DAI Assignment-2

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1 Mathemagic

1.1 Task A

For a Bernoulli random variable $X \sim Ber(p), P[X=0]=1-p, P[X=1]=p$ and $P[X=n]=0 \forall n \geq 2.$

$$G_{Ber}(z) = \sum_{n=0}^{\infty} P[X = n] z^n$$

= $P[X = 0] + P[X = 1] z$
 $G_{Ber} = (1 - p) + pz$

1.2 Task B

When $X \sim Bin(n, p)$, $P[X = k] = \binom{n}{k} p^k (1 - p)^{n-k} \ \forall k \in \mathbb{Z}, 0 \le k \le n$. $P[X = k] = 0 \ \forall k > n$.

$$G_{Bin}(z) = \sum_{k=0}^{\infty} P[X = n] z^{k}$$

$$= \sum_{k=0}^{n} \binom{n}{k} p^{k} (1-p)^{n-k} z^{k}$$

$$= \sum_{k=0}^{n} \binom{n}{k} (pz)^{k} (1-p)^{n-k}$$

$$= (pz + (1-p))^{n}$$

$$G_{Bin}(z) = G_{Ber}(z)^{n}$$

1.3 Task C

Let X_1 and X_2 be two random variables which take up non-negative integers and let $X = X_1 + X_2$.

$$G_X(z) = \sum_{n=0}^{\infty} P[X = n] z^n$$

$$= \sum_{n=0}^{\infty} \sum_{i=0}^{n} (P[X_1 = i] P[X_2 = n - i]) z^n$$

$$= \sum_{i=0}^{\infty} \sum_{n=i}^{\infty} (P[X_1 = i] P[X_2 = n - i]) z^n$$

$$= \sum_{i=0}^{\infty} \sum_{n=i}^{\infty} (P[X_1 = i] z^i) (P[X_2 = n - i] z^{n-i})$$

$$= \sum_{i=0}^{\infty} P[X_1 = i] z^i \sum_{n=i}^{\infty} P[X_2 = n - i] z^{n-i}$$

$$= \sum_{i=0}^{\infty} P[X_1 = i] z^i \sum_{n=i}^{\infty} P[X_2 = n] z^n$$

$$G_X(z) = G_{X_1}(z) G_{X_2}(z)$$

For k = 1, $G_{\Sigma}(z) = G(z)^k$. If for k = n - 1, $G_{\Sigma}(z) = G(z)^k$, then for k = n,

$$G_{\Sigma}(z) = G_{X_1 + X_2 + \dots + X_{n-1}}(z)G_{X_n}(x)$$

$$= \prod_{i=1}^n G_{X_i}(z)$$

$$G_{\Sigma}(z) = G(z)^n$$

From the Principle of Mathematical Induction, for all $k \in \mathbb{N}$, $G_{\Sigma}(z) = G(z)^k$

1.4 Task D

When $X \sim Geo(p)$, $P[X = n] = (1 - p)^{n-1}p$ and P[X = 0] = 0.

$$G_{Geo}(z) = \sum_{n=0}^{\infty} P[X = n] z^n$$

$$= \sum_{n=1}^{\infty} (1 - p)^{n-1} p z^n$$

$$= pz \sum_{n=1}^{\infty} ((1 - p)z)^{n-1}$$

$$= \frac{pz}{1 - (1 - p)z}$$

1.5 Task E

 $X \sim Bin(n,p)$ and $Y \sim NegBin(n,p)$. $Y = \sum_{k=1}^{n} Y_k$ where $Y_k \sim Geo(p)$.

$$\begin{split} G(Y) &= G(Y_1 + Y_2 + \dots + Y_n) \\ &= G(Y_1)G(Y_2) \dots G(Y_n) \\ &= \left(\frac{pz}{1 - (1 - p)z}\right)^n \\ G_Y^{(n,p)}(z) &= \left(\frac{pz}{1 - (1 - p)z}\right)^n \\ &= \left(\frac{1}{\frac{1}{pz} + (1 - \frac{1}{p})}\right)^n \\ G_X^{(n,p^{-1})}(z^{-1}) &= \left(\frac{1}{pz} + (1 - \frac{1}{p})\right)^n \\ \Longrightarrow G_Y^{(n,p)}(z) &= \left(G_X^{(n,p^{-1})}(z^{-1})\right)^{-1} \end{split}$$

1.6 Task F

 $P[Y = k] = {\binom{k-1}{n-1}} p^n (1-p)^{k-n}.$

$$G_Y^{(n,p)}(z) = \sum_{k=0}^{\infty} P[Y=k] z^k$$

$$= \sum_{k=n}^{\infty} \binom{k-1}{n-1} p^n (1-p)^{k-n} z^k$$

$$\left(\frac{1}{pz} + \left(1 - \frac{1}{p}\right)\right)^{-n} = \sum_{k=n}^{\infty} \binom{k-1}{n-1} p^n (1-p)^{k-n} z^k$$

Substitute x = z, p = 2

$$\left(\frac{1}{2x} + \frac{1}{2}\right)^{-n} = \sum_{k=n}^{\infty} {k-1 \choose n-1} 2^n (-1)^{k-n} x^k$$

$$x^n (1+x)^{-n} = \sum_{k=n}^{\infty} {k-1 \choose n-1} (-1)^{k-n} x^k$$

$$(1+x)^{-n} = \sum_{k=n}^{\infty} {k-1 \choose n-1} (-1)^{k-n} x^{k-n}$$

$$(1+x)^{-n} = \sum_{r=0}^{\infty} {n+r-1 \choose n-1} (-1)^r x^r$$

$$(1+x)^{-n} = \sum_{r=0}^{\infty} {n+r-1 \choose r} (-1)^r x^r$$

1.7 Task G

$$G(z) = \sum_{n=0}^{\infty} P[X = n] z^n$$

$$G'(z) = \sum_{n=0}^{\infty} n P[X = n] z^{n-1}$$

$$G'(1) = \sum_{n=0}^{\infty} n P[X = n]$$

$$G'(1) = E[X]$$

For Bernoulli random variable $X \sim Ber(p)$,

$$G_{Ber}(z) = (1 - p) + pz$$

$$G'_{Ber}(z) = p$$

$$G'_{Ber}(1) = p$$

$$E[X] = p$$

For Binomial random variable $X \sim Bin(n, p)$,

$$G_{Bin}(z) = (pz + (1-p))^n$$

 $G'_{Bin}(z) = np(pz + (1-p))^{n-1}$
 $G'_{Bin}(1) = np$
 $E[X] = np$

For Geometric random variable $X \sim Geo(p)$,

$$G_{Geo}(z) = \frac{pz}{1 - (1 - p)z}$$

$$G'_{Geo}(z) = \frac{p}{(1 - (1 - p)z)^2}$$

$$G'_{Geo}(1) = \frac{p}{p^2}$$

$$E[X] = \frac{1}{p}$$

For Negative Binomial random variable $X \sim NegBin(n, p)$,

$$G_{NegBin}(z) = \left(\frac{pz}{1 - (1 - p)z}\right)^n$$

$$G'_{NegBin}(z) = n\left(\frac{pz}{1 - (1 - p)z}\right)^{n-1} \frac{p}{(1 - (1 - p)z)^2}$$

$$G'_{NegBin}(1) = n\left(\frac{p}{p}\right)^{n-1} \frac{p}{p^2}$$

$$E[X] = \frac{n}{p}$$

2 Normal Sampling

2.1 Task A

To prove that Y is uniformly distributed in [0,1], we need to prove that for an arbitrary interval [a,b] in [0,1], Y lies in [a,b] with probability b-a.

The probability that $a \leq Y \leq b$ is the probability that the randomly chosen X is such that $a \leq F_X(x) \leq b$. The cumulative probability of x being in the interval $[a_0,b_0]$ such that $F_X(a_0)=a$ and $F_X(b_0)=b$ is b-a, since the cumulative probability that $x \leq b_0$ is b, and the cumulative probability of $x \leq a_0$ is a, hence probability that $a_0 \leq x \leq b_0$ is b-a. Our condition that $F_X()$ is invertible guarantees that the \leq can be replaced with <, since it cannot be that $\lim_{x\to a_0^+} \neq \lim_{x\to a_0^+}$, otherwise between a_0^- and a_0^+ , all elements would be mapped to by a_0 .

2.2 Task B

Given a sample y, let us return $F_X^{-1}(y)$. To prove that this has the same CDF as X, we need to prove that if $F_X(a_0) = a$ and $F_X(b_0) = b$, the probability of getting $a_0 \le F_X^{-1}(y) \le b_0$ is b-a. But this is fairly straightforward, since $a_0 \le F_X^{-1}(y) \le b_0 \iff a \le y \le b$, and the probability of getting $a \le y \le b$ is b-a.

Thus, this is the required function.

2.3 Task C

Code provided in 2c.ipynb.

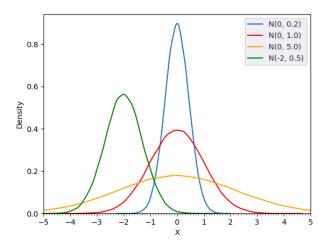


Figure 1: Histogram of the samples generated

2.4 Task D

Code provided in 2d.ipynb.

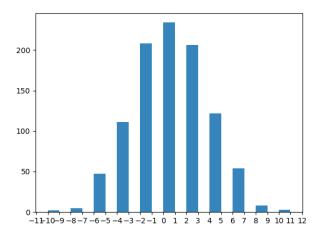


Figure 2: h=10

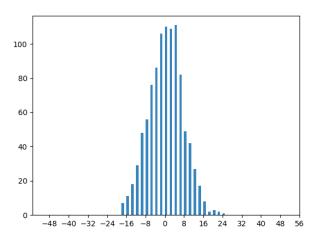


Figure 3: h=50

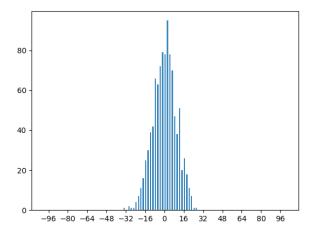


Figure 4: h=100

For h=100 alone, for better visibility, I have also written code with bins set to increment every 2 pockets. To observe the property of only even indexed bins being filled, simply replace range(-100, 102, 2) with range(-100, 102, 1).

2.5 Task E

$$P_h[X=2i] = \frac{\binom{2k}{i+k}}{2^{2k}}$$

Since the likelihood of a ball arriving at pocket i is equal to the likelihood of having i+k collisions which move it to the right within the 2k possible collisions overall. Applying Stirling's approximations,

$$P_{h}[X = 2j] \approx \frac{k^{2k}\sqrt{4\pi k}}{(k+j)^{k+j}\sqrt{2\pi(k+j)}(k-j)^{k-j}\sqrt{2\pi(k-j)}}$$

$$\approx \frac{k^{2k}\sqrt{2k}(k-j)^{j}}{(k^{2}-j^{2})^{k}\sqrt{2\pi(k^{2}-j^{2})}(k+j)^{j}}$$

$$\approx \frac{1}{\sqrt{\pi k}} \frac{k^{2k}(k-j)^{j}}{(k^{2})^{k}(k+j)^{j}} \qquad \text{(By approximating } k^{2}-j^{2}\approx j^{2})$$

$$= \frac{1}{\sqrt{\pi k}} \left(1 - \frac{2j}{k+j}\right)^{j}$$

Now, we independently approximate $\left(1 - \frac{2j}{k+j}\right)^j$ by taking its logarithm and using the property that for small x, $\ln(1-x) \approx -x$.

$$\ln\left(\left(1 - \frac{2j}{k+j}\right)^{j}\right) = j\ln\left(1 - \frac{2j}{k+j}\right)$$

$$\approx -\frac{2j^{2}}{k+j}$$

$$\left(1 - \frac{2j}{k+j}\right)^{j} \approx e^{-\frac{2j^{2}}{k}}$$

Setting i = 2j as required, we get

$$P_h[X=i] \approx \frac{1}{\sqrt{\pi k}} e^{-\frac{i^2}{2k}}$$
$$= \frac{1}{\sqrt{\pi k}} e^{-\frac{i^2}{\hbar}}$$

As required.

3 Fitting Data

3.1 Task A

$$\mu_1 = \frac{\sum x_i}{N} = 6.50$$

$$\mu_2 = \frac{\sum x_i^2}{N} = 46.55$$

3.2 Task B

The histogram of the data is shown in Figure 5.

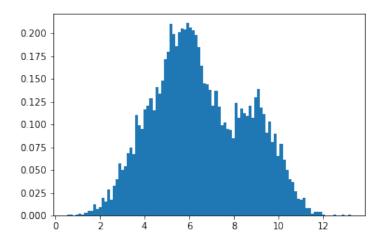


Figure 5: Histogram of the data

3.3 Task C

The binomial distribution is given by

$$\begin{split} \mu_1^{Bin} &= \sum_{r=0}^n \binom{n}{r} p^r (1-p)^{n-r} r \\ &= np \\ \mu_2^{Bin} &= \sum_{r=0}^n \binom{n}{r} p^r (1-p)^{n-r} r^2 \\ &= np \sum_{r=1}^n \binom{n-1}{r-1} p^{r-1} (1-p)^{n-r} r \\ &= np \sum_{r=0}^{n-1} \binom{n-1}{r} p^r (1-p)^{n-r-1} (r+1) \\ &= n(n-1)p^2 + np \sum_{r=0}^{n-1} \binom{n-1}{r} p^r (1-p)^{n-r-1} \\ &= n(n-1)p^2 + np \end{split}$$

The values of μ_1^{Bin} and μ_2^{Bin} are 6.50 and 46.55 respectively. These values are the same as the values of μ_1 and μ_2 calculated in Task A.

When the values of n and p are fit the data and the integers next to n are taken, the values of n and p are 20 and 0.33 respectively.

The histogram of the data and the binomial distribution with n = 20 and p = 0.33 is shown in Figure 6.

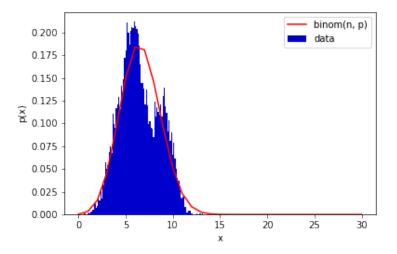


Figure 6: Histogram of the data and the binomial distribution with n=20 and p=0.33

3.4 Task D

The gamma distribution is given by

$$\begin{split} \mu_1^{Gamma} &= \int_0^\infty f(x;k,\theta) x dx \\ &= \int_0^\infty \frac{1}{\theta^k \Gamma(k)} x^{k-1} e^{-\frac{x}{\theta}} x dx \\ &= \frac{\theta^{k+1} \Gamma(k+1)}{\theta^k \Gamma(k)} \\ &= k\theta \end{split}$$

$$\begin{split} \mu_2^{Gamma} &= \int_0^\infty f(x;k,\theta) x^2 dx \\ &= \int_0^\infty \frac{1}{\theta^k \Gamma(k)} x^k e^{-\frac{x}{\theta}} x dx \\ &= \frac{\theta^{k+2} \Gamma(k+2)}{\theta^k \Gamma(k)} \\ &= k(k+1)\theta^2 \end{split}$$

The values of μ_1^{Gamma} and μ_2^{Gamma} are 6.50 and 46.55 respectively. These values are the same as the values of μ_1 and μ_2 calculated in Task A.

When the values of k and θ are fit the data, the values of k and θ are 9.70 and 0.67 respectively.

The histogram of the data and the gamma distribution with k=9.70 and $\theta=0.67$ is shown in Figure 7.

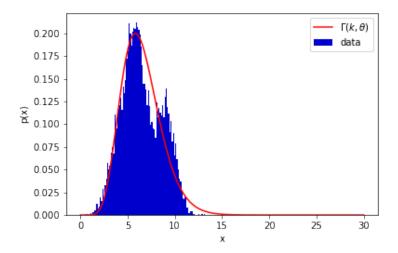


Figure 7: Histogram of the data and the gamma distribution with k = 9.70 and $\theta = 0.67$

3.5 Task E

Definition: Likelihood

Given a dataset S and a choice of parameter $\lambda = \lambda_0$ for a family of distributions $P[\lambda]$ parameterized by λ , define the likelihood of λ_0 by

$$\mathcal{L}(\lambda_0 \mid S) = P_{\lambda_0}[S] = \prod_{i=1}^{N} P_{\lambda_0}[x_i]$$

For large datasets, since each probability is a small number, the likelihood can underflow. Thus, the average log-likelihood is typically calculated:

$$\ell(\theta \mid S) = \frac{\log \mathcal{L}(\theta \mid S)}{N}$$

The average log-likelihood of the data under the binomial and gamma distributions are -2.157 and -2.161 respectively. Binomial distribution has a higher average log-likelihood than the gamma distribution.

3.6 Task F

$$\begin{split} \mu_1^{gmm} &= p_1 \mu_1 + p_2 \mu_2 \\ \mu_2^{gmm} &= p_1 (\sigma_1^2 + \mu_1^2) + p_2 (\sigma_2^2 + \mu_2^2) \\ \mu_3^{gmm} &= p_1 (\mu_1^3 + 3\mu_1 \sigma_1^2) + p_2 (\mu_2^3 + 3\mu_2 \sigma_2^2) \\ \mu_4^{gmm} &= p_1 (\mu_1^4 + 6\mu_1^2 \sigma_1^2 + 3\sigma_1^4) + p_2 (\mu_2^4 + 6\mu_2^2 \sigma_2^2 + 3\sigma_2^4) \end{split}$$

When the values of μ_1 , p_1 , μ_2 , and p_2 are fit to the data, with $\sigma_1 = \sigma_2 = 1$, the values of are 5.13, 0.61, 8.77, and 0.38 respectively.

The histogram of the data and the GMM with $\mu_1 = 5.13$, $p_1 = 0.61$, $\mu_2 = 8.77$, and $p_2 = 0.38$ is shown in Figure 8.

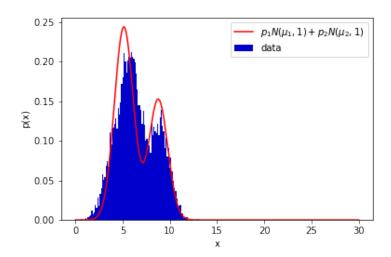


Figure 8: Histogram of the data and the GMM with $\mu_1 = 5.13$, $p_1 = 0.61$, $\mu_2 = 8.77$, and $p_2 = 0.38$

4 Quality in Inequalities

4.1 Task A

$$\begin{split} \mathbb{E}[X] &= \int_0^\infty X P(X) \, dX \\ &= \int_0^a X P(X) \, dX + \int_a^\infty X P(X) \, dX \\ &\geq \int_0^a X P(X) \, dX + \int_a^\infty a P(X) \, dX \\ &\geq \int_0^a X P(X) \, dX + a P[X \geq a] \\ &\geq a P[X \geq a] \end{split}$$

Here, we have used the fact that $\int_0^a XP(X)dX$ is ≥ 0 since $X \geq 0$ and $P(X) \geq 0$ everywhere in that interval.

From the above result, we get

$$\frac{\mathbb{E}[X]}{a} \ge P[X \ge a]$$

Intuitively, this corresponds to saying that if there is a high probability that $X \geq a$, then we can also expect that a random sample will give us a value greater than a.

4.2 Task B

First, we will prove a result which is useful later

$$\mathbb{E}\left[\left(X - \mathbb{E}[X] + \frac{\sigma^2}{\tau}\right)^2\right] = Var(X - \mathbb{E}[X] + \frac{\sigma^2}{\tau}) + \left(E[X - \mathbb{E}[X] + \frac{\sigma^2}{\tau}]\right)^2$$

Now, since the variance of a distribution does not change on adding a constant to the distribution, and since the expectation value of the sum of terms is the sum of expectations of the terms, we get that this is equal to

$$\mathbb{E}\left[\left(X - \mathbb{E}[X] + \frac{\sigma^2}{\tau}\right)^2\right] = \sigma^2 + \left(\frac{\sigma^2}{\tau}\right)^2$$

Now, we notice that whenever $X - \mu \ge \tau$, it also holds that $\left(X - \mu + \frac{\sigma^2}{\tau}\right)^2 \ge \left(\tau + \frac{\sigma^2}{\tau}\right)^2$. Since we are assuming that $x - \mu + \frac{\sigma^2}{\tau} \ge \tau + \frac{\sigma^2}{\tau} \ge 0$, this result is proven by the fact that x^2 is an increasing function over the non-negative reals.

Thus, $P[X - \mu \ge \tau] \le P[\left(X - \mu + \frac{\sigma^2}{\tau}\right)^2 \ge \left(\tau + \frac{\sigma^2}{\tau}\right)^2]$. Using Markov's inequality on the above result, we get

$$P[X - \mu \ge \tau] \le P\left[\left(X - \mu + \frac{\sigma^2}{\tau}\right)^2 \le \left(\tau + \frac{\sigma^2}{\tau}\right)^2\right] \le \frac{\sigma^2 + \left(\frac{\sigma^2}{\tau}\right)^2}{\left(\tau + \frac{\sigma^2}{\tau}\right)^2}$$
$$= \frac{\sigma^2}{\sigma^2 + \tau^2}$$

Thus proving our desired result.

4.3 Task C

$$\begin{split} P[X \geq x] &= \int_x^\infty P(X) \, dX \\ e^{-tx} E[e^{tx}] &= e^{-tx} \int_{-\infty}^\infty e^{tX} P(X) \, dX \\ &= e^{-tx} \int_{-\infty}^x e^{tX} P(X) \, dX + e^{-tx} \int_x^\infty e^{tX} P(X) \, dX \\ &\geq e^{-tx} \int_x^\infty e^{tX} P(X) \, dX \qquad \qquad \text{(since both P(X) and } e^{tX} \text{ are positive for all X)} \\ &\geq e^{-tx} e^{tx} \int_x^\infty P(X) \, dX \qquad \qquad \text{(since } e^{tX} \text{ is increasing if t$;0$)} \\ &= e^{-tx} e^{tx} P[X \geq x] \\ &= P[X > x] \end{split}$$

This proves that

$$P[X \ge x] \le e^{-tx} E[e^{tx}] \tag{1}$$

Similarly,

$$P[X \le x] = \int_{-\infty}^{x} P(X) \, dX$$

$$\begin{split} e^{-tx}E[e^{tx}] &= e^{-tx} \int_{-\infty}^{\infty} e^{tX} P(X) \, dX \\ &= e^{-tx} \int_{-\infty}^{x} e^{tX} P(X) \, dX + e^{-tx} \int_{x}^{\infty} e^{tX} P(X) \, dX \\ &\geq e^{-tx} \int_{-\infty}^{x} e^{tX} P(X) \, dX \qquad \qquad \text{(since both P(X) and } e^{tX} \text{ are positive for all X)} \\ &\geq e^{-tx} \cdot e^{tx} \cdot \int_{-\infty}^{x} P(X) \, dX \qquad \qquad \text{(since } e^{tx} \leq e^{tx'} \text{ where } x \geq x' \text{ when t is negative)} \\ &= \int_{-\infty}^{x} P(X) \, dX \\ &= P[X \leq x] \end{split}$$

as required.

4.4 Task D

1. We know, by the linearity of expectation, that

$$\mathbb{E}[Y] = \sum_{i=1}^{n} \mathbb{E}[X_i] = \sum_{i=1}^{n} p_i$$

2. For this question, we make use of the first result proved in Task C above1.

$$P[Y \ge (1+\delta)\mu] \le \frac{\mathbb{E}[e^{tY}]}{e^{(1+\delta)\mu t}}$$
$$= \frac{\mathbb{E}[e^{\left(t\sum_{i=1}^{n} X_i\right)}]}{e^{(1+\delta)\mu t}}$$

options We will now prove that for independent random variables, $\mathbb{E}[e^{\left(t\sum_{i=1}^{n}X_{i}\right)}] = \prod_{i=1}^{n}\mathbb{E}[e^{tX_{i}}]$.

To prove that, we prove the simpler result that for independent variables, $\mathbb{E}[XY] = \mathbb{E}[X]\mathbb{E}[Y]$.

$$\mathbb{E}[XY] = \int_{-\infty}^{\infty} (xy)p_{X,Y}(x,y) \, dy \, dx$$

$$= \int_{-\infty}^{\infty} (xy)p_X(x)p_Y(y) \, dy \, dx \qquad \text{(For independent RVs, p(xy)=p(x)p(y))}$$

$$= \int_{-\infty}^{\infty} xp_X(x) \int_{-\infty}^{\infty} yp_Y(y) \, dy \, dx$$

$$= \int_{-\infty}^{\infty} xp_X(x) \, dx \int_{-\infty}^{\infty} yp_Y(y) \, dy$$

$$= \mathbb{E}[X]\mathbb{E}[Y]$$

Since above, all X_i 's are independent, so are e^{tX_i} for all i. Thus,

$$E[e^{\left(t\sum_{i=1}^{n}X_{i}\right)}] = \mathbb{E}[e^{tX_{1}}]E[e^{\left(t\sum_{i=2}^{n}X_{i}\right)}]$$

$$= \mathbb{E}[e^{tX_{1}}]\mathbb{E}[e^{tX_{2}}]\mathbb{E}[e^{\left(t\sum_{i=3}^{n}X_{i}\right)}]$$

$$\cdots$$

$$= \prod_{i=1}^{n}\mathbb{E}[e^{tX_{i}}]$$

For each Bernoulli random variable X_i , it is easily observed that $E[e^{tX_i}] = 1 + (e^t - 1)p_i$. Using this result,

$$E[e^{(t\sum_{i=1}^{n} X_i)}] = \prod_{i=1}^{n} (1 + (e^t - 1)p_i)$$

We will now prove $\prod_{i=1}^{n} (1+x_i) \leq e^{\sum_{i=1}^{n} x_i}$ for positive x_i 's.

$$\ln\left(\prod_{i=1}^{n}(1+x_i)\right) = \sum_{i=1}^{n}\ln(1+x_i)$$

$$\leq \sum_{i=1}^{n}x_i \qquad \text{(Well known result that } \ln(1+x) \leq x\text{)}$$

$$= \ln\left(e^{\sum_{i=1}^{n}x_i}\right)$$

Since our initial assumption in using 1 required t > 0 and hence $e^t > 1$, all $(e^t - 1)p_i$ values are positive and we can use the above result. Hence, $\mathbb{E}[e^{tY}] \leq e^{\mu(e^t - 1)}$. Plugging this into ??, we get

$$P[Y \ge (1+\delta)\mu] \le \frac{e^{\mu(e^t-1)}}{e^{(1+\delta)t\mu}}$$

as required.

3. Since this holds for arbitrary positive t, we can find the t for which the value of $e^{\mu(e^t-1-(1+\delta)t)}$ is minimum. By differentiation, we observe the minima of the function is when $t = \ln(1+\delta)$. At the minima, we have

$$e^{t} - 1 - (1 + \delta)t = \delta - (1 + \delta)\ln(1 + \delta)$$

Using the expansion $\ln(1+\delta) = \sum_{i=1}^{\infty} \frac{\delta^{i}(-1)^{i-1}}{i}$

$$\delta - (1 + \delta) \ln(1 + \delta) = \delta + (1 + \delta) \left(\sum_{i=1}^{\infty} \frac{\delta^{i} (-1)^{i}}{i} \right)$$

$$= \delta + \sum_{i=1}^{\infty} \frac{\delta^{i} (-1)^{i}}{i} + \sum_{i=2}^{\infty} \frac{\delta^{i} (-1)^{i-1}}{i - 1}$$

$$= \sum_{i=2}^{\infty} \frac{\delta^{i} (-1)^{i-1}}{i(i - 1)}$$

This is approximately $-\frac{\delta^2}{2} + \frac{\delta^3}{6}$ for small delta. Thus we get a good upper bound of $e^{\frac{-\mu\delta^2(3-\delta)}{6}}$ for $P[Y \ge (1+\delta)\mu]$.

If $\delta > 0$, this can be simplified further, in a way that is more useful for us in question 5. Observe that for x > 0

$$\ln(1+x) \ge \frac{2x}{2+x}$$

This can be seen by observing that at x=0, $\ln(1+x) - \frac{2x}{2+x}$ evaluates to 0 and that its derivative, $\frac{4}{(1+x)(2+x)^2}$ is always non-negative. Thus,

$$\mu(\delta - (1+\delta)\ln(1+\delta)) \le \mu(\delta - \frac{2\delta(1+\delta)}{2+\delta})$$
$$= \frac{-\mu\delta^2}{2+\delta}$$

Thus, we get,

$$P[Y \ge (1+\delta)\mu] \le e^{\frac{-\mu\delta^2}{2+\delta}} \tag{2}$$

For the last few steps of simplification, https://math.mit.edu/ goemans/18310S15/chernoffnotes.pdf was used as a guide.

4.5 Task E

First, let us consider the corner cases where $\mu = 0$ or $\mu = 1$. In these cases, $A_n = 0$ or $A_n = 1$ always, and the result is trivial.

Now we can consider $0 < \mu < 1$.

Consider the random variable $B_n = nA_n$.

We need to prove $\lim_{n\to\infty} P[|B_n - n\mu| > n\epsilon] = 0$.

Note that we are done if we prove $\lim_{n\to\infty} P[B_n - n\mu > = n\epsilon] = 0$ and $\lim_{n\to\infty} P[n\mu - B_n > = n\epsilon] = 0$. We will prove the first part first, assuming μ is not 0.

$$P[B_n \ge n\epsilon + n\mu] \le e^{-\frac{n\epsilon^2}{2\mu + \epsilon}}$$

For a positive μ (since the distribution is the sum of Bernoulli RVs) and ϵ (given), $\frac{\epsilon^2}{2\mu+\epsilon}$ is a positive coefficient. e^{-n} tends to 0 as n tends to ∞ , which we will use henceforth.

Note that the result that $P[X \ge (1+\delta)\mu] \le e^{\frac{-\mu\delta^2}{2+\delta}}$ was not proven without referring to the internet, so I will prove the claim above using an alternative method which assumes only what has been proven by me.

$$P[X > (1+\delta)\mu] < e^{\mu(\delta - (1+\delta)\ln(1+\delta))}$$

We want to prove that $f(\delta) = \delta - (1 + \delta) \ln(1 + \delta)$ is strictly negative if δ is positive. For that, notice that f(0) = 0 and $f'(\delta) = -\ln(1 + \delta)$, which is non-positive for all non-negative δ . Thus, $f(\delta)$ is strictly negative for all positive δ .

This again means that $P[B_n \ge n\epsilon + n\mu] \le e^{-nk}$ for some positive coefficient k. Thus, $P[B_n \ge n\epsilon + n\mu]$ tends to 0 as n tends to ∞ .

Similarly, $P[B_n \le n\mu - n\epsilon]$ also tends to 0 as n tends to ∞ .

We can prove this by considering the RV $1-A_n$ instead of A_n and using the same logic as above. $P[n-B_n-n-n\mu\geq n\delta]=P[n\mu-nB_n\geq n\delta]\leq e^{-nk}$ for some positive coefficient k. Now, combining the two results,

$$\begin{split} P[|B_n - n\mu| > n\epsilon] &\leq P[B_n - n\mu \geq n\epsilon] + P[n\mu - B_n \geq n\epsilon] \\ &\leq e^{-nk} + e^{-nk'} \end{split} \qquad \text{(for some positive coefficients k and k')} \end{split}$$

If we get any positive value σ for the value of $P[|B_n - n\mu| > n\epsilon]$, simply set $n > \max(\frac{\ln(\frac{2}{\sigma})}{k}, \frac{\ln(\frac{2}{\sigma})}{k'})$ and we have

$$P[|B_n - n\mu| > n\epsilon] \le e^{-nk} + e^{-nk'}$$

$$< \frac{\sigma}{2} + \frac{\sigma}{2}$$

$$= \sigma$$

which is a contradiction.

Thus, we have proven that $\lim_{n\to\infty} P[|B_n - n\mu| > n\epsilon] = \lim_{n\to\infty} P[|A_n - \mu| > \epsilon] = 0$.

5 A Pretty "Normal" Mixture

5.1 Task A

Let A_i be the event in which you choose i in the first step of the algorithm. Then the set of all the A_i s is a valid partition, since they are completely disjoint and their union represents all the possibilities.

Hence we can use the law of Total Probability to get,

$$P[\mathcal{A} = x] = \sum_{i=1}^{K} P[A_i] \cdot P[X_i = x]$$
$$= \sum_{i=1}^{K} p_i \cdot P[X_i = x]$$
$$= P[X = x]$$
$$\implies f_{\mathcal{A}}(x) = f_X(x)$$

5.2 Task B

5.2.1 Part 1

$$\mathbb{E}[X] = \sum_{x \in \mathbb{Z}} \left(x \cdot \sum_{i=1}^{K} p_i \cdot P[X_i = x] \right)$$
$$= \sum_{i=1}^{K} \left(p_i \cdot \sum_{x \in \mathbb{Z}} x \cdot P[X_i = x] \right)$$
$$= \sum_{i=1}^{K} \left(p_i \cdot \mathbb{E}[X_i] \right) = \sum_{i=1}^{K} \left(p_i \cdot \mu_i \right)$$

5.2.2 Part 2

• Calculating $\mathbb{E}[X^2]$

$$\mathbb{E}[X^2] = \sum_{x \in \mathbb{Z}} \left(x^2 \cdot \sum_{i=1}^K p_i \cdot P[X_i = x] \right)$$

$$= \sum_{i=1}^K \left(p_i \cdot \sum_{x \in \mathbb{Z}} x^2 \cdot P[X_i = x] \right)$$

$$= \sum_{i=1}^K \left(p_i \cdot \mathbb{E}[X_i^2] \right) = \sum_{i=1}^K \left(p_i \cdot \left(Var(X_i) + (\mathbb{E}[X_i])^2 \right) \right)$$

$$= \sum_{i=1}^K \left(p_i \cdot \left(\sigma_i^2 + \mu_i^2 \right) \right)$$

• Calculating $(\mathbb{E}[X])^2$

$$(\mathbb{E}[X])^{2} = \left(\sum_{i=1}^{K} p_{i} \cdot \mathbb{E}[X_{i}]\right)^{2}$$

$$= \sum_{i=1}^{K} (p_{i} \cdot \mathbb{E}[X_{i}])^{2} + 2 \sum_{1 \leq i < j \leq K} (p_{i} \cdot \mathbb{E}[X_{i}] \cdot p_{j} \cdot \mathbb{E}[X_{j}])$$

$$= \sum_{i=1}^{K} (p_{i}\mu_{i})^{2} + 2 \sum_{1 \leq i < j \leq K} (p_{i}\mu_{i}p_{j}\mu_{j})$$

• Calculating Var[X]

$$Var[X] = (\mathbb{E}[X])^{2} - \mathbb{E}[X^{2}]$$

$$= \left(\sum_{i=1}^{K} (p_{i}\mu_{i})^{2} + 2\sum_{1 \leq i < j \leq K} (p_{i}\mu_{i}p_{j}\mu_{j})\right) - \sum_{i=1}^{K} (p_{i} \cdot (\sigma_{i}^{2} + \mu_{i}^{2}))$$

5.2.3 Part 3

$$M_X(t) = G(e^t) = \int_{-\infty}^{\infty} e^{xt} P[X = x] dx = \int_{-\infty}^{\infty} e^{xt} \sum_{i=1}^{K} p_i P[X_i = x] dx$$

$$= \int_{-\infty}^{\infty} \sum_{i=1}^{K} e^{xt} p_i P[X_i = x] dx = \sum_{i=1}^{K} p_i \int_{-\infty}^{\infty} e^{xt} P[X_i = x] dx$$

$$= \sum_{i=1}^{K} p_i M_{X_i} = \sum_{i=1}^{K} p_i e^{\mu_i t + \frac{1}{2}\sigma_i^2 t^2}$$

- 5.3 Task C
- 5.3.1 Part 1

$$\mathbb{E}[Z] = \mathbb{E}\left[\sum_{i=1}^{K} p_i X_i\right] = \sum_{i=1}^{K} p_i \cdot \mathbb{E}\left[X_i\right] = \sum_{i=1}^{K} (p_i \mu_i)$$

5.3.2 Part 2

$$Var(Z) = Var\left(\sum_{i=1}^{K} p_i X_i\right) = \sum_{i=1}^{K} Var\left(p_i X_i\right) = \sum_{i=1}^{K} p_i^2 Var\left(X_i\right) = \sum_{i=1}^{K} p_i^2 \sigma_i^2$$

5.3.3 Part 3

Since we already calculated μ_Z and $\sigma_Z = \sqrt{Var(Z)}$ We can just use the expression for Gaussian Random Variable (See Part 6 for the proof that Z is a Gaussian random variable)

$$f_Z(u) = \frac{1}{\sigma_Z \sqrt{2\pi}} e^{\frac{-(u-\mu_Z)^2}{2\sigma_Z^2}}$$

5.3.4 Part 4

Using the expression for Moment generating function of a Gaussian Random Variable (See Part 6 for the proof that Z is a gaussian random variable)

$$M_Z(t) = e^{\mu_Z t + \frac{1}{2}\sigma_Z^2 t^2}$$

5.3.5 Part 5

No, It can be seen that even though the expected value is the same in both the case, The values of both variance and MGF are different for both X and Z hence these are different random variables

5.3.6 Part 6

We can show that linear combination of independent random variables is also a random variable. For any random Gaussian variable X is given by $M_X(t) = \mathbb{E}[e^{tX}] = e^{\mu t + \frac{1}{2}\sigma^2 t^2}$ Let $Y = \sum_{i=1}^K a_i X_i$ where each X_i is an independent Gaussian random variable. Then

$$M_Y(t) = \mathbb{E}[e^{tY}] = \mathbb{E}\left[e^{t\sum_{i=1}^K a_i X_i}\right] = \mathbb{E}\left[\prod_{i=1}^K e^{ta_i X_i}\right]$$

$$= \prod_{i=1}^K \mathbb{E}\left[e^{ta_i X_i}\right] \qquad \text{,Since all the } X_i \text{ are mutually independent}$$

$$= \prod_{i=1}^K M_{a_i X_i}(t) = \prod_{i=1}^K \left(e^{a_i \mu_i t + \frac{1}{2}(a_i \sigma_i)^2 t^2}\right) = \exp\left(\sum_{i=1}^K (a_i \mu_i) t + \frac{1}{2}\sum_{i=1}^K ((a_i \sigma_i)^2) t^2\right)$$

If we define
$$\mu_y = \sum_{i=1}^K (a_i \mu_i)$$
 and $\sigma_y^2 = \sum_{i=1}^K ((a_i \sigma_i)^2)$ Then we get
$$M_y(t) = e^{\mu_y t + \frac{1}{2} \sigma_y^2 t^2}$$
(3)

This is just the Moment generating function of a Gaussian random variable, Hence Y is a Gaussian random variable. Hence Z is also a Gaussian random variable.

5.4 Task D

To prove that a PDF uniquely determines the MGF, we simply use the construction of the MGF. The MGF is defined as

$$M_X(t) = \mathbb{E}[e^{tX}] = \int_{-\infty}^{\infty} e^{tx} p_X(x) dx = \sum_{i=1}^{n} e^{tx_i} p(x_i)$$

where the support of the PDF is $\{x_1, \dots x_n\}$.

To prove the other direction, let us assume that we can find two distinct solutions $\{x_1, \dots x_n\}$ and $\{y_1, \dots y_m\}$ from an arbitrary MGF which we know is of a finite discrete PDF. Then, by the above result, we know that

$$\sum_{i=1}^{n} e^{tx_i} p(x_i) = \sum_{j=1}^{m} e^{ty_j} p(y_j)$$

If any of the x_i 's are not equal to any of the y_j 's, and given that all $p(x_i)$'s and $p(y_j)$'s are non-zero, then we can rearrange the terms to get that for some a_i 's

$$e^{a_0}tb_0 = \sum_{i=1}^k e^{a_i}tb_i$$

for some non-zero coefficients b_0 and for some k, and all a_i 's being distinct. Now, note that the LHS has an annihilating polynomial $D - a_0$, where D is the derivative operator w.r.t t. The RHS has an annihitating polynomial $(D - a_1) \cdots (D - a_k)$. Since the LHS equals the RHS, if an operator makes the LHS 0, it also makes the RHS zero. Thus, the LHS $(e^{a_0t}b_0)$ is also annihilated by the annihilator of the RHS.

$$(D - a_1) \cdots (D - a_k)(e^{a_0 t} b_0) = 0$$

$$\implies (D - a_1) \cdots (D - a_{k-1})((a_0 - a_k)e^{a_0 t} b_0) = 0$$

$$\implies (D - a_1) \cdots (D - a_{k-2})(a_0 - a_{k-1})(a_0 - a_k)e^{a_0 t} b_0 = 0$$

$$\cdots$$

$$\implies (a_0 - a_1) \cdots (a_0 - a_k)e^{a_0 t} b_0 = 0$$

Since e^{a_0t} is never 0, and we know $b_0 \neq 0$, $(a_0 - a_1) \cdots (a_0 - a_k) = 0$. But this is only possible if $a_0 = a_i$ for some $i \in \{1, \dots, k\}$. But that contradicts our assumption that the a_i 's are distinct.

Thus, it is not possible for two PDFs of finite discrete RVs to correspond to the same MGF.