

Math 104 HW 1

1. induction \rightarrow base case: $1^3 = 1$, and $1^2 = 1$; indeed, $1 = 1$

inductive step: assume $1^3 + 2^3 + \dots + n^3 = (1 + 2 + \dots + n)^2$

ACTUALLY, Lemma ①: $1 + 2 + \dots + n = \frac{n(n+1)}{2}$

proof: base case $n=1$: $\frac{1(2)}{2} = 1 = 1 \checkmark$

inductive step: assume $1 + 2 + \dots + n = \frac{n(n+1)}{2}$. adding $n+1$ gives

$$1 + 2 + \dots + n + (n+1) = \frac{n(n+1)}{2} + (n+1) = \frac{n(n+1) + 2(n+1)}{2} = \frac{(n+1)(n+2)}{2}, \text{ and}$$

we've shown $P(n) \rightarrow P(n+1)$. by induction, $1 + 2 + \dots + n = \frac{n(n+1)}{2}$ \blacksquare

using lemma ①, we can restate the problem to be

$$\text{proving that } 1^3 + 2^3 + \dots + n^3 = \left(\frac{n(n+1)}{2}\right)^2 = \frac{n^2(n+1)^2}{4} = \frac{n^2(n^2 + 2n + 1)}{4} = \frac{n^4 + 2n^3 + n^2}{4}.$$

we've already shown the base case, now for the

$$\text{inductive step: if } 1^3 + 2^3 + \dots + n^3 = \frac{n^4 + 2n^3 + n^2}{4} \text{ then}$$

$$1^3 + 2^3 + \dots + n^3 + (n+1)^3 = \frac{n^4 + 2n^3 + n^2}{4} + (n+1)^3 = \frac{n^4 + 2n^3 + n^2 + 4n^3 + 12n^2 + 12n + 4}{4}$$

$$= \frac{n^4 + 6n^3 + 13n^2 + 12n + 4}{4}. \text{ so we must show } n^4 + 6n^3 + 13n^2 + 12n + 4 = (n+1)^2(n+2)^2.$$

$$(n+1)^2(n+2)^2 = (n^2 + 2n + 1)(n^2 + 4n + 4) = n^4 + 4n^3 + 4n^2 + 2n^3 + 8n^2 + 8n + n^2 + 4n + 4 = n^4 + 6n^3 + 13n^2 + 12n + 4, \text{ as desired. } \blacksquare$$

2a. base case: $n=2$. $4 > 3$, as desired.

inductive step: if $n^2 > n+1$, then $n^2 + 2n + 1 > n+1 + 2n+1$

$$\rightarrow (n+1)^2 > (n+1) + 1 + 2n > (n+1) + 1, \text{ since } 2n > 0 \text{ (} n \geq 2 \text{)} \blacksquare$$

b. base case: $n=4$. $4! = 24 > 4^2 = 16$, as desired.

inductive step: if $n! > n^2$, $n!(n+1) > n^2(n+1)$

$\rightarrow (n+1)! > n^3 + n^2$. we want $(n+1)! > (n+1)^2 = n^2 + 2n + 1$, so we must

show that $n^3 + n^2 > n^2 + 2n + 1$ for $n \geq 4$:

base case: $n=4$, $64 > 9$, as desired. inductive step: if

$$n^3 > 2n + 1, \quad n^3 + 3n^2 + 3n + 1 > 3n^2 + 3n + 1 + 2n + 1, \text{ and } (n+1)^3 > 3n^2 + 5n + 2$$

$$> 5n + 2 = 4n + 2 + n > 4n + 2 + 1 \text{ (since } n \geq 4 \text{)} = 2(2n + 1) + 1 > 2(n+1) + 1. \blacksquare$$

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3. Let P_n be the proposition that $F_n = f(n)$ and $F_{n-1} = f(n-1)$.

base case $P_1: F_0 = 0$ and $F_1 = 1$ (given).

$$f(0) = \frac{1}{\sqrt{5}}(1-1) = 0 \text{ and } f(1) = \frac{1}{\sqrt{5}}\left(\frac{\sqrt{5}+1}{2} - \left(1 - \frac{\sqrt{5}+1}{2}\right)\right) \\ = \frac{1}{\sqrt{5}}\left(\frac{\sqrt{5}+1 - (2 - \sqrt{5} - 1)}{2}\right) = \frac{2\sqrt{5}+2-2}{2\sqrt{5}} = 1. \text{ Thus } P_1 \text{ is true.}$$

inductive step: assuming P_n (for $n \geq 1$), show P_{n+1} is true.

$$\text{if } F_n = \frac{1}{\sqrt{5}}(\varphi^n - (1-\varphi)^n) \text{ and } F_{n-1} = \frac{1}{\sqrt{5}}(\varphi^{n-1} - (1-\varphi)^{n-1}), \text{ then} \\ F_n + F_{n-1} = F_{n+1} \text{ (def. of sequence)} = \frac{1}{\sqrt{5}}(\varphi^n + \varphi^{n-1} - (1-\varphi)^n - (1-\varphi)^{n-1}) \\ = \frac{1}{\sqrt{5}}(\varphi^{n-1}(\varphi+1) - (1-\varphi)^{n-1}((\frac{1-\varphi}{\varphi})+1)).$$

For convenience, i will rename $1-\varphi$ to be α . We are given that φ solves $x^2 = x+1$. As it turns out, α does too: $(1-\varphi)^2 = \varphi^2 - 2\varphi + 1 = (1-\varphi+1) \Rightarrow \varphi^2 - \varphi - 1 = 0$, and $\varphi^2 = \varphi+1$, so $(\varphi+1) - \varphi - 1 = 0 = 0$, and $\alpha^2 = \alpha+1$, too. Continuing above:

$$F_{n+1} = \frac{1}{\sqrt{5}}(\varphi^{n-1}(\varphi^2) - \alpha^{n-1}(\alpha^2)) = \frac{1}{\sqrt{5}}(\varphi^{n+1} - (1-\varphi)^{n+1}). \text{ by induction,} \\ \text{the proof that } F_n = f(n) \text{ for } n \in \mathbb{N} \text{ (and } 0) \text{ is complete } \blacksquare$$

4. Using the same P_n as problem 3, ^{but with G_n instead of $f(n)$} base case:

if you can go up $1^{or 2}$ stairs at a time, there are no ^{is one} ways of going up 0 stairs, and 1 way of going up 1. (trivial exhaustion). indeed, $F_1 = 0$ and $F_2 = 1$.

inductive step: assume $G_{n-1} = F_{n+1}$ and $G_n = F_{n+2}$. To

get up $n+1$ stairs, you must either go up $n-1$ stairs and then go up 2, or up n steps then go up one.

Since you can take only 1 or 2 steps, these are

the only options, so $G_{n+1} = G_n + G_{n-1}$. Using the

inductive assumption, $G_{n+1} = F_{n+1} + F_{n+2} = F_{n+3}$ by definition of Fibonacci sequence, and by induction

$G_n = F_{n+1}$ for all $n \in \{0\} \cup \mathbb{N}$ \blacksquare

$$\sqrt{3} < 5$$

5. let $x = \sqrt[3]{5 - \sqrt{3}}$; $x^3 = 5 - \sqrt{3}$, $x^3 - 5 = -\sqrt{3}$

and $(x^3 - 5)^2 = 3$, $x^6 - 10x^3 + 22 = 0$. by the rational zeroes theorem, if x is rational it can only be $\pm 1, \pm 11, \pm 2, \pm 22$.

$5 > 5 - \sqrt{3} > \sqrt[3]{5 - \sqrt{3}} = x$, so 11 and 22 are eliminated.

Since $9 < 25$, $1 - \sqrt{3} < 5$ and $0 < 5 - \sqrt{3} < \sqrt[3]{5 - \sqrt{3}} = x$, leaving only $x = 1$ and $x = 2$. by exhaustion:

(1) $-10(1) + 22 = 1 - 10 + 22 = 13 \neq 0$

$64 - 80 + 22 = 86 - 80 = 6 \neq 0$.

Since no rational possibilities hold, x is irrational. \square

6. let $x = 30^{\frac{n}{m}}$. then $x^m = 30^n$.

Lemma ①: if $m \nmid n$ and $m, n \in \mathbb{N}$, and $m > n$, $30^{\frac{n}{m}}$ is irrational.

proof: by contradiction, assume x is rational, in which case it can be written as $\frac{a}{b}$ in lowest terms. then $(\frac{a}{b})^m = 30^n$, and $a^m = b^m \cdot 30^n$.

therefore a^m is even ($30 = 2 \cdot 15$), and so a is also even. we'll write $a = 2c$, where $c \in \mathbb{N}^+$, and

$$b^m = \frac{2^m \cdot c^m}{30^n} = \frac{2^m}{2^n} \cdot \frac{c^m}{15^n} = 2^{m-n} \cdot \frac{c^m}{15^n}. \text{ since } m > n, m-n > 0, \text{ and}$$

b^m is also even; b is even. but a and b cannot both be even, since lowest terms, so here $30^{\frac{n}{m}} \notin \mathbb{Q}$ \square

We've shown it's true when $m > n$. $m \neq n$, since $m \nmid n$, so we have to show for $m < n$:

since $n > m$, write $n = xm + y$, where $x, y \in \mathbb{N}$ and $y < m$

then $30^{\frac{n}{m}} = 30^{\frac{xm+y}{m}} = 30^x \cdot 30^{\frac{y}{m}}$. 30^x is rational, and since

$y < m$, we know $30^{\frac{y}{m}}$ is irrational. \square (lemma ①)

so $30^{\frac{n}{m}}$ is also irrational here. \square

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Lemma 2: rational \cdot irrational = irrational

rational can be written in lowest form as $\frac{a}{b}$, w/ $a, b \in \mathbb{Z}$

then rational \cdot irrational $\lambda = \frac{a}{b} \cdot \lambda = \frac{a\lambda}{b}$. so we want to show multiplication by an integer = irrational.

by contradiction: if $c \cdot \lambda = \frac{x}{y}$ (with $c, x, y \in \mathbb{Z}$), then $\lambda = \frac{x}{cy}$, which is impossible since λ is irrational.

similarly for division: $\frac{\lambda}{c} = \frac{x}{y} \rightarrow \lambda = \frac{cx}{y}$, impossible.

so, back to $\frac{a\lambda}{b}$, $a\lambda$ is irrational α , and $\frac{\alpha}{b}$ is also irrational. therefore, rational \cdot irrational = irrational \square

so we've shown if $n > m$, we can rewrite $30^{n/m}$ to $30^x \cdot 30^{y/m}$, where 30^x is rational and $30^{y/m}$ is not, by Lemma ①. By Lemma ②, this result is irrational.

Since we showed the result for $m > n$ and $m < n$ (and $m \neq n$ since $m \nmid n$) the proof is complete \square

7(1). First i'll show $N S' \rightarrow N S$:

Set $S \subseteq N$; $1 \in S$; $n \in S \rightarrow n+1 \in S$, by contradiction, assume $S \neq N$. Then the complement of S w/ respect to $N \neq \emptyset$, and by $N S'$ this complement \bar{S} contains a least element. The contrapositive of $n \in S \rightarrow n+1 \in S$ is $n+1 \notin S \rightarrow n \notin S$, or $n+1 \in \bar{S} \rightarrow n \in \bar{S}$. Therefore this smallest element of \bar{S} cannot exist, since for any $n \in \bar{S}$, there is a smaller $n-1 \in \bar{S}$. So our assumption that $S \neq N$ is false, and $N S' \rightarrow N S$.

so $S = \emptyset$,
contradicting
this

Next, that $N S \rightarrow N S'$:

if $(S \subseteq N \wedge 1 \in S \wedge n \in S \rightarrow n+1 \in S) \rightarrow S = N$, then we want to show that ① every nonempty subset has a least element and ② every $1 \neq n \in N$ is the successor of another number in N .

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② we know that $n \in N \rightarrow n+1 \in N$. by contradiction, assume $\exists w \in N$ s.t. w is not the successor of another number in N , and $w \neq 1$. That is, $w \in N$ and $w-1 \notin N$. i can construct a set S without w that is then N .

I start with $S = \{1\}$, and using $n \in S \rightarrow n+1 \in S$, build up the set as i go, adding each new successor: $S = \{1, 1+1, 1+1+1, \dots\}$. Since $w-1 \notin N$, w will never show up, and thus i have constructed a set $S \subseteq N$ ~~without~~ where $w \notin S$. $w \in N$.

① by contradiction, assume i CAN make such a subset S with no smallest element. The set S i made for ② contains only elements that are ordered — the bigger number is the one with more "+1"s, and no two different elements have a different # of "+1"s. Since $S \neq N$, i know this property is thus also true of N .

7(2). Consider the set $S = \{1, 1+1, 1+1+1, \dots\} \cup \{\omega, \omega+1, \omega+1+1, \dots\}$.
clearly, $\{1, \omega\}$ is a subset of S , but since they are
incomparable, there is no "smallest number" between
1 and ω . therefore, $\{1, \omega\} \subseteq S$ s.t. there is no smallest element