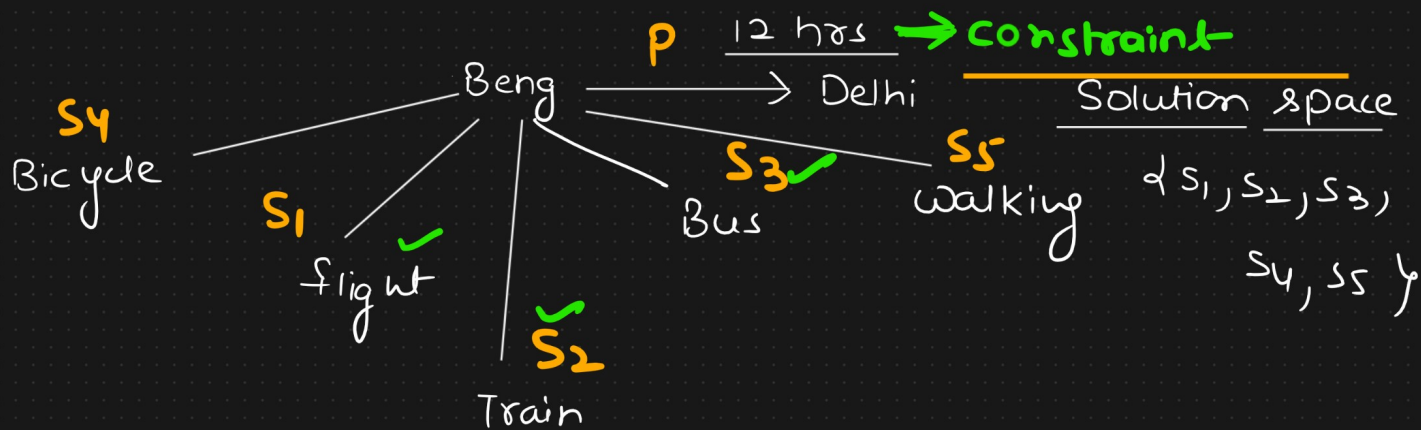


Greedy Algorithms



min. cost

feasible solution
 $\{S_1, S_2, S_3\}$

Under a given
constraint

← optimized solution

$\{S_1\}$ — AirIndia

Minimization/Maximization
→ given constraint

→ Optimized Problems

Greedy
Algorithms

Dynamic
Programming

Pseudocode

Greedy (arr, n):

for $i = 0$ to $n-1$:

$x = \text{select}(\text{arr})$

if $\text{feasible}(x)$:

$\text{solution} = \text{solution} + x$

Application

Fractional Knapsack

Items	1	2	3	4	5	6	7
Profit	25	75	100	50	45	90	30
Weight	5	10	12	4	7	9	3
P/W	5	7.5	8.33	12.5	6.42	10	10

$M = 37$

Optimization Problem

Sort
Data (Desc)
P/W

Maximize the Profit

Constraint $\leq M$

Item	4 ✓	6 ✓	7 ✓	3 ✓	2 ✓	5	1
Profit	50	90	30	100	75	45	25
Wt	4	9	3	12	<u>10</u>	7	5
<u>P/W</u>	12.5	10	10	8.33	7.5	6.42	5

Net wt

$$37 - 4 = \underline{\underline{33}}$$

$$33 - 9 = 24$$

$$24 - 3 = 21$$

$$21 - 12 = 9$$

→ $9 - \frac{9}{10} \times 10 = 0$

Net Profit

$$50$$

$$50 + 90 = 140$$

$$140 + 30 = 170$$

$$170 + 100 = 270$$

Profit

$$270 + \frac{9}{10} \times 75 =$$

$$\underline{\underline{337.5}}$$

Pseudocode

1) P_i/w_i — $\Theta(n)$

2) Sort P_i/w_i Desc $\Theta(n \log n)$

3) for $i = 0$ to $n-1$:

$\Theta(n)$ { decrease net weight & increase the profit

$$\underline{\Theta(n)} + \underline{\Theta(n \log n)} + \underline{\Theta(n)} = \underline{\Theta(n \log n)}$$