

equivalently

$$\begin{aligned}
 & \text{minimize} \quad \left\| \sum_{j=1}^q \mu_j v_j - \sum_{i=1}^p \lambda_i u_i \right\|_2 \\
 & \text{subject to} \\
 & \quad \sum_{i=1}^p \lambda_i = 1, \quad \lambda \geq 0 \\
 & \quad \sum_{j=1}^q \mu_j = 1, \quad \mu \geq 0.
 \end{aligned}$$

Geometrically,  $\sum_{i=1}^p \lambda_i u_i$  with  $\sum_{i=1}^p \lambda_i = 1$  and  $\lambda \geq 0$  is a convex combination of the  $u_i$ s, and  $\sum_{j=1}^q \mu_j v_j$  with  $\sum_{j=1}^q \mu_j = 1$  and  $\mu \geq 0$  is a convex combination of the  $v_j$ s, so the dual program is to minimize the distance between the polyhedron  $\text{conv}(u_1, \dots, u_p)$  (the convex hull of the  $u_i$ s) and the polyhedron  $\text{conv}(v_1, \dots, v_q)$  (the convex hull of the  $v_j$ s). Since both polyhedra are compact, the shortest distance between them is achieved. In fact, there is some vertex  $u_i$  such that if  $P(u_i)$  is its projection onto  $\text{conv}(v_1, \dots, v_q)$  (which exists by Hilbert space theory), then the length of the line segment  $(u_i, P(u_i))$  is the shortest distance between the two polyhedra (and similarly there is some vertex  $v_j$  such that if  $P(v_j)$  is its projection onto  $\text{conv}(u_1, \dots, u_p)$  then the length of the line segment  $(v_j, P(v_j))$  is the shortest distance between the two polyhedra).

If the two subsets are separable, in which case Problem (SVM<sub>h1</sub>) has an optimal solution  $\delta > 0$ , because the objective function is convex and the convex constraint  $\|w\|_2 \leq 1$  is qualified since  $\delta$  may be negative, by Theorem 50.17(2) the duality gap is zero, so  $\delta$  is half of the minimum distance between the two convex polyhedra  $\text{conv}(u_1, \dots, u_p)$  and  $\text{conv}(v_1, \dots, v_q)$ ; see Figure 50.19.

It should be noted that the constraint  $\|w\| \leq 1$  yields a formulation of the dual problem which has the advantage of having a nice geometric interpretation: finding the minimal distance between the convex polyhedra  $\text{conv}(u_1, \dots, u_p)$  and  $\text{conv}(v_1, \dots, v_q)$ . Unfortunately this formulation is not useful for actually solving the problem. However, if the equivalent constraint  $\|w\|^2 (= w^\top w) \leq 1$  is used, then the dual problem is much more useful as a solving tool.

In Chapter 54 we consider the case where the sets of points  $\{u_1, \dots, u_p\}$  and  $\{v_1, \dots, v_q\}$  are not linearly separable.

## 50.12 Some Techniques to Obtain a More Useful Dual Program

In some cases, it is advantageous to reformulate a primal optimization problem to obtain a more useful dual problem. Three different reformulations are proposed in Boyd and Van-