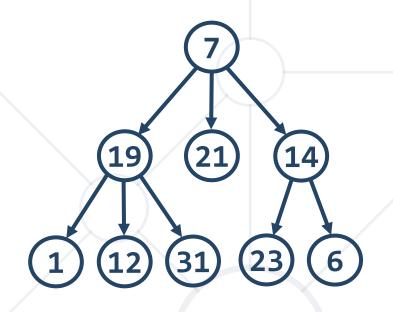
## Trees Representation and Traversal (BFS, DFS)

Trees Related Terminology and Traversal Algorithms



**SoftUni Team Technical Trainers** 







#### **Table of Contents**



#### 1. Why Trees?

- Definition and use cases of trees
- 2. Trees and Related Terminology
  - Node, Edge, Root, etc.
- 3. Implementing Trees
  - Recursive Tree Data Structure
- 4. Traversing Tree-Like Structures
  - BFS and DFS traversal





Definition and Use Cases of Trees

#### Summary





- So far, we have learned how to implement linear data structures like: List, Queue, Stack, LinkedList etc...
- We did a great job and learned how to take the best complexity we can, was that enough?
- Actually more of the operations we want to do like search, insert or remove are linear for unordered structures (sometimes we can do O(1)) but not for search

#### **Linear Data Structures - Types**



Atop an array



- Removing and searching with O(n)
- Inside sorted arrays remove, search O(log(n)), but sorting is needed on every new element added
- By using Node implementation
  - Ability to add and remove with O(1)
  - Every other operation is O(n)
  - Even sorted O(log(n)) is not achievable. But why?

#### Why Trees?





- We want not only to store data add or remove elements in efficient manner but also to search for elements, but can we do better than O(n)?
- Lets try to get down to O(log(n)) by using trees and see if we can

#### **Other Tree Benefits**



By learning how to work with trees you actually learn how to work with:

- Hierarchical structures like: file system, project structures and code branching, NoSQL data storage etc...
- Markup languages:
  - HTML
  - XML
- DFS and BFS algorithms



## **Trees and Related Terminology**

Node, Edge, Root, Etc.

#### **Tree Definition**

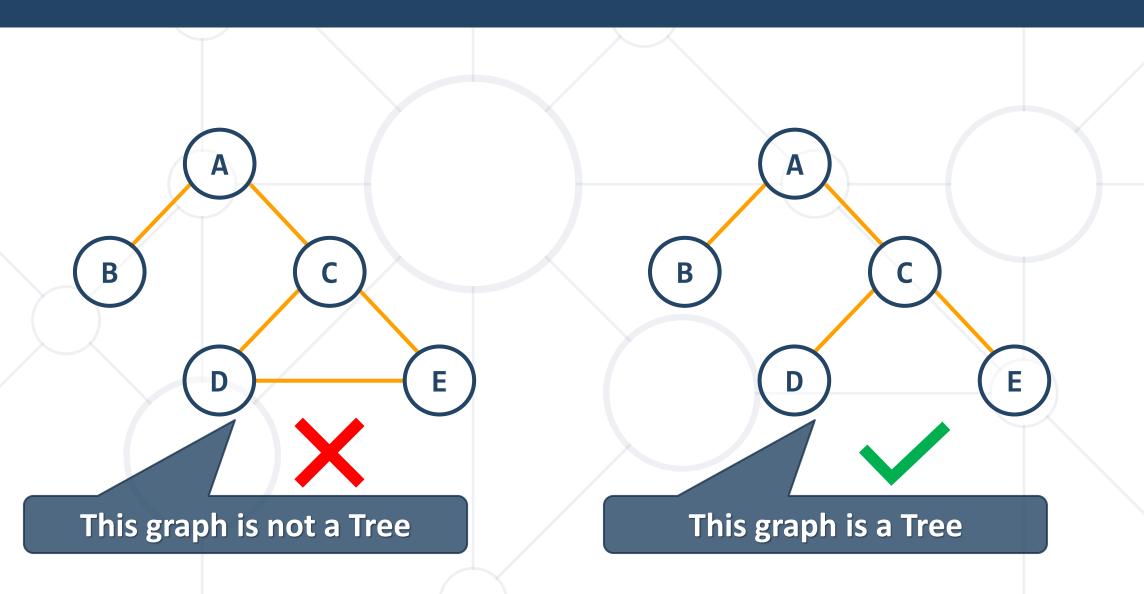




- Tree is a non-linear data structure which organizes data in a hierarchical structure, and this is a recursive definition.
- A tree is a connected graph without any circuits.
- If in a graph, there is one and only one path between every pair of vertices, then the graph is called as a tree.

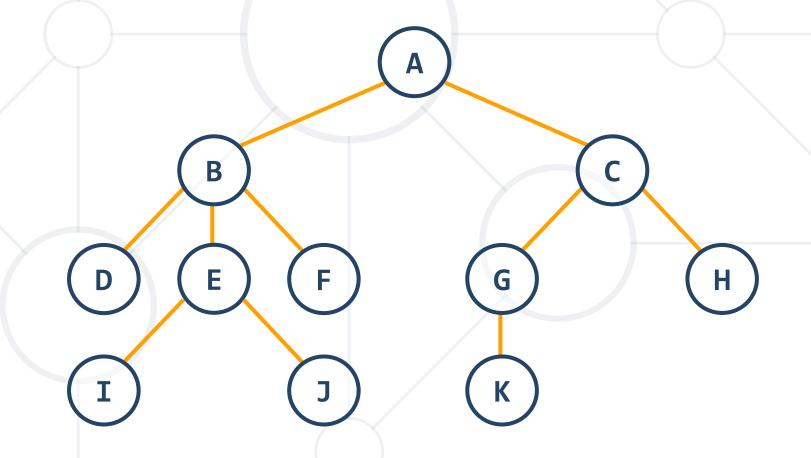
### **Tree Definition - Example**





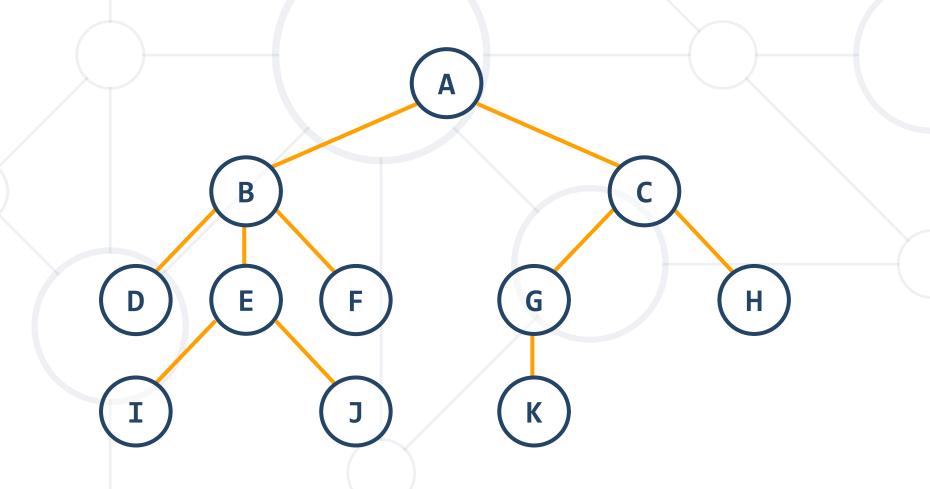


Root – the top node in a tree, the prime ancestor, there must be only one root node



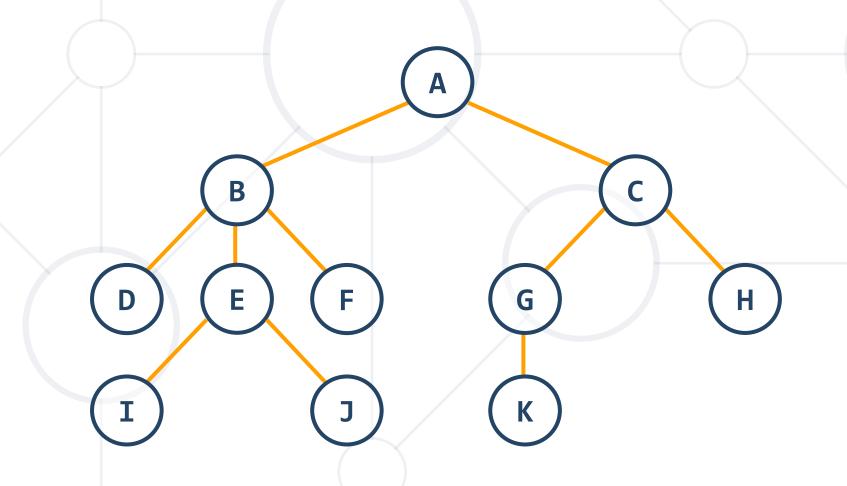


Edge – the connection between one node and another



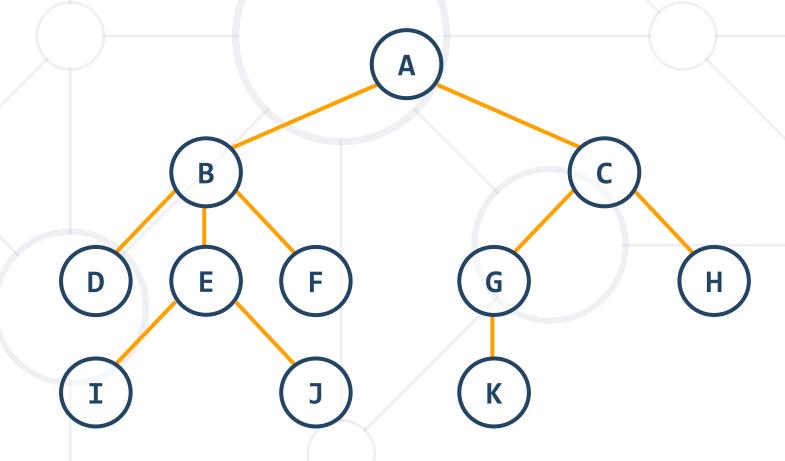


Parent – node, which has one or more children



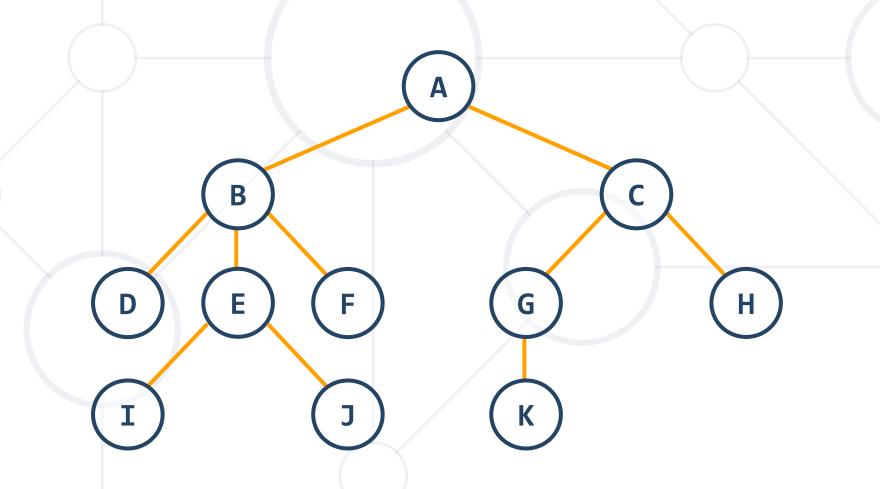


 Child – node directly connected to another node when moving away from the root, an immediate descendant



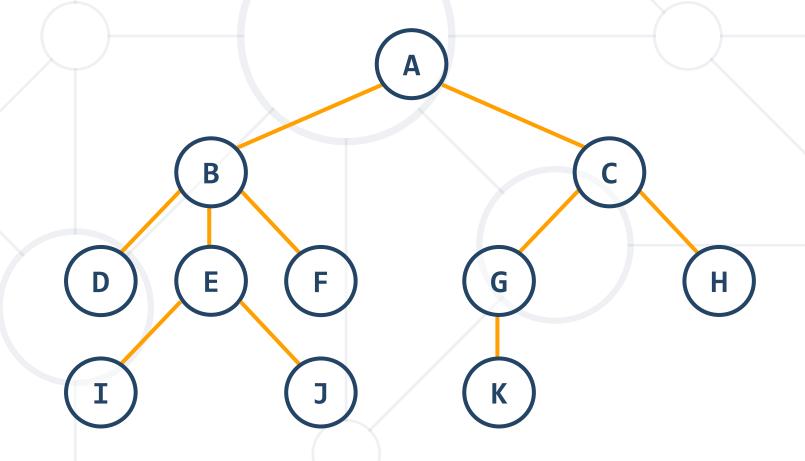


Siblings – nodes with the same parent



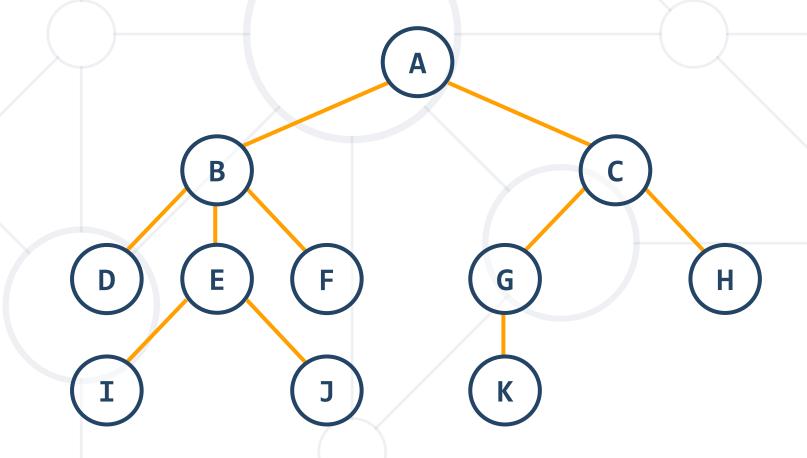


 Ancestor – node reachable by repeated proceeding from child to parent



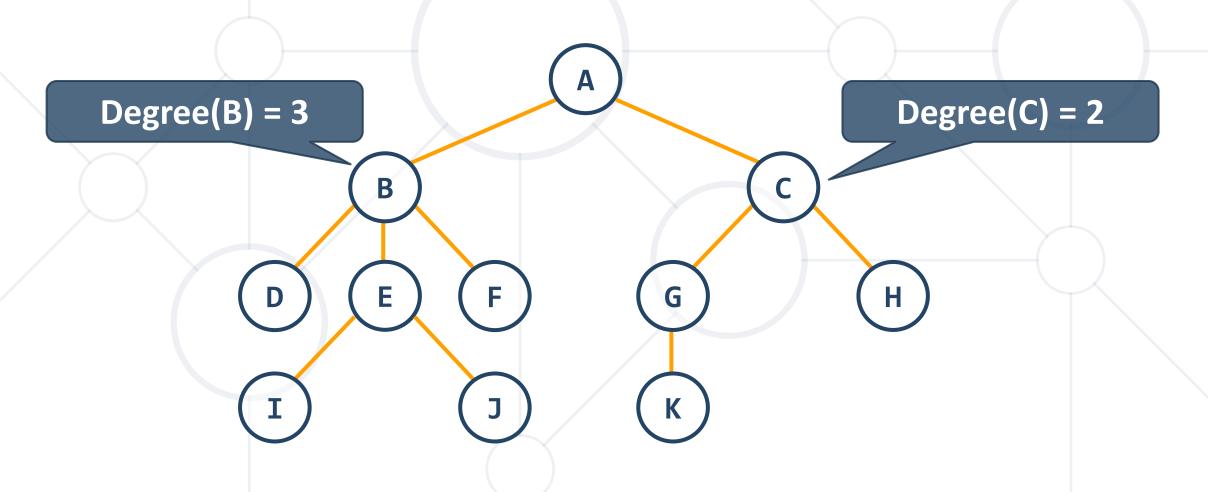


 Descendant – node reachable by repeated proceeding from parent to child



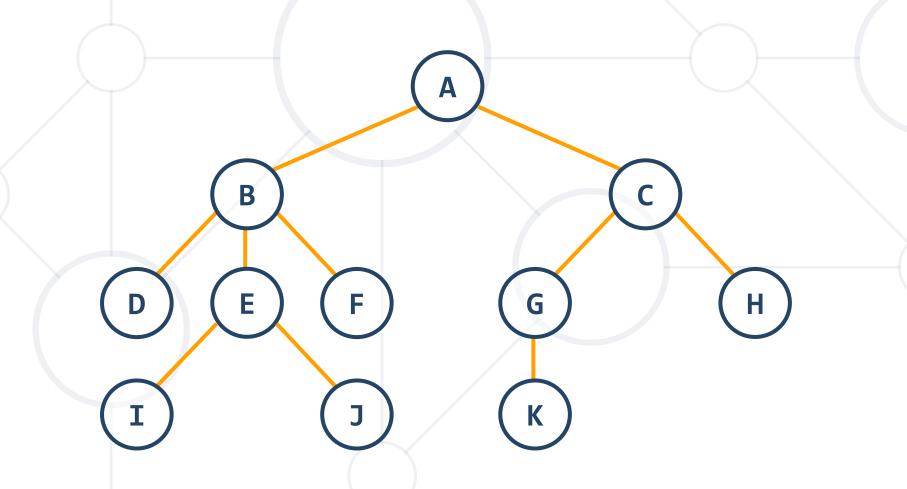


Degree – total number of a children for a node



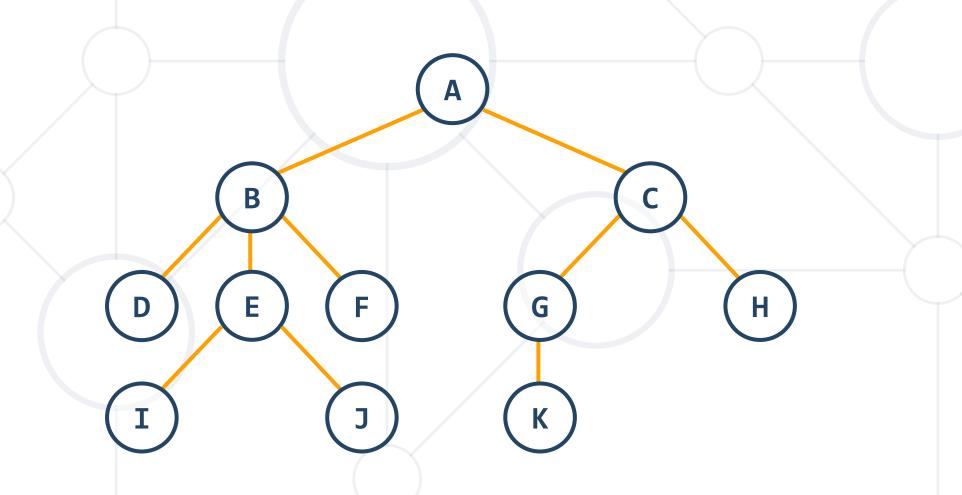


Leaf Node – node without any children



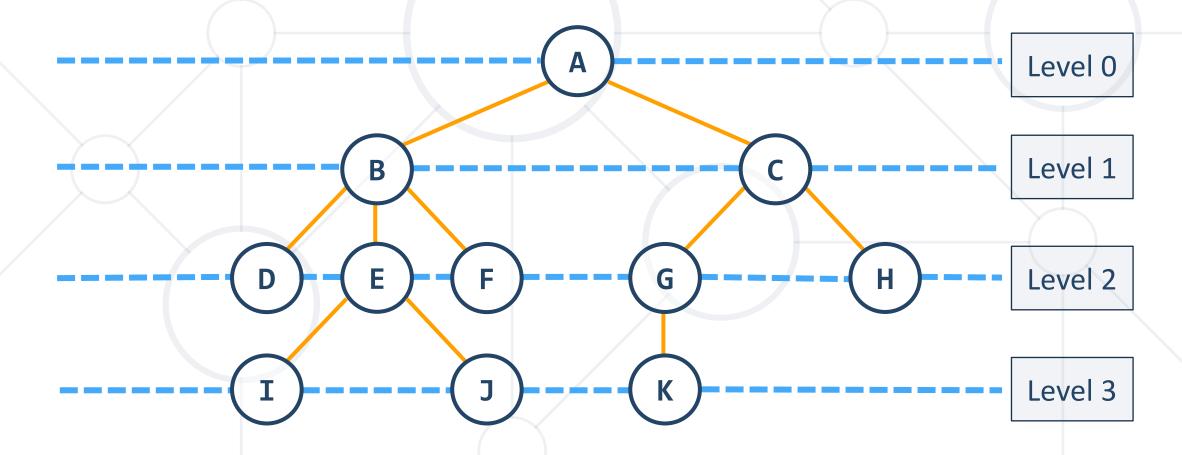


Internal Node – every non-leaf node, also know as branch





 Level – each step from top to bottom starting from 0 is called level of a tree



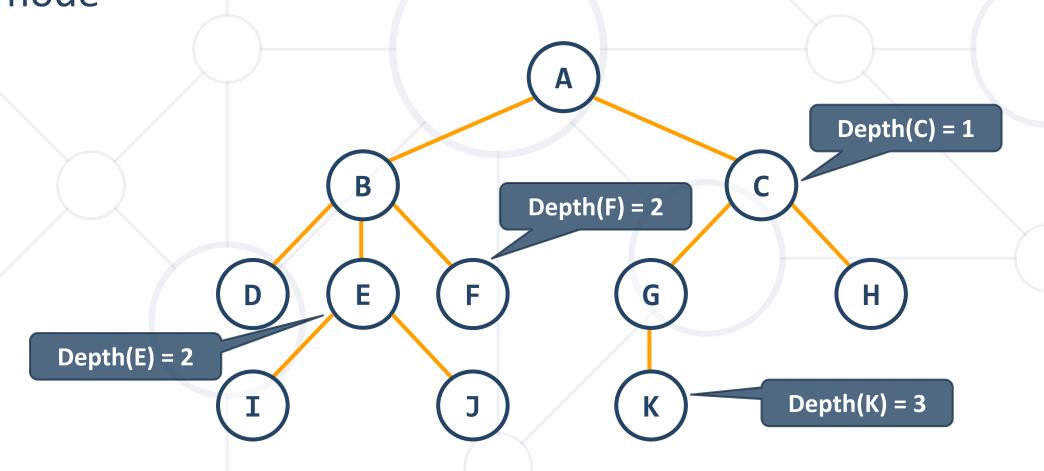


 Height – longest number of edges from a leaf to a specific node is called height of a node. Height of a tree is the height of the

root node Height(A) = 3Height(C) = 2 Height(F) = 0H G

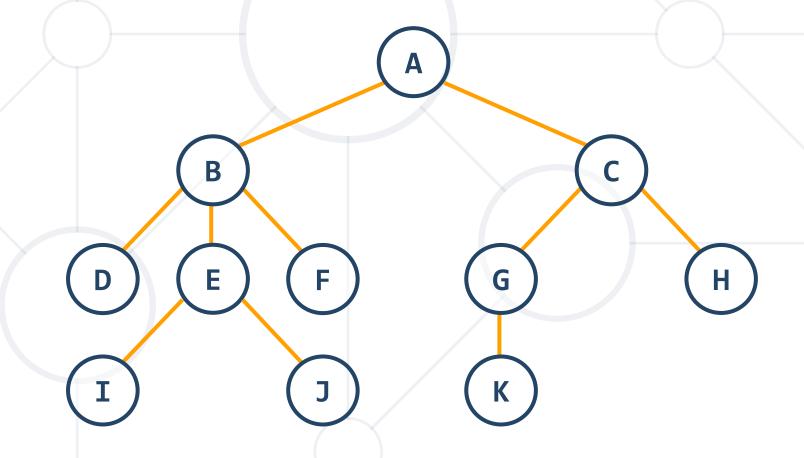


Depth – the total number of edges from the root to a particular node



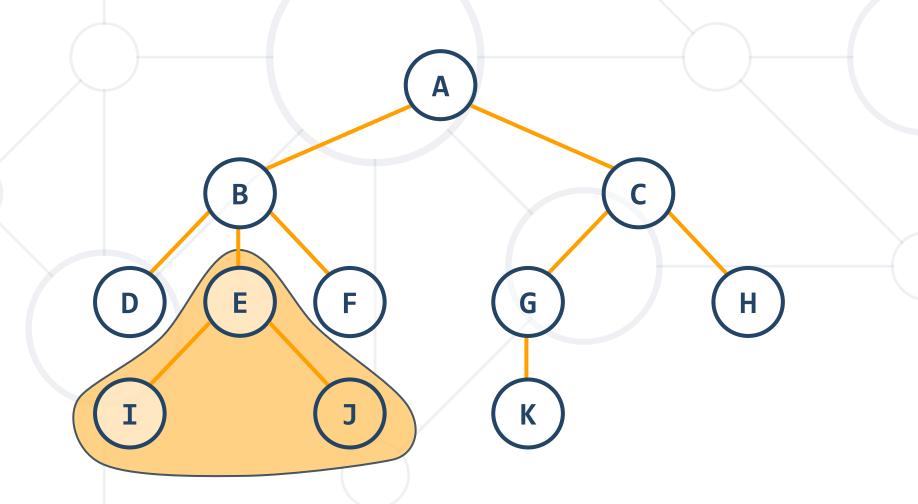


 Path – the sequence of nodes and edges from one node to another. Length of the path is the number of nodes



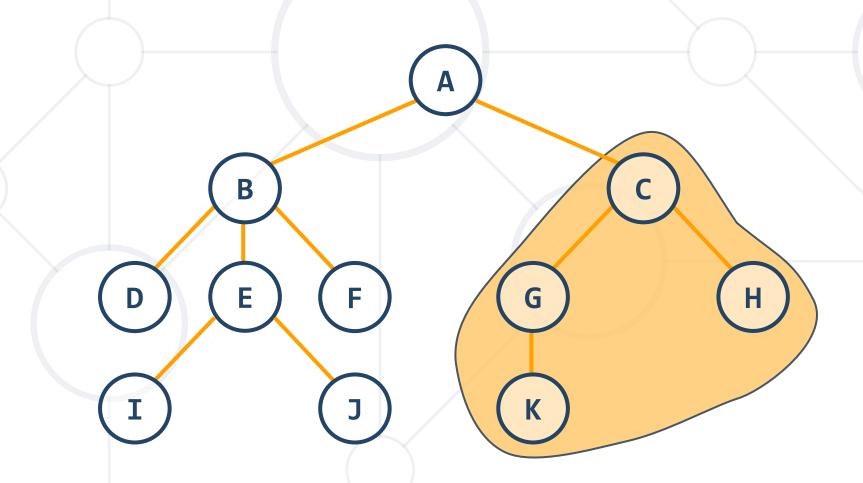


Subtree – each child from a node forms a subtree recursively



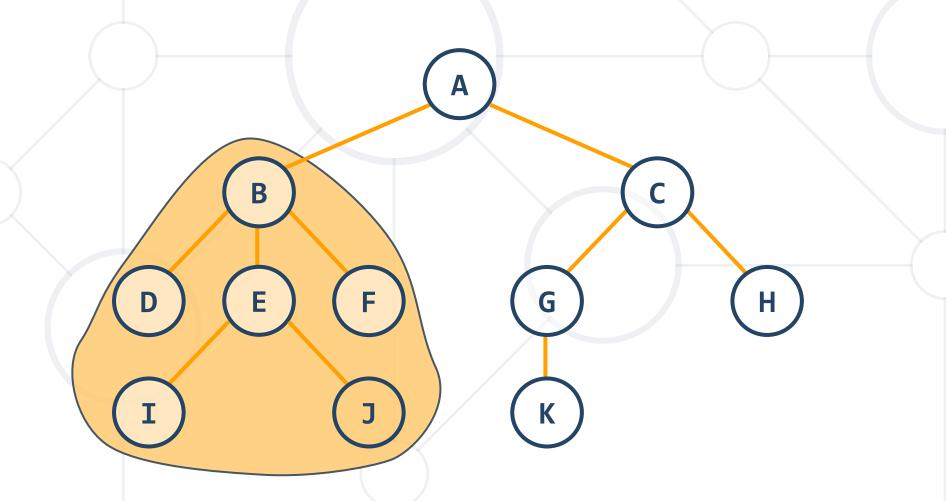


Subtree – each child from a node forms a subtree recursively



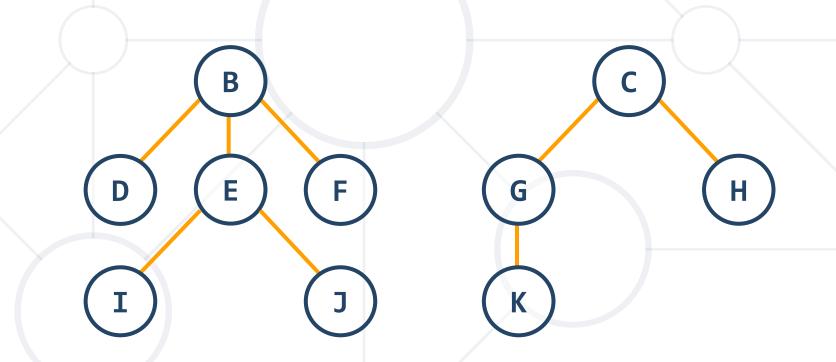


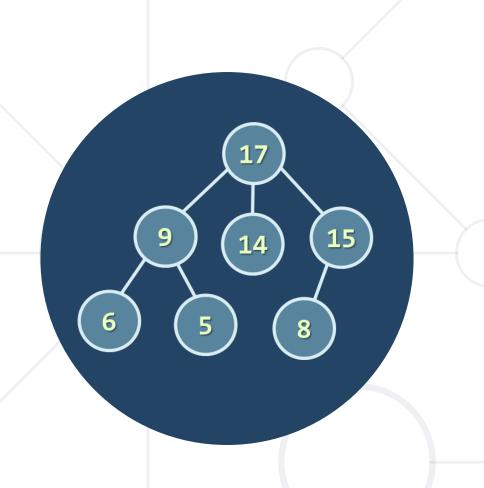
Subtree – each child from a node forms a subtree recursively





■ Forest – Set of disjoint trees





## Implementing Trees

Recursive Tree Data Structure

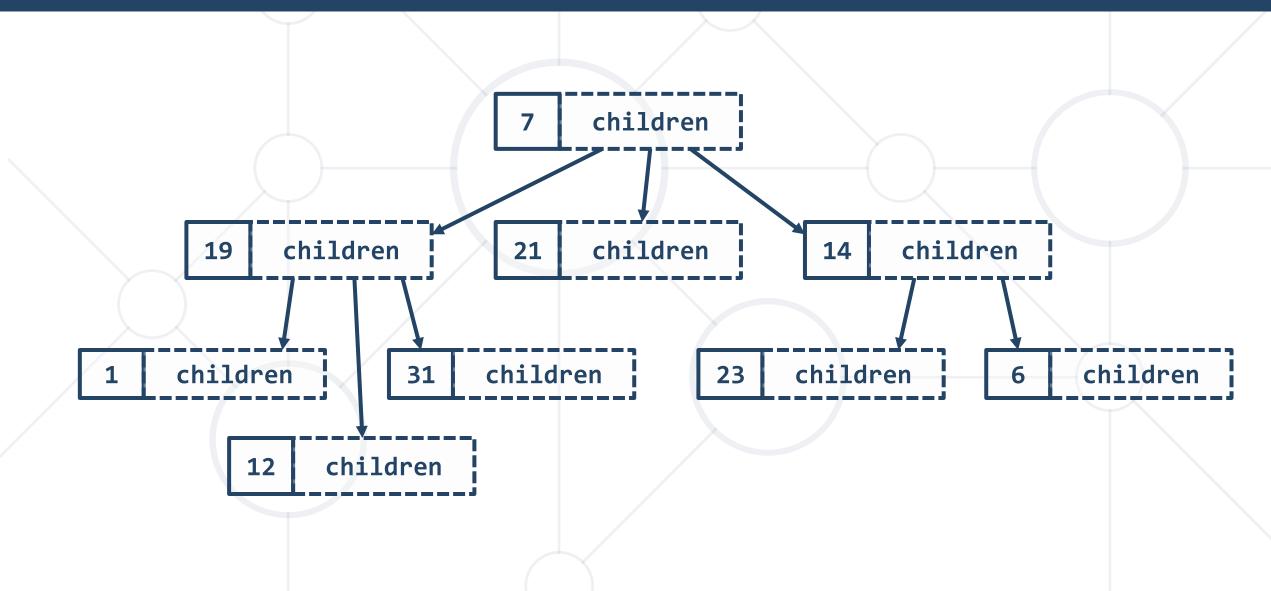
#### **Recursive Tree Definition**



- The recursive definition for tree data structure:
  - A single node is a tree
  - Nodes have zero or multiple children that are also trees

## Tree<int> Structure - Example



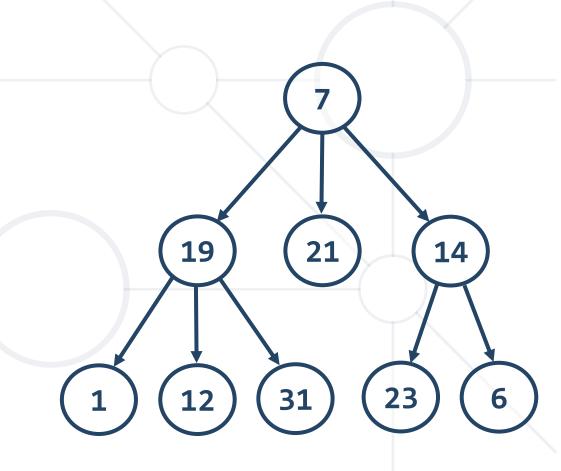


#### **Problem: Implement Tree Constructor**



Create a Tree constructor

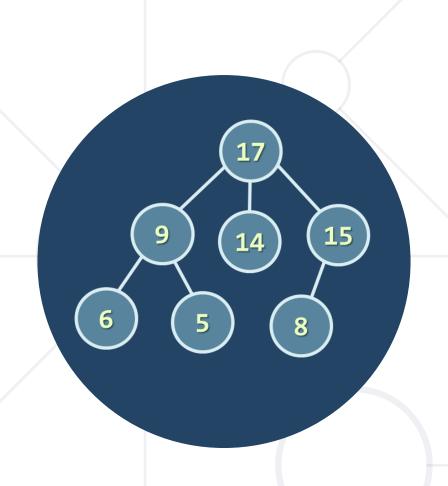
```
Tree<int> tree =
   new Tree<int>(7,
      new Tree<int>(19,
         new Tree<int>(1),
         new Tree<int>(12),
         new Tree<int>(31)),
      new Tree<int>(21),
      new Tree<int>(14,
         new Tree<int>(23),
         new Tree<int>(6))
```



#### **Solution: Implement Tree Constructor**



```
public class Tree<T> : IAbstractTree<T>
    private List<Tree<T>> children;
    private T Value;
    public Tree(T value, params Tree<T>[] children)
        this.Value = value;
        this.children = new List<Tree<T>>();
        for (Tree<T> child : children)
            this.children.Add(child);
```



# **Traversing Tree-Like Structures**

**DFS and BFS Traversals** 

#### **Tree Traversal Algorithms**



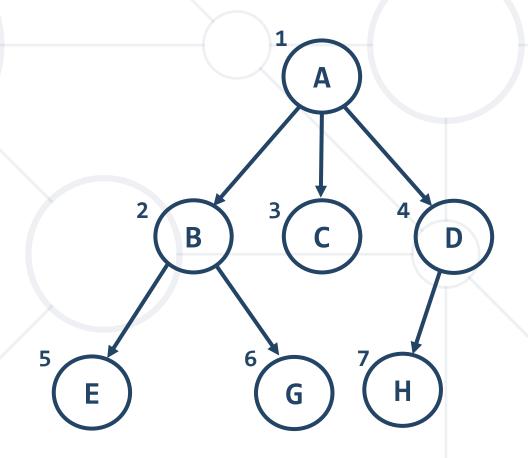
- Traversing a tree means to visit each of its nodes exactly once
  - The order of visiting nodes may vary on the traversal algorithm
  - Breadth-First Search (BFS)
    - Nearest nodes visited first
    - Implemented with a queue
  - Depth-First Search (DFS)
    - Visit node's successors first
    - Usually implemented by recursion

#### **Breadth-First Search (BFS)**

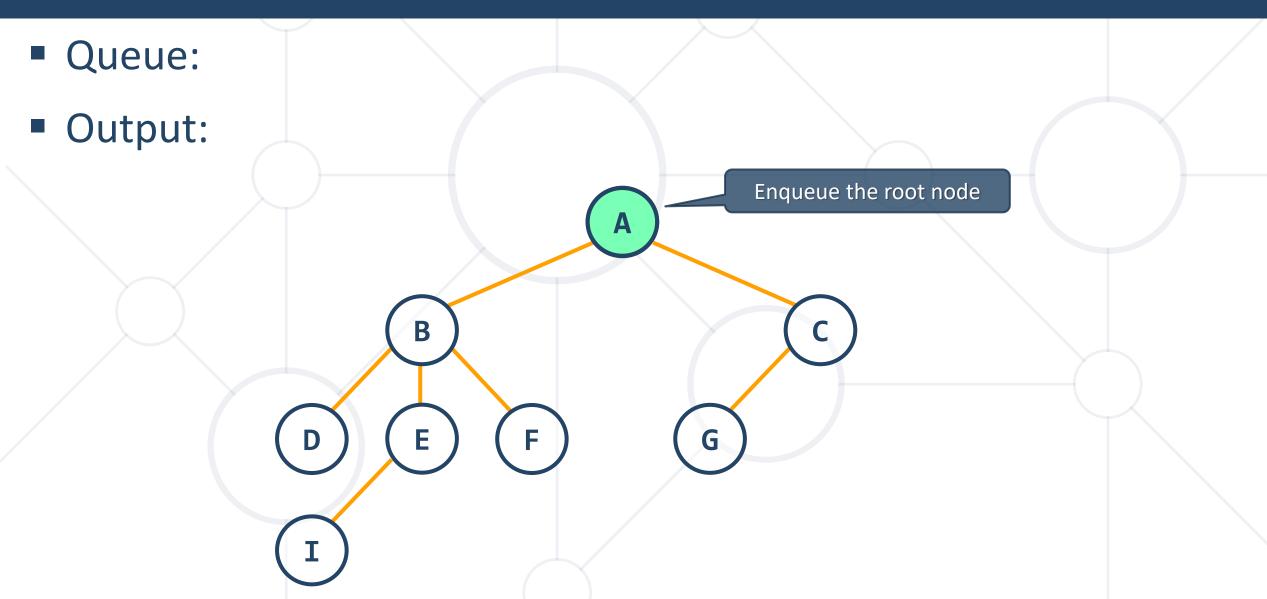


- Breadth-First Search (BFS) first traverses the tree level by level
- BFS algorithm pseudo code:

```
BFS (node) {
 queue ← node
 while queue not empty
   v ← queue
    print v
    for each child c of v
     queue ← c
```





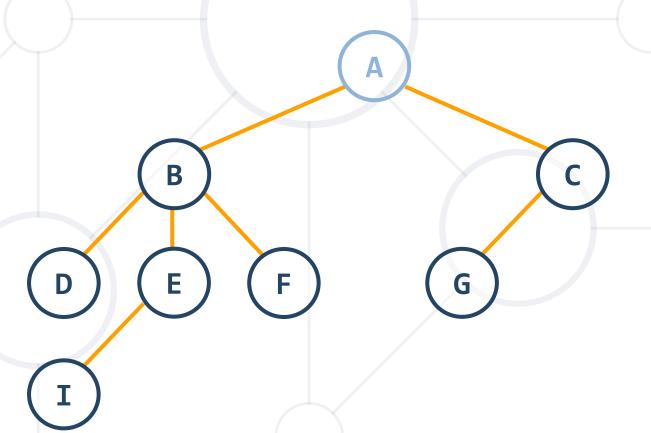




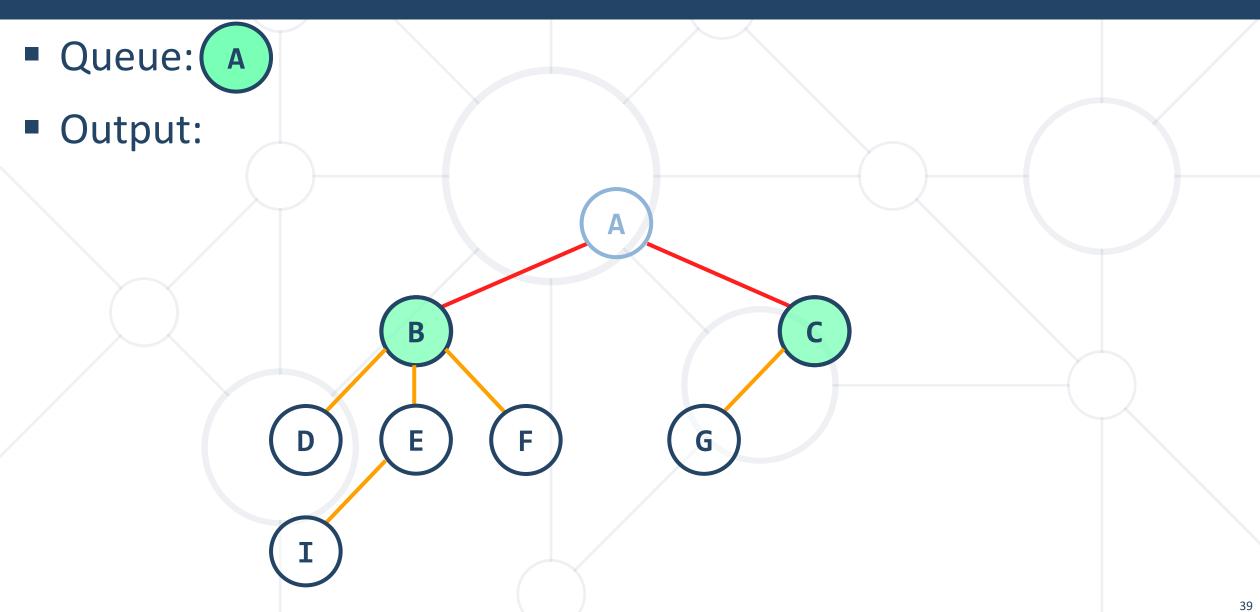
Queue:

Dequeue and print the current node and add its children to the queue

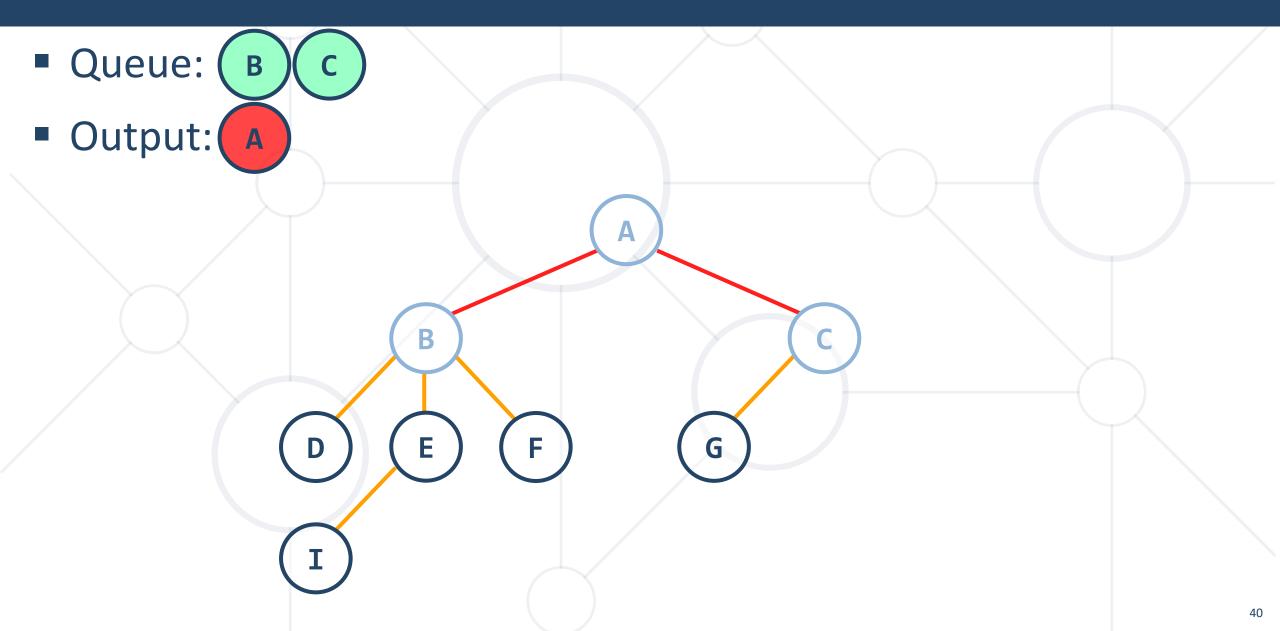
Output:



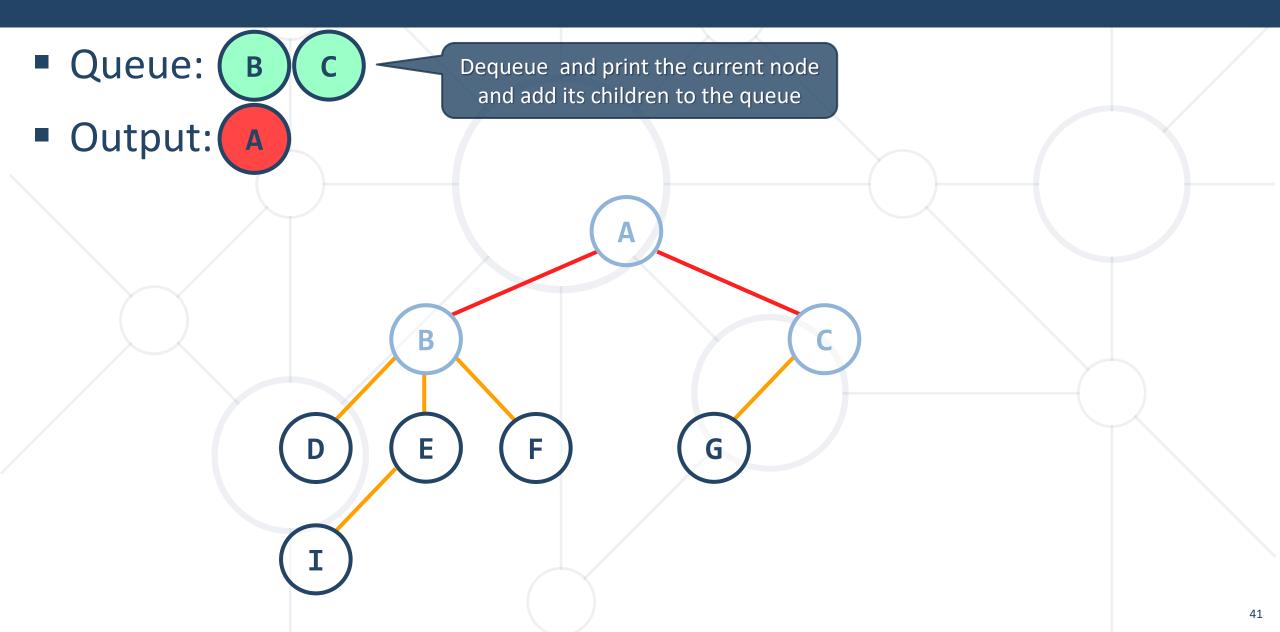




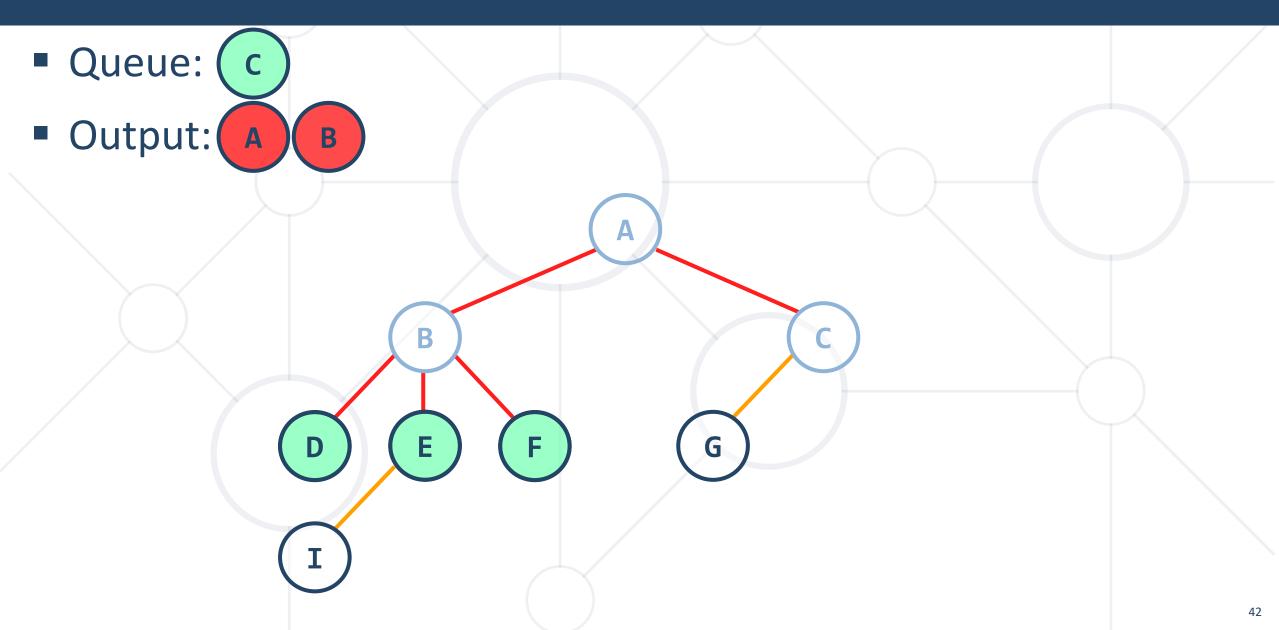




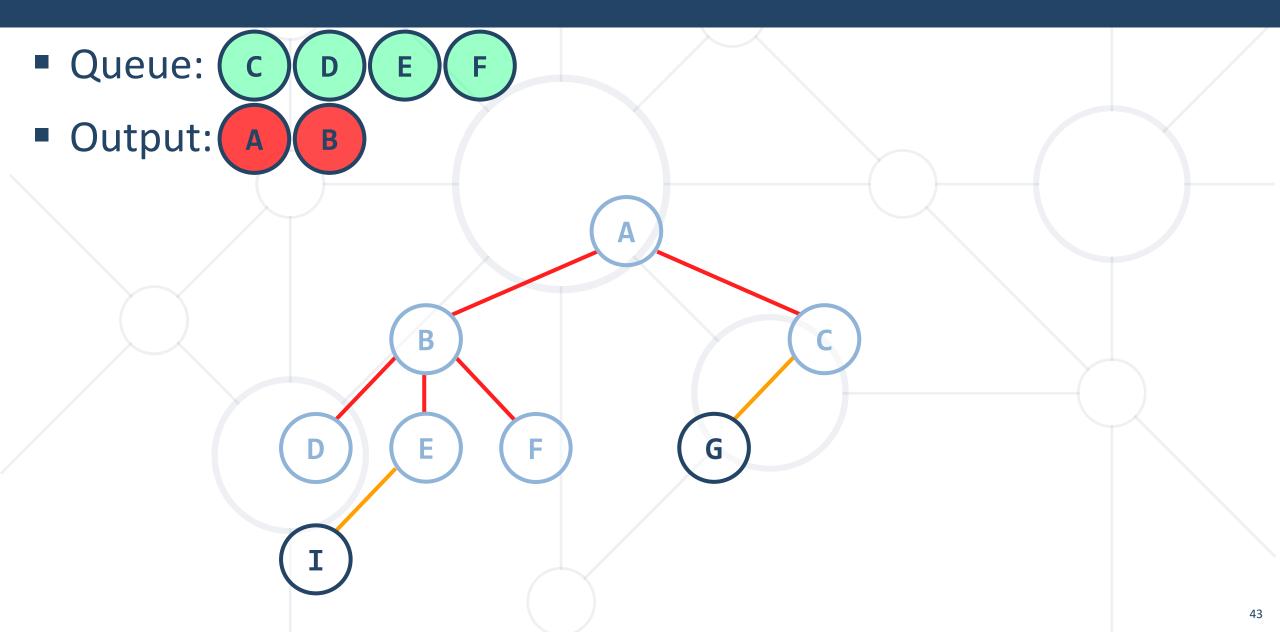




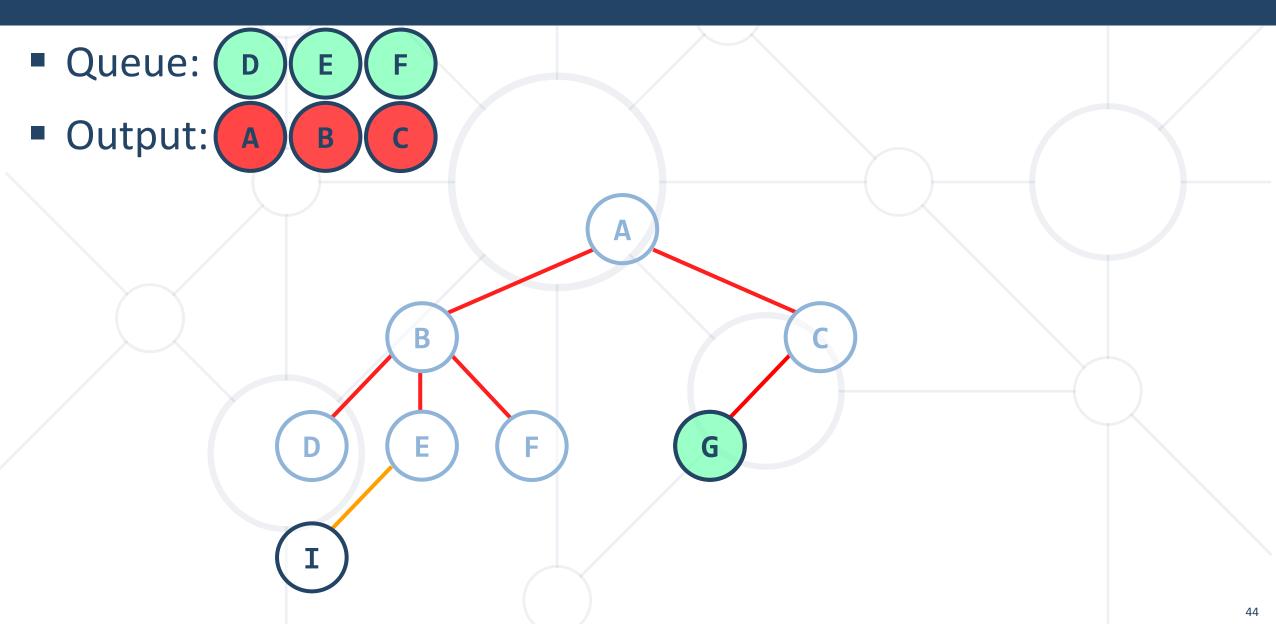




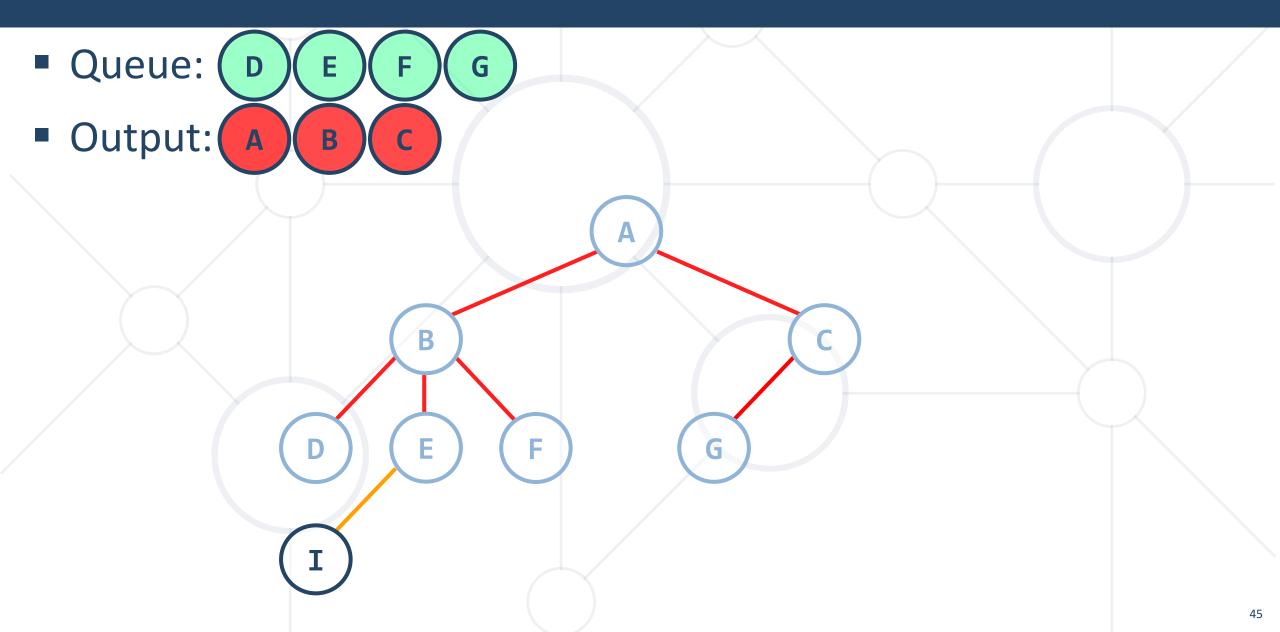




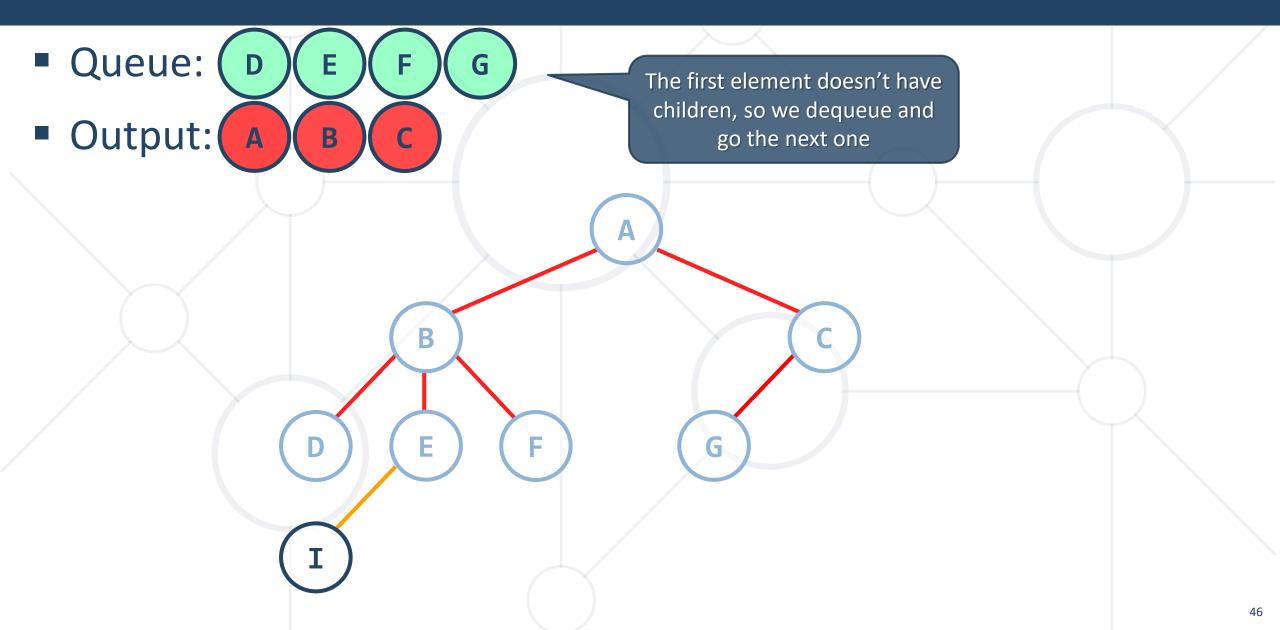




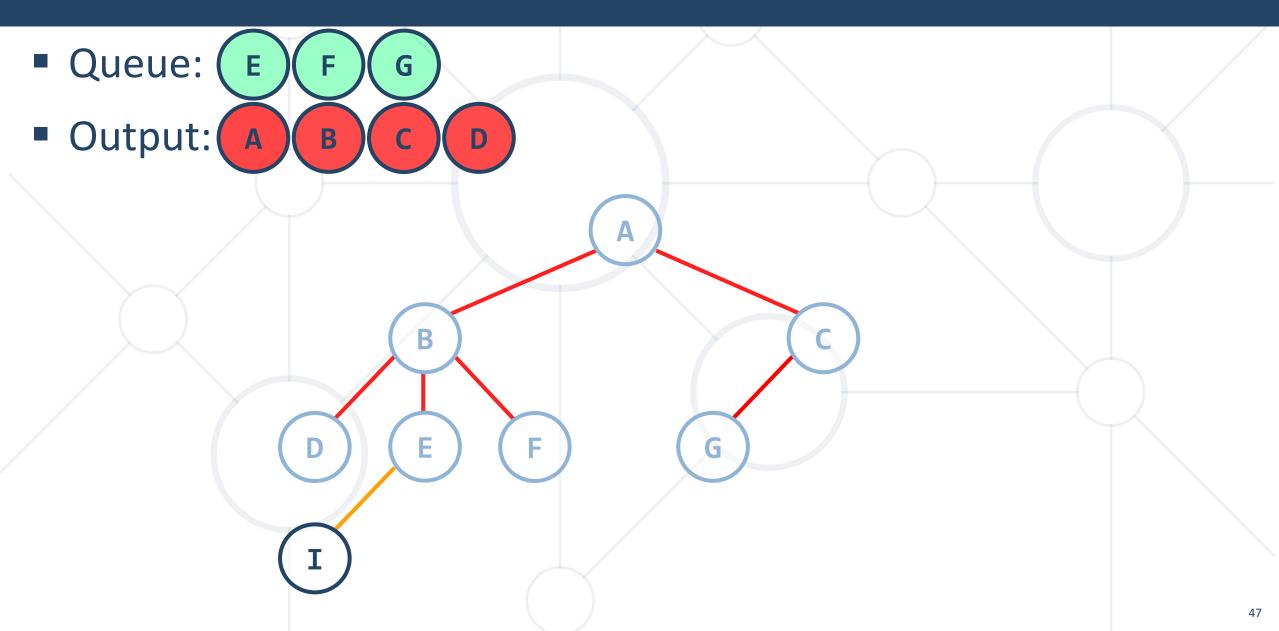




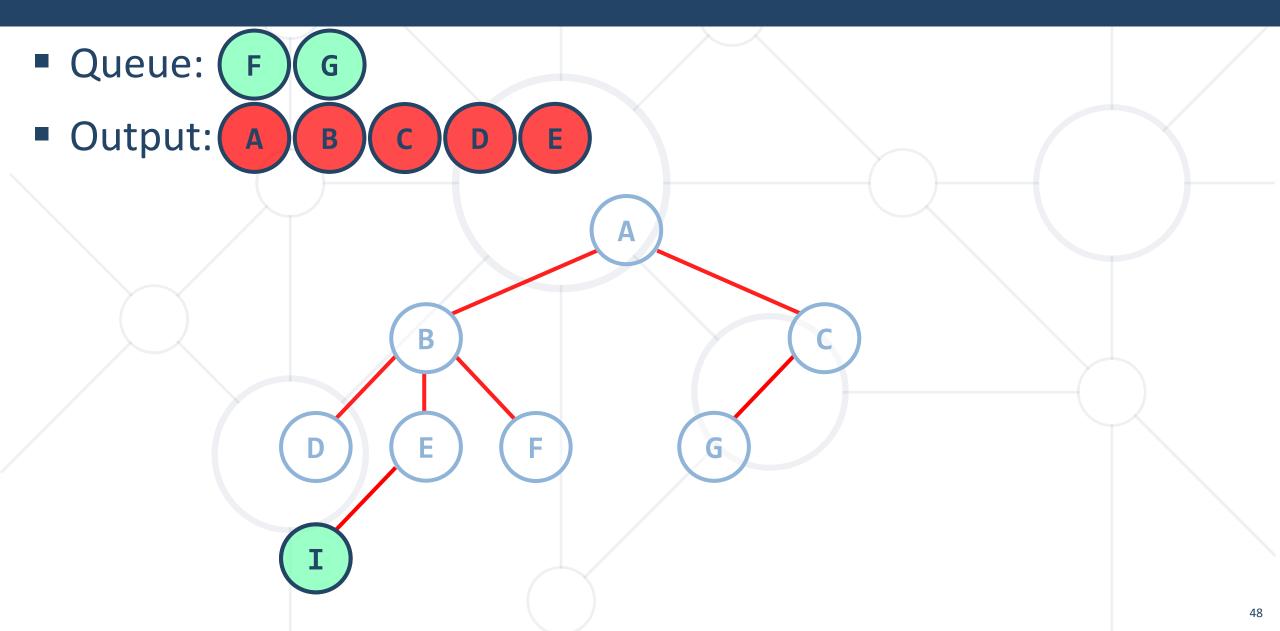




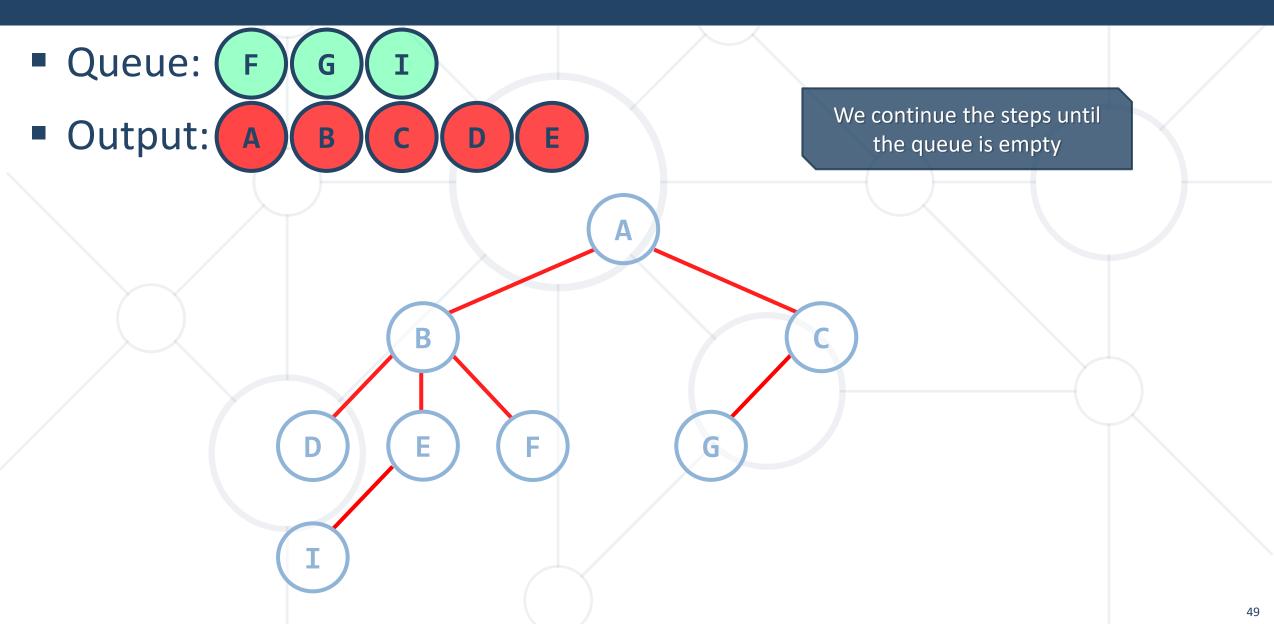




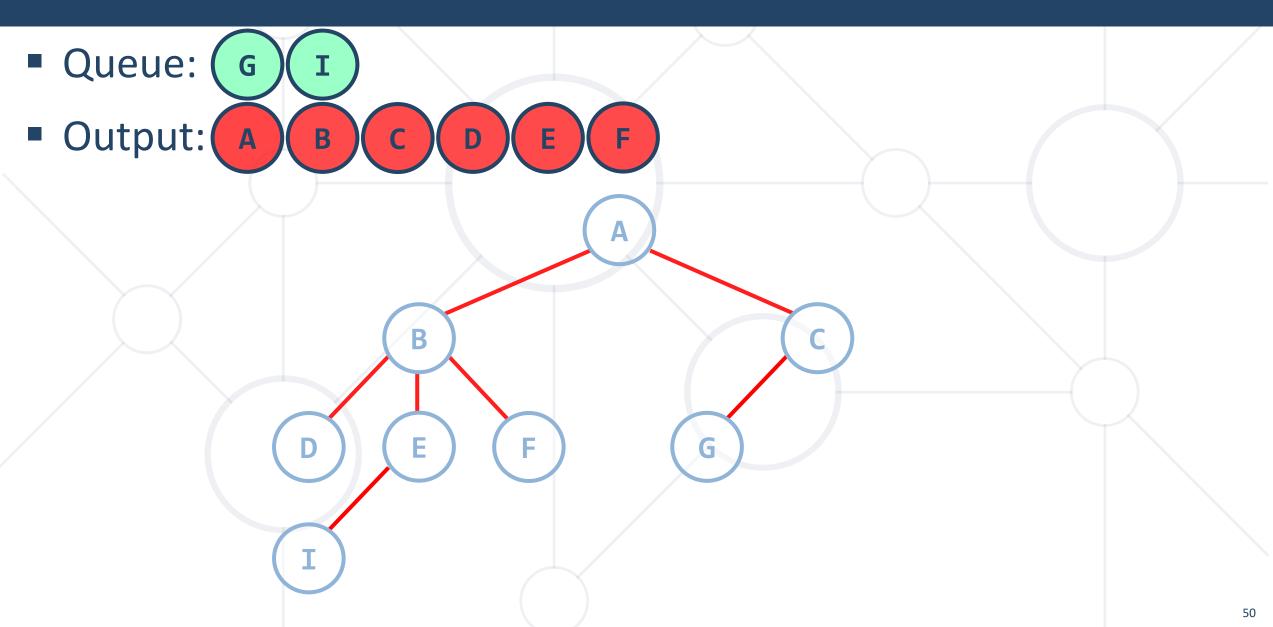




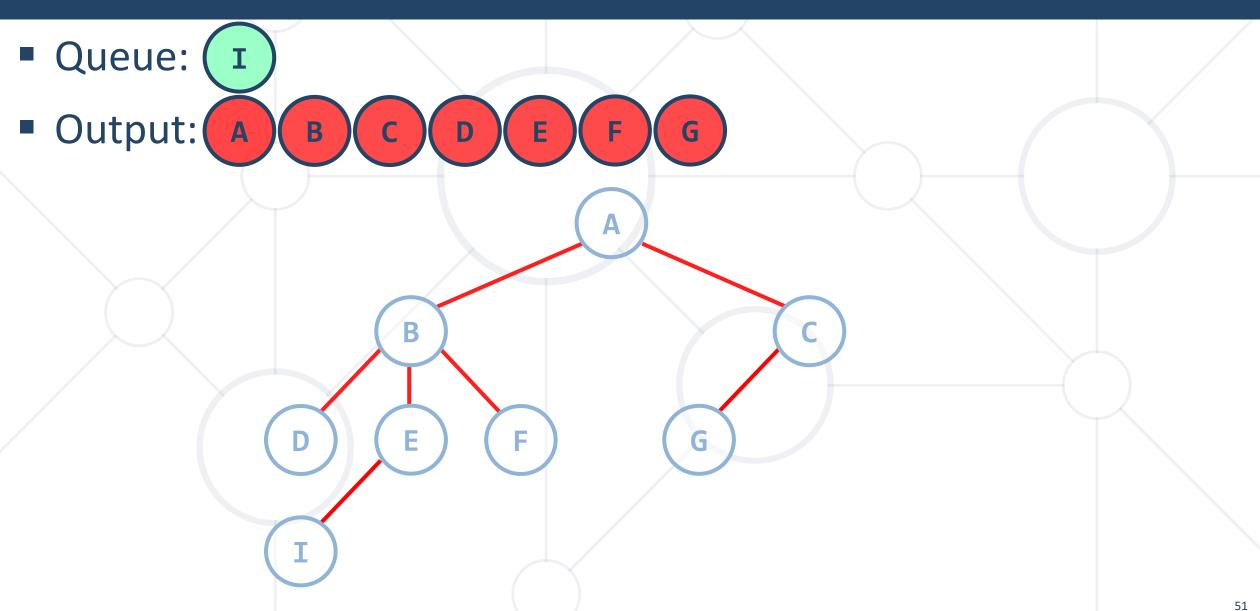




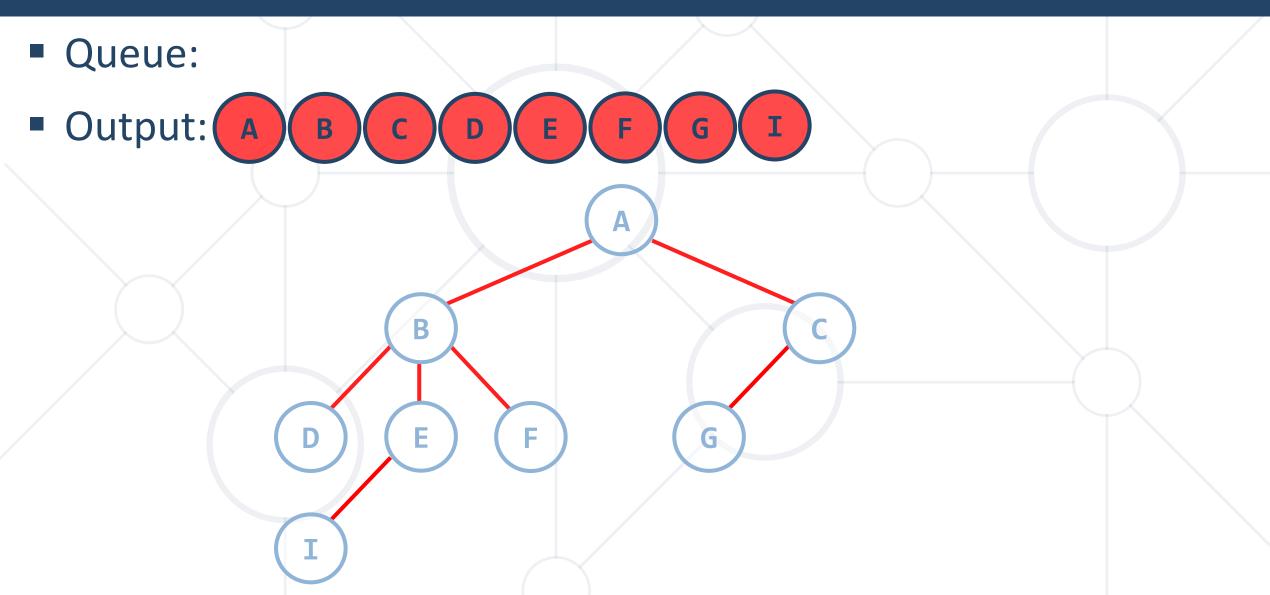








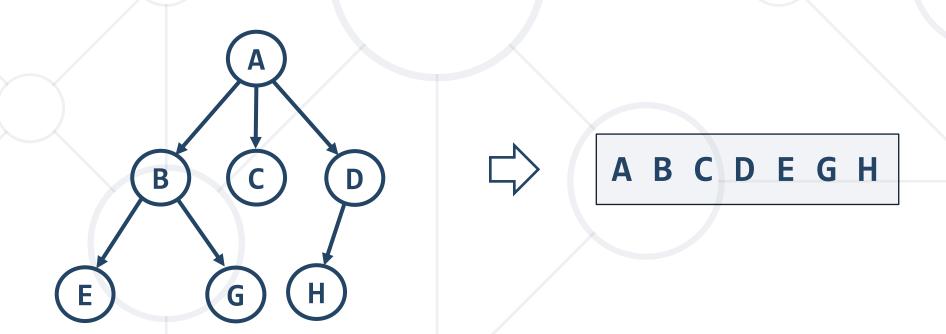




### **Problem: Order BFS**



- Given the Tree<T> structure, define a method
  - IEnumerable<T> OrderBfs()



#### **Solution: Order BFS**



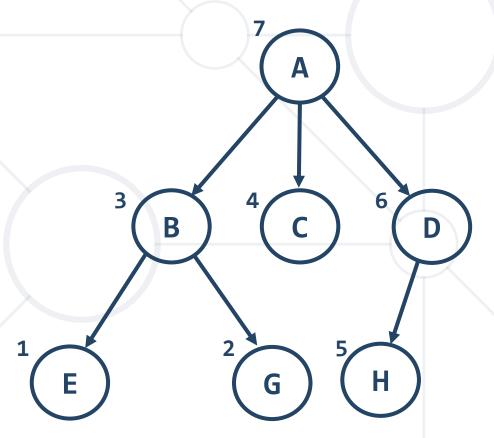
```
public IEnumerable<T> OrderBfs()
 var result = new List<T>();
 var queue = new Queue<Tree<T>>();
 queue.Enqueue(this);
 while (queue.Count > 0)
   // To Do: Implement this part
  return result;
```

# Depth-First Search (DFS)

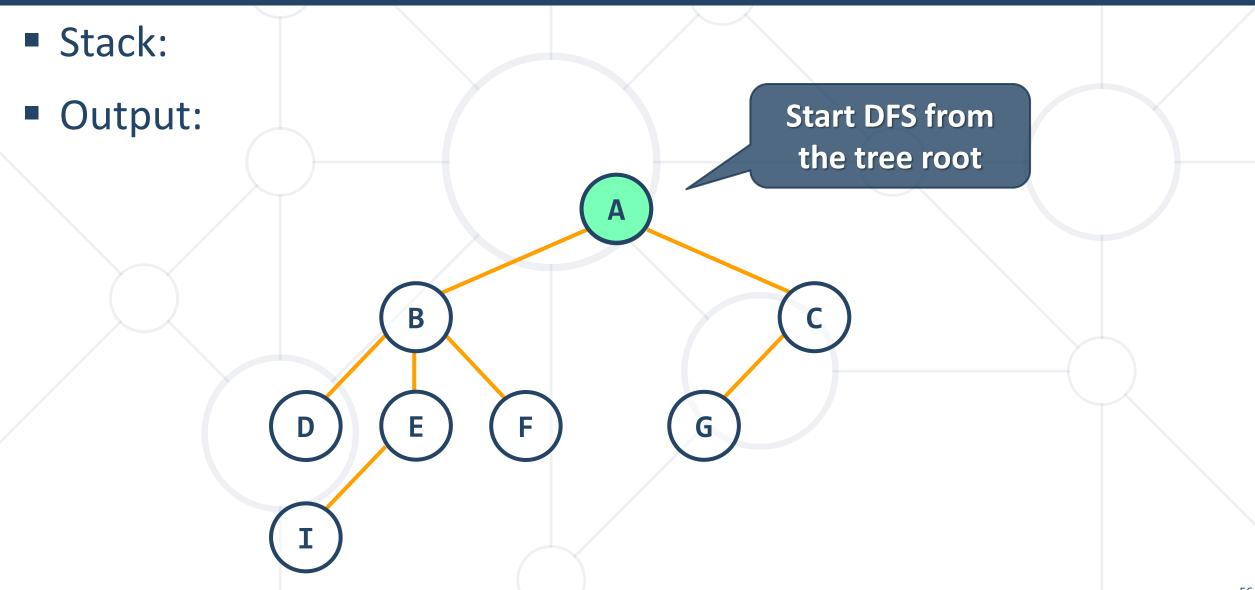


- Depth-First Search (DFS) first visits all descendants of given node recursively, finally visits the node itself or vice-versa
- DFS algorithm pseudo code:

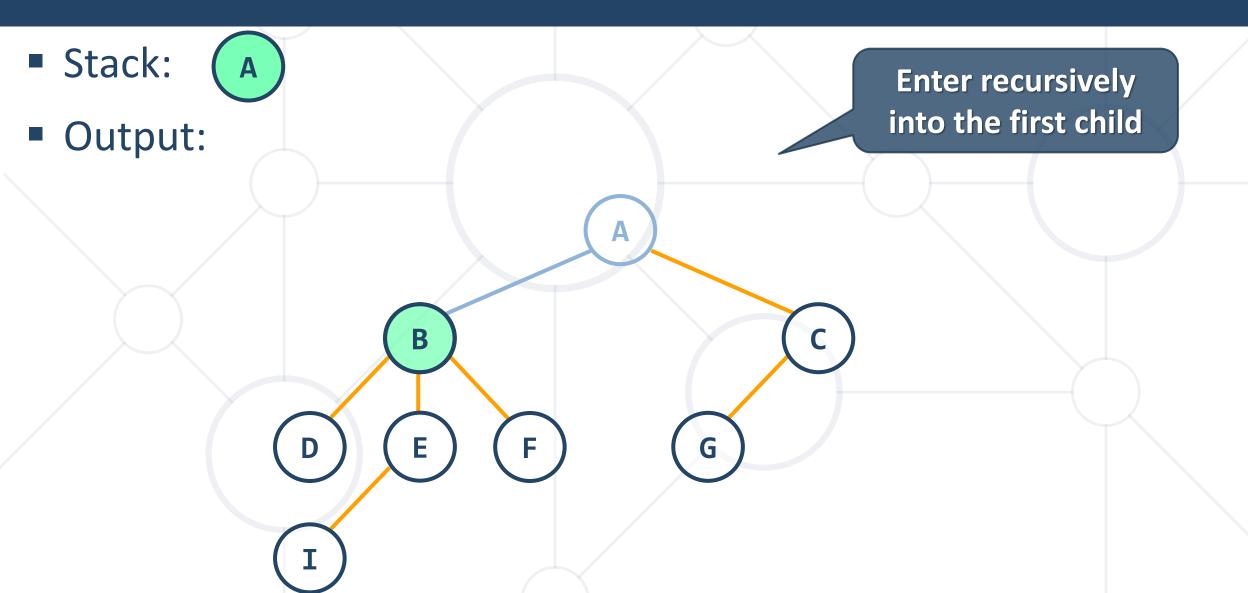
```
DFS (node) {
  for each child c of node
    DFS(c);
  print node;
}
```



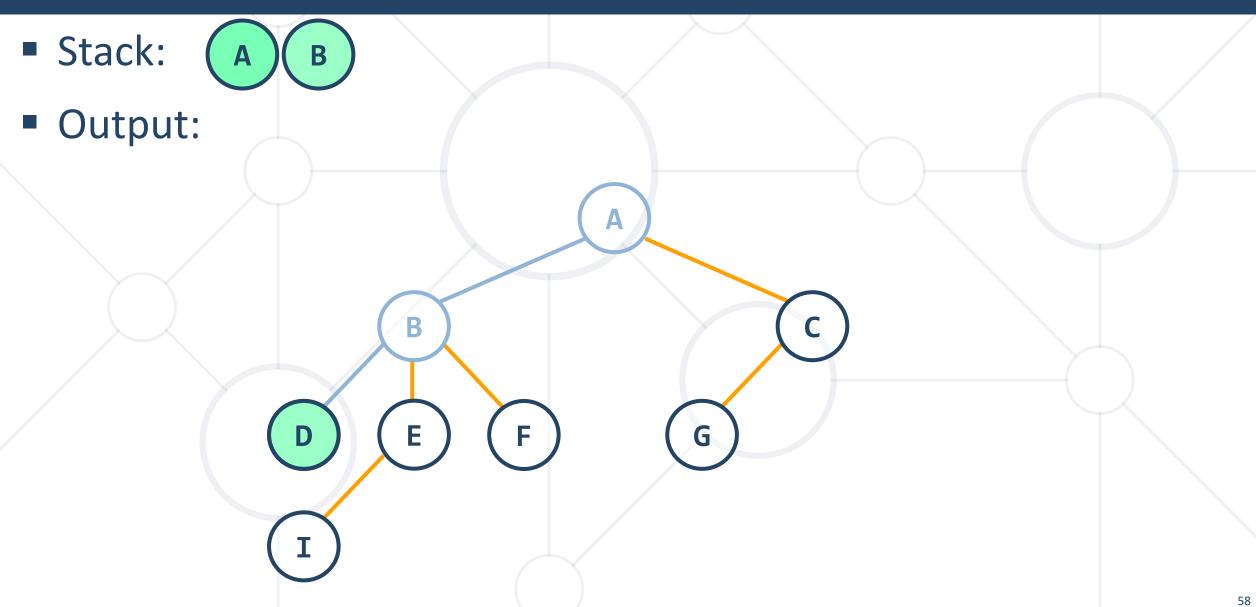




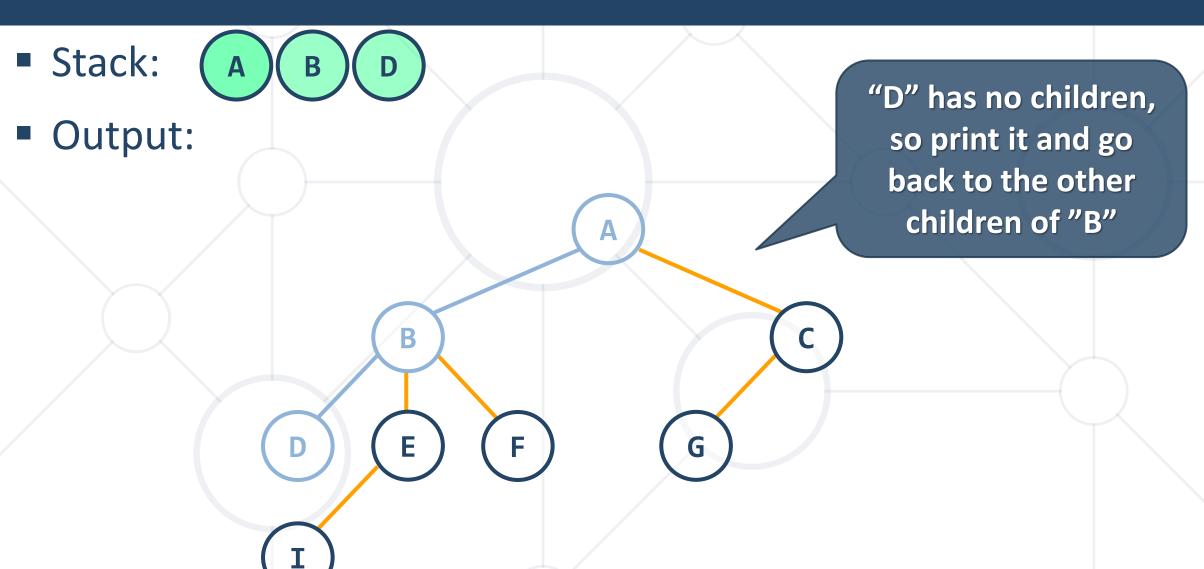




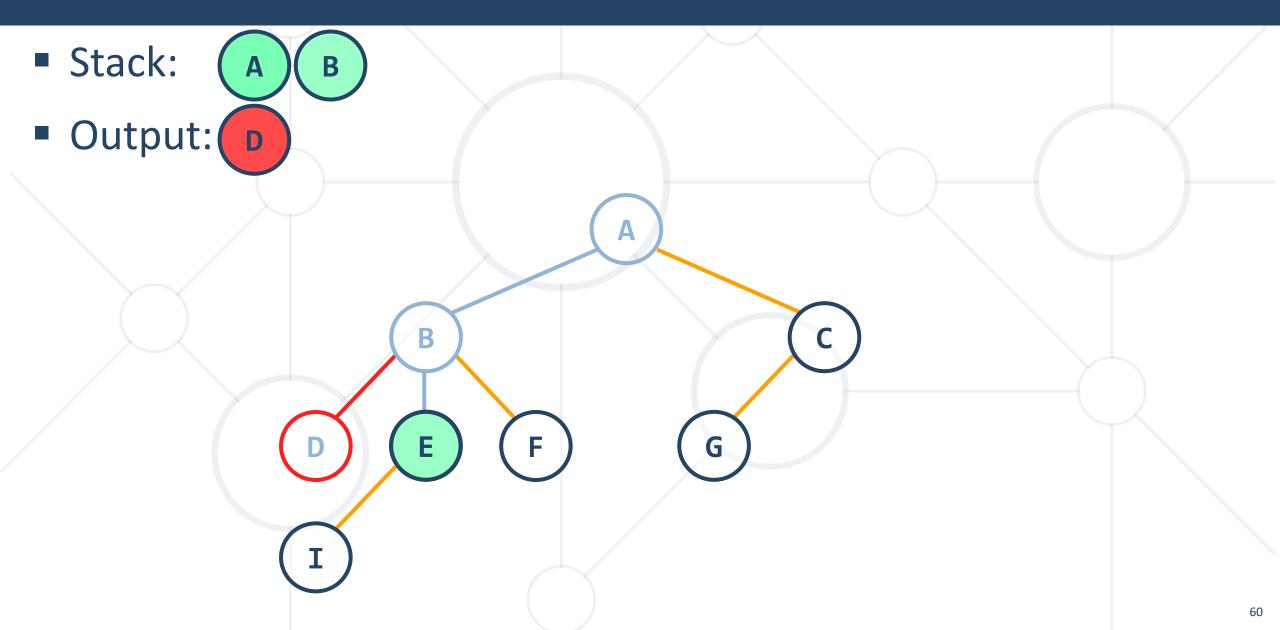




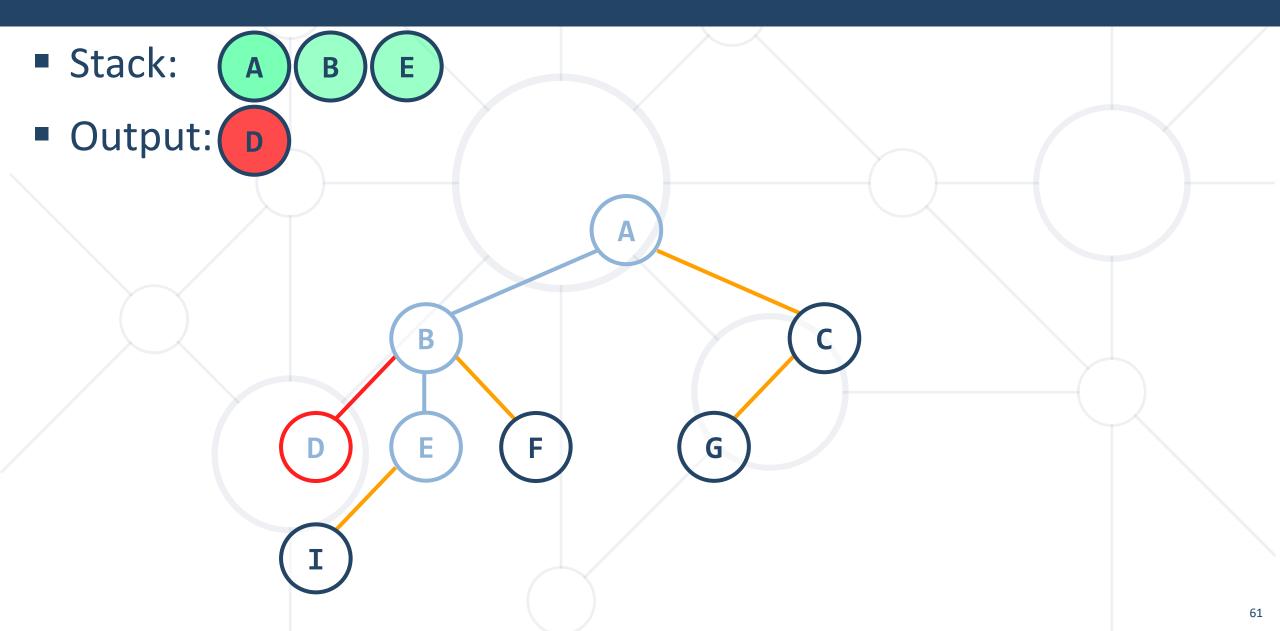




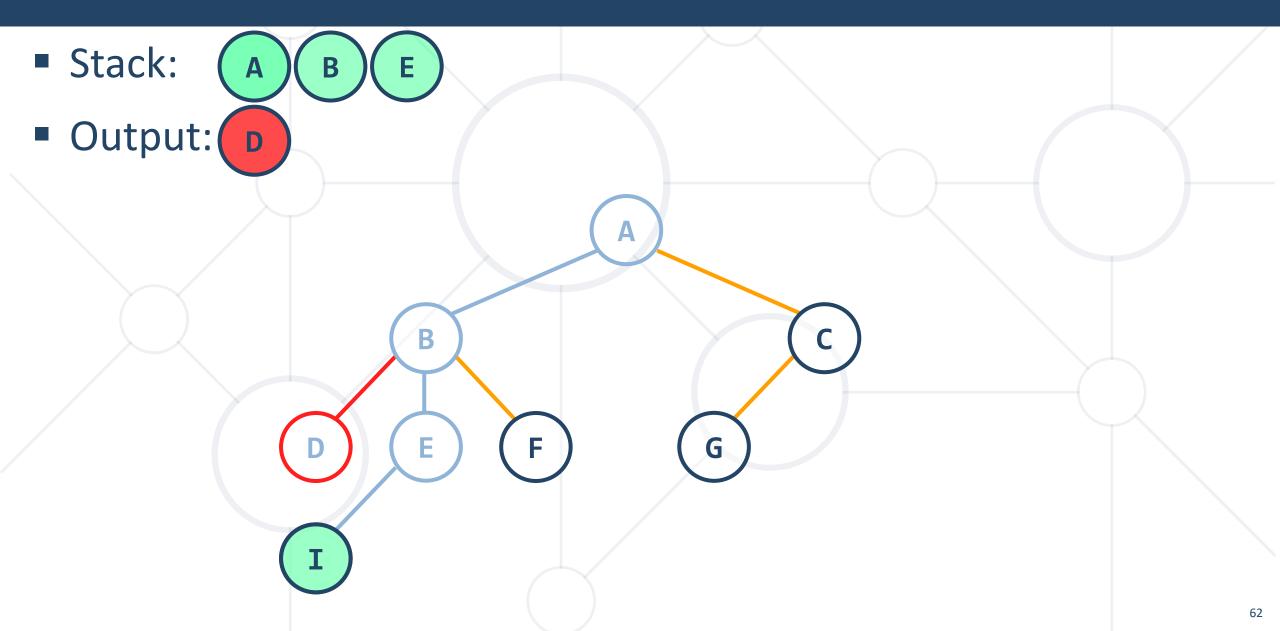




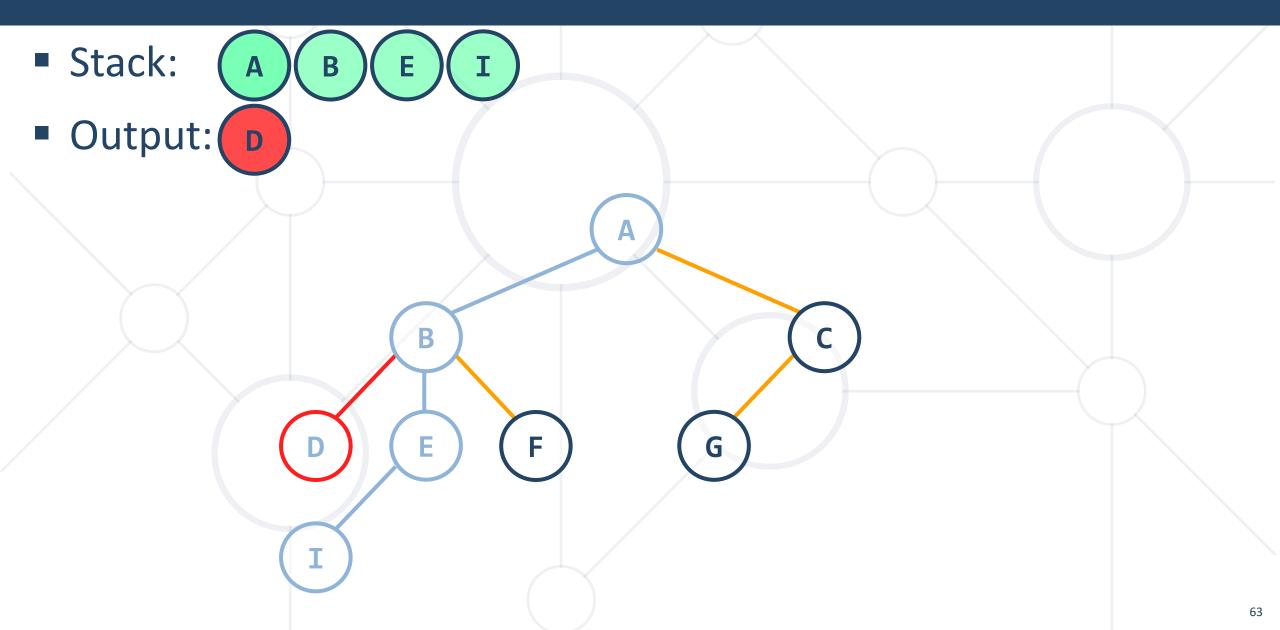




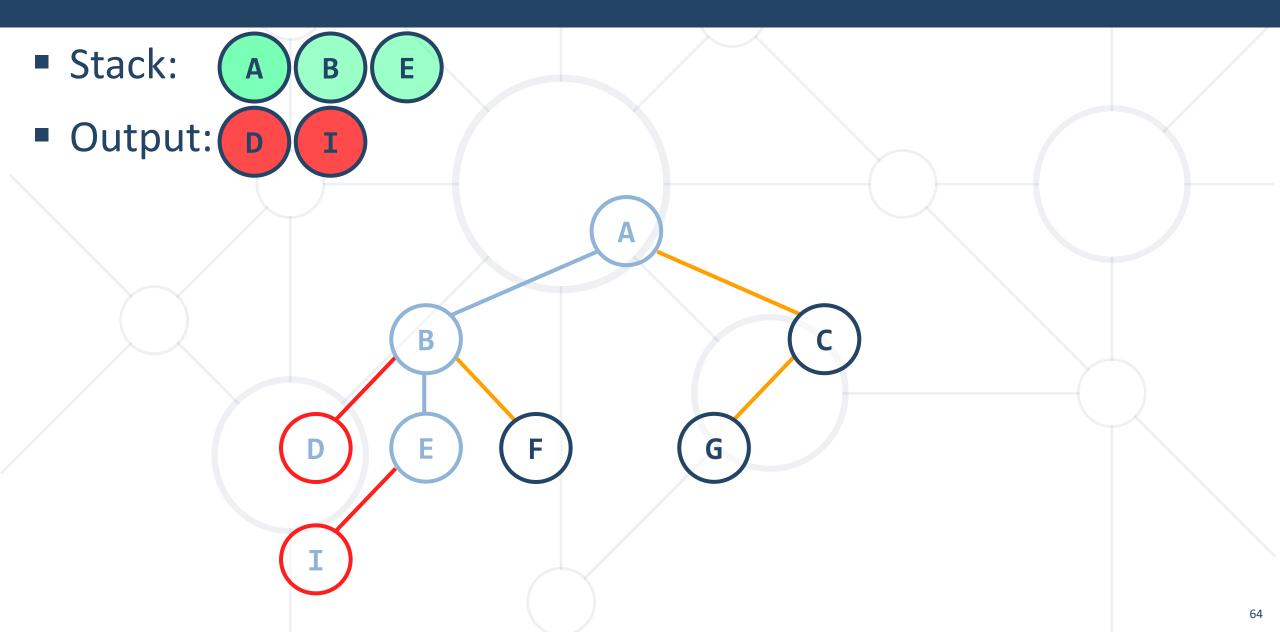




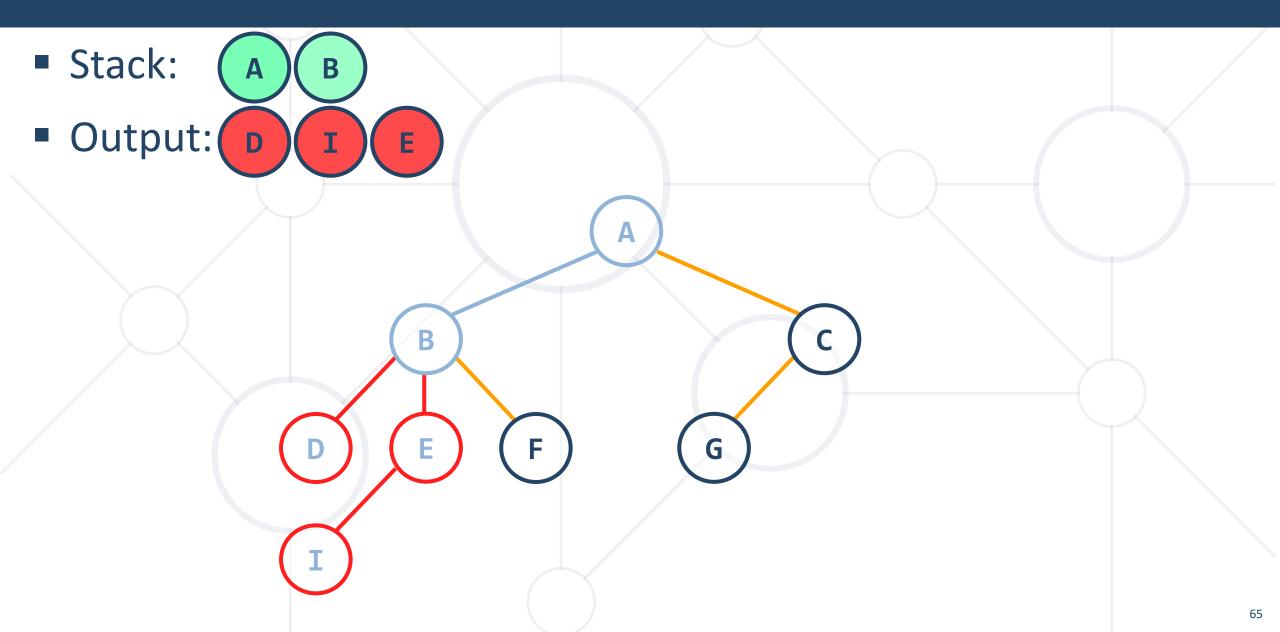




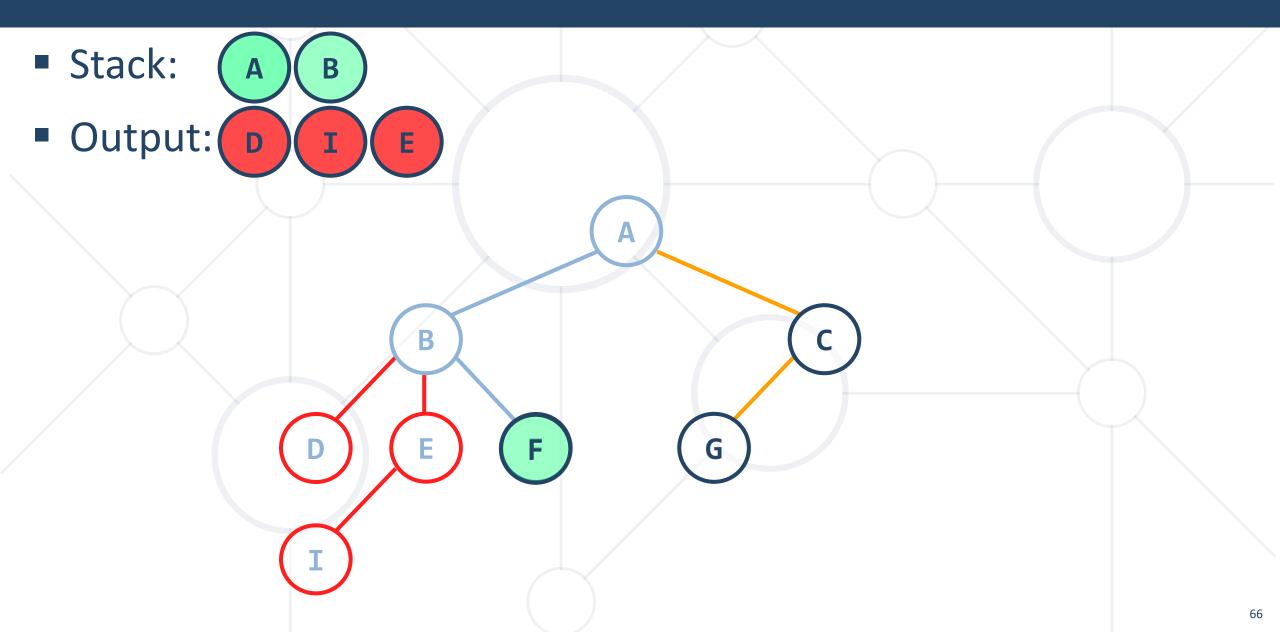




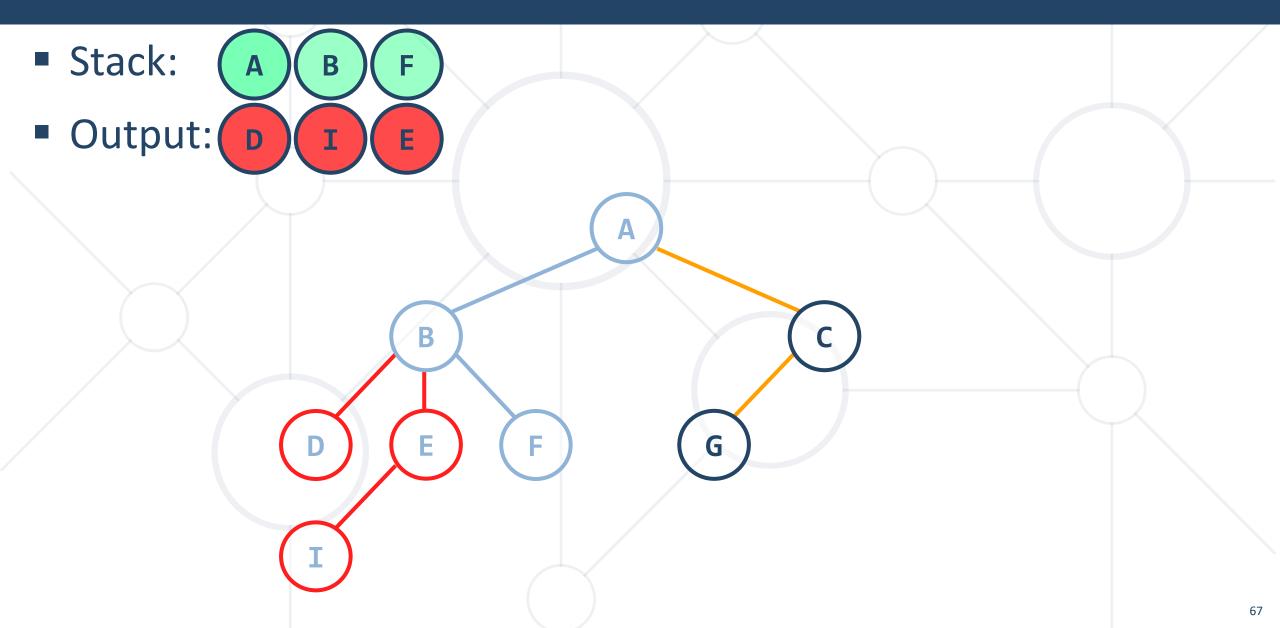




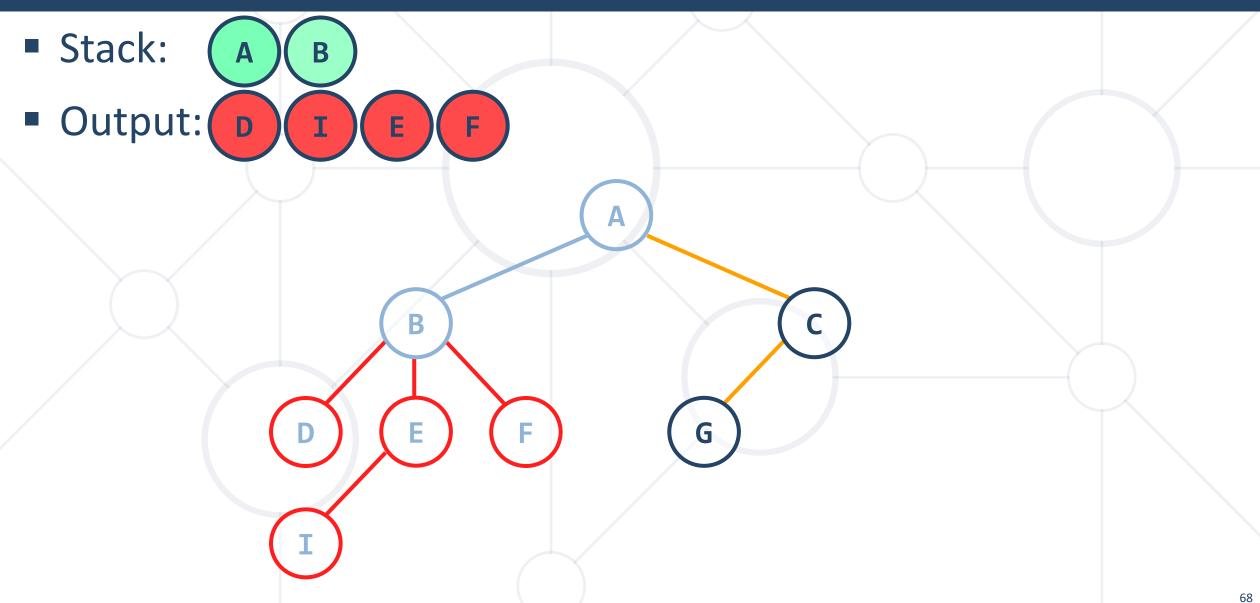




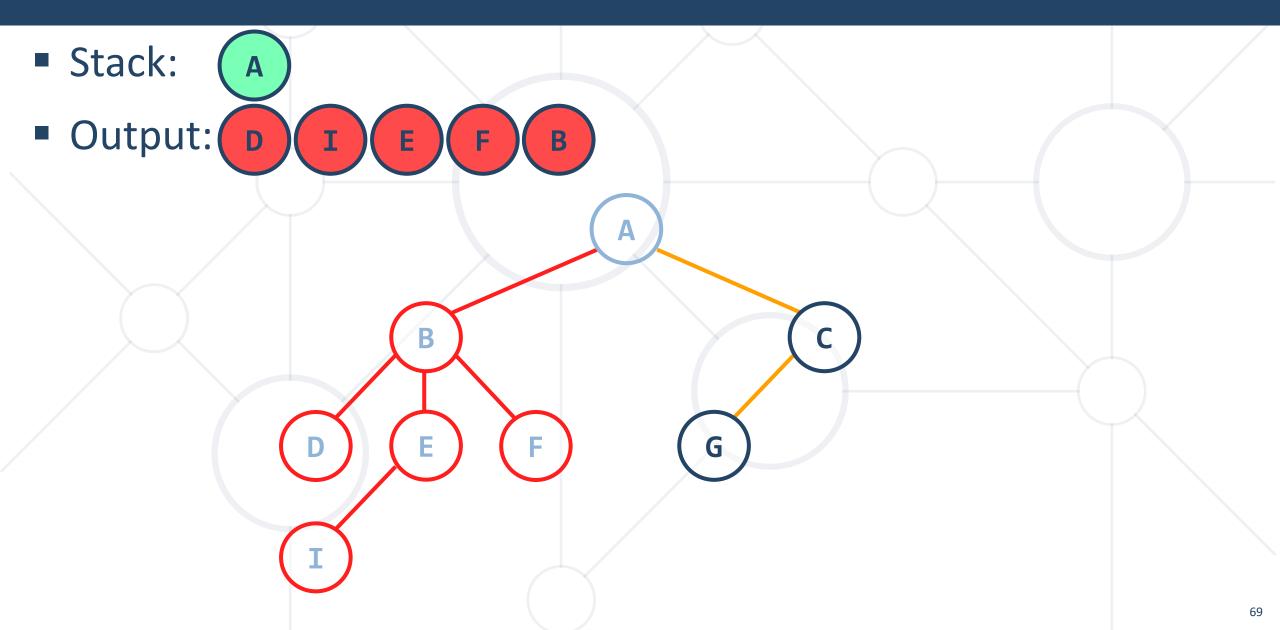




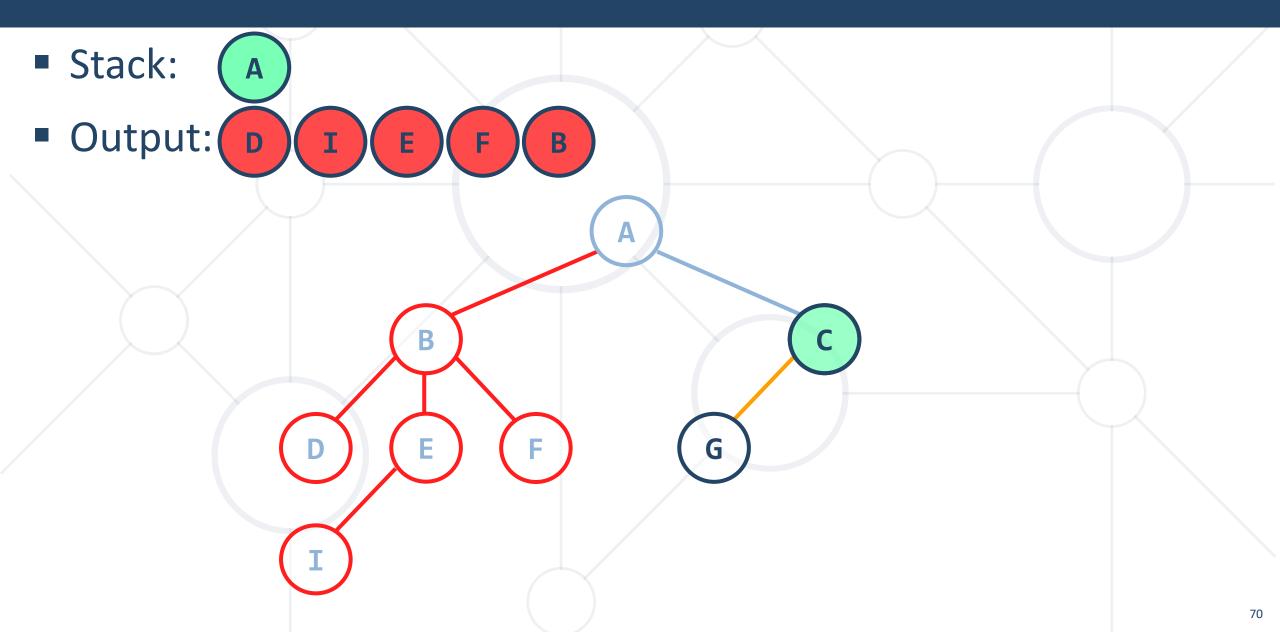




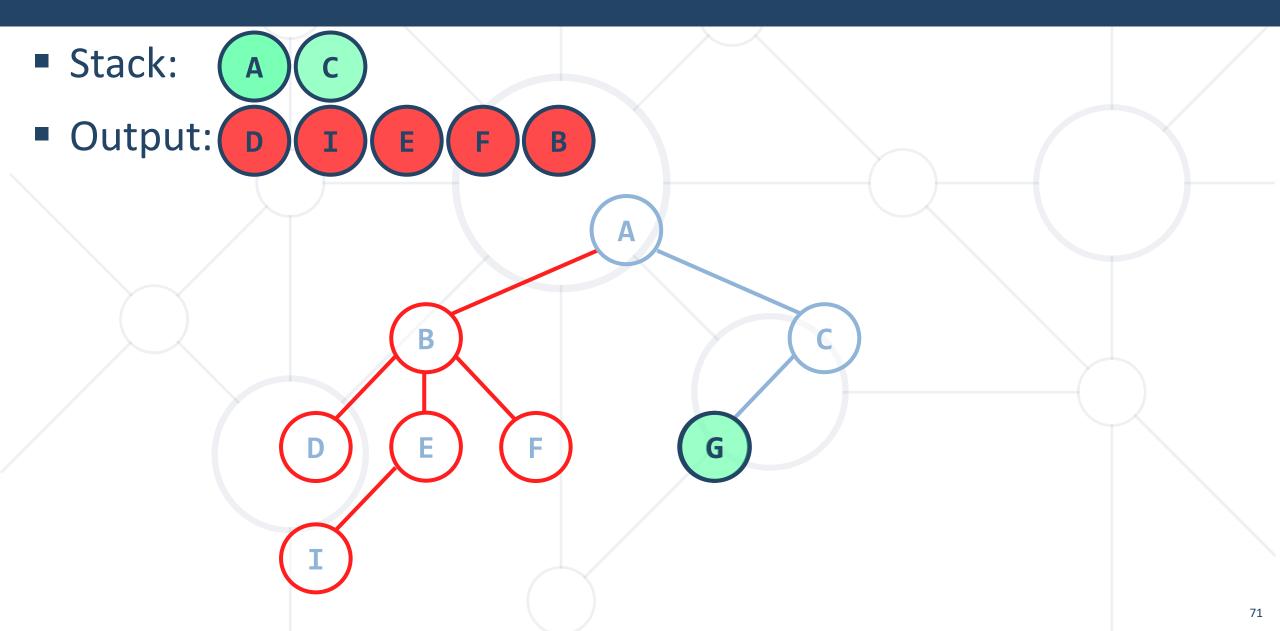




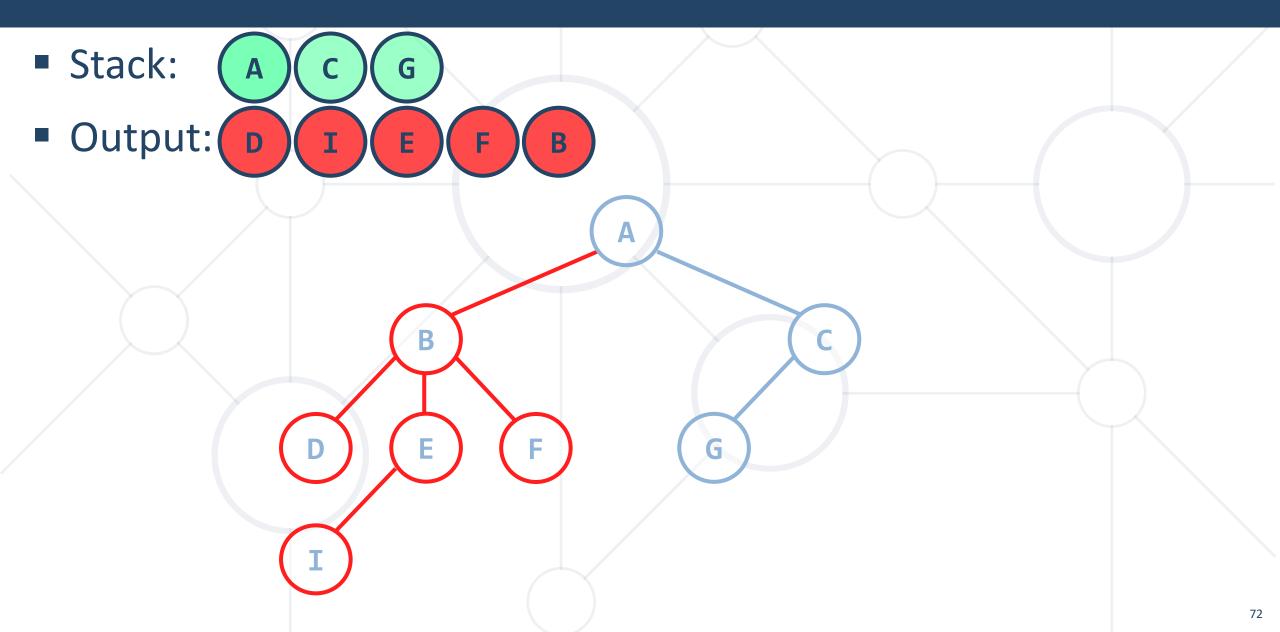






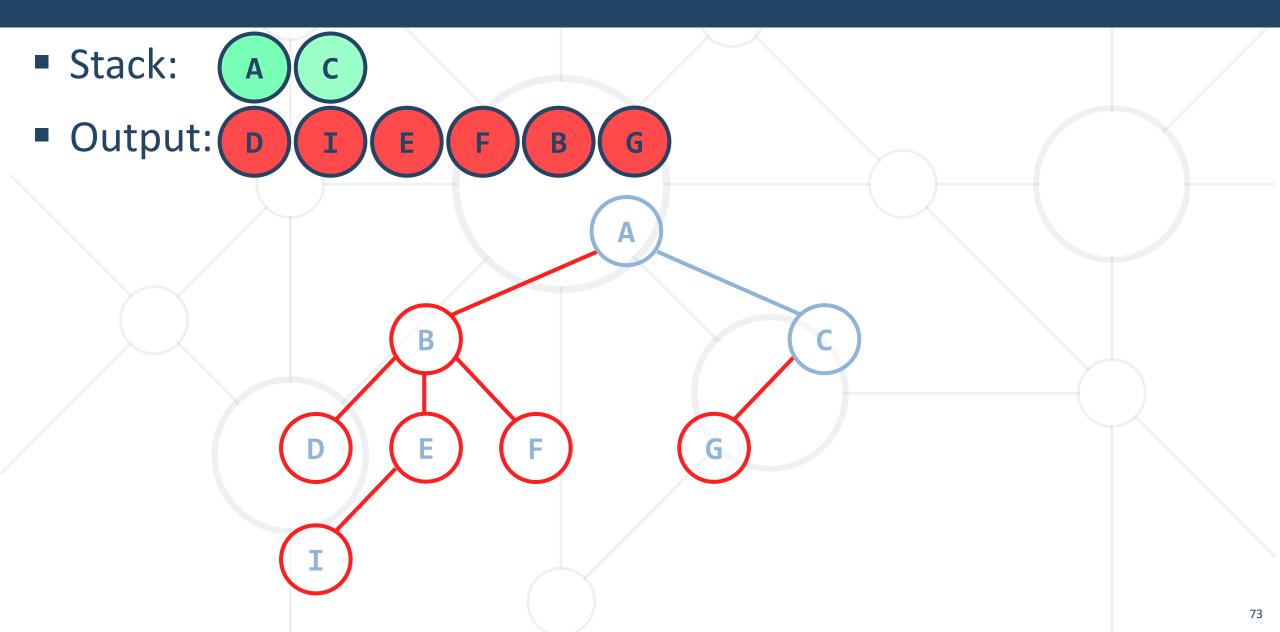






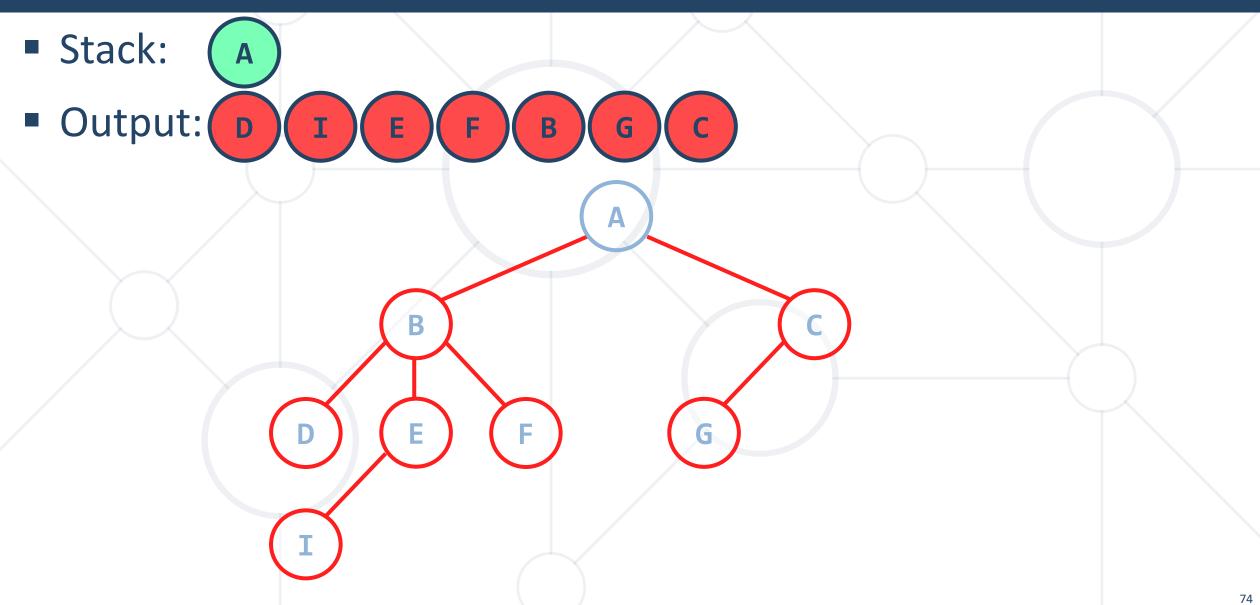
## **DFS Visualization**





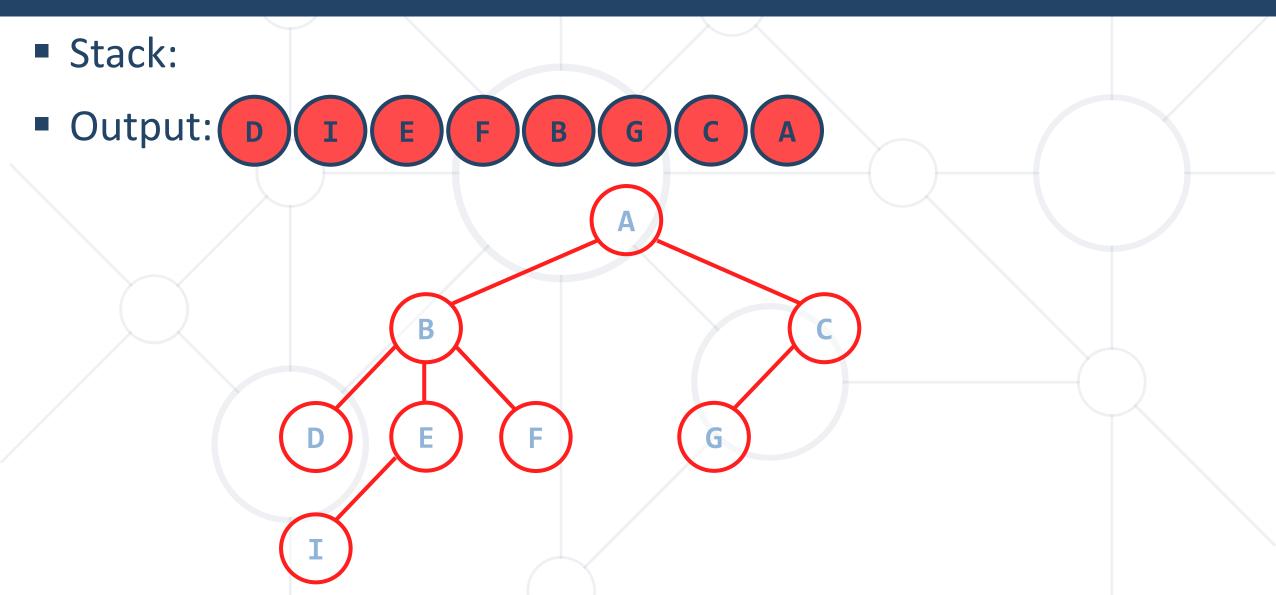
## **DFS Visualization**





## **DFS Visualization**

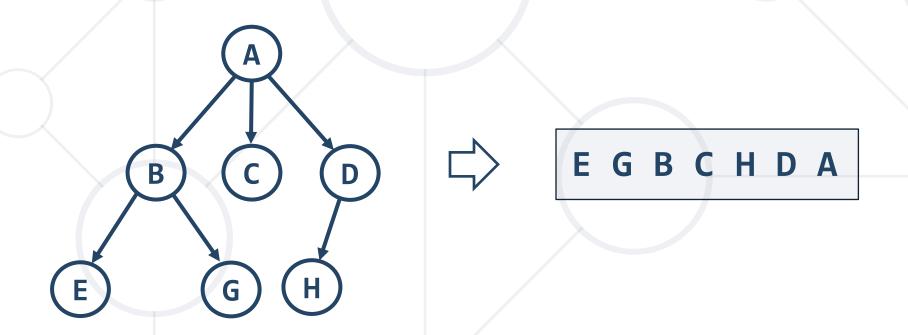




## **Problem: Order DFS**



- Given the Tree<T> structure, define a method
  - IEnumerable<T> OrderDfs()



#### **Solution: Order DFS with Stack**



```
public IEnumerable<T> OrderDfs()
 var result = new Stack<T>();
 var stack = new Stack<Tree<T>>();
  stack.Push(this);
 while (stack.Count > 0)
    // To Do: Implement this part
  return result;
```

#### **Solution: Recursive Order DFS**



```
public IEnumerable<T> OrderDfs()
 var order = new List<T>();
 this.Dfs(this, order);
  return order;
private void Dfs(Tree<T> tree, List<T> order)
 // To Do: Implement
```

#### Conclusion



What did we get so far?



- Are we working with O(log(n))?
- Well, the answer is...
  - No!
  - Why? Still, we are stuck at linear complexity for searching operations
- We will try to solve that with a Binary Search Tree



## Summary



- Trees are recursive data structures
  - A tree is a node holding a set of children (which are also nodes)
  - Edges connect Nodes
- DFS → children first
- BFS → root first





# Questions?

















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