

## ***Counting Sort***

- We will consider three algorithms that are faster and work by not making comparisons.
- Counting sort assumes that the numbers to be sorted are in the range 1 to  $k$  where  $k$  is small.
- The basic idea is to determine the rank of each number in final sorted array.

## ***Counting Sort***

- Recall that the rank of an item is the number of elements that are less than or equal to it.
- Once we know the ranks, we simply copy numbers to their final position in an output array.
- The algorithm sorts in  $\Theta(n + k)$ .
- If  $k$  is  $\Theta(n)$  then counting sort is an  $\Theta(n)$  time algorithm.

## ***Counting Sort***

The algorithm uses three arrays.

1.  $A[1..n]$ : Holds the initial input.
2.  $B[1..n]$ : Array that holds the sorted output.
3.  $C[1..k]$ : Array of integers.  $C[x]$  is the rank of  $x$  in  $A$ , where  $x \in [1..k]$ .

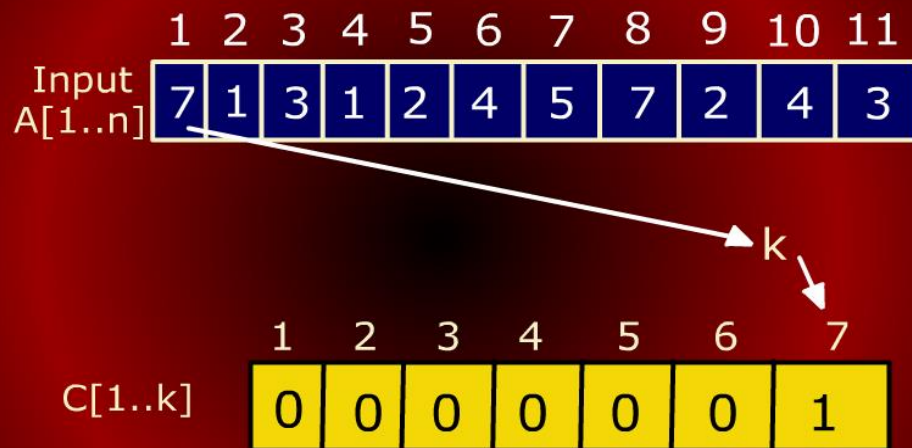
## Counting Sort

	1	2	3	4	5	6	7	8	9	10	11
Input A[1..n]	7	1	3	1	2	4	5	7	2	4	3

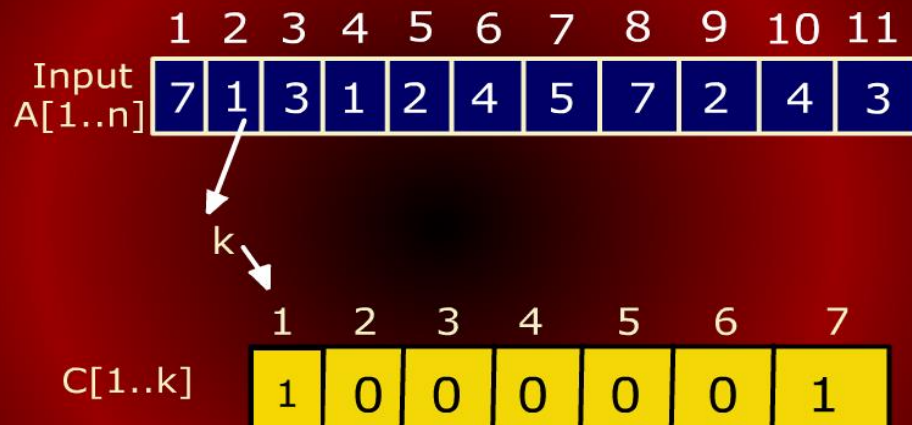
$k = 7$

	1	2	3	4	5	6	7
C[1..k]	0	0	0	0	0	0	0

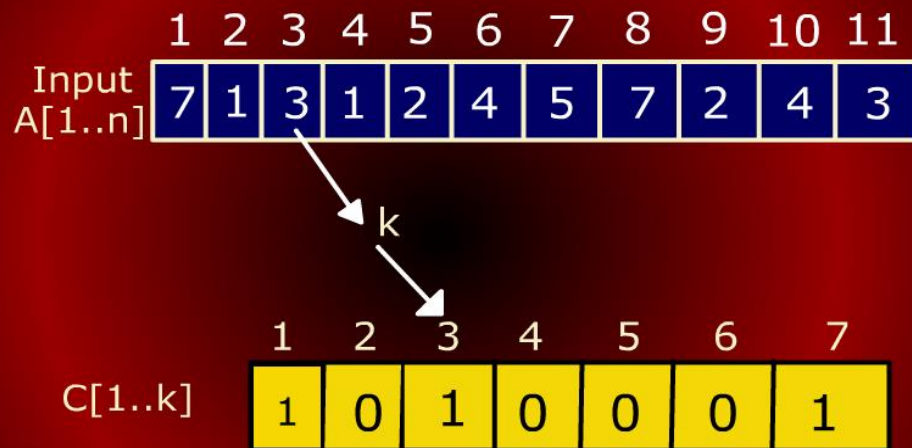
## Counting Sort



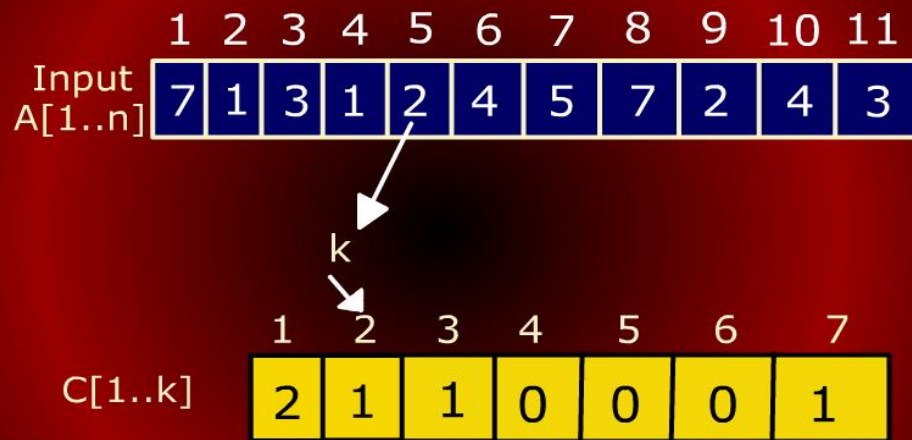
## Counting Sort



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## Counting Sort

	1	2	3	4	5	6	7	8	9	10	11
Input A[1..n]	7	1	3	1	2	4	5	7	2	4	3

Finally

	1	2	3	4	5	6	7
C[1..k]	2	2	2	2	1	0	2

## Counting Sort

	1	2	3	4	5	6	7	8	9	10	11
Input A[1..n]	7	1	3	1	2	4	5	7	2	4	3

	1	2	3	4	5	6	7
C[1..k]	2	2	2	2	1	0	2

**for i= 2 to 7**

**do** C[i] = C[i] + C[i- 1]

	1	2	3	4	5	6	7
C	2	4	6	8	9	9	11

6 elements  $\leq 3$

## Counting Sort

Input  
A[1..n]

1	2	3	4	5	6	7	8	9	10	11
7	1	3	1	2	4	5	7	2	4	3

Output  
B[1..n]

1	2	3	4	5	6	7	8	9	10	11
					3					

$B[6] = B[C[3]] = B[C[A[11]]] = A[11] = 3$

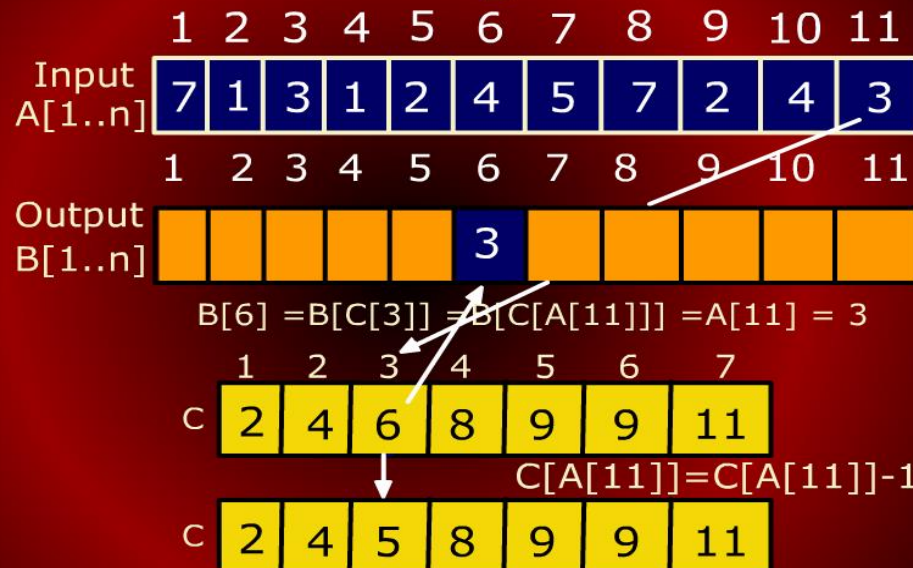
	1	2	3	4	5	6	7
C	2	4	6	8	9	9	11

$C[A[11]] = C[A[11]] - 1$

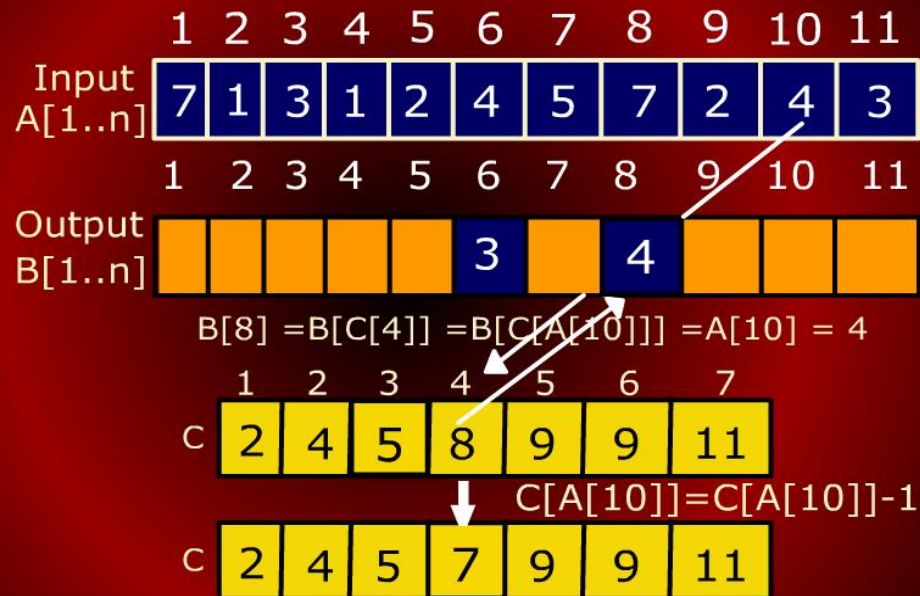
C	2	4	5	8	9	9	11
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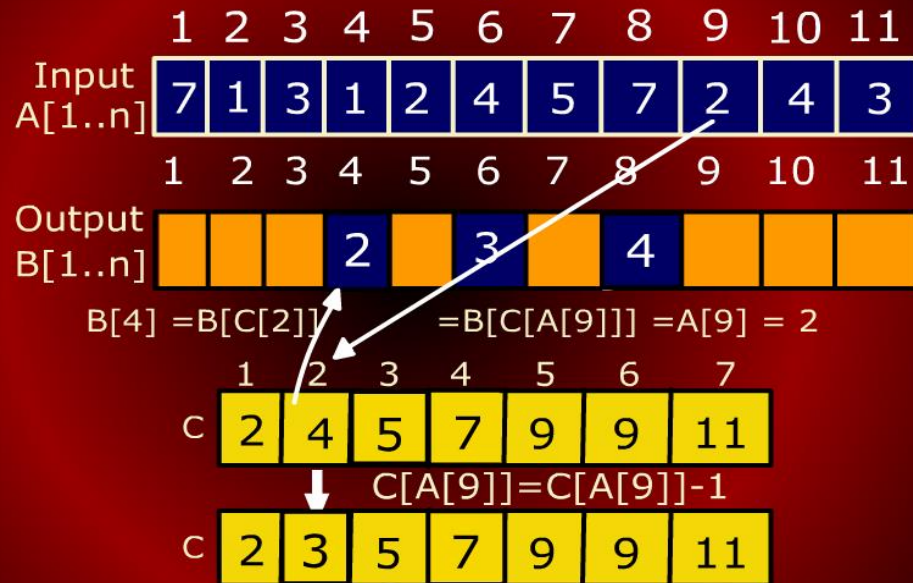
## Counting Sort



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## Counting Sort



## Counting Sort

Input  
A[1..n]

1	2	3	4	5	6	7	8	9	10	11
7	1	3	1	2	4	5	7	2	4	3

Output  
B[1..n]

1	2	3	4	5	6	7	8	9	10	11
			2		3		4			7

$B[11] = B[C[7]] = B[C[A[8]]] = A[8] = 7$

C

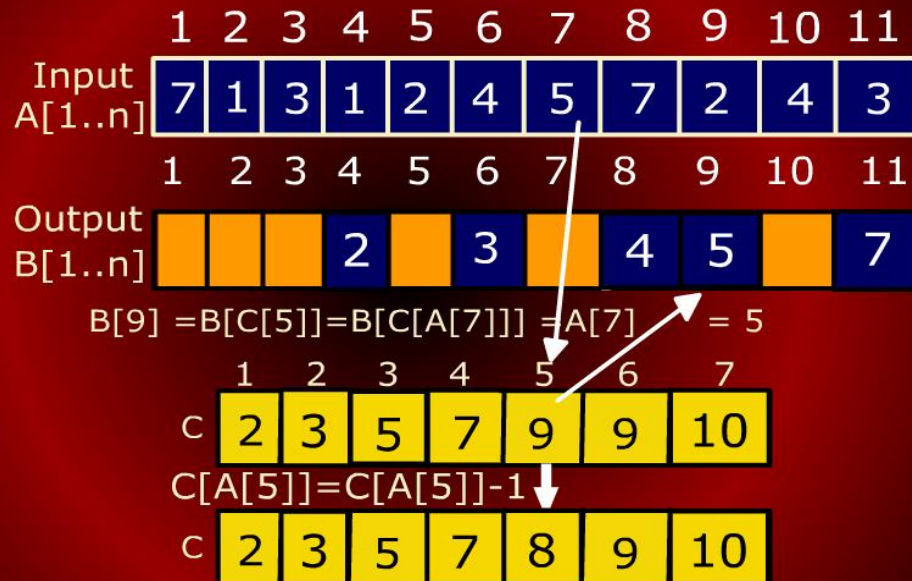
1	2	3	4	5	6	7
2	3	5	7	9	9	11

$C[A[8]] = C[A[8]] - 1$

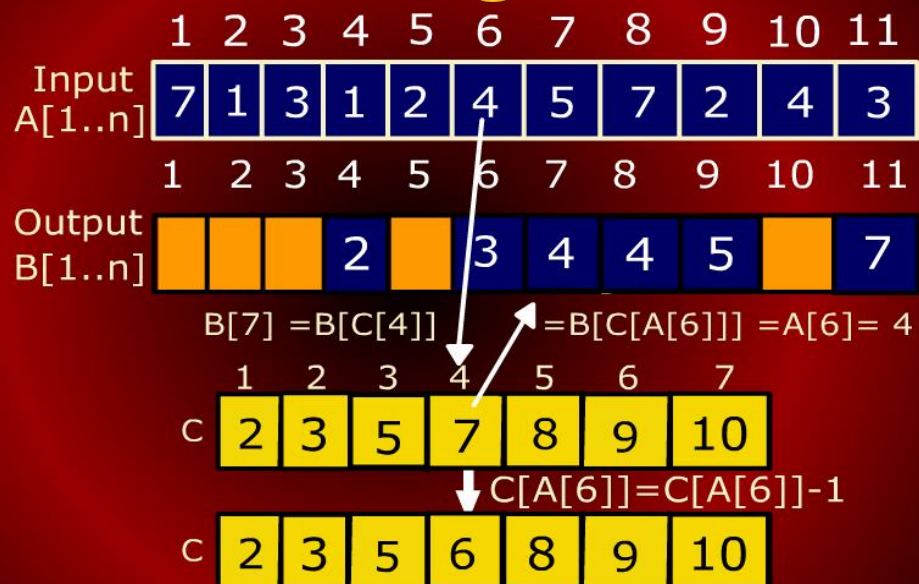
C

2	3	5	7	9	9	10
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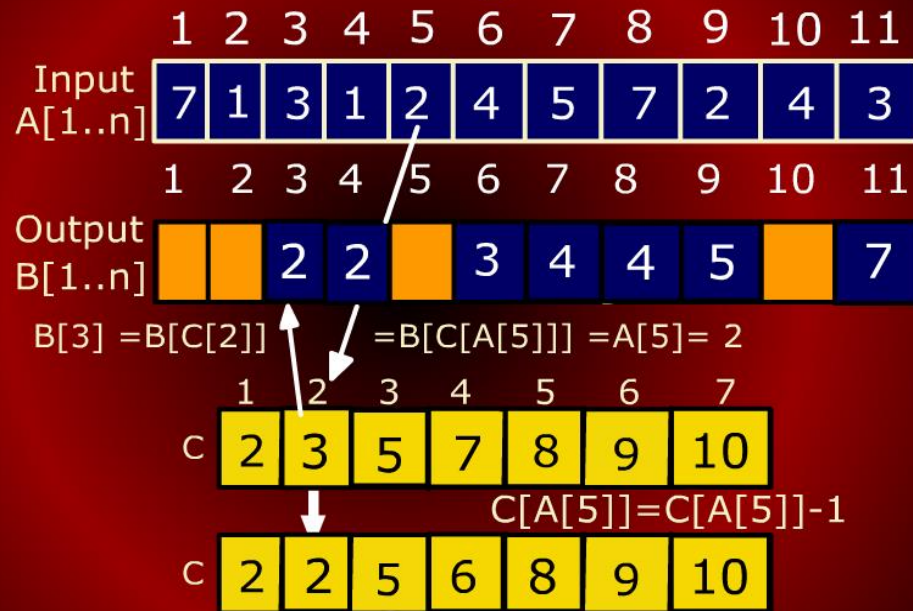
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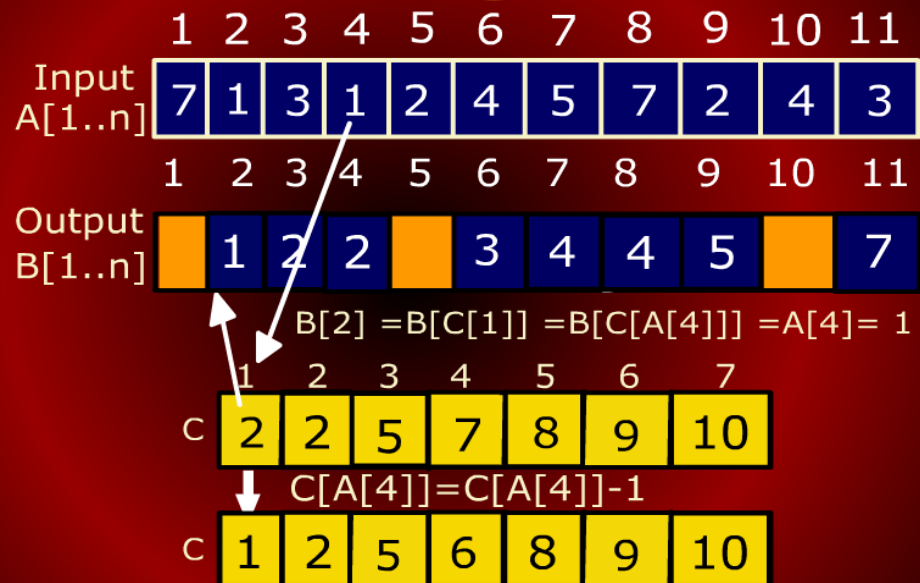


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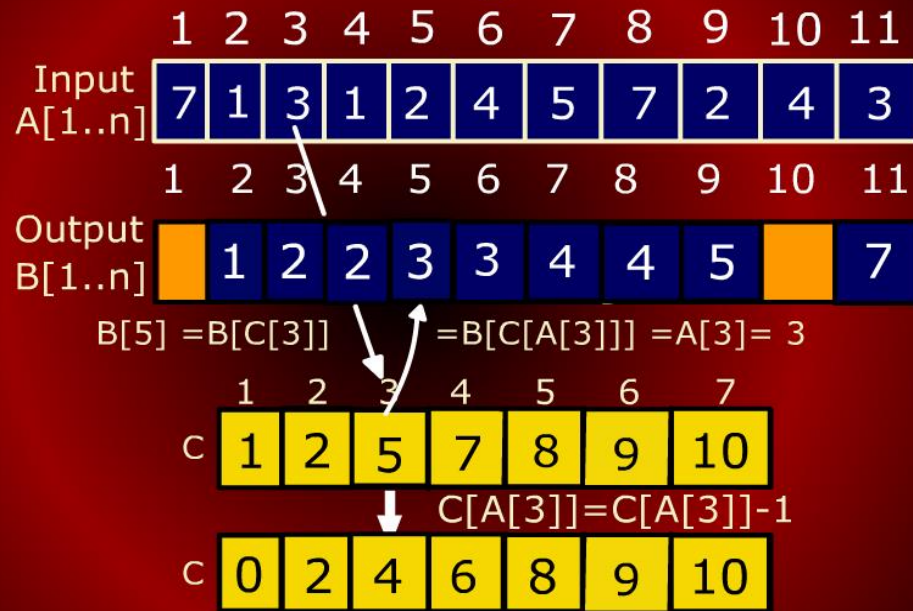




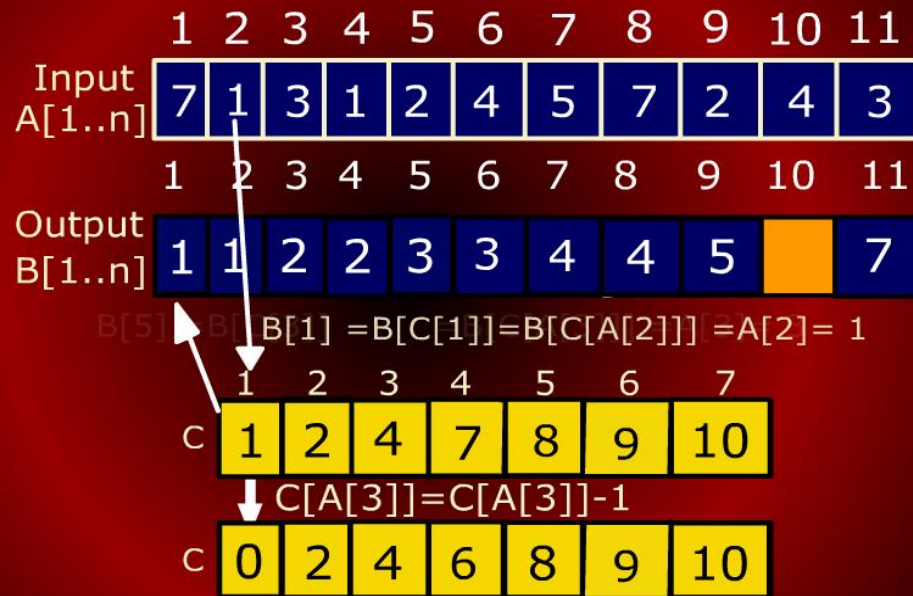
## Counting Sort



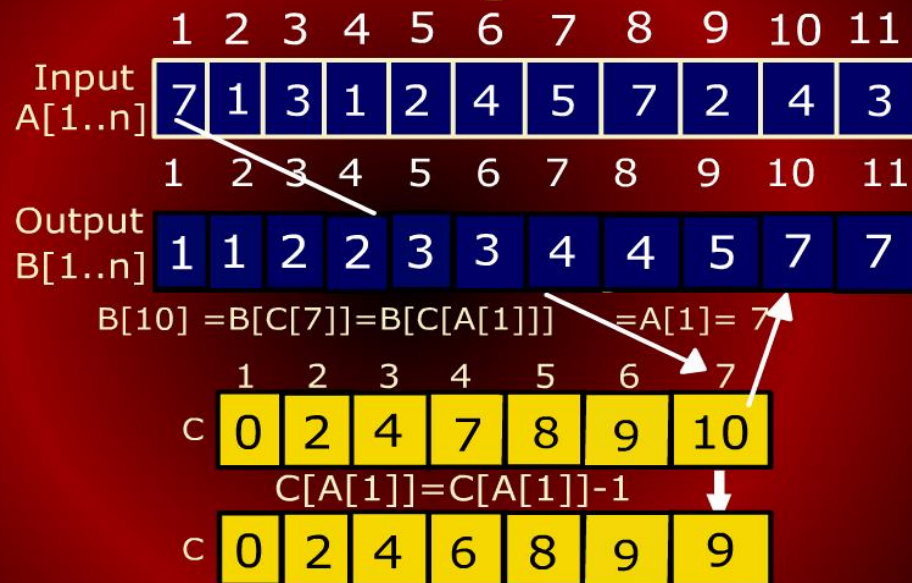
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## Counting Sort

COUNTING-SORT

( array A, array B, int k)

- 1    **for**  $i \leftarrow 1$  **to**  $k$
- 2    **do**  $C[i] \leftarrow 0$     k times
- 3    **for**  $j \leftarrow 1$  **to**  $\text{length}[A]$
- 4    **do**  $C[A[j]] \leftarrow C[A[j]] + 1$     n times
- 5    *// C[i] now contains the number of elements = i*
- 6    **for**  $i \leftarrow 2$  **to**  $k$

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7   do C[i] ← C[i] + C[i - 1]  k times
8   // C[i] now contains the number
   // of elements ≤ i
9   for j ← length[A] downto 1
10  do B[C[A[j]]] A[j]
11  C[A[j]] C[A[j]] - 1  n times
```