

MAXIMUM CONTIGUOUS SUM

① INTRO!

Sps we have a list of length n containing numbers

ex $\rightarrow A = [-3, 5, 1, -2, 8, -10, 1]$

If we look at all contiguous sums, what is the max?

ex For, here are some CS:

$$-3 + 5 + 1 = 3$$

$$5 + 1 - 2 + 8 = 12$$

$$-10 = -10$$

must use at least one elt.

Q: What's the max?

ex in ours, it's $5 + 1 - 2 + 8 = 12$ \leftarrow Cannot get more!

not acs:

$$5 - 2 + 8 \text{ (misses +1)}$$

② Application! If list contains daily inc/dec in stock price,
MCS = max contig. increase

③ Why are we looking at it?!

mostly b/c it introduces us to several (three!)

different alg. approaches which we'll use a lot.

- Brute Force
- Divide and Conquer
- Dynamic Programming (saw in coin-changing!)

④ BRUTE FORCE!

Try EVERY possible contiguous sum!

Here's the pseudocode:

\\ PRE: A is a list of length n.

max = A[0]

for i = 0 to n-1

sum = 0

for j = i to n-1

sum = sum + A[j]

if sum > max

max = sum

end

end

end

\\ POST: max is the maximum sum.

to start w/ something

$c_1 \leftarrow \theta(1)$

$\rightarrow n$ times $c_1 \leftarrow \theta(1)$

$\rightarrow n-1-i+1$ times

$c_1 \leftarrow \theta(1)$

$c_2 \leftarrow \theta(1)$

if larger, inc. overall MCS

i, j iterate through all possible index pairs w/ $i \leq j$

for each i, this checks all MCS starting at index i

How much time does this take?

$$T(n) = c_1 + \sum_{i=0}^{n-1} \left[c_1 + \sum_{j=i}^{n-1} (c_1 + c_2) \right]$$

note: Each for loop creates a sum Σ .
Sums to know!

now we need to simplify!

$$T(n) = c_1 + \sum_{i=0}^{n-1} \left[c_1 + \sum_{j=i}^{n-1} (c_1 + c_2) \right]$$

$$= c_1 + \sum_{i=0}^{n-1} \left[c_1 + (c_1 + c_2)(n-1-i+1) \right]$$

var. on ①

$$= c_1 + \sum_{i=0}^{n-1} \left[c_1 + (c_1 + c_2)n - (c_1 + c_2)i \right]$$

$$= c_1 + (c_1 + (c_1 + c_2)n) \sum_{i=0}^{n-1} 1 - (c_1 + c_2) \sum_{i=0}^{n-1} i$$

var on ②

$$= c_1 + (c_1 + (c_1 + c_2)n)(n) - (c_1 + c_2) \frac{(n-1)(n)}{2}$$

$$= \theta(n^2) \text{ b/c } n^2 \text{ is the highest power of } n \text{ in } \dots$$

$$\begin{aligned} \sum_{i=1}^n 1 &= n & \text{①} \\ \sum_{i=1}^n i &= \frac{n(n+1)}{2} & \text{②} \end{aligned}$$

END $T(n) = \theta(n^2)$

↳ loose Meaning \rightarrow as n (list length) inc linearly
the time req. inc. quadratically

⑤ Can we do better... faster?!