DAY-8

1) Implement Floyd's Algorithm to find the shortest path between all pairs of cities. Display the distance matrix before and after applying the algorithm. Identify and print the shortest

```
Input: n = 4, edges = [[0,1,3],[1,2,1],[1,3,4],[2,3,1]], distanceThreshold = 4
Output: 3
The neighboring cities at a distanceThreshold = 4 for each city are:
City 0 -> [City 1, City 2]
City 1 -> [City 0, City 2, City 3]
City 2 -> [City 0, City 1, City 3]
City 3 -> [City 1, City 2]
Cities 0 and 3 have 2 neighboring cities at a distanceThreshold = 4, but we have to return
city 3 since it has the greatest number.
```

CODE:

```
import sys
def floyd warshall(n, edges, distanceThreshold):
  # Initialize the distance matrix
  dist = [[sys.maxsize] * n for in range(n)]
     for i in range(n):
     dist[i][i] = 0
  for edge in edges:
     u, v, w = edge
     dist[u][v] = w
     dist[v][u] = w # Since the graph is undirected
  print("Distance matrix before applying Floyd's algorithm:")
  print matrix(dist)
  for k in range(n):
     for i in range(n):
       for j in range(n):
          if dist[i][k] != sys.maxsize and dist[k][j] != sys.maxsize:
             dist[i][j] = min(dist[i][j], dist[i][k] + dist[k][j])
  print("\nDistance matrix after applying Floyd's algorithm:")
  print matrix(dist)
  neighboring cities = []
  for i in range(n):
     count = 0
```

```
for j in range(n):
       if dist[i][j] <= distanceThreshold and i != j:
          count += 1
     neighboring cities.append((i, count))
     city with max neighbors = \max(\text{neighboring cities, key=lambda x: } (x[1], -x[0]))[0]
  print("\nCity with the most neighbors within distance threshold =", distanceThreshold, ":",
city_with_max_neighbors)
 return city_with_max_neighbors
def print_matrix(matrix):
  for row in matrix:
     print(row)
n = 4
edges = [[0, 1, 3], [1, 2, 1], [1, 3, 4], [2, 3, 1]]
distanceThreshold = 4
result = floyd warshall(n, edges, distanceThreshold)
print("Output:", result)
```

Distance matrix after applying Floyd's algorithm:

[0, 3, 4, 5]

[3, 0, 1, 2]

[4, 1, 0, 1]

[5, 2, 1, 0]

City with the most neighbors within distance threshold = 4:3

Output: 3

2) Write a Program to implement Floyd's Algorithm to calculate the shortest paths between all pairs of routers. Simulate a change where the link between Router B and Router D fails. Update the distance matrix accordingly. Display the shortest path from Router A to Router F before and after the link failure.

Input as above

Output: Router A to Router F = 5

CODE:

```
import sys
def floyd warshall(n, edges):
  # Initialize the distance matrix
  dist = [[sys.maxsize] * n for _ in range(n)]
     for i in range(n):
     dist[i][i] = 0
     for edge in edges:
     u, v, w = edge
     dist[u][v] = w
     dist[v][u] = w
  print("Distance matrix before applying Floyd's algorithm:")
  print matrix(dist)
  for k in range(n):
     for i in range(n):
       for j in range(n):
          if dist[i][k] != sys.maxsize and dist[k][j] != sys.maxsize:
             dist[i][j] = min(dist[i][j], dist[i][k] + dist[k][j])
  return dist
def simulate link failure(dist, routerB, routerD):
  dist[routerB][routerD] = sys.maxsize
  dist[routerD][routerB] = sys.maxsize
  print("\nSimulated link failure between Router B and Router D.")
def print matrix(matrix):
  for row in matrix:
     print(row)
def find shortest path(dist, routerA, routerF):
  if dist[routerA][routerF] == sys.maxsize:
```

```
return "No path available"
  return dist[routerA][routerF]
n = 6
edges = [
  [0, 1, 2],
  [0, 2, 4],
  [1, 2, 1],
  [1, 3, 7],
  [2, 4, 3],
  [3, 4, 2],
  [3, 5, 1]
  [4, 5, 5]
]router A = 0
routerF = 5
routerB = 1
routerD = 3
dist = floyd warshall(n, edges)
print("\nShortest path from Router A to Router F before link failure:")
shortest path before = find shortest path(dist, routerA, routerF)
print("Router A to Router F =", shortest path before)
simulate link failure(dist, routerB, routerD)
for k in range(n):
  for i in range(n):
     for j in range(n):
       if dist[i][k] != sys.maxsize and dist[k][j] != sys.maxsize:
          dist[i][j] = min(dist[i][j], dist[i][k] + dist[k][j])
print("\nShortest path from Router A to Router F after link failure:")
shortest path after = find shortest path(dist, routerA, routerF)
print("Router A to Router F =", shortest path after)
```

Distance matrix before applying Floyd's algorithm:

[0, 2, 4, 9, 7, 10]

[2, 0, 1, 7, 4, 8]

[4, 1, 0, 6, 3, 7]

[9, 7, 6, 0, 2, 1]

[7, 4, 3, 2, 0, 5]

[10, 8, 7, 1, 5, 0]

Shortest path from Router A to Router F before link failure:

Router A to Router F = 5

Simulated link failure between Router B and Router D.

Shortest path from Router A to Router F after link failure:

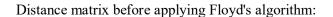
Router A to Router F = 5

the distance matrix before and after applying the algorithm. Identify and print the shortest Input: n = 5, edges = [[0,1,2],[0,4,8],[1,2,3],[1,4,2],[2,3,1],[3,4,1]], distanceThreshold = 2 Output: 0 Explanation: The figure above describes the graph. The neighboring cities at a distanceThreshold = 2 for each city are: City 0 -> [City 1] City 1 -> [City 0, City 4] City 2 -> [City 3, City 4] City 3 -> [City 2, City 4] City 4 -> [City 1, City 2, City 3] The city 0 has 1 neighboring city at a distance Threshold = 2. **CODE:** import sys def floyd warshall(n, edges): dist = [[sys.maxsize] * n for in range(n)]for i in range(n): dist[i][i] = 0for edge in edges: u, v, w = edgedist[u][v] = wdist[v][u] = w # The graph is undirectedprint("Distance matrix before applying Floyd's algorithm:") print matrix(dist) for k in range(n): for i in range(n): for j in range(n): if dist[i][k] != sys.maxsize and dist[k][i] != sys.maxsize: dist[i][j] = min(dist[i][j], dist[i][k] + dist[k][j])return dist def print matrix(matrix): for row in matrix: print(row) def count neighbors(dist, distanceThreshold): neighbors count = [0] * len(dist)

for i in range(len(dist)):

3) Implement Floyd's Algorithm to find the shortest path between all pairs of cities. Display

```
for j in range(len(dist)):
       if i != j and dist[i][j] <= distanceThreshold:
          neighbors count[i] += 1
  return neighbors count
def find city with fewest neighbors(neighbors count):
  min neighbors = sys.maxsize
  city with min neighbors = -1
  for i, count in enumerate(neighbors count):
     if count < min neighbors:
       min neighbors = count
       city with min neighbors = i
  return city with min neighbors
n = 5
edges = [
  [0, 1, 2], # City 0 -> City 1
  [0, 4, 8], # City 0 -> City 4
  [1, 2, 3], # City 1 -> City 2
  [1, 4, 2], # City 1 -> City 4
  [2, 3, 1], # City 2 -> City 3
  [3, 4, 1] # City 3 -> City 4
1
distanceThreshold = 2
dist = floyd warshall(n, edges)
print("\nDistance matrix after applying Floyd's algorithm:")
print matrix(dist)
neighbors count = count neighbors(dist, distanceThreshold)
print("\nNumber of neighboring cities within distance threshold:")
for i in range(n):
  print(f"City {i} -> {neighbors count[i]} neighbors")
city with min neighbors = find city with fewest neighbors(neighbors count)
print(f"\nCity with the fewest neighboring cities within the distance threshold: City
{city with min neighbors}")
```



[0, 2, inf, inf, 8]

[2, 0, 3, inf, 2]

[inf, 3, 0, 1, inf]

[inf, inf, 1, 0, 1]

[8, 2, inf, 1, 0]

Distance matrix after applying Floyd's algorithm:

[0, 2, 5, 6, 4]

[2, 0, 3, 4, 2]

[5, 3, 0, 1, 2]

[6, 4, 1, 0, 1]

[4, 2, 2, 1, 0]

Number of neighboring cities within distance threshold:

City 0 -> 1 neighbors

City 1 -> 2 neighbors

City 2 -> 2 neighbors

City 3 -> 2 neighbors

City 4 -> 3 neighbors

City with the fewest neighboring cities within the distance threshold: City 0

4) Implement the Optimal Binary Search Tree algorithm for the keys A,B,C,D with frequencies 0.1,0.2,0.4,0.3 Write the code using any programming language to construct the OBST for the given keys and frequencies. Execute your code and display the resulting OBST and its cost. Print the cost and root matrix.

```
Input N =4, Keys = \{A,B,C,D\} Frequencies = \{01.02.,0.3,0.4\}
Output : 1.7
Cost Table
01234
1 0 0.1 0.4 1.1 1.7
2 0 0.2 0.8 0.4
3 0 0.4 1.0
4 0 0.3
50
Root table
1234
11233
2233
333
44
CODE:
import sys
def optimal bst(keys, freq, n):
  # Initialize the cost and root tables
  cost = [[0 for in range(n)] for in range(n)]
  root = [[0 for in range(n)] for in range(n)]
     for i in range(n):
     cost[i][i] = freq[i]
  for L in range(2, n + 1): # L is the chain length
     for i in range(n - L + 1):
       j = i + L - 1
       cost[i][j] = sys.maxsize
       sum freq = sum(freq[i:j+1]) # Sum of frequencies from i to j
       for r in range(i, j + 1):
          # Calculate cost when r is the root
          c = (cost[i][r - 1] if r > i else 0) + (cost[r + 1][j] if r < j else 0) + sum freq
          if c < cost[i][j]:
             cost[i][j] = c
            root[i][j] = r
  return cost, root
```

```
def print_matrix(matrix, name):
    print(f"\n{name} Table:")
    for row in matrix:
        print(row)
    keys = ['A', 'B', 'C', 'D']
    freq = [0.1, 0.2, 0.4, 0.3]
    n = len(keys)
    cost, root = optimal_bst(keys, freq, n)
    print_matrix(cost, "Cost")
    print_matrix(root, "Root")
    print(f"\nThe minimum cost of the OBST is: {cost[0][n-1]}")
```

Cost Table:

[0.1, 0.4, 1.1, 1.7]

[0, 0.2, 0.8, 1.4]

[0, 0, 0.4, 1.0]

[0, 0, 0, 0.3]

Root Table:

[0, 1, 2, 2]

[0, 1, 2, 2]

[0, 0, 2, 3]

[0, 0, 0, 3]

The minimum cost of the OBST is: 1.7

```
5) Consider a set of keys 10,12,16,21 with frequencies 4,2,6,3 and the respective
probabilities. Write a Program to construct an OBST in a programming language of your
choice. Execute your code and display the resulting OBST, its cost and root matrix.
Input N = 4, Keys = \{10,12,16,21\} Frequencies = \{4,2,6,3\}
Output: 26
0123
0 4 80 202 262
1 2 102 162
2612
33
a) Test cases
Input: keys[] = \{10, 12\}, freq[] = \{34, 50\}
Output = 118
b) Input: keys[] = \{10, 12, 20\}, freq[] = \{34, 8, 50\}
Output = 142
CODE:
import sys
def optimal_bst(keys, freq, n):
  # Initialize the cost and root tables
  cost = [[0 \text{ for } in range(n)] \text{ for } in range(n)]
  root = [[0 for in range(n)] for in range(n)]
     for i in range(n):
     cost[i][i] = freq[i]
     for L in range(2, n + 1): # L is the chain length
     for i in range(n - L + 1):
       i = i + L - 1
       cost[i][j] = sys.maxsize
       sum freq = sum(freq[i:j+1]) # Sum of frequencies from i to j
               for r in range(i, j + 1):
          c = (cost[i][r-1] if r > i else 0) + (cost[r+1][j] if r < j else 0) + sum freq
          if c < cost[i][j]:
            cost[i][j] = c
            root[i][j] = r
  return cost, root
def print matrix(matrix, name):
  print(f"\n{name} Table:")
```

for row in matrix:

```
print(row)
keys = [10, 12, 16, 21]
freq = [4, 2, 6, 3]
n = len(keys)

cost, root = optimal_bst(keys, freq, n)
print_matrix(cost, "Cost")
print_matrix(root, "Root")
print(f"\nThe minimum cost of the OBST is: {cost[0][n-1]}")
```

Cost Table:

[4, 10, 26, 46]

[0, 2, 14, 28]

[0, 0, 6, 15]

[0, 0, 0, 3]

Root Table:

[0, 0, 2, 2]

[0, 1, 2, 2]

[0, 0, 2, 3]

[0, 0, 0, 3]

The minimum cost of the OBST is: 26

6) A game on an undirected graph is played by two players, Mouse and Cat, who alternate turns. The graph is given as follows: graph[a] is a list of all nodes b such that ab is an edge of the graph. The mouse starts at node 1 and goes first, the cat starts at node 2 and goes second, and there is a hole at node 0. During each player's turn, they must travel along one edge of the graph that meets where they are. For example, if the Mouse is at node 1, it must travel to any node in graph[1]. Additionally, it is not allowed for the Cat to travel to the Hole (node 0). Then, the game can end in three ways:

If ever the Cat occupies the same node as the Mouse, the Cat wins.

If ever the Mouse reaches the Hole, the Mouse wins.

If ever a position is repeated (i.e., the players are in the same position as a previous turn, and it is the same player's turn to move), the game is a draw.

Given a graph, and assuming both players play optimally, return

1 if the mouse wins the game,

2 if the cat wins the game, or

0 if the game is a draw.

Example 1:

```
Input: graph = [[2,5],[3],[0,4,5],[1,4,5],[2,3],[0,2,3]]
```

Output: 0

CODE:

```
from collections import deque
```

```
def catMouseGame(graph):
  n = len(graph)
  dp = [[[0] * 2 \text{ for in range}(n)] \text{ for in range}(n)]
     queue = deque()
  for cat in range(1, n):
     dp[0][cat][0] = 1 # Mouse's turn, Mouse wins
     dp[0][cat][1] = 1 \# Cat's turn, Mouse wins
     queue.append((0, cat, 0))
     queue.append((0, cat, 1))
  for mouse in range(1, n):
     dp[mouse][mouse][0] = 2 \# Mouse's turn, Cat wins
     dp[mouse][mouse][1] = 2 \# Cat's turn, Cat wins
     queue.append((mouse, mouse, 0))
     queue.append((mouse, mouse, 1))
  while queue:
     mouse, cat, turn = queue.popleft()
     result = dp[mouse][cat][turn]
     if turn == 0:
```

for prev cat in graph[cat]:

```
if prev cat == 0:
            continue
         if dp[mouse][prev cat][1] == 0:
            if result == 2
               dp[mouse][prev cat][1] = 2
               queue.append((mouse, prev cat, 1))
            elif all(dp[mouse][next cat][0] == 1 for next cat in graph[mouse]):
               # If every possible move for the Cat leads to Mouse winning
               dp[mouse][prev cat][1] = 1
               queue.append((mouse, prev cat, 1))
else:
       for prev mouse in graph[mouse]:
         if dp[prev mouse][cat][0] == 0:
            # If the game hasn't been decided yet for this state
            if result == 1: # Mouse wins this state
               dp[prev\_mouse][cat][0] = 1
               queue.append((prev mouse, cat, 0))
            elif all(dp[next mouse][cat][1] == 2 for next mouse in graph[prev mouse]):
               dp[prev mouse][cat][0] = 2
               queue.append((prev mouse, cat, 0))
     return dp[1][2][0]
graph = [[2, 5], [3], [0, 4, 5], [1, 4, 5], [2, 3], [0, 2, 3]]
result = catMouseGame(graph)
print(result)
```

7) You are given an undirected weighted graph of n nodes (0-indexed), represented by an edge list where edges[i] = [a, b] is an undirected edge connecting the nodes a and b with a probability of success of traversing that edge succProb[i]. Given two nodes start and end, find the path with the maximum probability of success to go from start to end and return its success probability. If there is no path from start to end, return 0. Your answer will be accepted if it differs from the correct answer by at most 1e-5.

Example 1:

```
Input: n = 3, edges = [[0,1],[1,2],[0,2]], succProb = [0.5,0.5,0.2], start = 0, end = 2
Output: 0.25000
```

CODE:

```
import heapq
def maxProbability(n, edges, succProb, start, end):
  graph = [[] for _ in range(n)]
  for (a, b), prob in zip(edges, succProb):
    graph[a].append((b, prob))
     graph[b].append((a, prob))
  \max \text{ prob} = [0.0] * n
  max prob[start] = 1.0 # Start node has probability 1 to itself
    pq = [(-1.0, start)] # We use -1.0 because heapq is a min-heap, and we want to maximize the
probability
    while pq:
    current prob, node = heapq.heappop(pq)
    current prob = -current prob # Convert back to positive
if node == end:
       return current prob
          for neighbor, edge prob in graph[node]:
       new prob = current prob * edge prob
       if new prob > max prob[neighbor]:
          max prob[neighbor] = new prob
         heapq.heappush(pq, (-new prob, neighbor))
    return 0.0
n = 3
edges = [[0, 1], [1, 2], [0, 2]]
succProb = [0.5, 0.5, 0.2]
start = 0
end = 2
result = maxProbability(n, edges, succProb, start, end)
print(f"Output: {result:.5f}")
```

OUTPUT:

0.25000

8) grid[0][0]). The robot tries to move to the bottom-right corner (i.e., grid[m - 1][n - 1]). The robot can only move either down or right at any point in time. Given the two integers m and n, return the number of possible unique paths that the robot can take to reach the bottom-right corner. The test cases are generated so that the answer will be less than or equal to 2 * 10 9.

Example 1:

START FINISH

Input: m = 3, n = 7

Output: 28

CODE:

```
def uniquePaths(m, n):
    dp = [[1] * n for _ in range(m)]
    for i in range(1, m):
    for j in range(1, n):
        dp[i][j] = dp[i-1][j] + dp[i][j-1]
    return dp[m-1][n-1]

m = 3
n = 7
result = uniquePaths(m, n)
print(f''Output: {result}'')
```

OUTPUT:

1 1 1 1 1 1 1 1 2 3 4 5 6 7

1 3 6 10 15 21 28

9) Given an array of integers nums, return the number of good pairs. A pair (i, j) is called good if nums[i] == nums[j] and i < j.

Example 1:

```
Input: nums = [1,2,3,1,1,3]
```

Output: 4

CODE:

```
def numIdenticalPairs(nums):
    freq = {}
    good_pairs = 0
    for num in nums:
    if num in freq:
        good_pairs += freq[num]
        freq[num] += 1
    else:
        freq[num] = 1
    return good_pairs
nums = [1, 2, 3, 1, 1, 3]
result = numIdenticalPairs(nums)
print(f"Output: {result}")
```

OUTPUT:

4

10) There are n cities numbered from 0 to n-1. Given the array edges where edges[i] = [fromi, toi, weighti] represents a bidirectional and weighted edge between cities fromi and toi, and given the integer distanceThreshold. Return the city with the smallest number of cities that are reachable through some path and whose distance is at most distance Threshold, If there are multiple such cities, return the city with the greatest number. Notice that the distance of a path connecting cities i and j is equal to the sum of the edges' weights along that path. Example 1:

```
Input: n = 4, edges = [[0,1,3],[1,2,1],[1,3,4],[2,3,1]], distanceThreshold = 4
Output: 3
```

CODE:

```
import heapq
def findTheCity(n, edges, distanceThreshold):
  graph = [[] for in range(n)]
  for u, v, w in edges:
    graph[u].append((v, w))
    graph[v].append((u, w))
def dijkstra(start):
    distances = [float('inf')] * n
    distances[start] = 0
    min heap = [(0, start)] # (distance, node)
     while min heap:
       current distance, current node = heapq.heappop(min heap)
       if current distance > distances[current node]:
          continue
       for neighbor, weight in graph[current node]:
          distance = current distance + weight
          if distance < distances[neighbor]:
            distances[neighbor] = distance
 heapq.heappush(min heap, (distance, neighbor))
    return distances
  min reachable count = float('inf')
```

```
city_with_min_reachable = -1
for city in range(n):
    distances = dijkstra(city)
    reachable_count = sum(1 for dist in distances if dist <= distanceThreshold)
    if (reachable_count < min_reachable_count) or (
        reachable_count == min_reachable_count and city > city_with_min_reachable):
        min_reachable_count = reachable_count
        city_with_min_reachable = city
return city_with_min_reachable
```

```
n = 4 edges = [[0, 1, 3], [1, 2, 1], [1, 3, 4], [2, 3, 1]] distanceThreshold = 4 result = findTheCity(n, edges, distanceThreshold) print(f''Output: {result}'')
```

3

11) You are given a network of n nodes, labeled from 1 to n. You are also given times, a list of travel times as directed edges times[i] = (ui, vi, wi), where ui is the source node, vi is the target node, and wi is the time it takes for a signal to travel from source to target. We will send a signal from a given node k. Return the minimum time it takes for all the n nodes to receive the signal. If it is impossible for all the n nodes to receive the signal, return -1. Example 1:

```
Input: times = [[2,1,1],[2,3,1],[3,4,1]], n = 4, k
Output: 2
```

CODE:

```
import heapq
def networkDelayTime(times, n, k):
  # Step 1: Create the graph as an adjacency list
  graph = [[] for in range(n + 1)]
  for u, v, w in times:
     graph[u].append((v, w)) # u -> (v, w)
  distances = [float('inf')] * (n + 1)
  distances[k] = 0
  min heap = [(0, k)] # (time, node)
    while min heap:
    current time, current node = heapq.heappop(min heap)
          if current time > distances[current node]:
       continue
          for neighbor, travel time in graph[current node]:
       new time = current time + travel time
              if new time < distances[neighbor]:
         distances[neighbor] = new time
         heapq.heappush(min heap, (new time, neighbor))
  max time = max(distances[1:]) # Ignore index 0 as nodes are 1-indexed
  return max time if max time != float('inf') else -1
times = [[2, 1, 1], [2, 3, 1], [3, 4, 1]]
n = 4
k = 2
result = networkDelayTime(times, n, k)
print(f"Output: {result}")
```

OUTPUT: