

RECURRENCE TREES



① Intro: Recurrence trees are another way to understand recurrence relations:

- (A) To calculate specific values
- (B) To construct closed fmlas in (perhaps) more organized manner than digging down.
- (C) Serves as background for proof of Master Theorem!

Note: used for RR of the form

↪ not things like $T(n-1)$, $T(n-2)$, etc.

$$T(n) = a T\left(\frac{n}{b}\right) + f(n)$$

↪ yes, like $T(n) = 2T\left(\frac{n}{3}\right) + 6n$

Let's look at:

$$T(n) = 2T\left(\frac{n}{3}\right) + 5n + 1 \quad w/ \quad T(1) = 7$$

② Calc. Specific Values

Sps we want $T(9)$

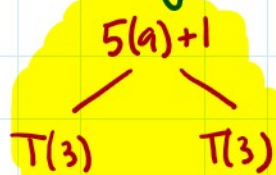
we'll construct the tree in stages! $n=9$

start

$$T(9) \otimes$$

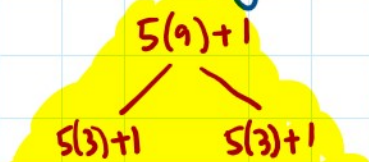
we obs. that $T(9) = 2T(3) + 5(9) + 1$

here's a tree which diagrams ↗

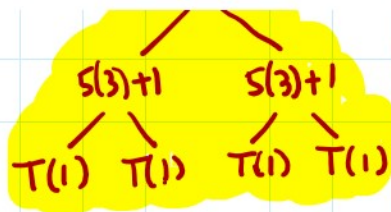


\otimes identical total sum!

repeat(ish) using the fact that $T(3) = 2T(1) + 5(3) + 1$
replace the two $T(3)$ by subtrees:

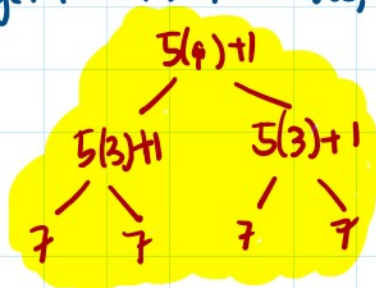


\otimes and again,



⊗ and again,
same total sum!

b/c we get to $T(1)$ in leaves, and $T(1) = 7$:



⊗ and again!

This total is $T(9)$!

Thus

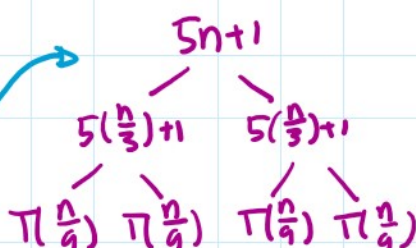
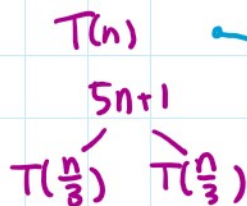
$$T(9) = S(9)+1 \\ + 2(S(3)+1) \\ + 4(7) = 106$$

Level by level

③ Finding a closed formula!

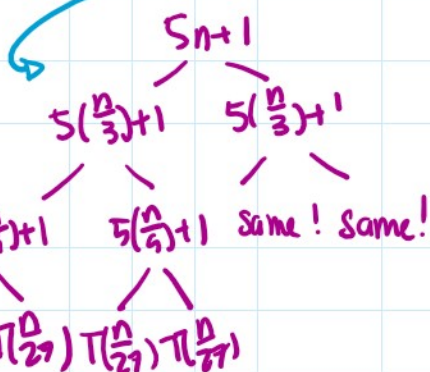
Same as ↑ but w/out specific n ! $T(n) = 2T(\frac{n}{3}) + S(n)+1$ w/ $T(1) = 7$

start
next



$$T(\frac{n}{3}) = 2T(\frac{n}{9}) + S(\frac{n}{3})+1$$

$$T(\frac{n}{9}) = 2T(\frac{n}{27}) + S(\frac{n}{9})+1$$



does this go forever?!

technically it stops at the base case
which happens when the leaves
are $T(1)$

(idealistic - assume n is nice,
and we reach $T(1)$!)

Say level 0 = root level, next level is level 1, etc.

At level k , as we built it, we had $T(\frac{n}{3^k})$
before proceeding!

We get to $T(1)$ when $\frac{n}{3^k} = 1$ which is when $k = \log_3 n$

• non-leaf levels are $0, 1, 2, \dots, (\log_3 n) - 1$

We get to 1(1) when $3^k = 1$ which is when $k = \log_3 n$

• non-leaf levels are $0, 1, 2, \dots, (\log_3 n) - 1$

• leaf level is $\log_3 n$

This next table yields the total in each level!

Level	#nodes	value/node	total = #nodes * value/node
0	1	$5n+1$	$(1)(5n+1)$
1	2	$5(\frac{n}{3})+1$	$(2)(5(\frac{n}{3})+1)$
2	4	$5(\frac{n}{3^2})+1$	$(4)(5(\frac{n}{3^2})+1)$
\vdots	\vdots	\vdots	\vdots
$k-1$	2^{k-1}	$5(\frac{n}{3^{k-1}})+1$	$2^{k-1}(5(\frac{n}{3^{k-1}})+1)$
k	2^k	7	$2^k(7)$

↳ Add 'em up!

b/c $T(1)=7$

Find the pattern!

$$T(n) = 7 \cdot 2^k + \sum_{i=0}^{k-1} (2^i) (5(\frac{n}{3^i}) + 1)$$

Let's evaluate! Eventually we'll sub. $k = \log_3 n$

$$\begin{aligned} T(n) &= 7 \cdot 2^k + \sum_{i=0}^{k-1} \left[5n \left(\frac{2}{3}\right)^i + 2^i \right] \\ &= 7 \cdot 2^k + 5n \sum_{i=0}^{k-1} \left(\frac{2}{3}\right)^i + \sum_{i=0}^{k-1} 2^i \\ &= 7 \cdot 2^k + 5n \left[\frac{\left(\frac{2}{3}\right)^k - 1}{\frac{2}{3} - 1} \right] + 2^k - 1 \\ &= 7 \cdot 2^k - 15n \left[\left(\frac{2}{3}\right)^k - 1 \right] + 2^k - 1 \end{aligned}$$

$$\begin{aligned} \sum_{i=0}^J r^i &= \frac{r^{J+1} - 1}{r - 1} \\ &= \frac{1 - r^{J+1}}{1 - r} \end{aligned}$$

now sub $k = \log_3 n$

$$T(n) = 8 \cdot 2^{\log_3 n} - 15n \left[\left(\frac{2}{3}\right)^{\log_3 n} - 1 \right] - 1$$

technically we're done. However, note that (using c.o.B)

$$\log_3 n \quad \frac{\log n}{\log 3} \quad \frac{1}{\log 3} \quad \frac{1}{\log 3} \quad \log_3 2 \quad \log_3 7$$

↳ Could sub in to get a more detailed

technicality were true. However, note that (using C.B.O.)

$$2^{\log_3 n} = 2^{\frac{\lg n}{\lg 3}} = (2^{\lg n})^{\frac{1}{\lg 3}} = n^{1/\lg 3} = n^{\log_3 2}$$

and

$$\left(\frac{2}{3}\right)^{\log_3 n} = \frac{2^{\log_3 n}}{3^{\log_3 n}} = \frac{n^{\log_3 2}}{n} = n^{\log_3 2 - 1}$$

} to get a nicer value!