# S2021 COMS311 Recitation Week 2 Sections 1 4

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Please refer to, Recitation Notes by Weisi, for the first portion of the recitation notes. I will almost always try to use the common recitation materials first and then may augment them with more material.

This recitation is about Big-O and time complexity. Make sure you know it and can use it both for analyzing an algorithm given to you and in making your own.

### 1 Problem1

#### 1.1 a

$$f(n) = n$$
. Is  $f(n) = O(\log(n))$ ? Why or why not?

#### 1.2

$$f(n) = n$$
. Is  $f(n) = O(n^2)$ ? Why or why not?

#### 1.3 b

$$f(n) = n$$
. Is  $f(n) = O(n^{999})$ ? Why or why not?

#### 1.4 d

Prove that  $x^2 + 10x + 1000$  is  $O(x^2)$ .

#### $1.5 \epsilon$

Disprove that  $x^3$  is  $O(x^2)$ .

## 2 Problem2

Rank the following functions by order of growth; that is, find an arrangement  $g_1, g_2, ..., g_{30}$  of the functions satisfying  $g_1 = O(g_2), g_2 = O(g_3), ..., g_{29} = O(g_{30})$ . Partition your list into equivalence classes such that functions f(n) and g(n) are in the same class if and only if  $f(n) = \Theta(g(n))$ . Scroll down to the next page for the answer.

$$2^{lgn}, \ 2^{2^n}, \ e^n, \ n^2, \ 1, \ 4^{lgn}, \ 9n, \ 2^n, \ 2^{2^n+1}, \ n!,$$

$$1 < 9n = 2^{\lg n} < n^2 = 4^{\lg n} < 2^n < e^n < n! < 2^{2^n} < 2^{2^n + 1}$$

Remember that  $lg(n) \neq log(n)$  and  $lg(n) = log_2(n)$  and thus  $2^{lg(n)} = n$