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Code updating history of Jump Code Experiment

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Notes about the Matlab code used in the Jump Code Experiments.

BACKGROUND (APR 07, 2015)

All the codes are available in SVN '[https://github.com/bannisilloyd/dawei-qip-matlab/trunk/Jump Code](https://github.com/bannisilloyd/dawei-qip-matlab/trunk/Jump%20Code)'

The flipping error channel and erasing error channel are

$$A_0 = \begin{pmatrix} 1 & 0 \\ 0 & \sqrt{1-r} \end{pmatrix}, A_1 = \begin{pmatrix} 0 & 0 \\ 0 & \sqrt{r} \end{pmatrix}, \quad (1)$$

Since $A_0 = [(1 + \sqrt{1-r})/2]I + [(1 - \sqrt{1-r})/2]Z$, we know in the Hadamard basis the qubit will be flipped with the probability

$$p_F = [(1 - \sqrt{1-r})/2]^2, \quad (2)$$

and no error probability

$$p_0 = [(1 + \sqrt{1-r})/2]^2. \quad (3)$$

And $A_1|+\rangle = \sqrt{r/2}|1\rangle$, $A_1|-\rangle = -\sqrt{r/2}|1\rangle$, so A_1 will erase all the information in the Hadamard basis with the probability $p_E = r/2$.

To make a unitary transformation, we need one ancilla. Assume the ancilla starts at $|0\rangle$, so we can use the following way to construct unitary. The channel Φ on these 2 qubits (ancilla A and system B) should give us

$$\Phi(\rho_{AB}) = |0\rangle\langle 0| \otimes A_0\rho_B A_0^\dagger + |1\rangle\langle 1| \otimes A_1\rho_B A_1^\dagger. \quad (4)$$

Φ is not unitary but we can use a unitary channel U with the aid of post selection to realize it. It means if we start from $|00\rangle$ and $|01\rangle$ (Note we have assumed the ancilla A is from $|0\rangle$), by applying U we will have

$$U|00\rangle = |00\rangle, U|01\rangle = \sqrt{1-r}|01\rangle + \sqrt{r}|11\rangle. \quad (5)$$

So the unitary can be chosen in this way

$$U = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & \sqrt{1-r} & 0 & \sqrt{r} \\ 0 & 0 & -1 & 0 \\ 0 & \sqrt{r} & 0 & -\sqrt{1-r} \end{pmatrix}. \quad (6)$$

In summary, we have three probabilities

$$\begin{aligned} p_0 &= [(1 + \sqrt{1-r})/2]^2, \\ p_F &= [(1 - \sqrt{1-r})/2]^2, \\ p_E &= r/2. \end{aligned} \quad (7)$$

CLASSICAL CASE WITHOUT CORRECTION (APR 07, 2015)

First consider the classical case. Without correction, one needs one bit with the two states $|+\rangle$ or $|-\rangle$. Suppose it is $|+\rangle$.

We can use 2 qubits (ancilla and system) starting from $|0+\rangle$. Applying U and post-select the ancilla. If we get $|0\rangle$, we know A_0 happens and with probability p_0 no error happens, with probability p_F $|+\rangle$ is flipped to $|-\rangle$. If we get $|1\rangle$, we know A_1 happens (the probability of this case is p_E) and we will get $|1\rangle$ on the system. We have half chance to recover it to the correct state $|+\rangle$. Therefore, the fidelity for this classical case will be

$$F^C = p_0 + \frac{1}{2}p_E = \frac{1}{2} + \frac{1}{2}\sqrt{1-r} \quad (8)$$

The matlab code is (for a given r)

```

1 % projective measurement
2 M0 = kron(ST0, I); M1 = kron(ST1, I);
3
4 rho_projective0 = M0*rho_channel*M0'; % if ST0 is detected on the ancilla
5 % equal to rho_echo = Gz_echo(rho_channel,2), simulated by Gradient Echo
6 rho_sys0 = ptrace(rho_projective0, 1, [2 2]); % trace out ancilla
7 Fidelity0 = trace(System*rho_sys0);
8
9 rho_projective1 = M1*rho_channel*M1'; % if ST1 is detected on the ancilla
10 % equal to rho_echo = Gz_echo(rho_channel,2), simulated by Gradient Echo
11 rho_sys1 = ptrace(rho_projective1, 1, [2 2]); % trace out ancilla
12 Fidelity1 = trace(System*rho_sys1);

```

The total fidelity is Fidelity0+Fidelity1 and equals to F^C .

CLASSICAL CASE WITH CORRECTION (APR 07, 2015)

Without correction, one needs two bits with the two states $|++\rangle$ or $|--\rangle$ to encode one bit information. Suppose it is $|++\rangle$.

We have to use 4 qubits (ancilla, system, system, ancilla) to simulate this case. Applying U and post-select the 2 ancillas.

Case 1: If we get $|00\rangle$, we know A_0A_0 happens. We have p_0^2 probability for no error and p_F^2 for flip error on both qubits. The latter one cannot be corrected. Besides, we have $2p_0p_F$ probability that only one flip error happens. So we can project the system to $\{|++\rangle, |+-\rangle, |-+\rangle, |--\rangle\}$ basis. If we get $|++\rangle$ or $|--\rangle$, do nothing. If we get $|+-\rangle$ or $|-+\rangle$, discard the 2nd qubit and use the 1st qubit state as the real one. It means we have p_0p_F probability to succeed, and p_Fp_0 probability to fail. In total, for Case 1, the succeeding probability is $p_0^2 + p_0p_F$.

Case 2: If we get $|01\rangle$, we know A_0A_1 happens. Discard the 2nd qubit and use the 1st qubit state. The succeeding probability is thus p_0p_E .

Case 3: If we get $|10\rangle$, we know A_1A_0 happens. Discard the 1st qubit and use the 2nd qubit state. The succeeding probability is thus p_Ep_0 .

Case 4: If we get $|11\rangle$, we know A_1A_1 happens. The final result would be $|11\rangle$ on the system and we have half chance to get the correct state $|++\rangle$. The probability is thus $\frac{1}{2}p_E^2$.

The total fidelity with error correction in the classical case will be

$$F_{code}^C = p_0^2 + p_0p_F + p_0p_E + p_Ep_0 + \frac{1}{2}p_E^2. \quad (9)$$

Comparing the classical case with and without error correction, we will have the following figure. And we know for any r , the error correction performs better than the original case.

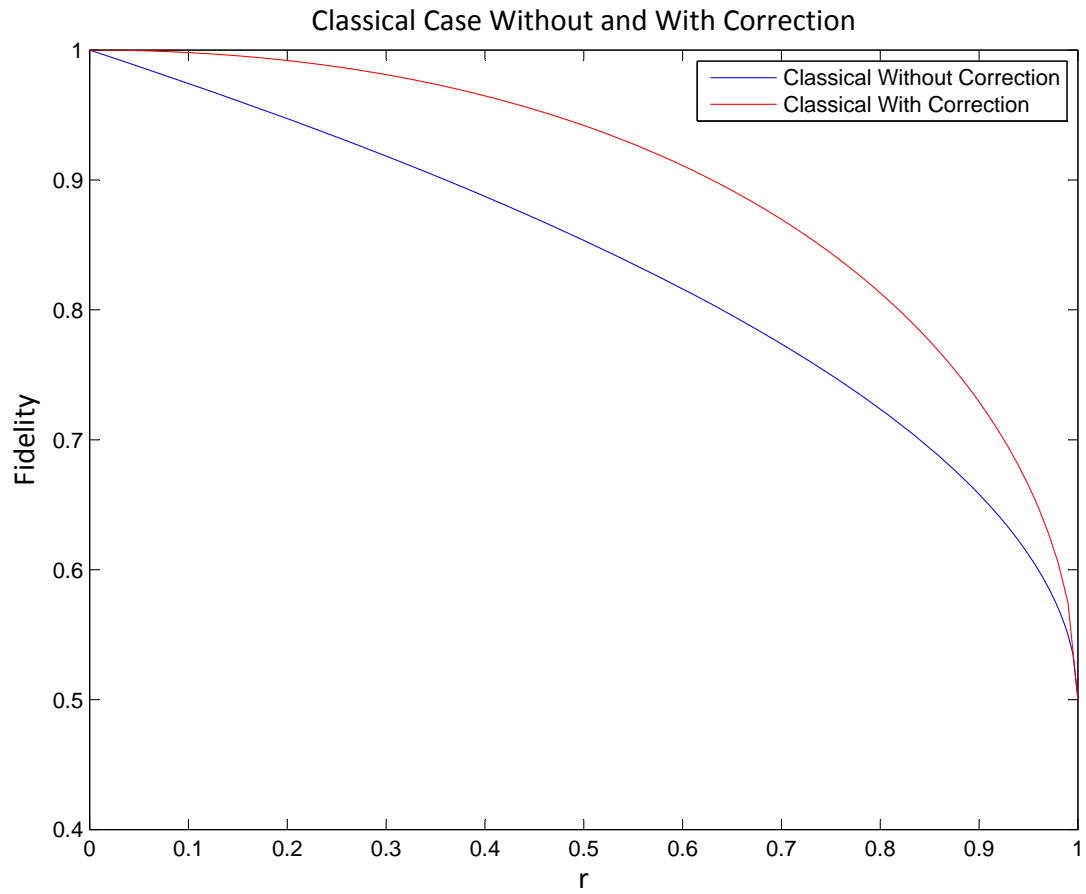


FIG. 1. Classical case with and without error correction.