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http://www.cs.princeton.edu/~wayne/kleinberg-tardos

PRIORITY QUEUES

binary heaps

Priority queue data type

A min-oriented priority queue supports the following core operations:

- Make-Heap(): create an empty heap.
- INSERT(H, x): insert an element x into the heap.
- EXTRACT-MIN(H): remove and return an element with the smallest key.
- DECREASE-KEY(H, x, k): decrease the key of element x to k.

The following operations are also useful:

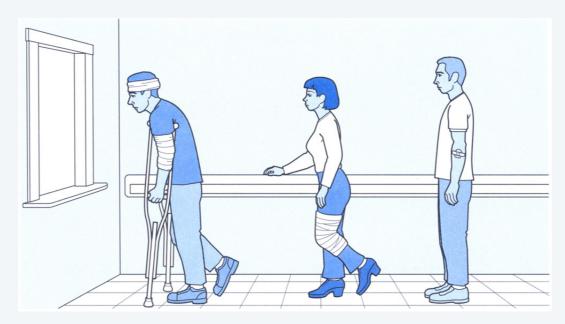
- Is-EMPTY(H): is the heap empty?
- FIND-MIN(*H*): return an element with smallest key.
- DELETE(H, x): delete element x from the heap.
- MELD(H_1, H_2): replace heaps H_1 and H_2 with their union.

Note. Each element contains a key (duplicate keys are permitted) from a totally-ordered universe.

Priority queue applications

Applications.

- A* search.
- Heapsort.
- · Online median.
- · Huffman encoding.
- Prim's MST algorithm.
- Discrete event-driven simulation.
- Network bandwidth management.
- · Dijkstra's shortest-paths algorithm.
- ...

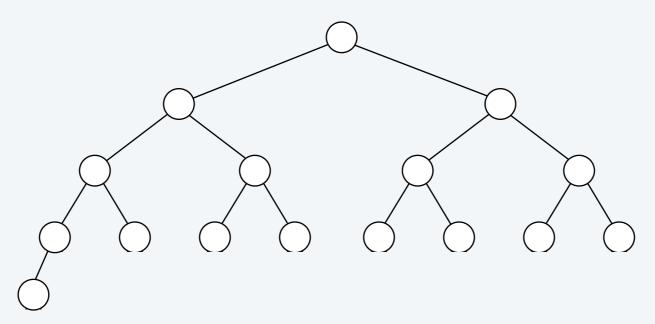


http://younginc.sitell.com/source/5895/fos0092.html

Complete binary tree

Binary tree. Empty or node with links to two disjoint binary trees.

Complete tree. Perfectly balanced, except for bottom level.



complete tree with n = 16 nodes (height = 4)

Property. Height of complete binary tree with n nodes is $\lfloor \log_2 n \rfloor$. Pf. Height increases (by 1) only when n is a power of 2. \blacksquare

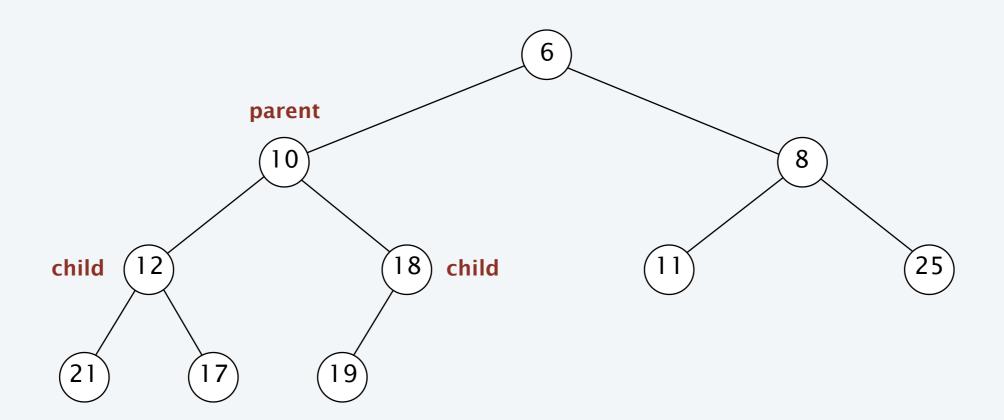
A complete binary tree in nature



Binary heap

Binary heap. Heap-ordered complete binary tree.

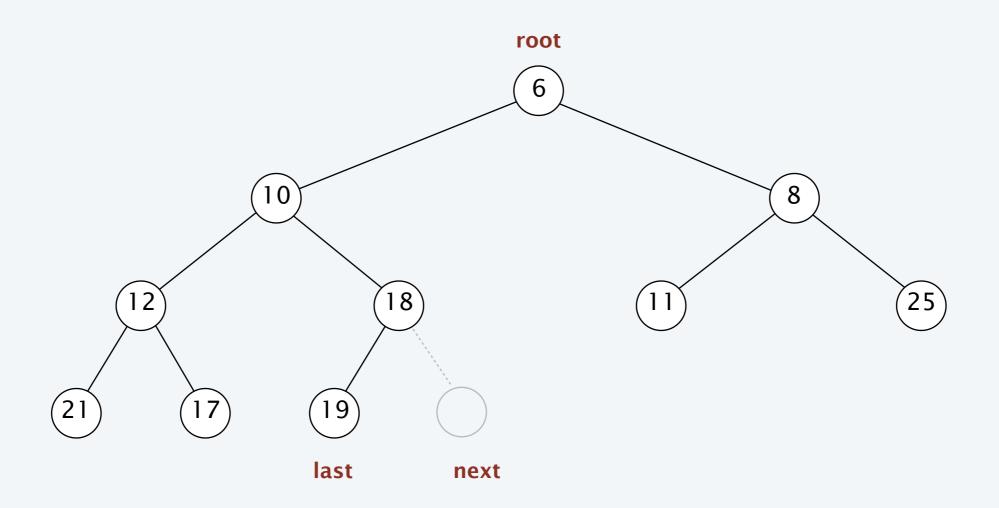
Heap-ordered tree. For each child, the key in child \geq key in parent.



Explicit binary heap

Pointer representation. Each node has a pointer to parent and two children.

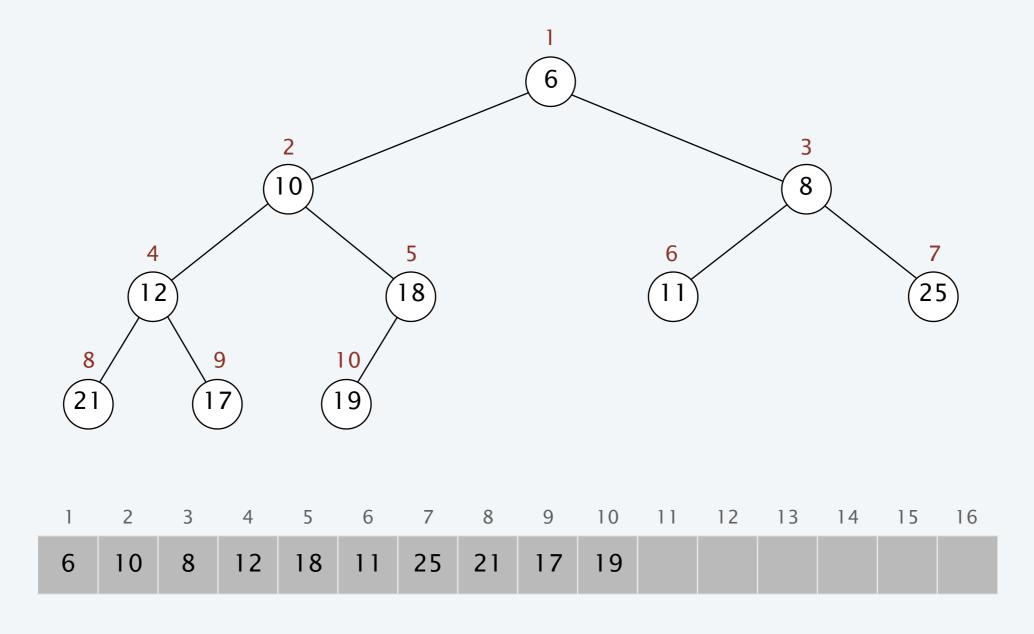
- Maintain number of elements *n*.
- Maintain pointer to root node.
- Can find pointer to last node or next node in $O(\log n)$ time.



Implicit binary heap

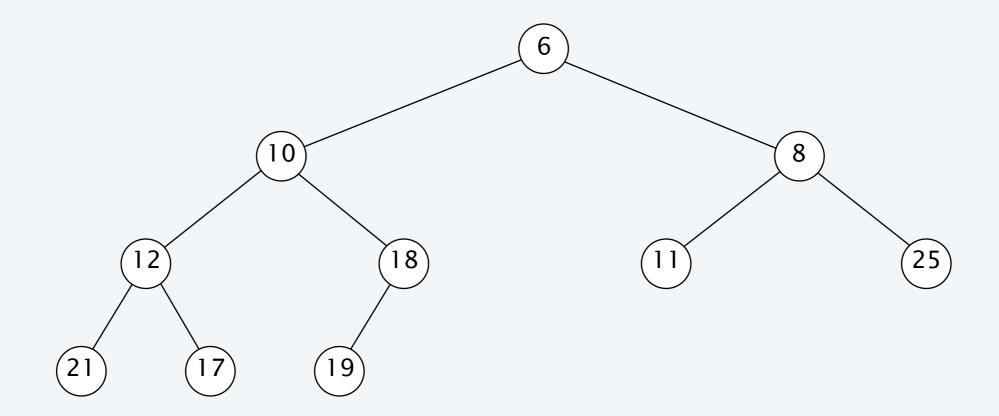
Array representation. Indices start at 1.

- Take nodes in level order.
- Parent of node at k is at $\lfloor k/2 \rfloor$.
- Children of node at k are at 2k and 2k + 1.



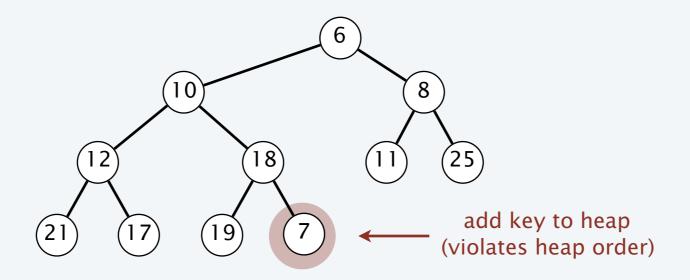


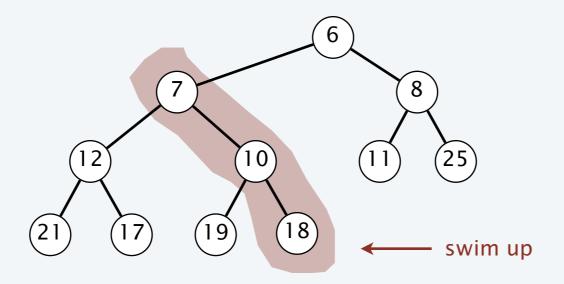
heap ordered



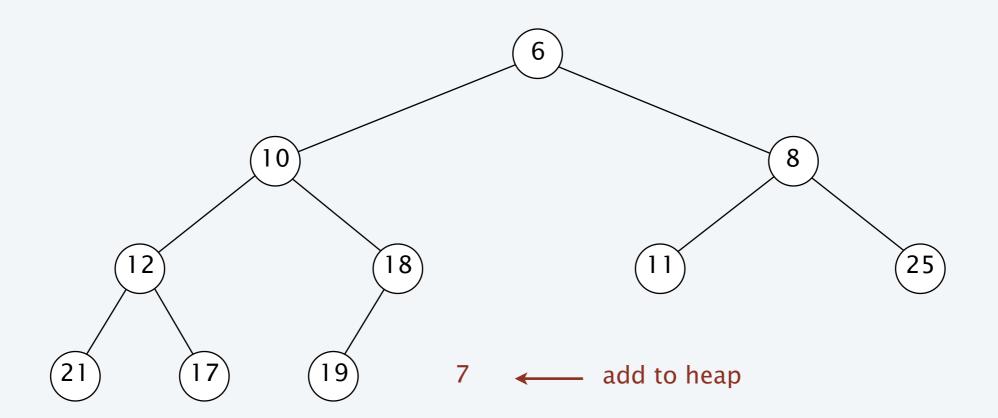
Binary heap: insert

Insert. Add element in new node at end; repeatedly exchange new element with element in its parent until heap order is restored.

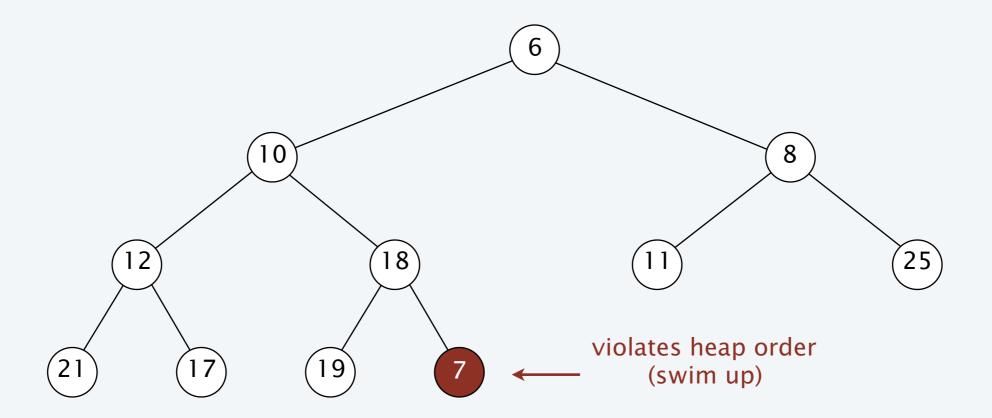




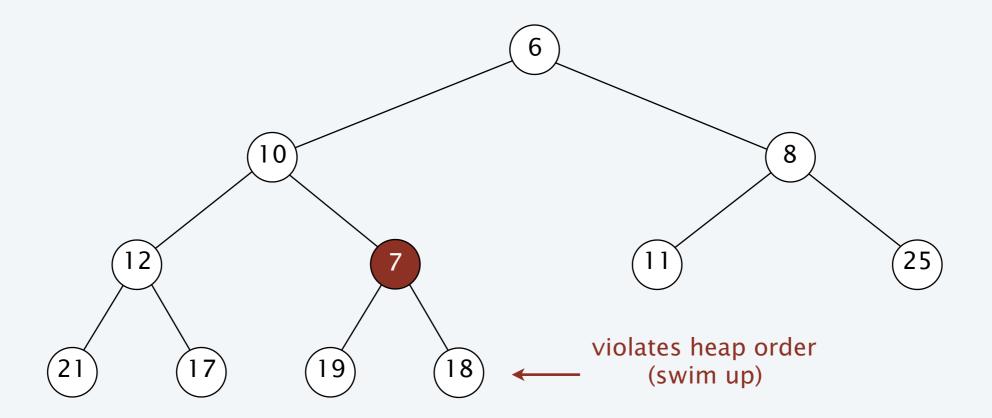
Insert. Add node at end; repeatedly exchange element in child with element in parent until heap order is restored.



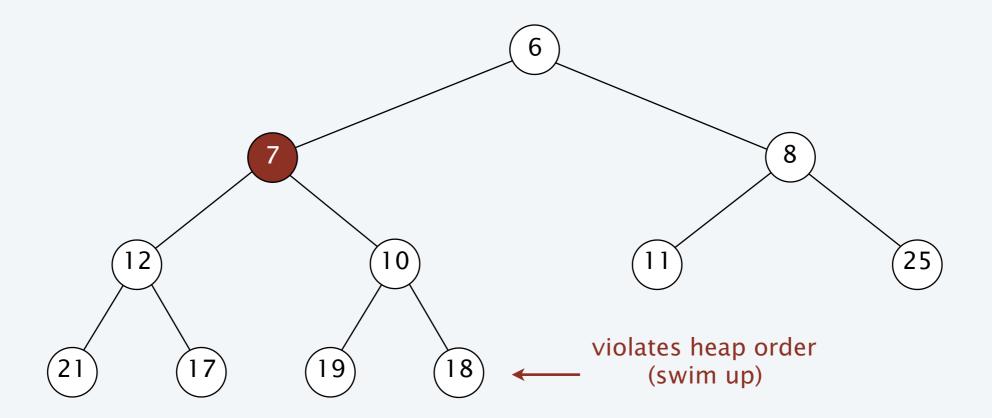
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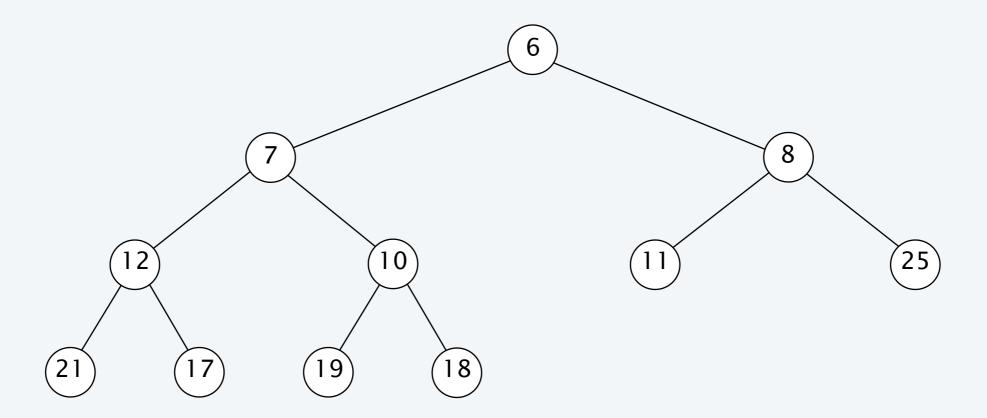
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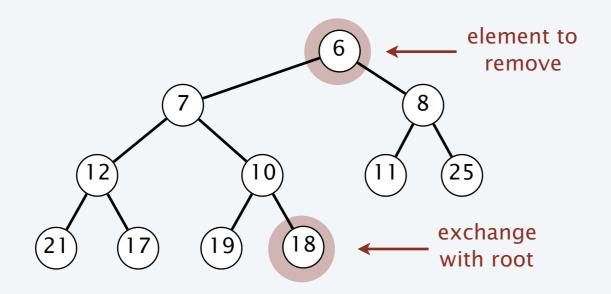


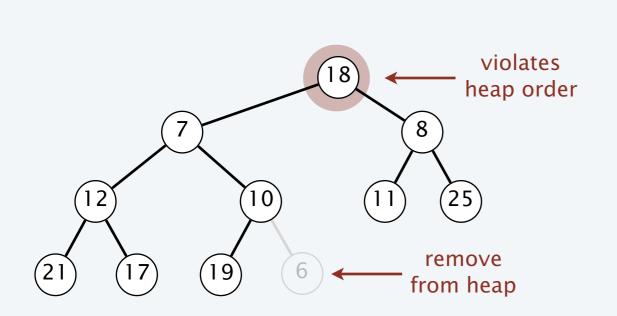
heap ordered

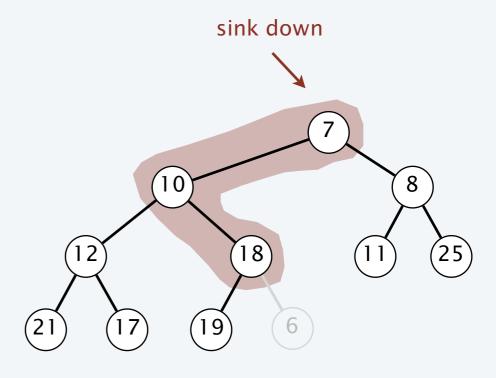


Binary heap: extract the minimum

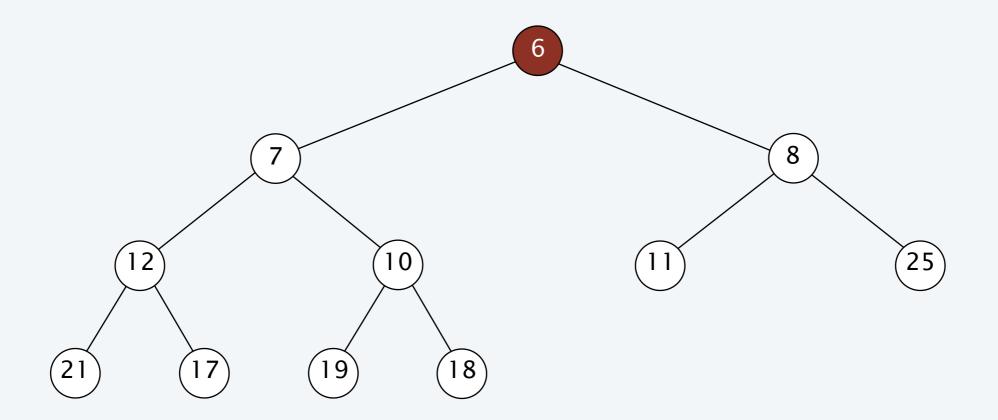
Extract min. Exchange element in root node with last node; repeatedly exchange element in root with its smaller child until heap order is restored.



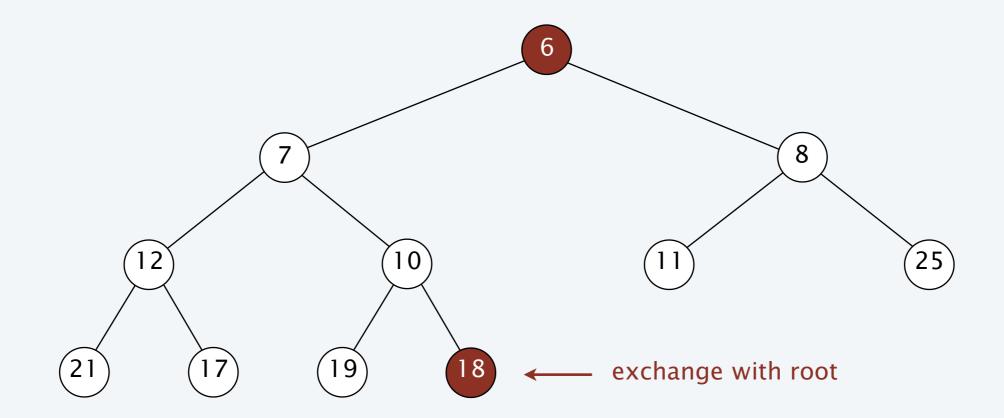




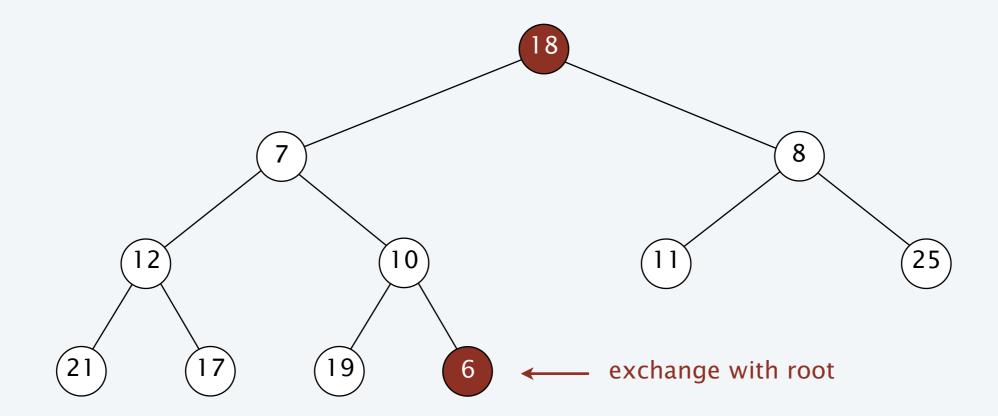
Extract min. Exchange root node with last node; repeatedly exchange element in parent with element in larger child until heap order is restored.



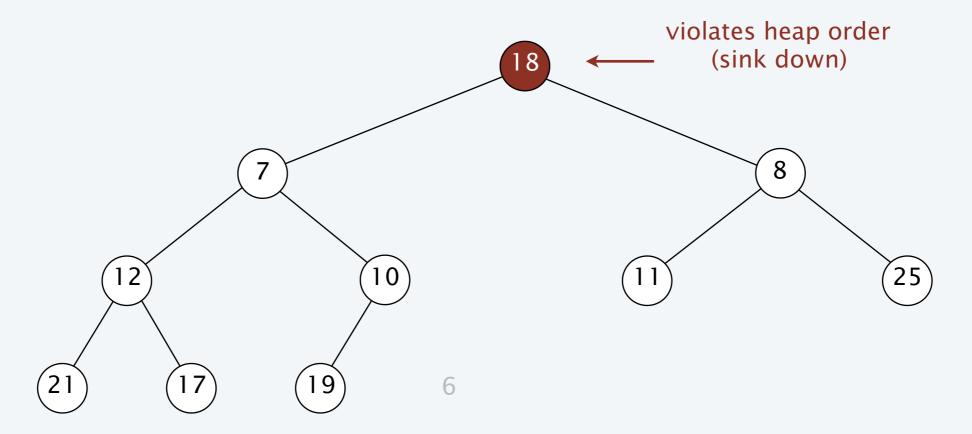
Extract min. Exchange root node with last node; repeatedly exchange element in parent with element in larger child until heap order is restored.



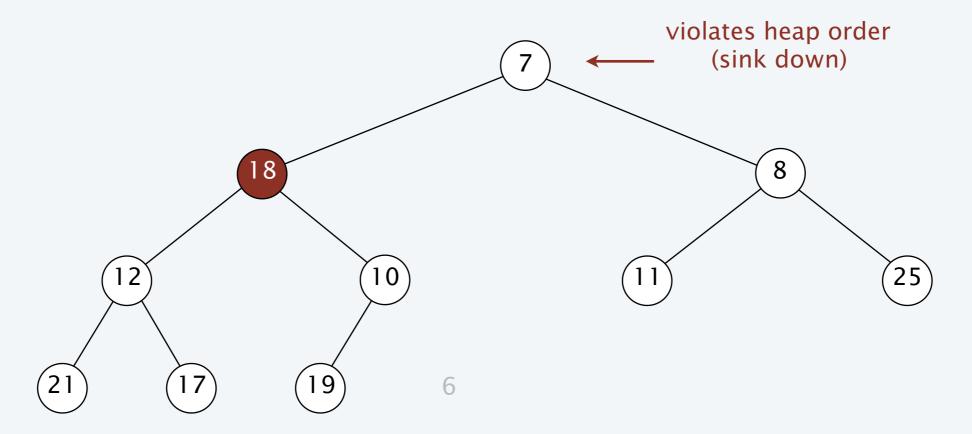
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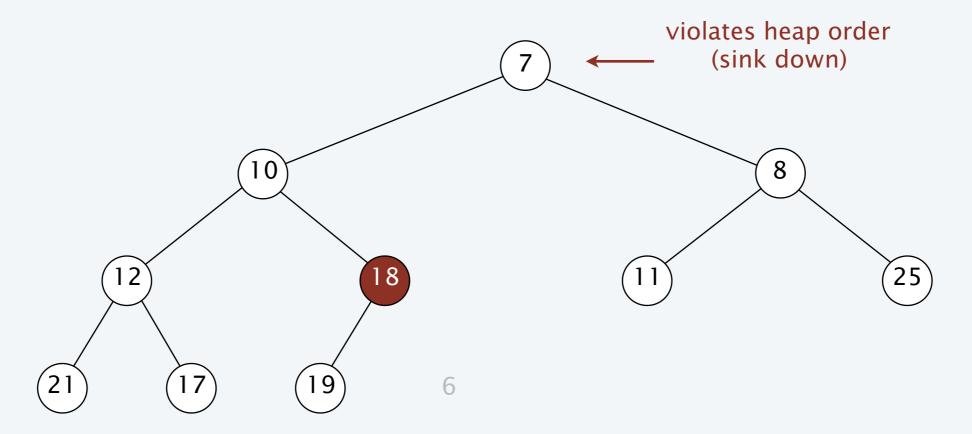
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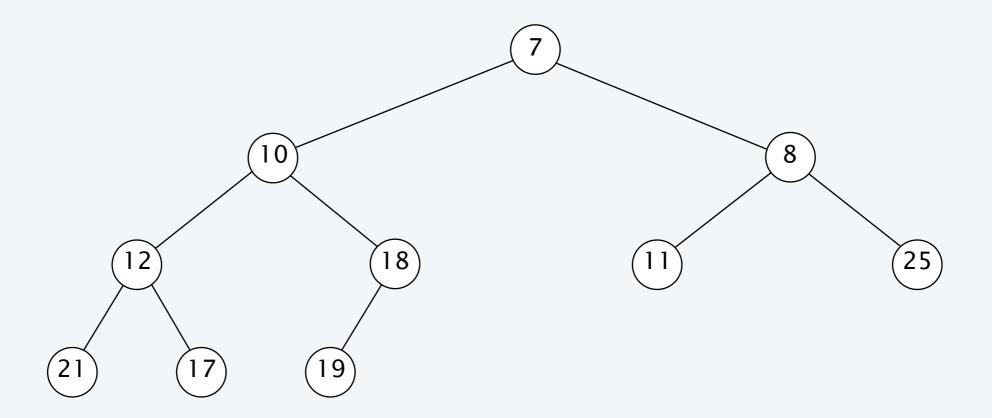
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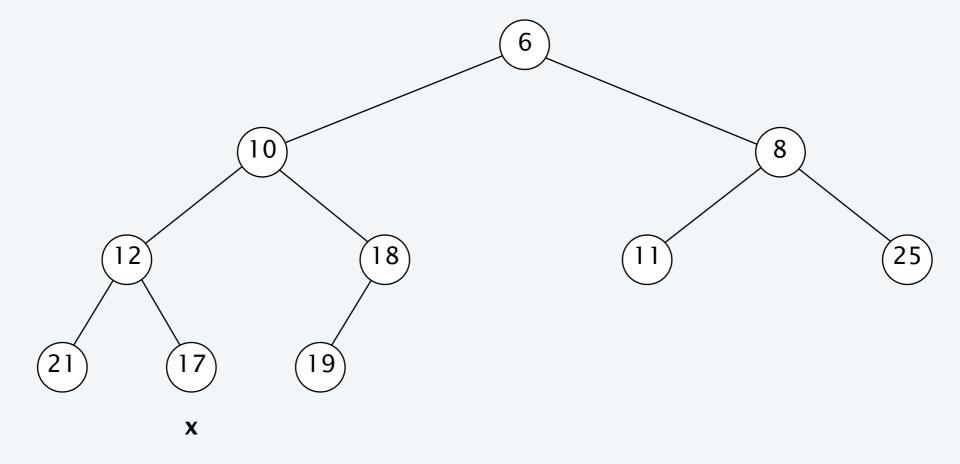
heap ordered



Binary heap: decrease key

Decrease key. Given a handle to node, repeatedly exchange element with its parent until heap order is restored.

decrease key of node x to 11



Binary heap: analysis

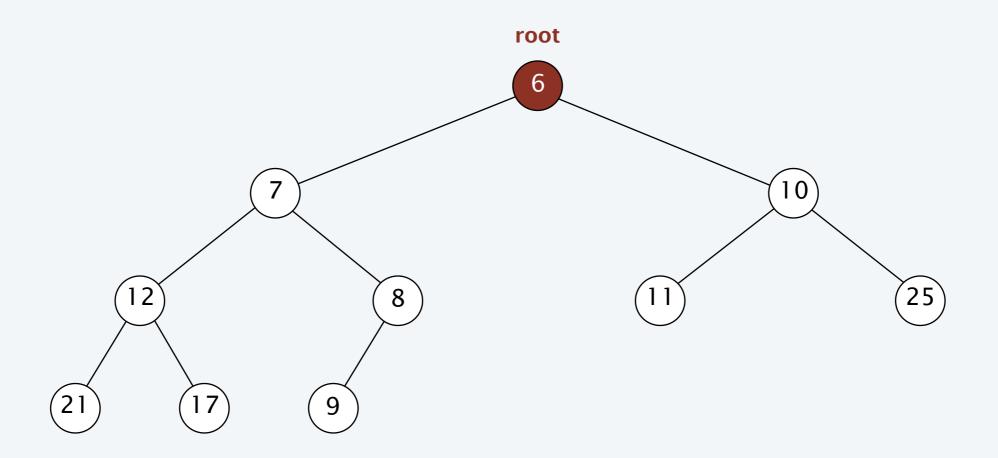
Theorem. In an implicit binary heap, any sequence of m INSERT, EXTRACT-MIN, and DECREASE-KEY operations with n INSERT operations takes $O(m \log n)$ time. Pf.

- Each heap op touches nodes only on a path from the root to a leaf; the height of the tree is at most $\log_2 n$.
- The total cost of expanding and contracting the arrays is O(n).

Theorem. In an explicit binary heap with n nodes, the operations INSERT, DECREASE-KEY, and EXTRACT-MIN take $O(\log n)$ time in the worst case.

Binary heap: find-min

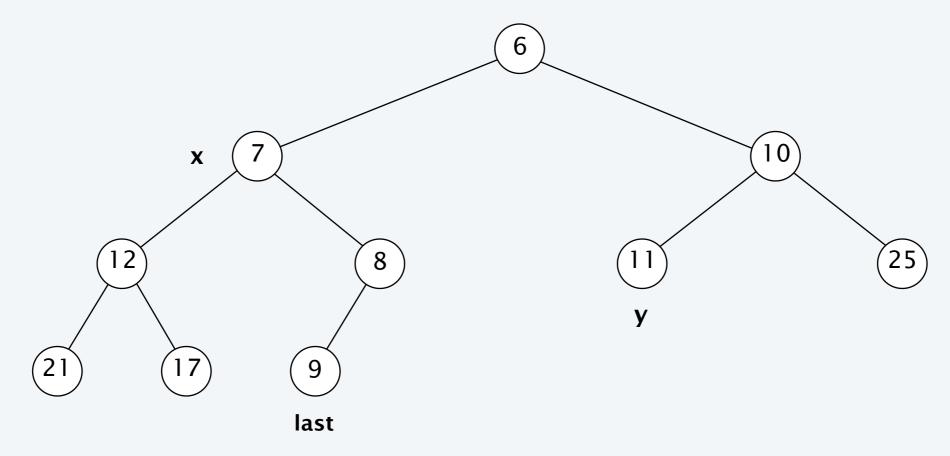
Find the minimum. Return element in the root node.



Binary heap: delete

Delete. Given a handle to a node, exchange element in node with last node; either swim down or sink up the node until heap order is restored.

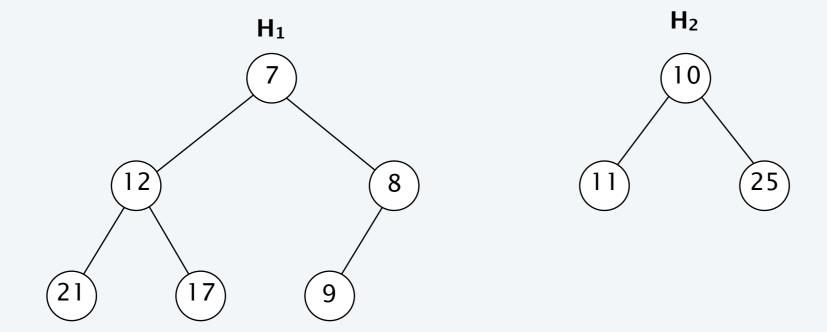
delete node x or y



Binary heap: meld

Meld. Given two binary heaps H_1 and H_2 , merge into a single binary heap.

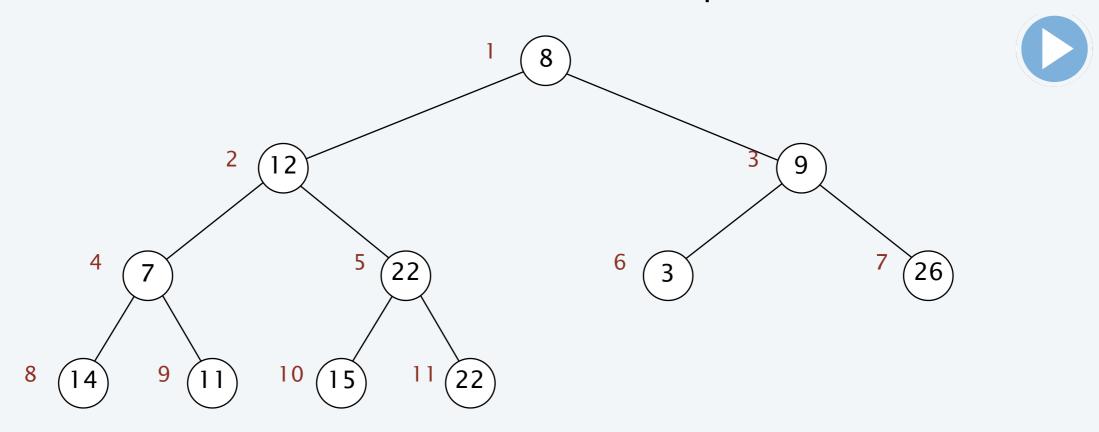
Observation. No easy solution: $\Omega(n)$ time apparently required.



Binary heap: heapify

Heapify. Given n elements, construct a binary heap containing them. Observation. Can do in $O(n \log n)$ time by inserting each element.

Bottom-up method. For i = n to 1, repeatedly exchange the element in node i with its smaller child until subtree rooted at i is heap-ordered.

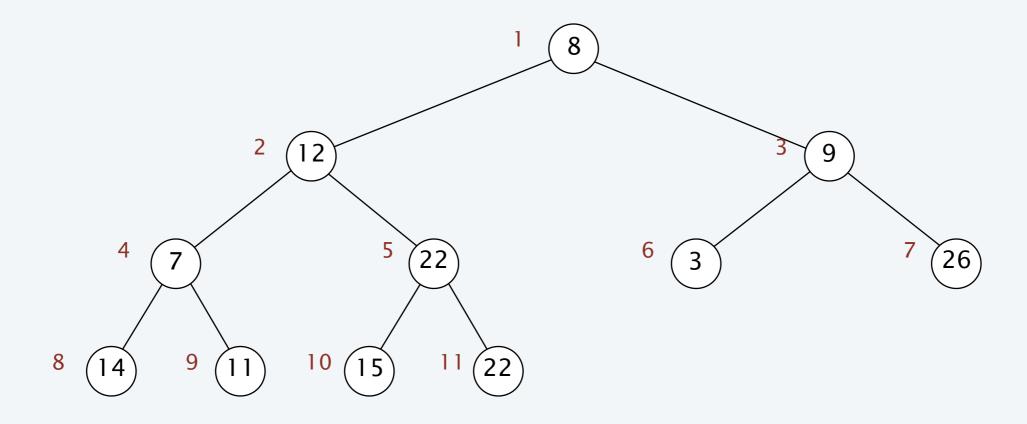


8	12	9	7	22	3	26	14	11	15	22
1	2	3	4	5	6	7	8	9	10	11

Heapify. For each element in reverse-array order, sink it down.

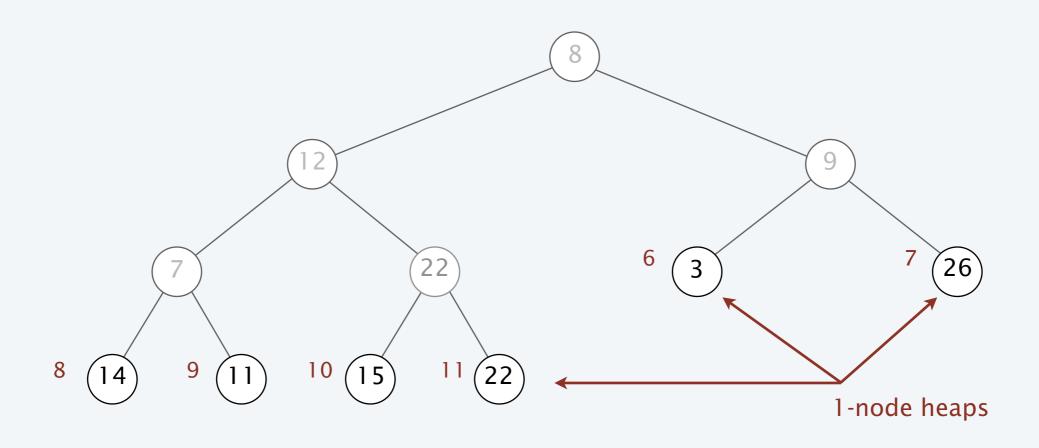
we assume array entries are indexed 1 to n

array in arbitrary order

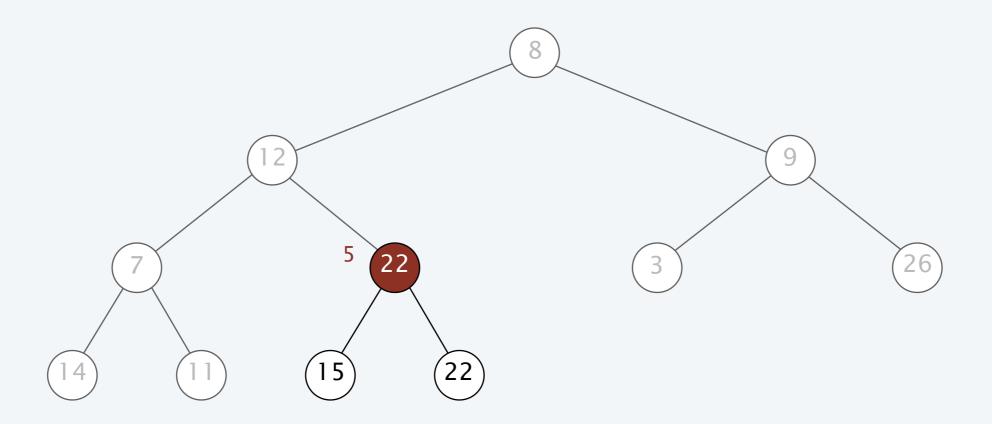


8	12	9	7	22	3	26	14	11	15	22
1	2	3	4	5	6	7	8	9	10	11

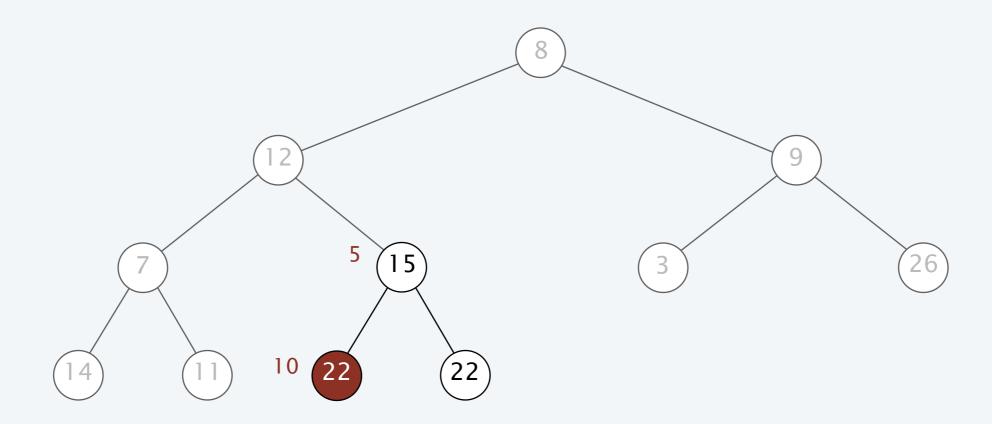
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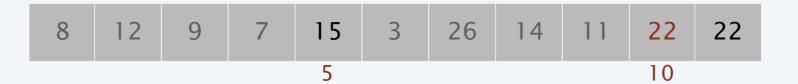


Heapify. For each element in reverse-array order, sink it down.



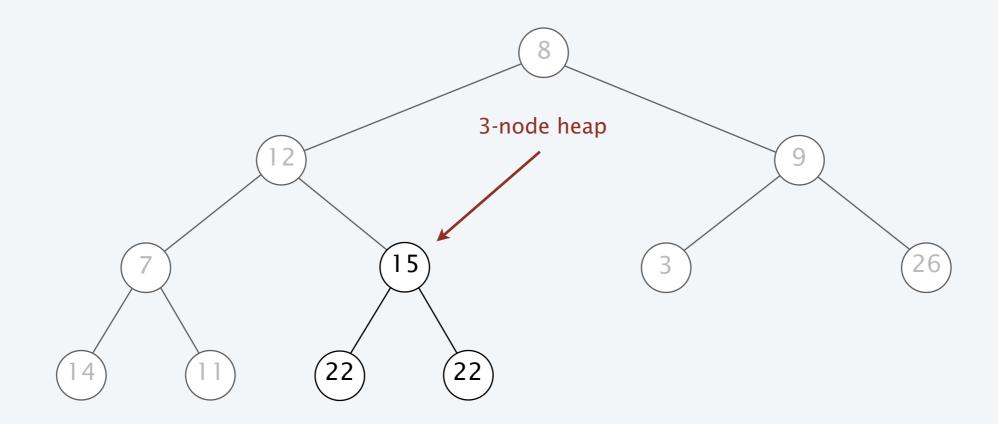
Heapify. For each element in reverse-array order, sink it down.





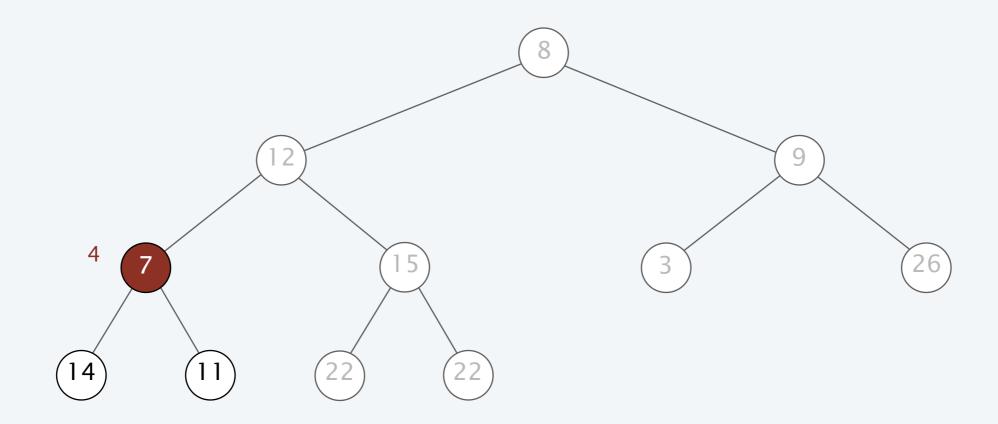
Heapify. For each element in reverse-array order, sink it down.

sink 5



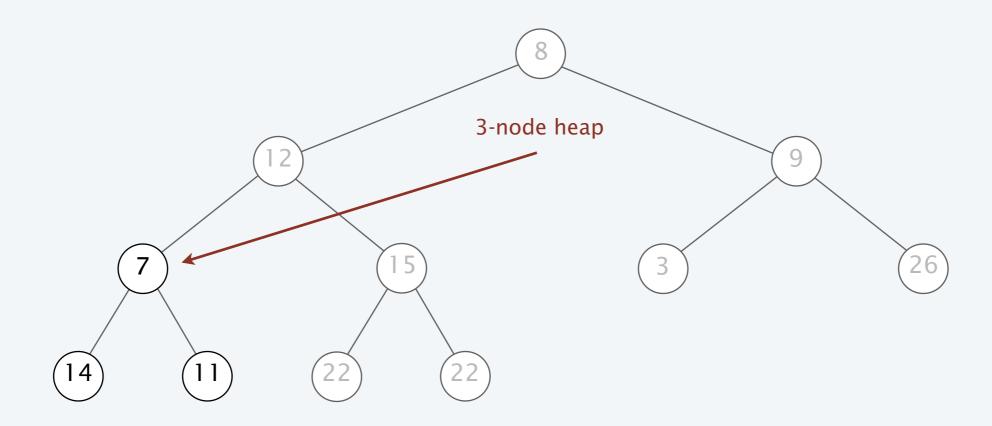
8 12 9 7 15 3 26 14 11 22 22

Heapify. For each element in reverse-array order, sink it down.



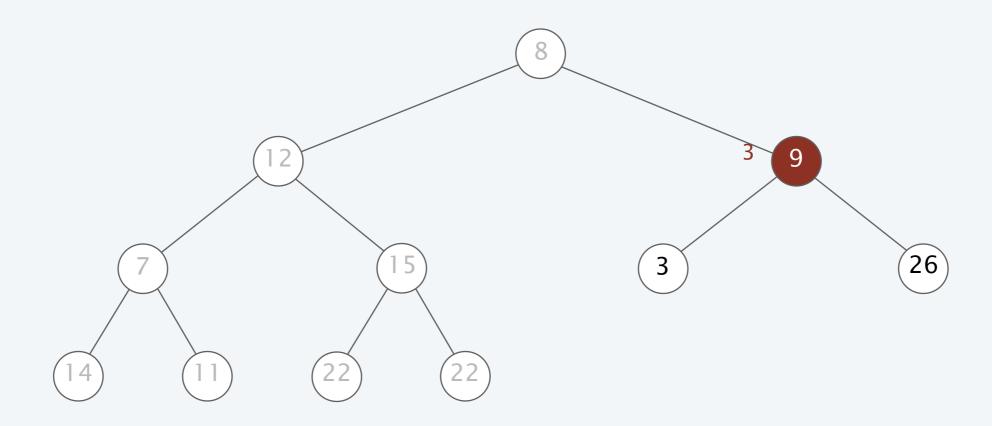
Heapify. For each element in reverse-array order, sink it down.

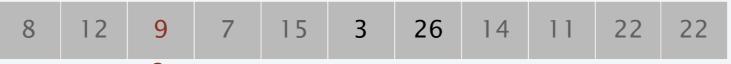
sink 4



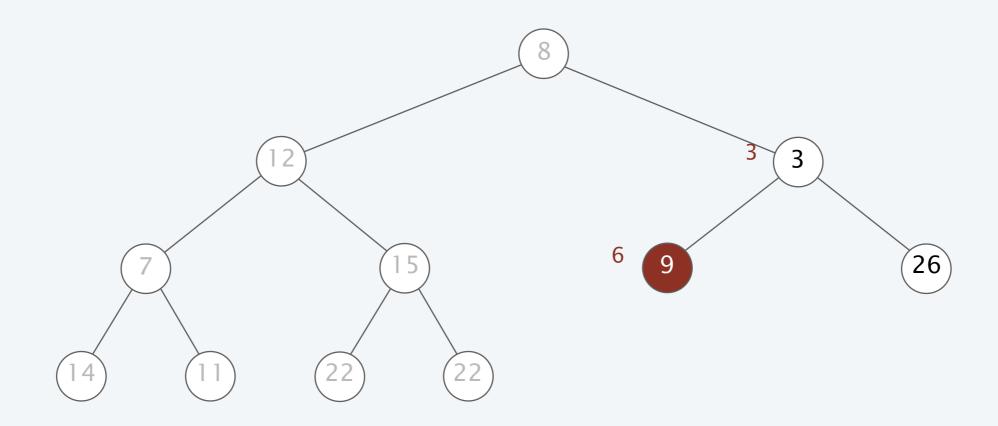
8 12 9 **7** 15 3 26 **14 11** 22 22

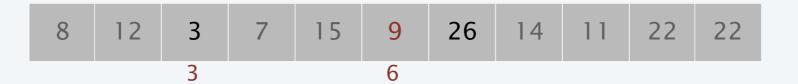
Heapify. For each element in reverse-array order, sink it down.





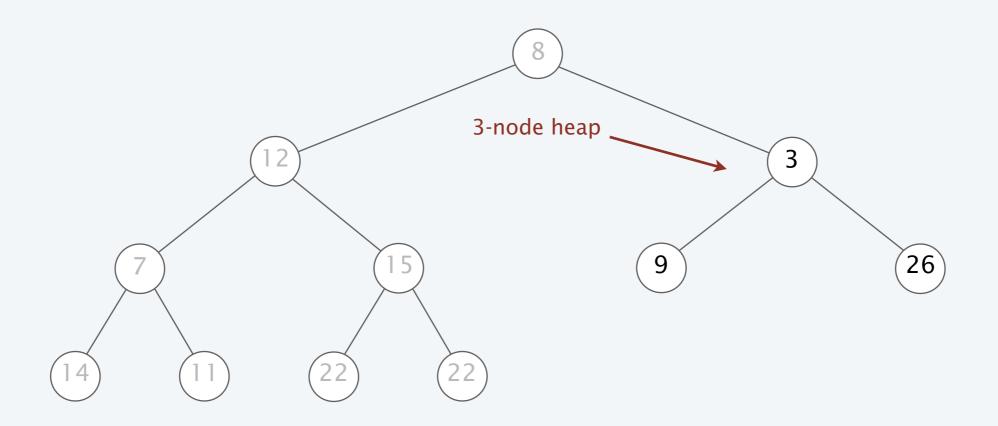
Heapify. For each element in reverse-array order, sink it down.





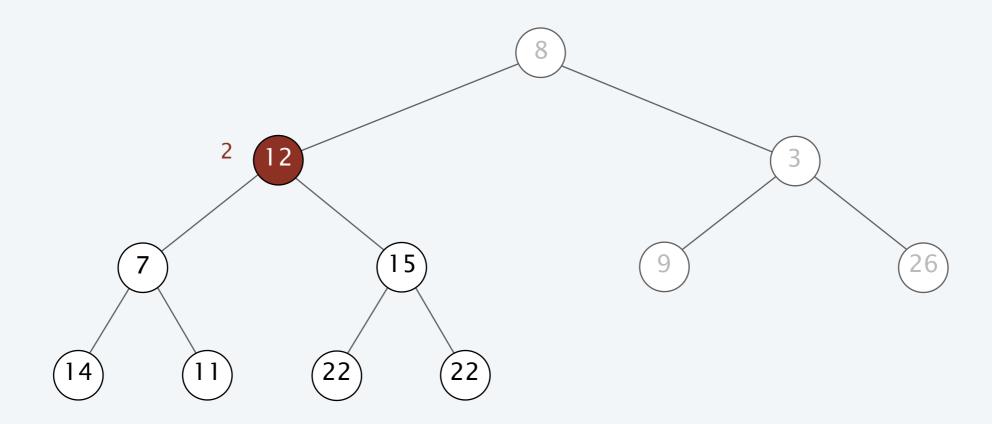
Heapify. For each element in reverse-array order, sink it down.

sink 3



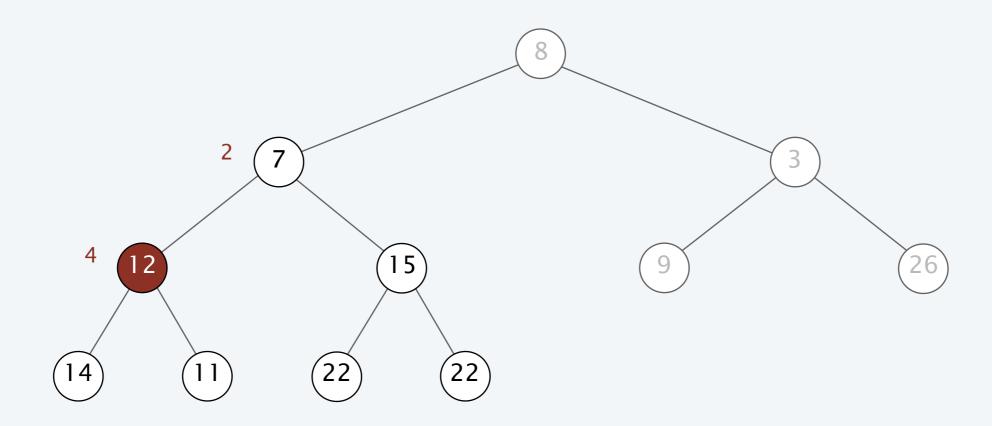
8 12 3 7 15 26 26 14 11 22 22

Heapify. For each element in reverse-array order, sink it down.



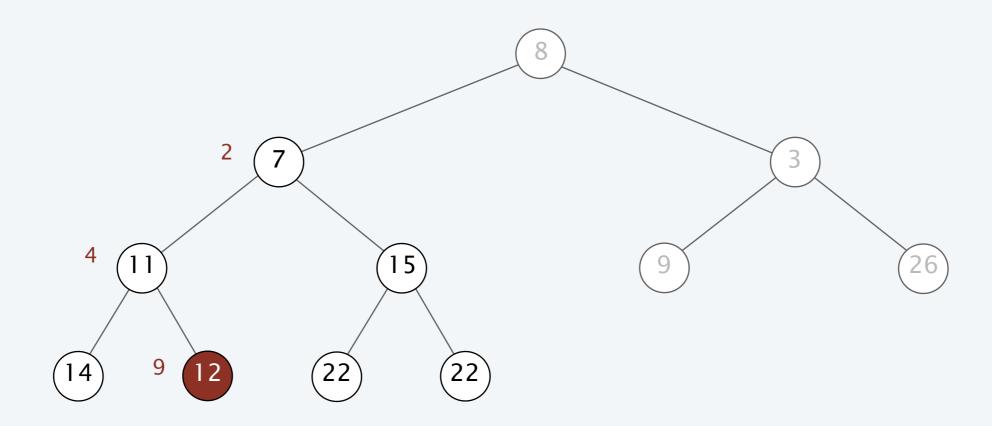


Heapify. For each element in reverse-array order, sink it down.





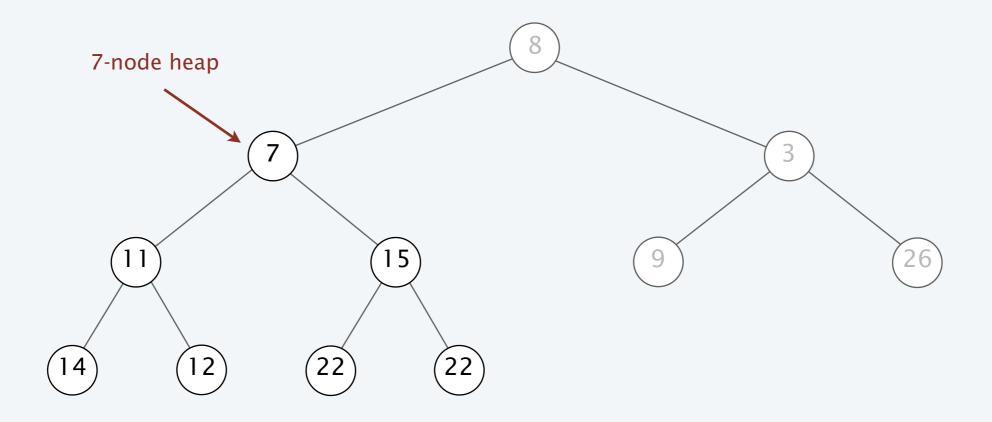
Heapify. For each element in reverse-array order, sink it down.





Heapify. For each element in reverse-array order, sink it down.

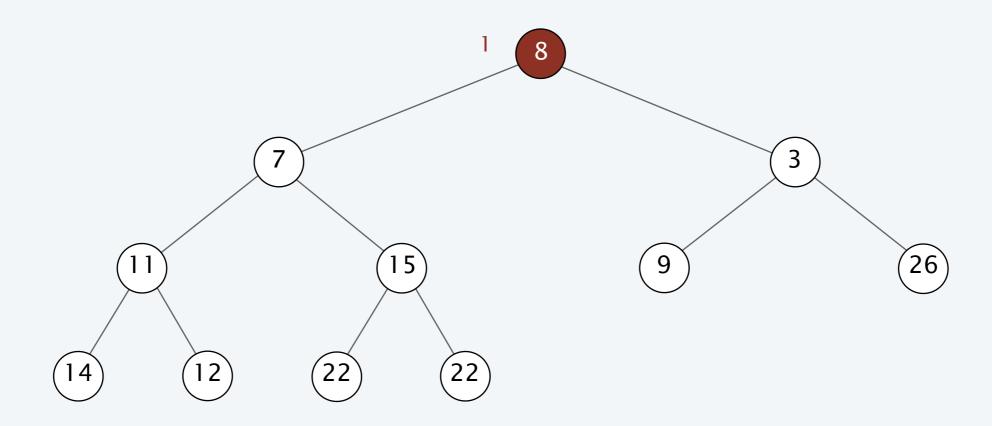
sink 2



8 7 3 11 15 9 26 14 12 22 22

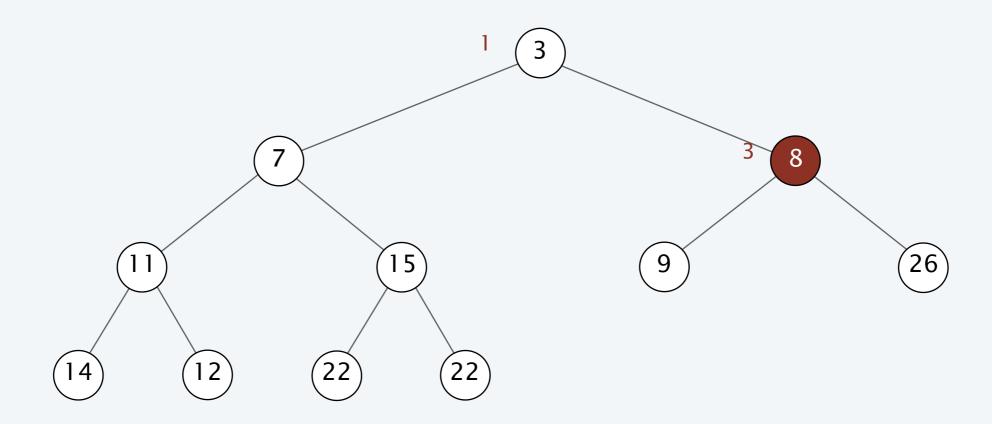
Heapify. For each element in reverse-array order, sink it down.

sink 1



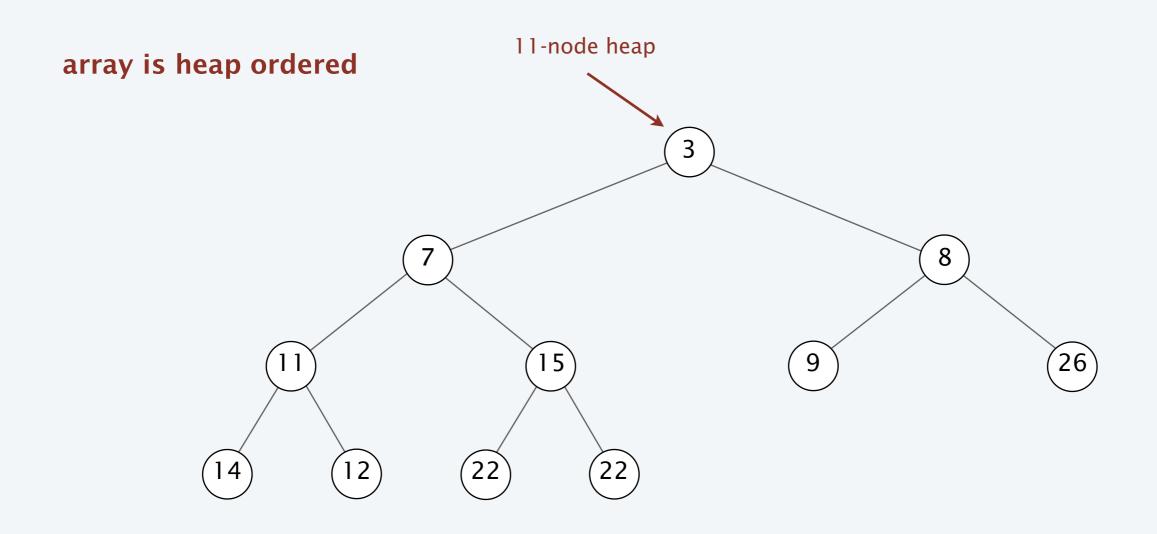
 8
 7
 3
 11
 15
 9
 26
 14
 12
 22
 22

Heapify. For each element in reverse-array order, sink it down.





Heapify. For each element in reverse-array order, sink it down.



Priority queues performance cost summary

operation	linked list	binary heap		
Маке-Неар	<i>O</i> (1)	<i>O</i> (1)		
ISEMPTY	<i>O</i> (1)	<i>O</i> (1)		
INSERT	<i>O</i> (1)	$O(\log n)$		
Extract-Min	O(n)	$O(\log n)$		
Decrease-Key	<i>O</i> (1)	$O(\log n)$		
DELETE	<i>O</i> (1)	$O(\log n)$		
Meld	<i>O</i> (1)	O(n)		
FIND-MIN	O(n)	<i>O</i> (1)		