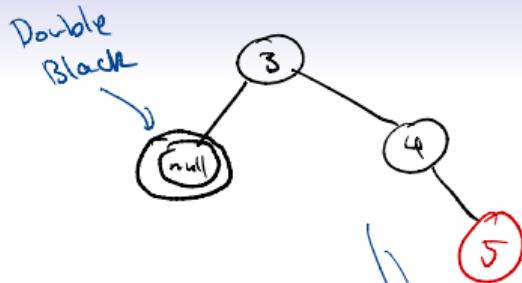
- ② if both u and v are Black.
cause Double Black



②

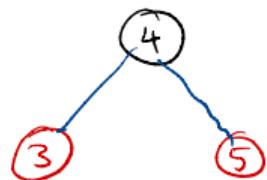


Delete (2)

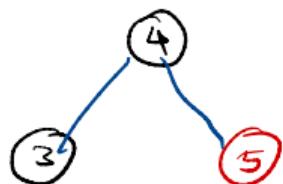


case 1) if sibling node is black and at least one of its children is Red
→ rotation and recoloring.

rotation

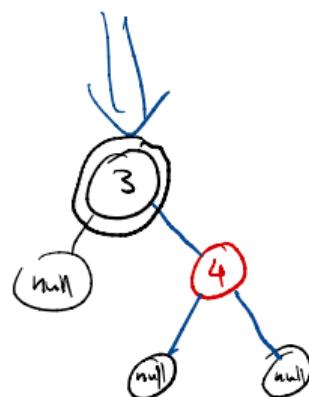
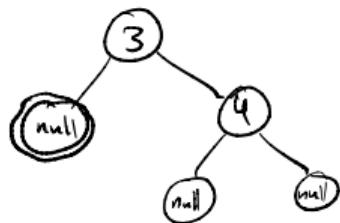
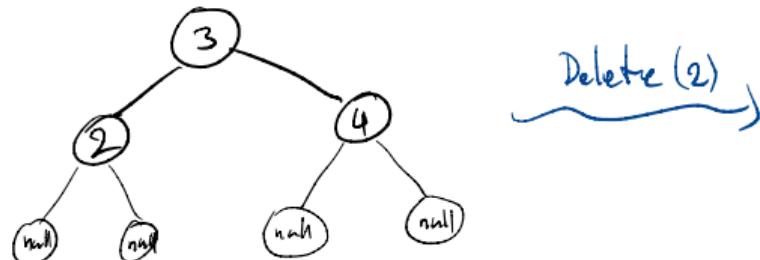


recolor



* Consider four cases for Rotation.

case 2) if the sibling is Black and both of its children are Black . \Rightarrow Perform recoloring and check the parent.



Case 3) if sibling is red \Rightarrow perform Rotation to move the sibling up and recolor the sibling and parent.

