45.2 Basic Feasible Solutions and Vertices

If the system Ax = b has a solution and if some row of A is a linear combination of other rows, then the corresponding equation is redundant, so we may assume that the rows of A are linearly independent; that is, we may assume that A has rank m, so $m \le n$.

Definition 45.6. If A is an $m \times n$ matrix, for any nonempty subset K of $\{1, \ldots, n\}$, let A_K be the submatrix of A consisting of the columns of A whose indices belong to K. We denote the jth column of the matrix A by A^j .

Definition 45.7. Given a Linear Program (P_2)

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maximize cx
subject to Ax = b and x \ge 0,
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where A has rank m, a vector $x \in \mathbb{R}^n$ is a basic feasible solution of (P) if $x \in \mathcal{P}(A, b) \neq \emptyset$, and if there is some subset K of $\{1, \ldots, n\}$ of size m such that

- (1) The matrix A_K is invertible (that is, the columns of A_K are linearly independent).
- (2) $x_i = 0$ for all $j \notin K$.

The subset K is called a basis of x. Every index $k \in K$ is called basic, and every index $j \notin K$ is called nonbasic. Similarly, the columns A^k corresponding to indices $k \in K$ are called basic, and the columns A^j corresponding to indices $j \notin K$ are called nonbasic. The variables corresponding to basic indices $k \in K$ are called basic variables, and the variables corresponding to indices $j \notin K$ are called nonbasic.

For example, the linear program

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maximize x_1 + x_2
subject to x_1 + x_2 + x_3 = 1 and x_1 \ge 0, x_2 \ge 0, x_3 \ge 0, (*)
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has three basic feasible solutions; the basic feasible solution $K = \{1\}$ corresponds to the point (1,0,0); the basic feasible solution $K = \{2\}$ corresponds to the point (0,1,0); the basic feasible solution $K = \{3\}$ corresponds to the point (0,0,1). Each of these points corresponds to the vertices of the slanted purple triangle illustrated in Figure 45.3. The vertices (1,0,0) and (0,1,0) optimize the objective function with a value of 1.

We now show that if the Standard Linear Program (P_2) as in Definition 45.7 has some feasible solution and is bounded above, then some basic feasible solution is an optimal solution. We follow Matousek and Gardner [123] (Chapter 4, Section 2, Theorem 4.2.3).

First we obtain a more convenient characterization of a basic feasible solution.

Proposition 45.2. Given any Standard Linear Program (P_2) where Ax = b and A is an $m \times n$ matrix of rank m, for any feasible solution x, if $J_{>} = \{j \in \{1, ..., n\} \mid x_j > 0\}$, then x is a basic feasible solution iff the columns of the matrix $A_{J_{>}}$ are linearly independent.