

lists are implemented as
linked lists in OCaml,
must be Homogenous

Let bindings

We use **let** to bind name (identifier) to a value:

```
#let x = 100;; (* x is an immutable binding 100 *)  
val x : int = 100
```

Since functions are values, just like **ints** or **strings**, **let** is also used to define functions:

```
#let add x y = x + y;;  
val add : int -> int -> int
```

Type Annotations

- OCaml compiler infers the types. But type inference is tricky. It gives vague error messages. We can annotate types manually.

- The syntax `(e : t)` asserts that “`e` has type `t`”.

```
let (x : int) = 3
let z = (x : int) + 5
```

- Define functions' parameter and return types

```
let add (x:int) (y:int):int = x + y
let id x = x (* 'a → 'a *)
let id (x:int) = x (* int → int *)
```

- Checked by compiler: Very useful for debugging.



Lists in OCaml

- The basic data structure in OCaml
 - Lists can be of *arbitrary length*
 - Implemented as a linked data structure
 - Lists must be *homogeneous*
 - All elements have the same type
- Operations
 - Construct lists
 - Destruct them via pattern matching

Constructing Lists: Syntax

Syntax

$$e1 :: e2 \rightarrow [e1; e2]$$

- `[]` is the empty list (pronounced “nil”)
- $e1 :: e2$ prepends element $e1$ to list $e2$
 - $e1$ is the head, $e2$ is the tail
- $[e1; e2; \dots; en]$ is *syntactic sugar* for $e1 :: e2 :: \dots :: en :: []$

Examples

<code>3 :: []</code>	<code>(* [3] *)</code>
<code>2 :: (3 :: [])</code>	<code>(* [2; 3] *)</code>
<code>[1; 2; 3]</code>	<code>(* 1 :: (2 :: (3 :: [])) *)</code>

Constructing Lists: Evaluation

Evaluation

- `[]` is a value
- `[e1; ...; en]` evaluates to a list of `[v1; ...; vn]`
 - Where
 - *e1* \Rightarrow *v1*,
 - ...,
 - *en* \Rightarrow *vn*

each member of list is evaluated to a value

Constructing Lists: Examples

```
# let y = [1; 1+1; 1+1+1] ;;  
val y : int list = [1; 2; 3]
```

```
# let x = 4::y ;;  
val x : int list = [4; 1; 2; 3]
```

↓ pre pended

```
# let z = 5::y ;;  
val z : int list = [5; 1; 2; 3]
```

↓ pre pended, no change saved from prev, because functional language.

```
# let m = "hello"::"bob"::[] ;;  
val m : string list = ["hello"; "bob"]
```

Constructing Lists: Typing

Nil:

`[]: 'a list (* empty list *)`

Cons:

If $e1 : t$ and $e2 : t \text{ list}$ then $e1 :: e2 : t \text{ list}$

is of \uparrow
type

also of \uparrow
type

$[e1; e2] : t \text{ list}$

Examples

```
# let x = [1; "world"] ;;
```

X

This expression has type string but an expression was expected of type int

```
# let m = [[1];[2;3]];;
```

a list of int lists

```
val y : int list list = [[1]; [2; 3]]
```

```
# let y = 0::[1;2;3] ;;
```

```
val y : int list = [0; 1; 2; 3]
```

list of ints

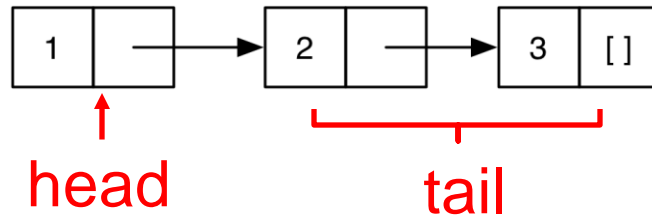
```
# let w = [1;2]::y ;;
```

This expression has type int list but is here used with type int list list

int list cannot be prepended to
list of ints b/c
int list + int

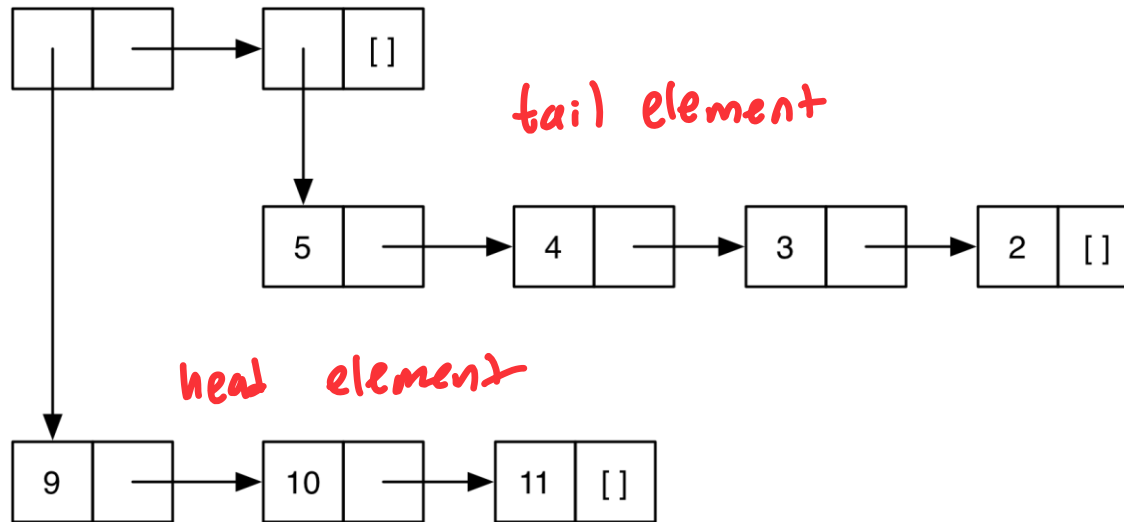
Lists in Ocaml are Linked

[1;2;3] is represented as:



Lists of Lists

- Lists can be nested arbitrarily
 - Example: `[[9; 10; 11]; [5; 4; 3; 2]]`
 - Type `int list list`, also written as `(int list) list`

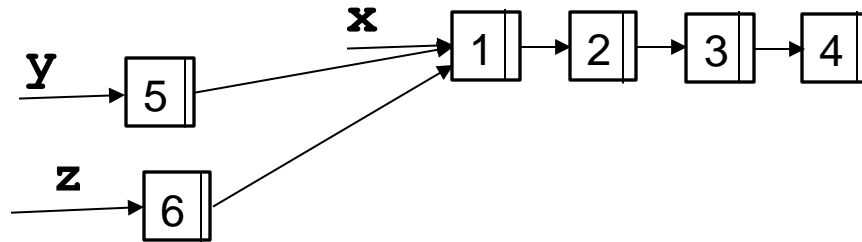


Lists are Immutable

- No way to *mutate* (change) an element of a list
- Instead, build up new lists out of old, e.g., using `::`

```
let x = [1;2;3;4]  
let y = 5::x  
let z = 6::x
```

y = 5 ; 1 ; 2 ; 3 ; 4
z = 6 ; 1 ; 2 ; 3 ; 4



Quiz 1

What is the type of the following expression?

[1.0; 2.0; 3.0; 4.0]

Float list

A. array

B. list

☒ C. float list

D. int list

Quiz 1

What is the type of the following expression?

`[1.0; 2.0; 3.0; 4.0]`

A. array

B. list

C. float list ✓

D. int list

Quiz 2

What is the type of the following expression?

`10 :: [20]`

A. `int`

int list

☒ B. `int list`

C. `int list list`

D. `error`

Quiz 2

What is the type of the following expression?

`10 :: [20]`

A. `int`

B. `int list`

C. `int list list`

D. `error`

Quiz 3

What is the type of the following definition?

```
let f a = "umd" :: [a]
```

- A. `string -> string`
- B. `string list`
- C. `string list -> string list`
- ☒ D. `string -> string list`

string \rightarrow string list

Quiz 3

What is the type of the following definition?

```
let f a = "umd" :: [a]
```

A. `string -> string`

B. `string list`

C. `string list -> string list`

D. `string -> string list`

Pattern Matching

- To pull lists apart, use the **match** construct

- Syntax

```
match e with  
| p1 -> e1  
| ...  
| pn -> en
```

analogous to switch
case:
case:
default,

- *p1...pn* are *patterns*
- *e1...en* are *branch expressions*

Pattern Matching Example

```
let is_empty l =  
  match l with  
    [] -> true  
  | (h::t) -> false
```

► Example runs

- `is_empty []` `(* true *)`
- `is_empty [1]` `(* false *)`
- `is_empty [1;2]` `(* false *)`

Pattern Matching Example (cont.)

```
let hd l =  
  match l with  
    (h::t) -> h
```

*gets head of a list
basically saying
if (l == h::t), return h, ~*

- Example runs

```
- hd [1;2;3] (* 1 *)  
- hd [2;3]   (* 2 *)  
- hd [3]     (* 3 *)  
- hd []      (* Exception: Match_failure *)
```

Pattern Matching Example (cont.)

```
let neg n =  
  match n with  
  | true -> false  
  | _ -> true
```

```
let is_empty l =  
  match l with  
  [] -> true  
  | _ -> false
```


- An underscore is a wildcard pattern. It matches anything

↳ analogous to 'default'
in switch case
statements

Quiz 4

To what does the following expression evaluate?

```
match [1;2;3] with
  [] -> [0]
| h::t -> t
```

- A. []
- B. [0]
- C. [1]
-  D. [2;3]

Quiz 4

To what does the following expression evaluate?

```
match [1;2;3] with  
  [] -> [0]  
| h::t -> t
```

- A. []
- B. [0]
- C. [1]
- D. [2;3]

"Deep" pattern matching

- $a :: b$ matches lists with **at least one** element
- $a :: []$ matches lists with **exactly one** element
- $a :: b :: []$ matches lists with **exactly two** elements
- $a :: b :: c :: d$ matches lists with **at least three** elements

Quiz 5

To what does the following expression evaluate?

```
match [1;2;3] with
| 1 :: []      -> [0]
| _ :: _        -> [1]
| 1 :: _ :: [] -> []
```

- A. []
- B. [0]
- ☒ C. [1]
- D. [2;3]

Quiz 5

To what does the following expression evaluate?

```
match [1;2;3] with
| 1::[]      -> [0]
| _::_      -> [1]
| 1::_:[]   -> []
```

- A. []
- B. [0]
- C. [1]
- D. [2;3]

Pattern Matching – An Abbreviation

- `let f p = e`, where `p` is a pattern
 - is shorthand for `let f x = match x with p -> e`

`let hd l =
 match l with
 | h::_ -> h` (* get head of list *)

- Examples

– `let hd (h::_) = h`

– `let tl (_::t) = t`

← shortened to

- Useful if there's only one acceptable input

`let tl l =
 match l with
 | _::t -> t`

o.g: `let f(x::y::_) = x + y`

`let f list =
 match list with
 | x::y::_ -> x + y`

Polymorphic Types

head

- The `hd` function works for *any type of list*
 - `hd [1; 2; 3]` `(* 1 *)`
 - `hd ["a"; "b"; "c"]` `(* "a" *)`
- OCaml gives such functions **polymorphic** types
 - `hd : 'a list -> 'a`
- These are basically generic types in Java
 - `'a list` is like `List<T>`

Examples Of Polymorphic Types

```
let t1 (_::t) = t
```

```
# t1 [1; 2; 3];;
```

```
- : int list = [2; 3]
```

```
# t1 [1.0; 2.0];;
```

```
- : float list = [2.0]
```

```
(* t1 : 'a list -> 'a list *)
```

Examples Of Polymorphic Types

```
let eq x y = (x = y)
```

-

- ```
eq 1 2;;
- : bool = false
```

```
eq "hello" "there";;
- : bool = false
```

```
eq "hello" 1 -- type error
(* eq : 'a -> 'a -> bool *)
```

## Quiz 6

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What is the type of the following function?

```
let f x y =
 if x = y then 1 else 0
```

- ~~A.~~ 'a -> 'b -> int
- B. 'a -> 'a -> bool
- ☒ C. 'a -> 'a -> int *because x=y  
comparison; must be same type*
- D. int

## Quiz 6

---

What is the type of the following function?

```
let f x y =
 if x = y then 1 else 0
```

- A. `'a -> 'b -> int`
- B. `'a -> 'a -> bool`
- C. `'a -> 'a -> int`
- D. `int`



# Missing Cases

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- Exceptions for inputs that don't match any pattern
  - OCaml will warn you about non-exhaustive matches
- Example:

```
let hd l = match l with (h::_) -> h;;
```

```
Warning: this pattern-matching is not exhaustive.
```

```
Here is an example of a value that is not matched:
```

```
[]
```

```
hd [];;
```

```
Exception: Match_failure ("", 1, 11).
```

# Pattern matching is *AWESOME*

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1. You can't forget a case
  - Compiler issues inexhaustive pattern-match warning
2. You can't duplicate a case
  - Compiler issues unused match case warning
3. You can't get an exception
  - Can't do something like `List.hd []`
4. Pattern matching leads to elegant, concise, beautiful code

# Lists and Recursion

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- Lists have a recursive structure
  - And so most functions over lists will be recursive

```
let rec length l = match l with
 [] -> 0
 | (_::t) -> 1 + (length t)
```

- This is just like an inductive definition
  - *The length of the empty list is zero*
  - *The length of a nonempty list is 1 plus the length of the tail*
- Type of `length`?
  - `'a list -> int`

# More Examples

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- `sum l (* sum of elts in l *)`  
`let rec sum l = match l with`  
    `[] -> 0`  
    `| (x::xs) -> x + (sum xs)`
- `negate l (* negate elements in list *)`  
`let rec negate l = match l with`  
    `[] -> []`  
    `| (x::xs) -> (-x) :: (negate xs)`
- `last l (* last element of l *)`  
`let rec last l = match l with`  
    `[x] -> x`  
    `| (x::xs) -> last xs`

## More Examples (cont.)

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(\* return a list containing all the elements in the list l followed by all the elements in list m \*)

- `append l m`

```
let rec append l m = match l with
 [] -> m
 | (x::xs) -> x::(append xs m)
```

- `rev l` (\* reverse list; hint: use append \*)

```
let rec rev l = match l with
 [] -> []
 | (x::xs) -> append (rev xs) (x::[])
```

- `rev` takes  $O(n^2)$  time. Can you do better?