

A Play on Regular Expressions

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ICFP 2010

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- intuitive method for regular expression matching
- automata construction with elegant Haskell implementation
- can be generalized in surprising ways

$((a|b) * c (a|b) * c) * (a|b) *$

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$((a|b) * c (a|b) * c) * (a|b) *$

$((a|b) * c (a|b) * c) * (a|b) *$

$((a|b)^*c(a|b)^*c)^*(a|b)^*$

"abc"

$((a|b)^*c(a|b)^*c)^*(a|b)^*$

"abc"

$((a|b)^*c(a|b)^*c)^*(a|b)^*$

"abc"

$((a|b)^*c(a|b)^*c)^*(a|b)^*$

"abc"

$((a|b)^*c(a|b)^*c)^*(a|b)^*$

"abcc"

```
data Reg = Eps
        | Sym Bool Char
        | Alt Reg Reg
        | Seq Reg Reg
        | Rep Reg
```

```
match :: Reg -> String -> Bool
match r "" = empty r
...
```

```
empty :: Reg -> Bool
empty Eps          = True
empty (Sym _ _)    = False
empty (Alt p q)     = empty p || empty q
empty (Seq p q)     = empty p && empty q
empty (Rep r)       = True
```



```
...  
match r (c:cs) =  
    final $ foldl (shift False)  
                (shift True r c)  
                cs
```

```
final :: Reg -> Bool
```

```
final Eps          = False
```

```
final (Sym m _) = m
```

```
final (Alt p q) = final p || final q
```

```
final (Seq p q) =  
    final q || final p && empty q
```

```
final (Rep r)      = final r
```

```
match r (c:cs) =  
    final $ foldl (shift False)  
                (shift True r c)  
                cs
```

```
shift :: Bool -> Reg -> Char -> Reg
```

```
shift _ Eps      _ = Eps
shift m (Sym _ x) c = Sym (m && x==c)
...
```

...
shift m (Alt p q) c =
 Alt (shift m p c) (shift m q c)
...

```
...  
shift m (Seq p q) c =  
    Seq (shift m p c)  
        (shift (m && empty p || final p)  
              q c)  
...
```

```
...  
shift m (Rep r) c =  
  Rep (shift (m || final r) r c)
```

- **False** $\mapsto 0$
- **True** $\mapsto 1$
- **(||)** $\mapsto (+)$
- **(&&)** $\mapsto (*)$

match :: Reg -> String -> Int


```
match (a|a*) "a" == 2
match ((a|a*)(b|b*)) "ab" == 4
```

```
match :: Semiring s  
      => Reg -> String -> s
```

- position of leftmost matching
- length of longest matching
- ...

Laziness \rightsquigarrow infinite regular expressions!

non-regular languages like:

$$\{a^n b^n \mid n \in \mathbb{N}\}$$

$$\{a^n b^n c^n \mid n \in \mathbb{N}\}$$

and more.

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curious? read the play!

cabal install weighted-regexp

github.com/sebfisch/haskell-regexp