(Viet Discrete : Good Stuff. D Recall . a s b moder if n la-b. STAF · congruence mod is an equivalence relation on Z. • arith rules: $a \equiv b \mod n$ => $\begin{cases} a \pm a \equiv b \pm b \mod n \\ aa' \equiv bb' \mod n \end{cases}$ In part, a = 6 mad = > ak = 6 mad 1 ha = klo mad 1. Example Lit dide ... dy = N . dy de digita mod 2 N 10"d, + 10"dz + + 10 dn -, + du = d, nod 2.

omod 2 mod 2 / /same for mod 5. mod 3 . N = 10"d, +10"-2 dz + ... + lody + du = d, + dz + ... + dn -1 + du mod 3. (10 = 1 mod 5) => (10 x = 1 mod 5) 2 * model : N = 10 dy + 10 2 dz tout Loday + du (10 = -1 mod 11) => (0 = (-1) k mod n); t content of mod n)

N = ...+ d_{n-2} -d_{n-1} + d_n mod 11 Ex: 100 mod ?? reductions address have order exponent. (60se) too = 2 mod 7 . => 1000 = 2 mod 7 2 (2) = 2 (0) = 1024 mod 7 = 2 mod 7 flows = flow (15 = 01(015) weathyworks of prime mad (exponent) recall Formal's little theorem. a' = a mod > Cate a modit das Z april = 1 modip (00° = (00°)+ + ged (100,7)=1 OK. => 100 (1)+4 = 104 mod 7 = 24 mod 7 = 2 mod 7. Z/nZ Again congruence mod u , n & Zt , is an equivalence relation on Z => equivalence classes. [1] = {1, n+1, 2n+1, -n+1, -2n+1...} when representatives are the

From (3) a | e => c = ac' ab | aq bq2 | = ax + by
ble => ble = ac' ab | c = aq = bq2 ged(a, o)=1 (6/c) = ablac'=c 1 0=aq - lage & because by a A blac' so more common parts between Theorem: ged (o, b) . I em (a, b) = a · b . my: d= ged(e,b). d|a=> a=da' } , ged(a',b')=1 dla= b=db' J[[k=ged(a'b') >1 . Hun, K|a' = 5 Kd|a'd = a Now, show law (a, b) = da' b' (then god e fem = d. da'b' = ab) . k|b' => kd |b'd = 5 da'b' is a common multiple of a and b, so hid greater than good of 1 (/da'b' => /6 | 6a') at milite of a ord, then do o m. relatively prime. Bod's prog . alm] dlm => m = m'd a = da' | m = dm' => a'lm' .) => a'b' | m' Same for la : => lo! lou! da'b' | m'd => da'b | m , 1 recall gives a, b, - prime fectorisations, "joint" 0 = 10, 12 1 2, ... xx > 0 => gcd(a,b) = mi(a, A,) p. mi(k, A,) . mi(a, A,) len (a,b) = p mar(x, s,) | mar(az 182) | max(ax, px) ged (a, b) · lem (a, b) = To milax, bx) . To medax. bx) = To px x + Bx

| Discrete , Cost'. Congruences "The same w/o being equal". Almost equality ? equi relation.) fa 1>0 integer. a = 6 mod n -> a and 6 have the same remainder upon trustin by n i.e. a a -b -> makes sense. Equivalence relation: 1 Reflexive: a = a mod n. a has some remainder os a. (Try to prose this). @ Sym: a = b mod n () b za nod n. Than; a = b wodn, b = c mod n => a = c mod n - Congruence mad in is an equivalence relation on I. Rules: (arith), a= bmoden A a' = b' maden => a ± a' = b ± b' moden => aa' = bb' mad n