

Concurrent Programming (02158) - resubmission of assignment 2

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1 Problem 1

In order to prevent race conditions that were caused by several threads reading and writing into the `HashMap` at the same time. The map is wrapped using the `Collections.synchronizedMap` method at creation time. When a car enters a tile it firstly checks whether the current position key exists in the map if not it creates a new element with a semaphore for the particular position by using `putIfAbsent` method. Then it calls `P` function for the corresponding semaphore to decrease the value by 1, and its value remains until the car leaves so that other cars waiting to enter can continue.

2 Problem 3

Semaphore `dirMutex` is used to protect variable `isDirDownward` from being read and written at the same time within function `enter`. In order to protect variable `isFirst` properly, the operation of changing value of variable `isFirst` is put in the critical region right after the declaration of local variable `isFirstDownwards` which is protected by mutex `isFirstMutex`. Therefore only one car can have boolean variable `isFirstDownwards` as true.

3 Problem 4

The changes made to the problem 4 mostly rely on the solution from problem 3. The main issue was the inconsistency between the *Java* implementation and the *Promela* code. These include the use of the same semaphores as in the *Java* case, which were fixed in the *Promela* implementation. Statements as the *counter* increment were protected by a semaphore (as in *Java*) and not by atomic statement in the *Promela* code and finally all changes from problem 3 were translated into the *Promela* script.