CS 344: Operating Systems Lab

ASSIGNMENT 0A

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Question 1

ex1.c code

Output

Question 2

1st instruction: [f000:fff0] 0xffff0: ljmp $0x3630,$0xf000e05b

- Jump to CS = $0xf000 & IP = 0xe05b

- 0x3630 is jump to this CS (earlier in the BIOS)

- 0xf000e05b is the IP which is different from the lab because it is 32 bits rather than 16 bits and that is all the way into the top of the extended memory location but before the memory mapped PCI device location reserved by the BIOS

2nd Instruction: [f000:e05b] 0xfe05b: cmpw $0xffc8,%cs:(%esi)

- Compare content at 0xffc8 & with content at code segment offset with value at esi.

- esi:- 32-bit source index register

3rd Instruction: [f000:e062] 0xfe062: jne 0xd241d0b0

- Jump to 0xd241d0b0 if the above comparison does not set ZF

4th instruction: [f000:e066] 0xfe066: xor %edx,%edx

- ZF was set thus jump of previous instruction doesn't occur

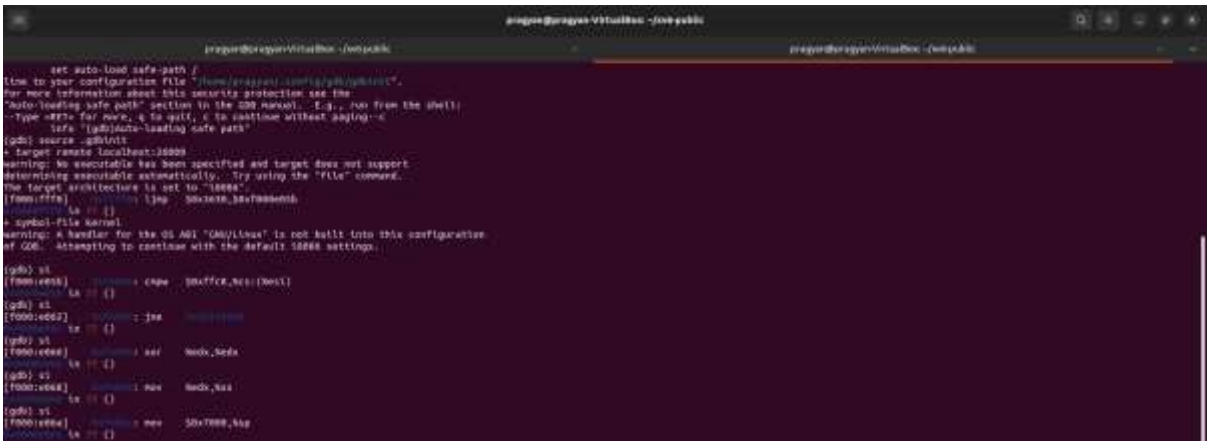
- It set edx to zero, edx is 32-bit general-purpose register.

5th instruction: [f000:e068] 0xfe068: mov %edx,%ss

- Move content of stack segment register(ss) to edx

6th instruction: [f000:e06a] 0xfe06a: mov $0x7000,%sp

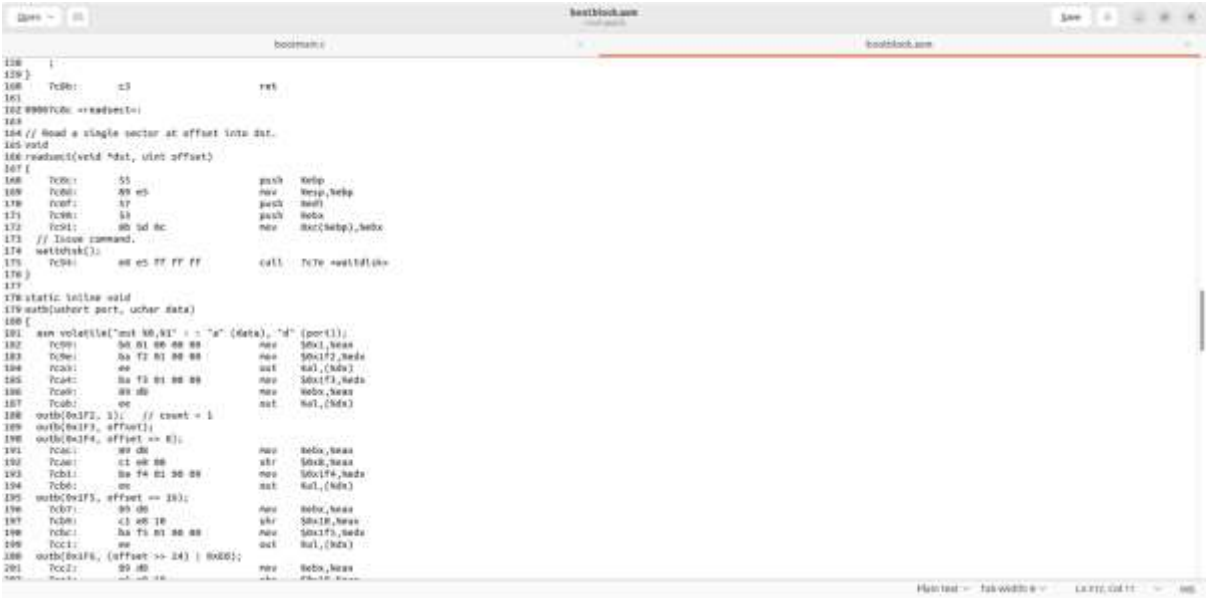
- Move content at the location pointed 16-bit stack pointer(sp) to $0x7000

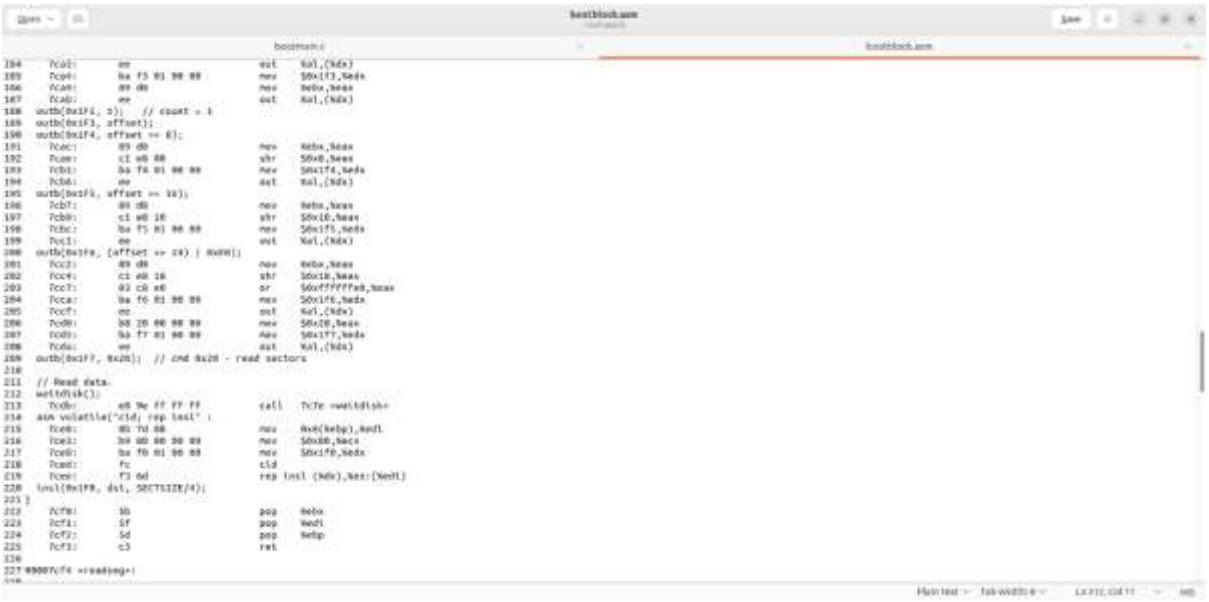


Question 3

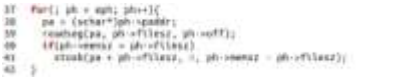
The code for readsect() is given below

The assembly code for readsect() is given below





The for loop that reads the sectors of the kernel from the disk is given below



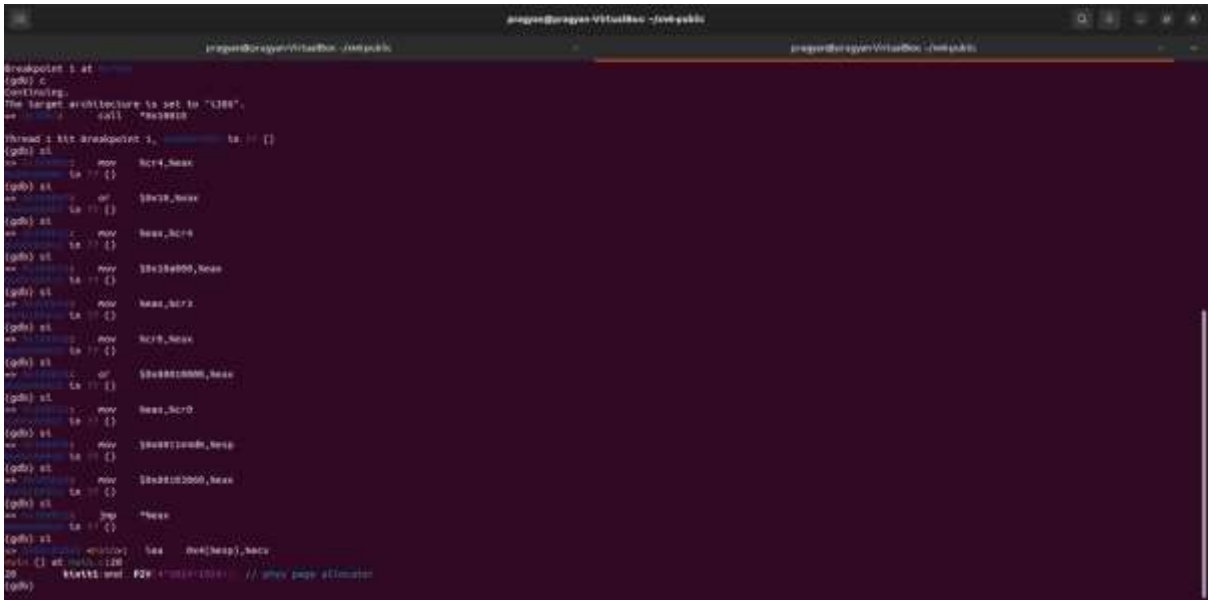
The first instruction of the loop

The last instruction of the loop

The explanation for the first instruction is that the first operation on entering the for loop will be comparison between the values of ph and eph because the loop will run only when ph < eph. The explanation of last instruction is that the loop ends when the values of ph and eph become equal and hence the loop jumps to the next instruction at 0x7d91. Hence the jump instruction will be the last instruction of the for loop. The next instruction after the for loop is

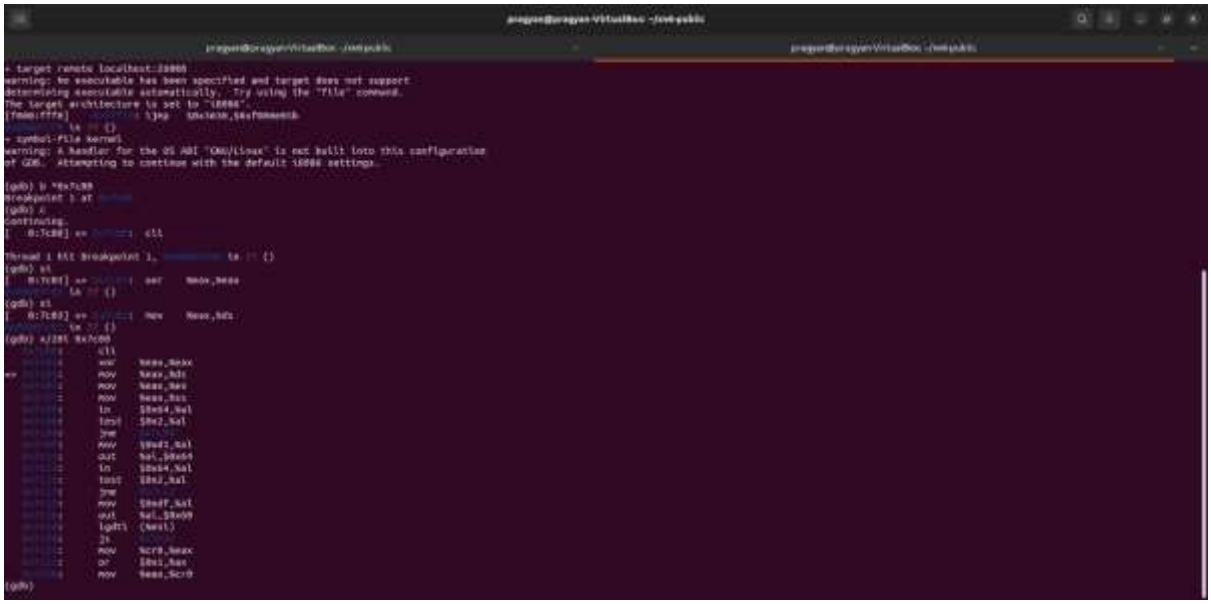
```
313     7d81:       ff 15 10 00 01 00       call   *0x10010
```

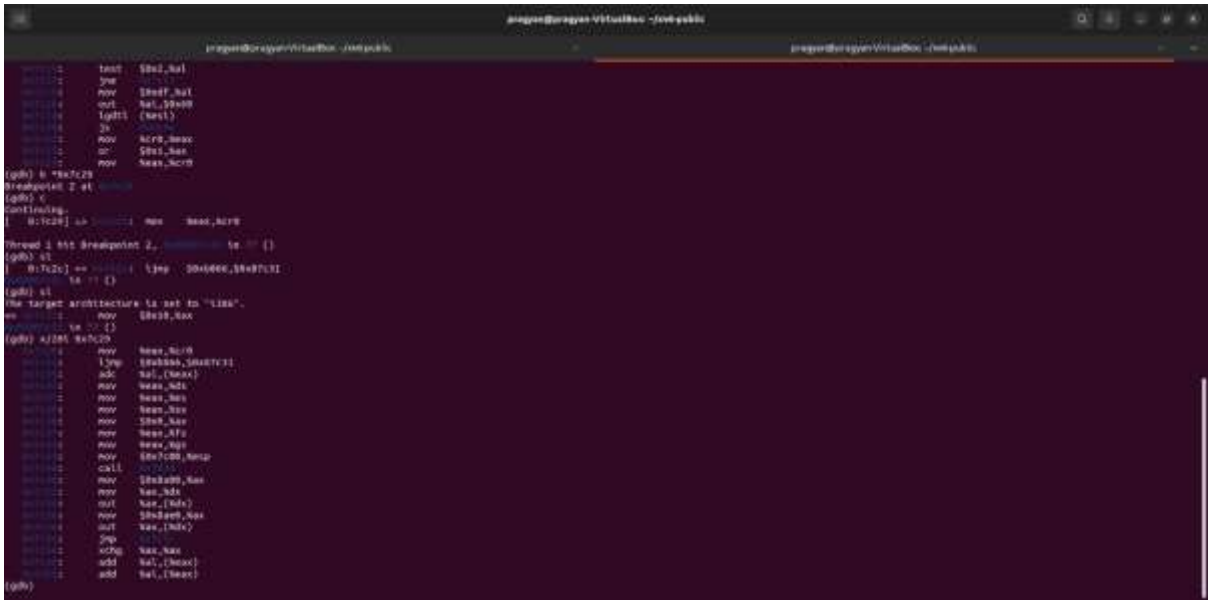
Making a breakpoint at that address and then stepping into further instructions gives the following output.



(a)

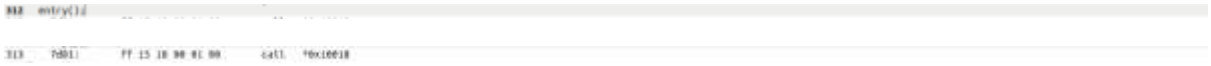
The command line $(SEG_KCODE<<3), $start32 causes the switch from 16 to 32-bit mode in bootasm.S





(b)

By analysing the contents of bootasm.S, bootmain.c and bootblock.asm, we conclude that bootasm.S switches the OS into 32-bit mode and then calls bootmain.c which first loads the kernel using ELF header and the enters the kernel using entry(). Hence the last instruction of bootloader is entry(). Looking for the same in bootblock.asm, we find out the instruction to be



which is a call instruction which shifts control to the address stored at 0x10018 since dereferencing operator (*) has been used. Now we need to know the starting address of the kernel. We can find this by two methods:

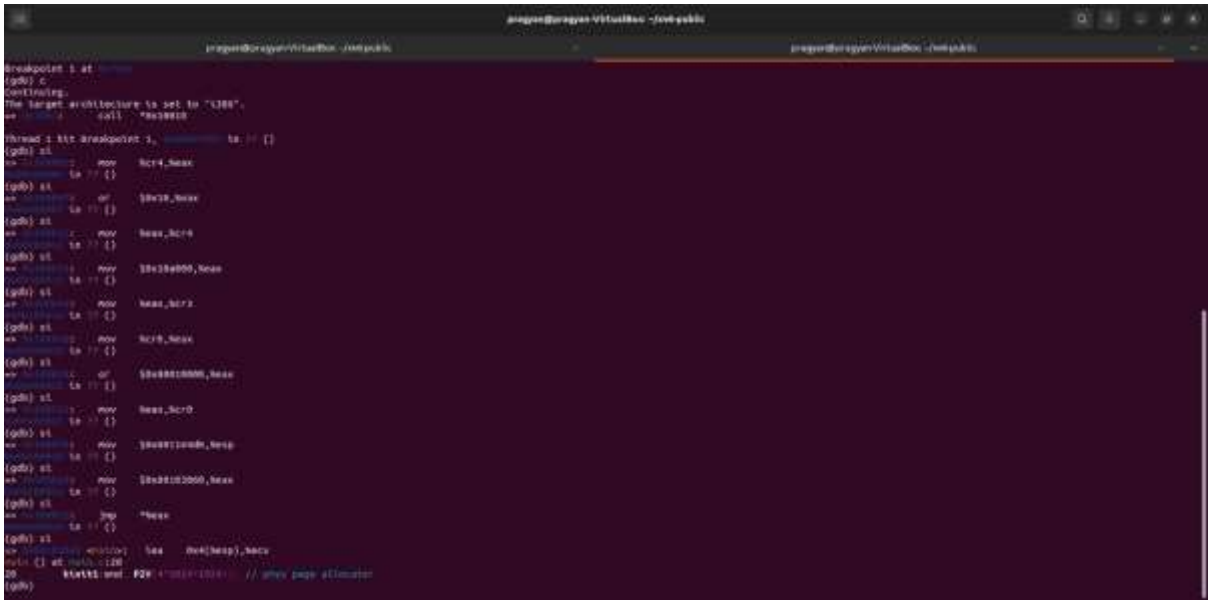
(i) By looking at the first word of memory stored at 0x10018 (by using the command "x/1x 0x10018")

(ii) By looking at the contents of "objdump -f kernel"

After getting the starting address of kernel, we need to see what is the instruction stored at that address to get the first instruction of kernel. We can do this by two methods:

(i) By using "x/1i 0x0010000c"

(ii) By looking into kernel.asm



Hence, the first instruction of kernel is

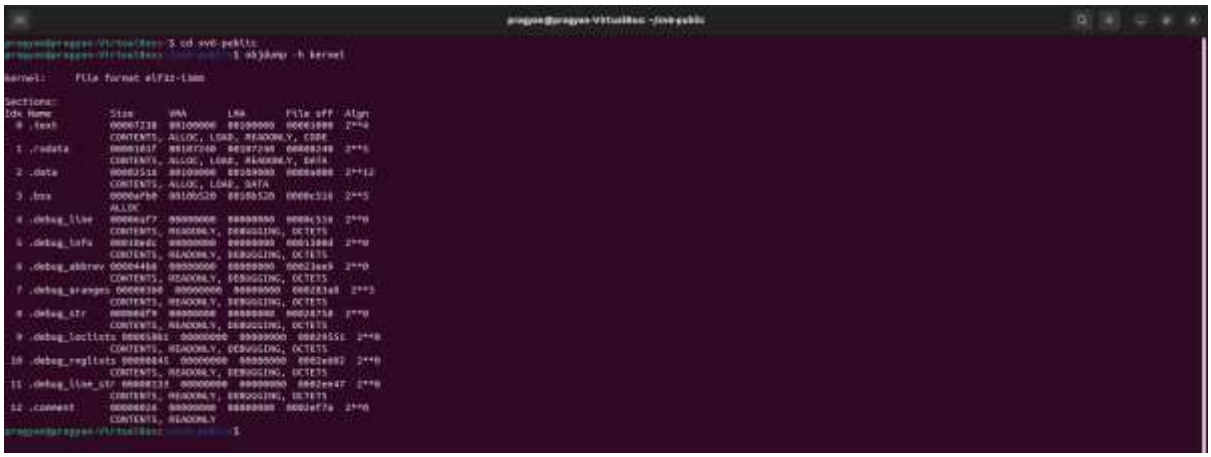
0x10000c: mov %cr4,%eax

(c)

Question 4

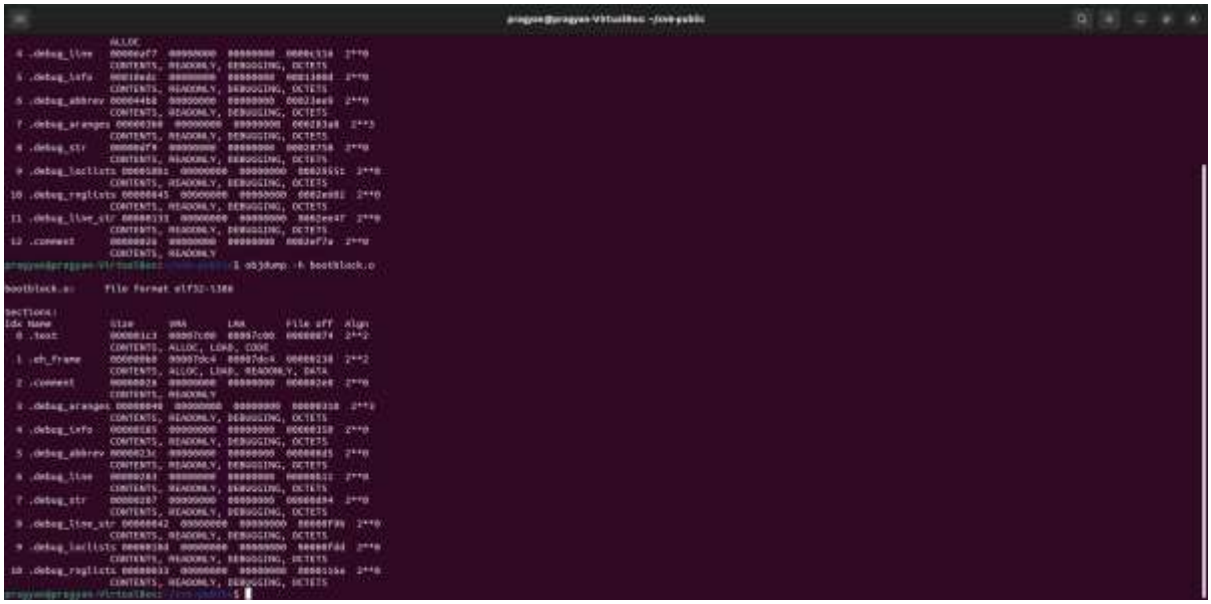
pointers.c

objdump -h kernel



As we can see in the above screenshot, VMA and LMA of .text section is different indicating that it loads and executes from different addresses.

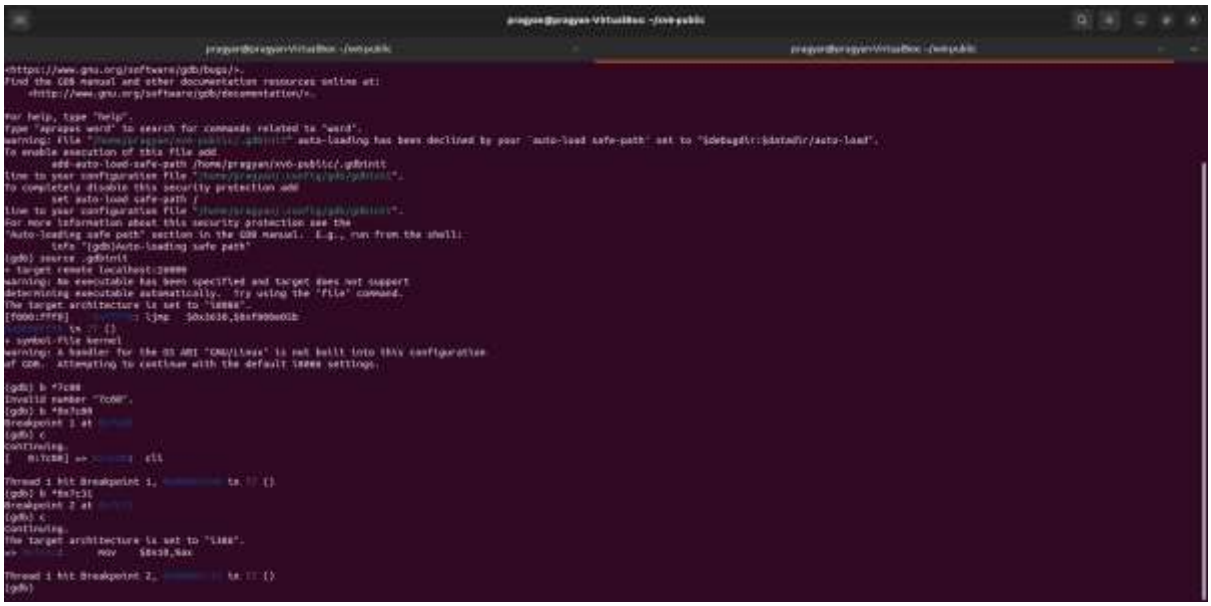
objdump -h bootblock.o



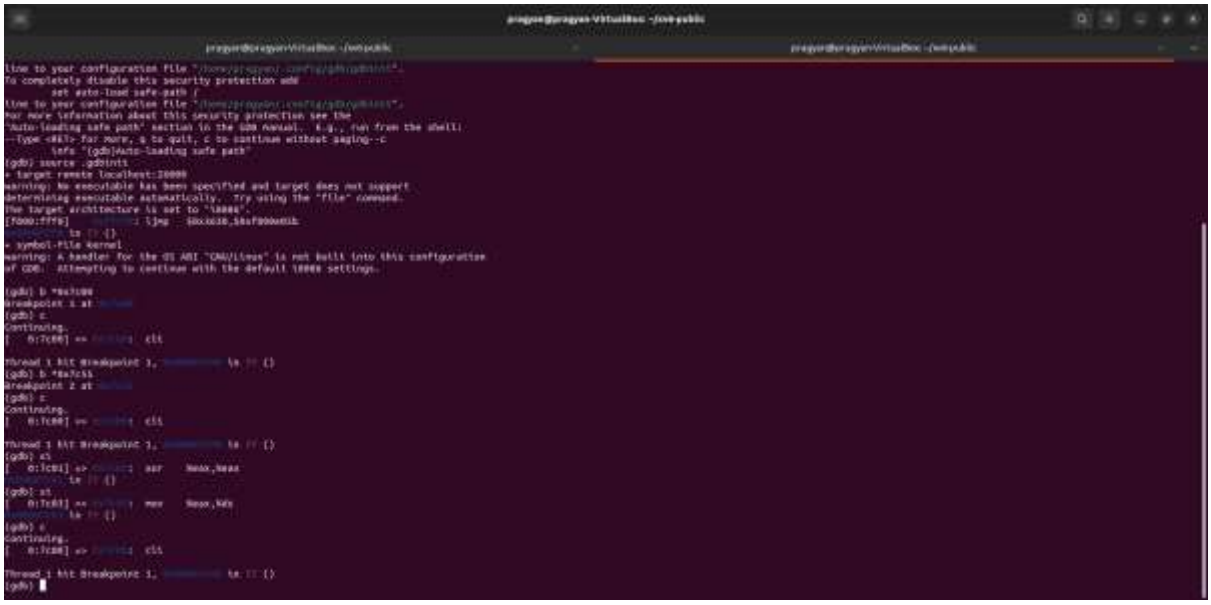
As we can see in the above screenshot, VMA and LMA of .text section is same indicating that it loads and executes from the same address.

Question 5

When boot loader's link address is 0x7C00 then commands are running properly and transition from 16 to 32 bit was occurring at 0x7C31 address location as seen below:



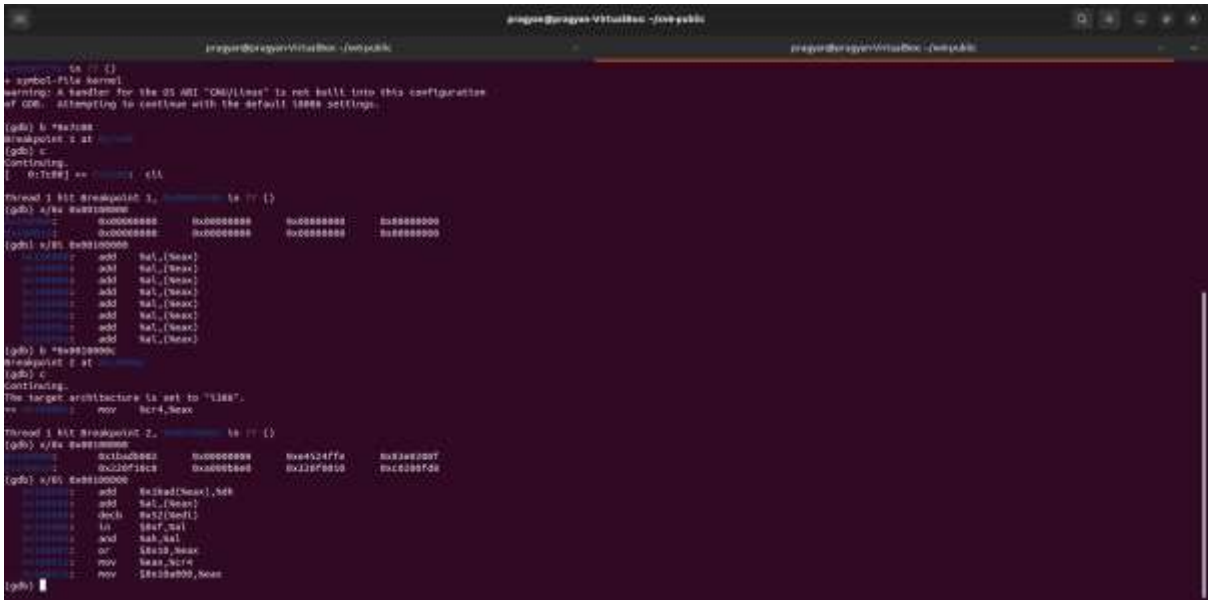
But when the boot loader's link address is changed to any other address (we took 0x7C24 in this case), after running make clean make and restarting gdb and continuing from address location 0x7C00, then the boot loader is restarting again and again after running some instructions in the gdb.



As seen in the image above, we tried to run commands after continuing from breakpoint at 0x7C00 address location and we always end up hitting the same breakpoint at 0x7C00. Also 16 to 32 bit architecture change didn't occurred as breakpoint b *0x7C55 is not hit which should be responsible for architecture change in this case.

Question 6

For this experiment, we have to examine the 8 words of memory at 0x00100000 at two different instances of time, the first when the BIOS enters boot loader and the second when the boot loader enters the kernel. For this, we will use the command "x/8x 0x00100000" but before that we will have to set our breakpoints. The first breakpoint will be at 0x7c00 because this is the point where the BIOS hands control over to the boot loader. The second breakpoint will be at 0x0010000c because this is the point when the kernel is passed control by the bootloader.



As we can see in the diagram, we get different values at both the breakpoints. The explanation to this is as follows. The address 0x00100000 is actually 1MB which is the address from where the kernel is loaded into the memory. Before the kernel is loaded into the memory, this address contains no data (i.e. garbage value). By default, all the uninitialized values are set to 0 in xv6. Hence, when we tried to read the 8 words of memory at 0x00100000 at the first breakpoint, we got all zeroes since no data had been loaded until that point. When we check the values at the second breakpoint, the kernel has already been loaded into the memory and thus this address now contains meaningful data instead of zeroes.