

# Customer Sharing in Economic Networks with Costs

**Bin Li<sup>1</sup> Dong Hao<sup>1</sup> Dengji Zhao<sup>2</sup> Tao Zhou<sup>1</sup>**

<sup>1</sup>University of Electronic Science and Technology of China, Chengdu, China

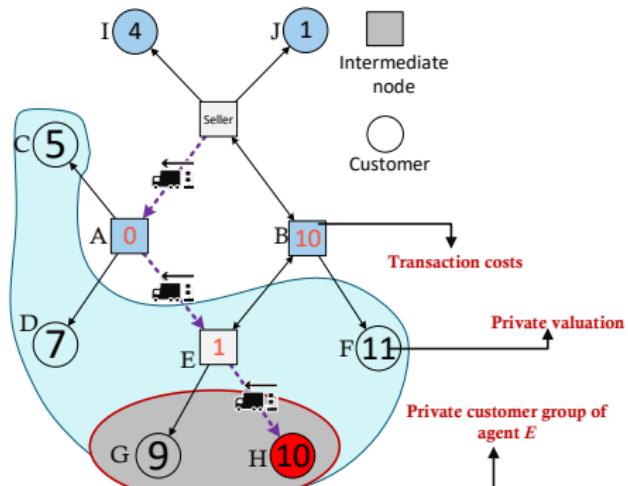
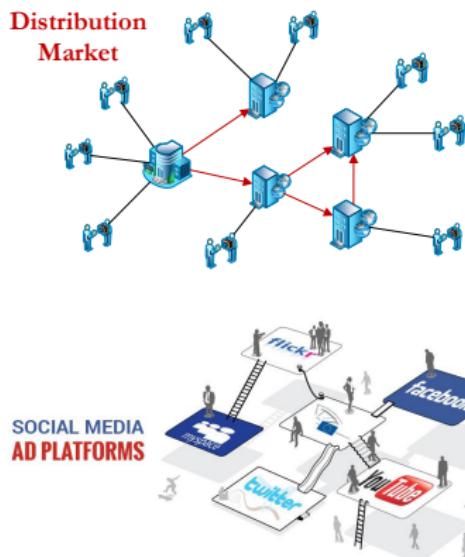
<sup>2</sup>ShanghaiTech University, Shanghai, China

July 18, 2018

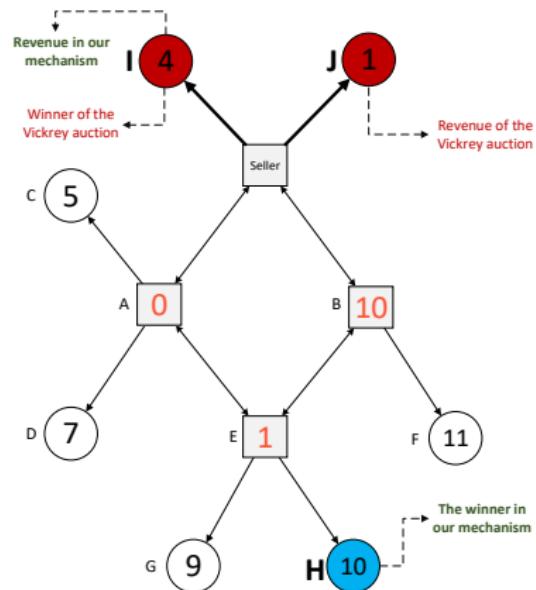
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# Economic Networks



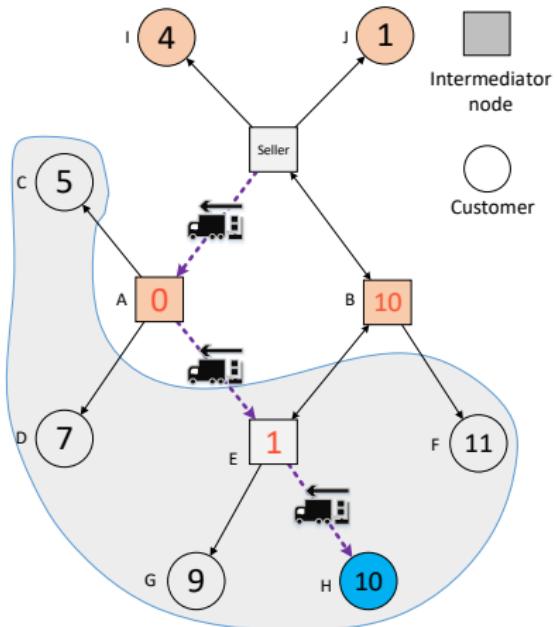
# Customer Sharing



## Benefits

- Sharing improves social welfare.
- Sharing increases revenue.

# Customer Sharing: Single Item Auction



- Unobservable Network.
- Private Customer Group.
- Strategic Neighbour-Sharing.
- Strategic Bids Report.
- Transaction Costs.
- Limited Communication.

# Customer Sharing: Single Item Auction

## The Goal

Design mechanisms that incentivize intermediate nodes to **share all their neighbors to the seller** and at the same time, motivate all buyers to **reveal their true valuations** on the commodity.

# Customer Sharing: Single Item Auction

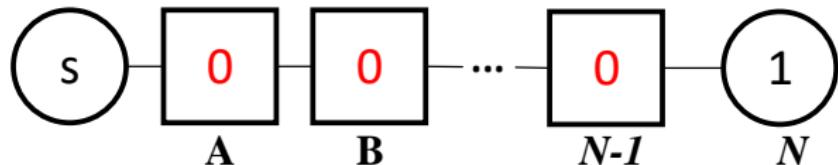
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## Efficiency & Revenue

More importantly, we also want **maximize the social welfare** (globally) and achieve a **higher revenue** compared with holding an auction in seller's neighbors only (the Vickrey auction).

# The VCG Mechanism is not Applicable



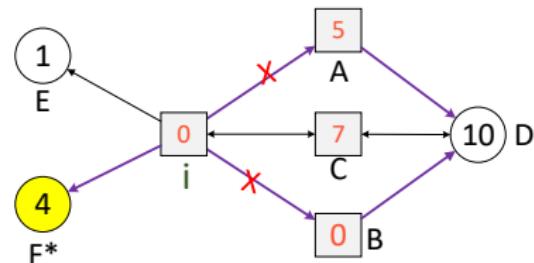
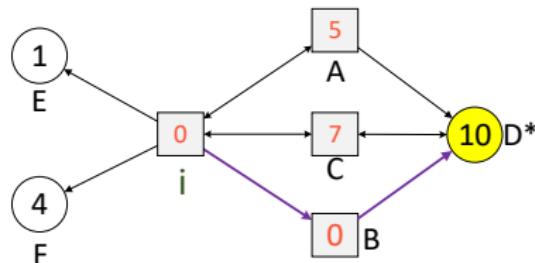
Claim.

The seller's revenue is  $-(N - 1)$  when applying the VCG mechanism in above line network.

# Customer Sharing Mechanism: Threshold Neighborhood

## Definition

Given the agents' type report profile  $t'$ , for each intermediate node  $i \in N$ , define  $i$ 's threshold neighbourhood  $r_i^{*''}$  as the minimum subset of  $r_i''$  that makes **the winner under efficient allocation changed if  $i$  does not share the sale information to  $r_i^{*''}$** , i.e.,  $r_i^{*''} = \arg \min_{r_i'' \subseteq r_i'} \{|r_i''|\}$  where  $\pi_{m'}(r_i' \setminus r_i'', t_{-i}'') = 1 \wedge \pi_m(r_i', t_{-i}') = 1 \wedge m' \neq m$  and  $\pi$  is an efficient allocation.



# Customer Sharing Mechanism(CSM)

- **Allocation policy:** Given a feasible type profile  $t'$ , allocate the commodity to buyer  $m = \arg \max_{j \in N} SW_j$  (i.e., efficient allocation) and trade the commodity along  $LCC_m$ .
- **Payment policy:** The payment policy is defined for each category of agents as follows.
  - for customer  $i \in N$ , her payment is defined as:

$$W_{-i}^* - W^*(t') + v_i(t'_i, \pi^{csm}).$$

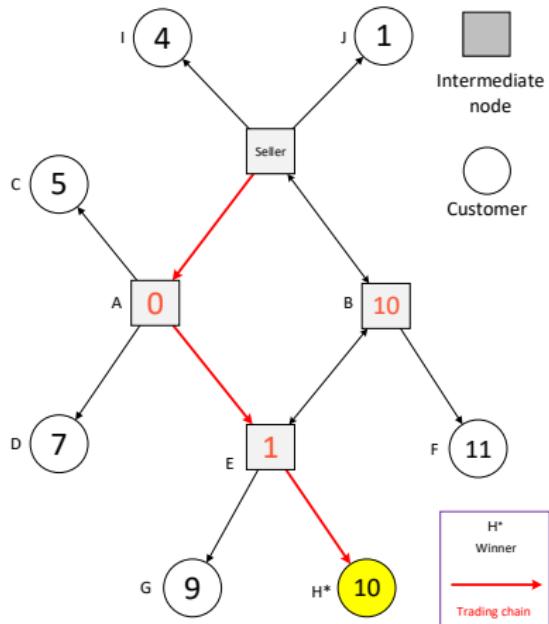
- for an intermediate node  $i$ , her payment is:

$$W_{-d_i}^* - W_{-r_i''}^* + v_i(t'_i, \pi^{csm}),$$

where  $W_{-r_i''}^*$ , denotes the maximum social welfare under feasible type profile  $(r'_i \setminus r_i'', t''_{-i})$ .

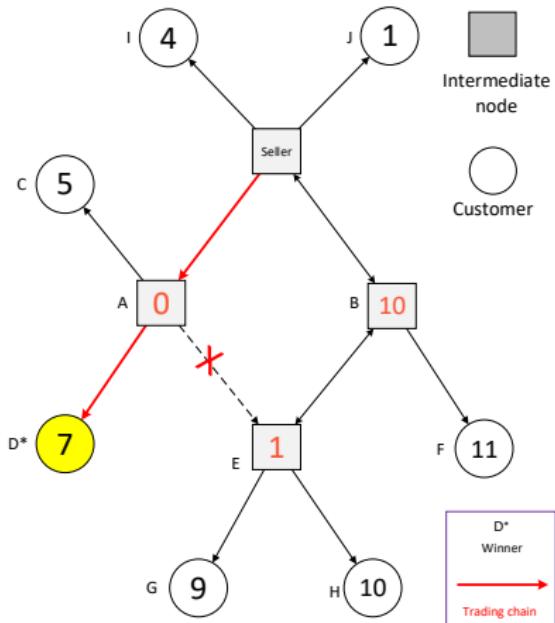
**PS:**  $LCC_m$  is the trading chain with least costs from  $s$  to  $m$ .

# Running Example



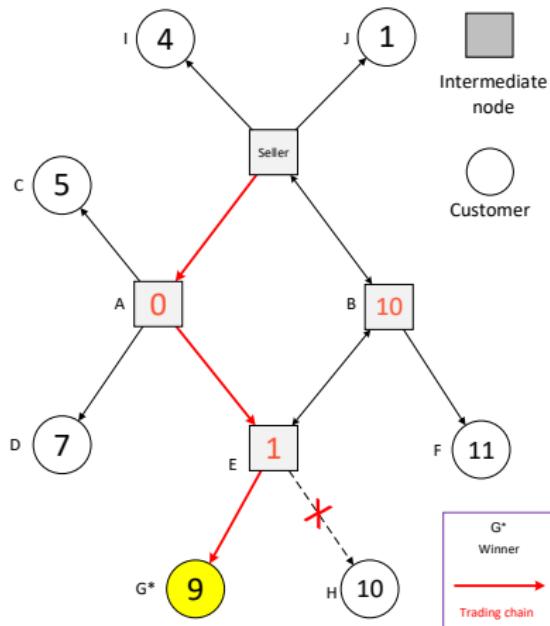
- ① The trading chain is **S-A-E-H** and the winner is **H**.

# Running Example



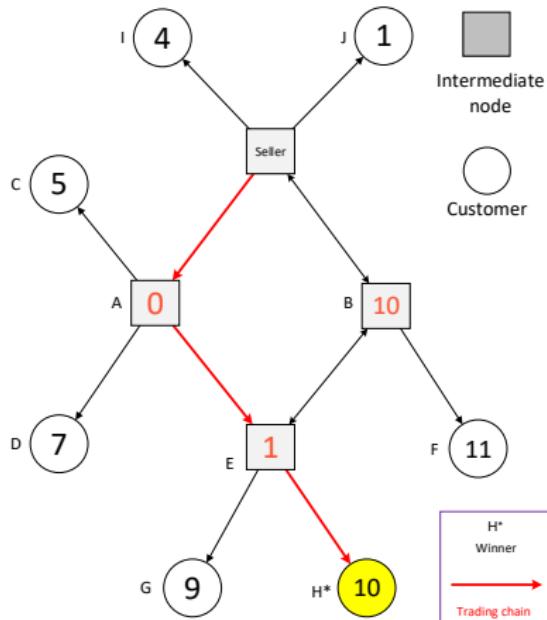
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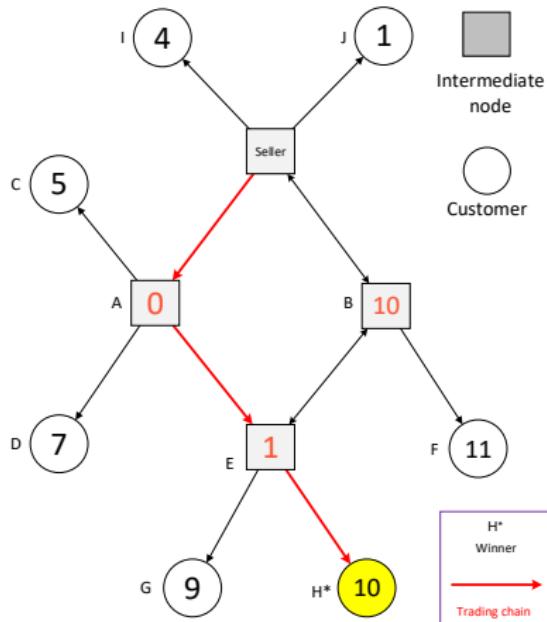
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- ③ E's threshold neighborhood is **H** and her payment is  $7-8+(-1)=-2$ .

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- ④  $H^*$ 's payment is  $8-9+10=9$ .

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- ④ H\*'s payment is  $8-9+10=9$ .
- ⑤ The total revenue of seller is  $-3+(-2)+9=4$ .

# Main Theorem

## Theorem

*CSM is efficient, individually rational, incentive compatible and budget balanced. In particular, the seller's revenue is no less than  $W_{-d_1^*}^*$  (no less than that given in the Vickrey auction) where agent  $1^*$  is the first agent in  $LCC_m$  with  $r_i^* \neq r_i$ .*

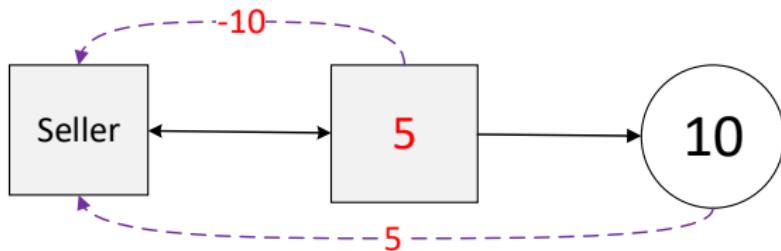
# Future Work

- What if intermediate nodes also have valuations on the commodity?
- Characterize all truthful customer sharing mechanisms.
- Revenue maximization problem in economic networks.

## Q & A

Thank you for listening!

## When Costs are Private



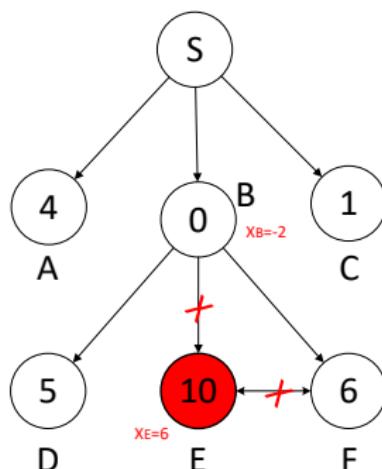
Myerson et al. 1983

There is no bilateral trade mechanism which satisfies the four properties mentioned above, i.e., IR, WBB, BNIC and EE.

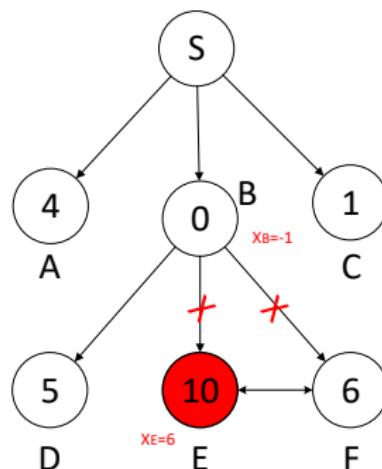
## Other Issues: When Costs Disappear

Claim.

When costs are disappear, IDM with  $-r_i*$  instead of  $-d_i$  achieves a higher revenue for the seller because of  $-r_i* \subseteq -d_i$ .



Original IDM with  $-d_i$



New IDM with  $-r_i^*$

# Simulations

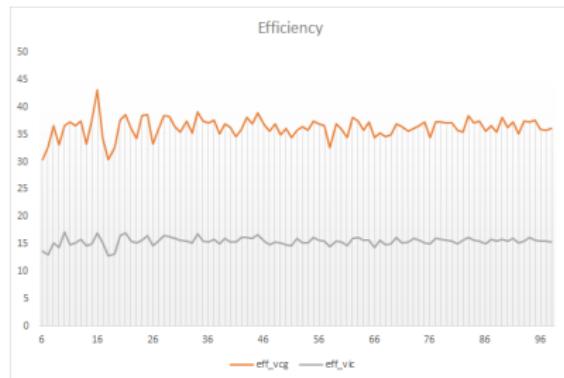


Figure: Efficiency comparison of three mechanisms

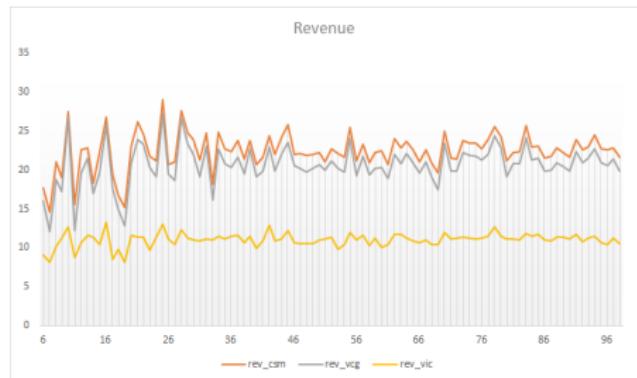


Figure: Revenue comparison of three mechanisms