



Figure 45.5: The cube centered at the origin with diagonal through  $(-1, -1, -1)$  and  $(1, 1, 1)$  has twelve edges. The edge from  $(1, 1, -1)$  to  $(1, 1, 1)$  is associated with the linear form  $x + y = 2$ .

Observe that a 0-dimensional face of  $\mathcal{P}$  is a vertex. If  $\mathcal{P}$  has dimension  $d$ , then the  $(d - 1)$ -dimensional faces of  $\mathcal{P}$  are called its *facets*.

If  $(P)$  is a linear program in standard form, then its basic feasible solutions are exactly the vertices of the polyhedron  $\mathcal{P}(A, b)$ . To prove this fact we need the following simple proposition

**Proposition 45.5.** *Let  $Ax = b$  be a linear system where  $A$  is an  $m \times n$  matrix of rank  $m$ . For any subset  $K \subseteq \{1, \dots, n\}$  of size  $m$ , if  $A_K$  is invertible, then there is at most one basic feasible solution  $x \in \mathbb{R}^n$  with  $x_j = 0$  for all  $j \notin K$  (of course,  $x \geq 0$ )*

*Proof.* In order for  $x$  to be feasible we must have  $Ax = b$ . Write  $N = \{1, \dots, n\} - K$ ,  $x_K$  for the vector consisting of the coordinates of  $x$  with indices in  $K$ , and  $x_N$  for the vector consisting of the coordinates of  $x$  with indices in  $N$ . Then

$$Ax = A_K x_K + A_N x_N = b.$$

In order for  $x$  to be a basic feasible solution we must have  $x_N = 0$ , so

$$A_K x_K = b.$$

Since by hypothesis  $A_K$  is invertible,  $x_K = A_K^{-1}b$  is uniquely determined. If  $x_K \geq 0$  then  $x$  is a basic feasible solution, otherwise it is not. This proves that there is at most one basic feasible solution  $x \in \mathbb{R}^n$  with  $x_j = 0$  for all  $j \notin K$ .  $\square$

**Theorem 45.6.** *Let  $(P)$  be a linear program in standard form, where  $Ax = b$  and  $A$  is an  $m \times n$  matrix of rank  $m$ . For every  $v \in \mathcal{P}(A, b)$ , the following conditions are equivalent:*