

**COMP9120 Relational Database Systems****Tutorial Week 9: Transaction Management****Exercise 1. Transaction Support in SQL**

The transaction concept is reflected in several parts of SQL. In SQL, you can group several statements into a transaction. Many SQL dialects, e.g. with SQL Server or also PostgreSQL, use an explicit **BEGIN TRANSACTION** (or **BEGIN**) command to start a transaction. A transaction is requested to finish successfully using **COMMIT** or aborted using **ROLLBACK**.

- a) Try out the following small script in PostgreSQL that tries to add three new lecture theatres in the new law building to our University database (using the University schema from previous tutorials): what classrooms starting with LS are shown before and after the **ROLLBACK** keyword?

```
BEGIN;
INSERT INTO Classroom VALUES ('LS101', 300, 'sloping');
INSERT INTO Classroom VALUES ('LS104', 100, 'sloping');
INSERT INTO Classroom VALUES ('LS106', 100, 'sloping');
-- check what we have so far:
SELECT * FROM Classroom WHERE ClassroomId LIKE 'LS%';
-- simulate a problem and abort our transaction
ROLLBACK;
-- check what is kept in the database
SELECT * FROM Classroom WHERE ClassroomId LIKE 'LS%';
```

- b) Now try the same script, but with the **ROLLBACK** statement replaced with **COMMIT**.

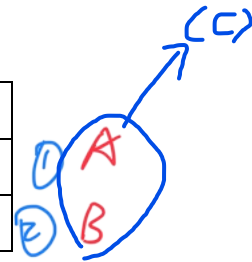
## Exercise 2. Serializability and Update Anomalies

Given the following relation

Offerings(uosCode, year, semester, lecturerId)

with this example instance:

uosCode	year	semester	lecturerId
COMP5138	2021	S1	4711
INFO2120	2021	S2	4711



Consider the following hypothetical interleaved execution of two transactions T1 and T2 in a DBMS where concurrency control (such as locking) is not done; that is, each statement is executed as it is submitted, using the most up-to-date values of the database contents.

T1	<b>SELECT * FROM</b> Offerings <b>WHERE</b> lecturerId = 4711	① ②
T2	<b>SELECT</b> year <b>INTO</b> :yr <b>FROM</b> Offerings <b>WHERE</b> uosCode = 'COMP5138'	①
T1	<b>UPDATE</b> Offerings <b>SET</b> year=year+1 <b>WHERE</b> lecturerId = 4711 <b>AND</b> uosCode = 'COMP5138'	update ① 2022
T2	<b>UPDATE</b> Offerings <b>SET</b> year=:yr+2 <b>WHERE</b> uosCode = 'COMP5138'	also update ① 2023
T1	<b>COMMIT</b>	store before update
T2	<b>COMMIT</b>	if it is year+2 then it's 2024

Indicate:

- the values returned by the **SELECT** statements and the final value of the database; → ... 2023 ...
- whether the execution produces any update anomalies *lost update*
- whether the execution is conflict serializable or not.

Not serializable:

T<sub>1</sub>: a, c, e first

T<sub>2</sub>: b, d, f second

and

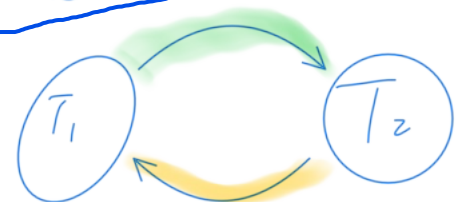
T<sub>1</sub> second

T<sub>2</sub> first

not equal current

not conflict

serializable



because year

draw based on conflict pairs

T<sub>1</sub>

same(A) no(B)  
R1(A) R1(B)

R1(A), W1(A)

T<sub>2</sub>

R2(A)

W2(A)

2  
4 conflict pairs

no read because : yr

### Exercise 3. Transaction Isolation experiment

In this exercise, we are going to give you some intuition of transaction management. You will begin by opening pgAdmin in two windows.

Execute the following two transactions line-by-line as seen in the following figure (press 'Execute' after each line) and write down the results of the SELECT COUNT() statements:

Window 1	Window 2	Count
BEGIN TRANSACTION;	BEGIN TRANSACTION; <u>SET TRANSACTION ISOLATION LEVEL</u> SERIALIZABLE;	<i>in this transaction</i> <i>Only see before this transaction</i>
SELECT COUNT(*) FROM student;		? N
	SELECT COUNT(*) FROM student;	? N
INSERT INTO student VALUES (); //insert test values		
SELECT COUNT(*) FROM student;		? N+1
	SELECT COUNT(*) FROM student;	? N
COMMIT;		
SELECT COUNT(*) FROM student;		? N+1
	SELECT COUNT(*) FROM student;	? N
	COMMIT;	
	SELECT COUNT(*) FROM student;	? N+1

↙ now can see outside