COMP2123
Data structures and Algorithms
Lecture 11: Divide and Conquer
[GT 11 and 9]

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Some content is taken from material provided by the textbook publisher Wiley.

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Divide and Conquer

Divide and Conquer algorithms can normally be broken into these three parts:

1. Divide If it is a base case, solve directly, otherwise break up the problem into several parts.

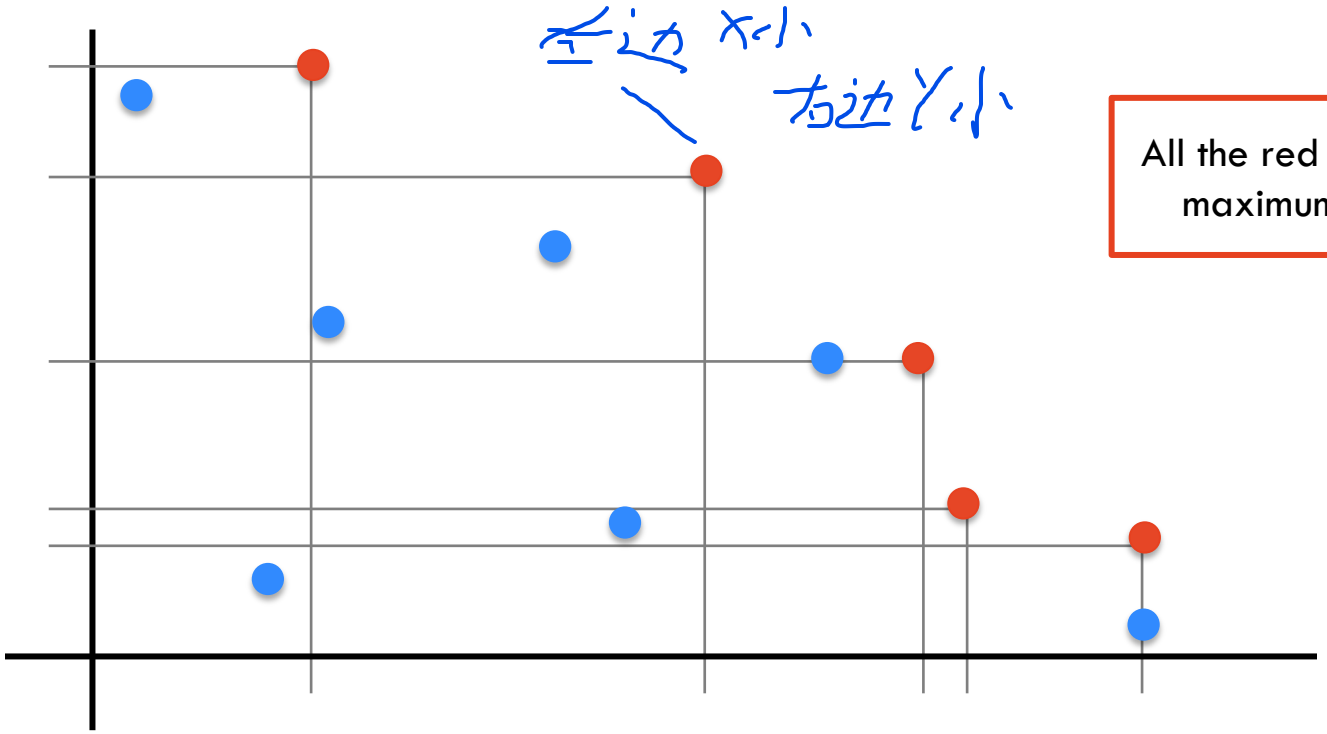
2. Recur/Delegate Recursively solve each part [each sub-problem].

3. Conquer Combine the solutions of each part into the overall solution.

Maxima-Set (Pareto frontier)

Definition A point is maximum in a set if all other points in the set have either a smaller x- or smaller y-coordinate.

Problem Given a set S of n distinct points in the plane (2D), find the set of all maximum points.



All the red points are maximum points.

Maxima-Set: Naïve Solution

Brute force $O(n^2)$

Idea: Check every point (one at a time) to see if it is a maximum point in the set S.

To check if point p is a maximum point in S:

```
for q in S do
   if q ≠ p and q.x ≥ p.x and q.y ≥ p.y then
      return "No"
return "Yes"
```

There is a point q that dominates p

Maxima-Set: Naïve Solution

Idea: Check every point (one at a time) to see if it is a maximum point in the set S.

To check if point p is a maximum point in S:

```
for q in S do
   if q ≠ p and q.x ≥ p.x and q.y ≥ p.y then
      return "No"
return "Yes"
```

Naïve algorithm to find the maxima-set of S:

```
maximaSet ← empty list
for p in S do
   if p is a maximum point in S then
      add p to the maximaSet
return maximaSet
```

Maxima-Set: Naïve Solution

Idea: Check every point (one at a time) to see if it is a maximum point in the set S.

To check if point p is a maximum point in S:

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O(n)

Naïve algorithm to find the maxima-set of S:

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O(n)

$O(n^2)$

Maxima-Set

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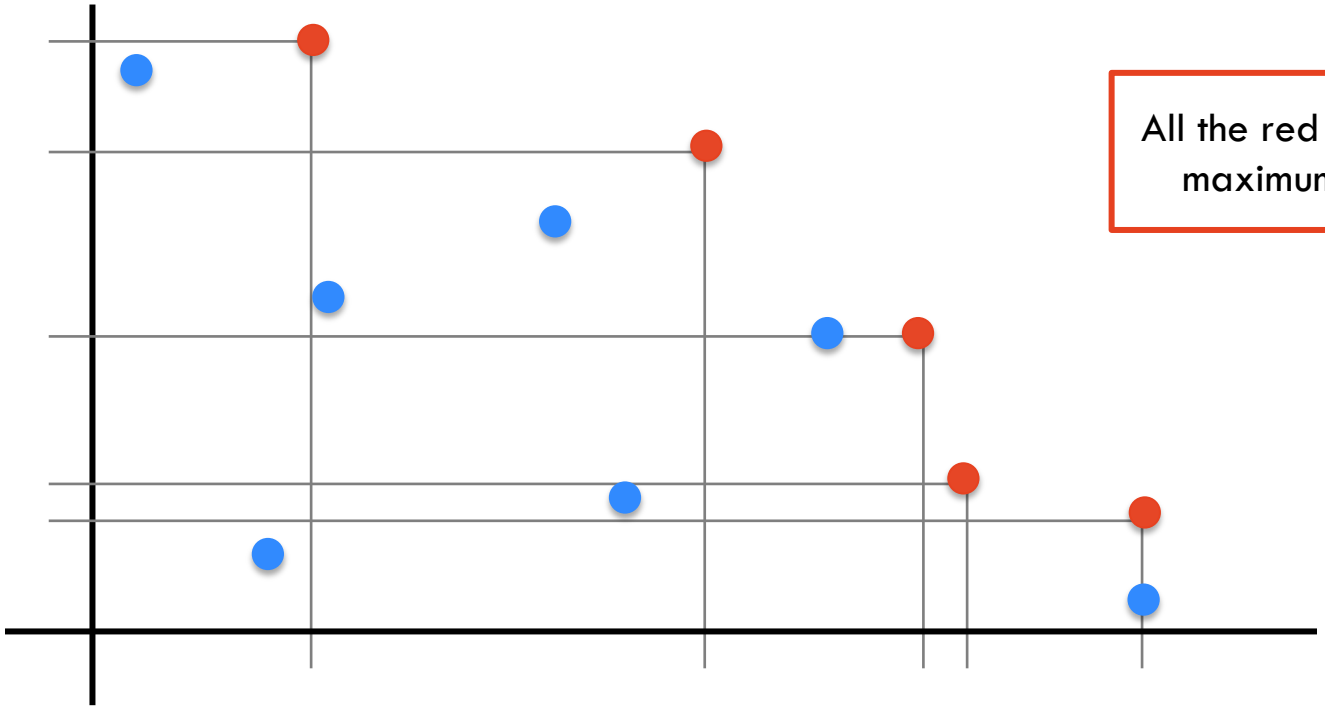


All the red points are maximum points.

Maxima-Set

Preprocessing Sort the points by increasing x coordinate and store them in an array. Note: we only do this once. Break ties in x by sorting by increasing y coordinate.

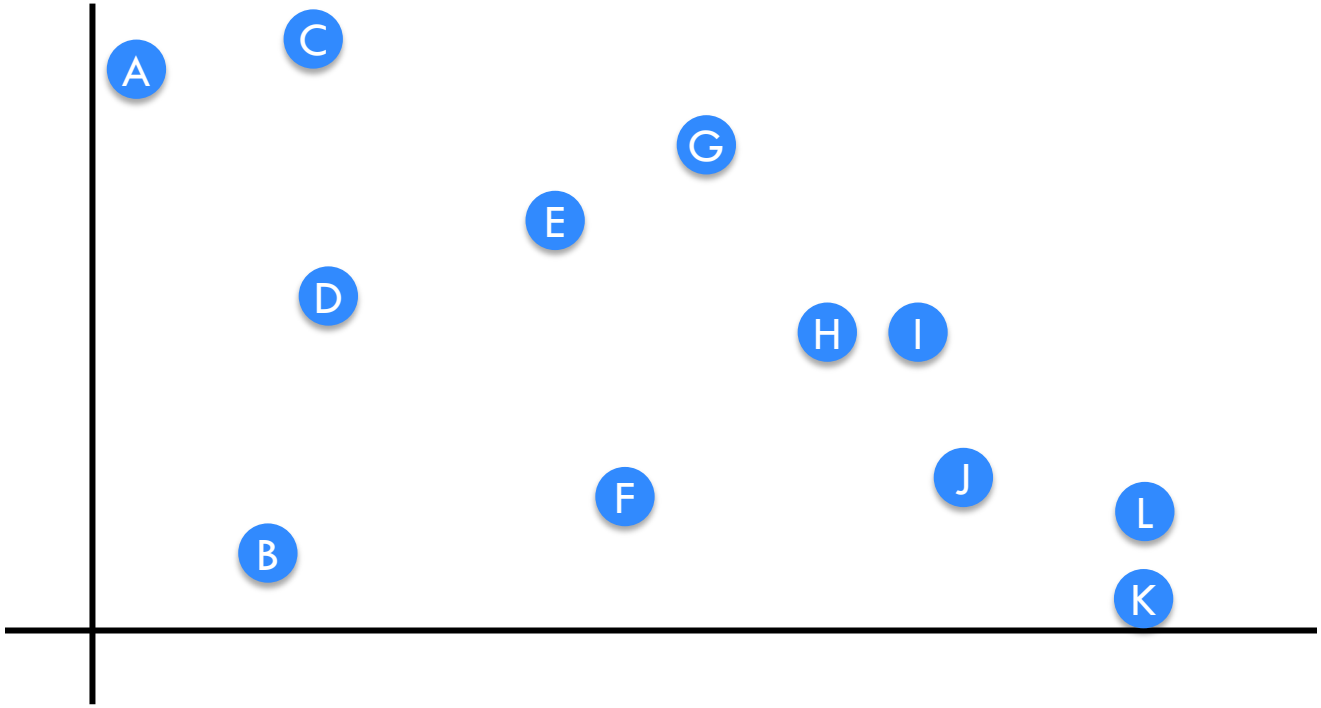
Divide sorted array into two halves.

Recur recursively find the MS of each half.

Conquer compute the MS of the union of Left and Right MS

Maxima-Set

Preprocessing Sort the points by increasing x coordinate and store them in an array. Note: we only do this once. Break ties in x by sorting by increasing y coordinate.

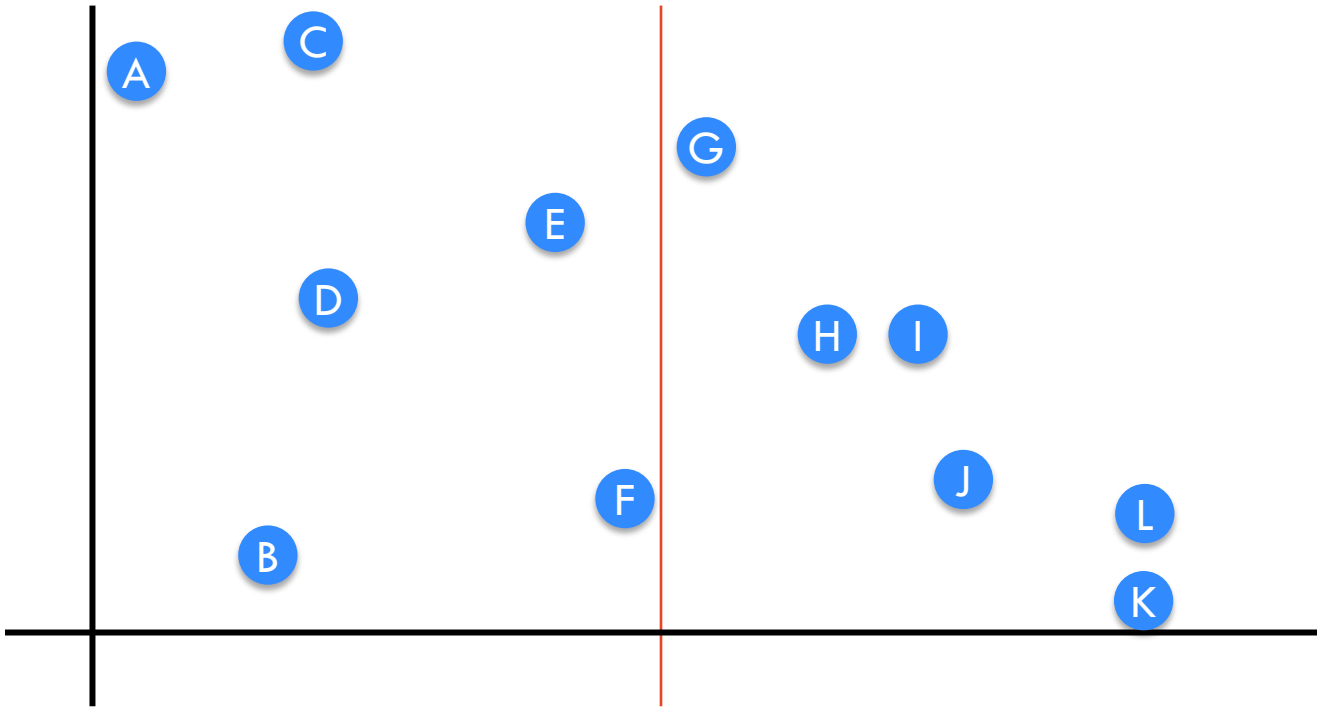


Sorted Points

A	B	C	D	E	F	G	H	I	J	K	L

Maxima-Set

Divide array into two halves.

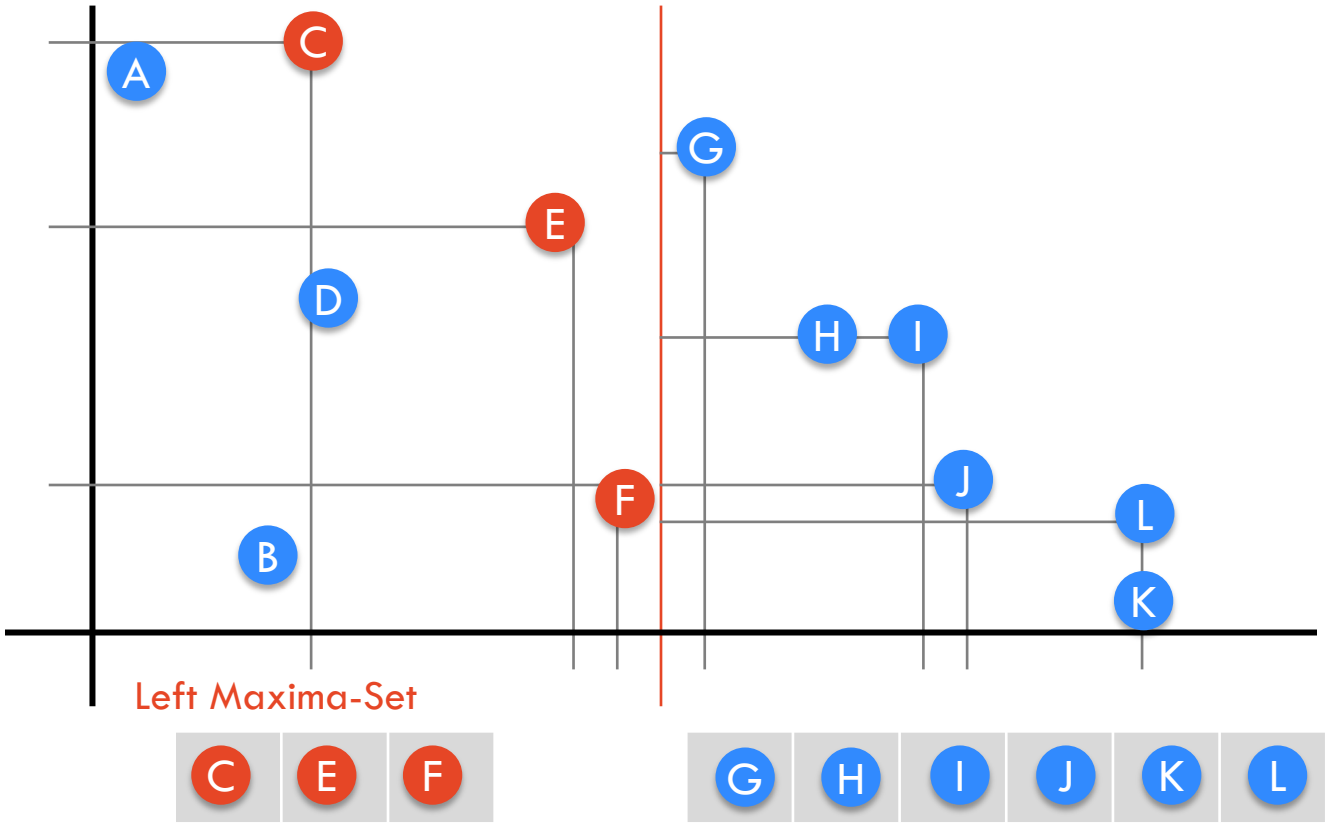


Sorted Points

A	B	C	D	E	F	G	H	I	J	K	L

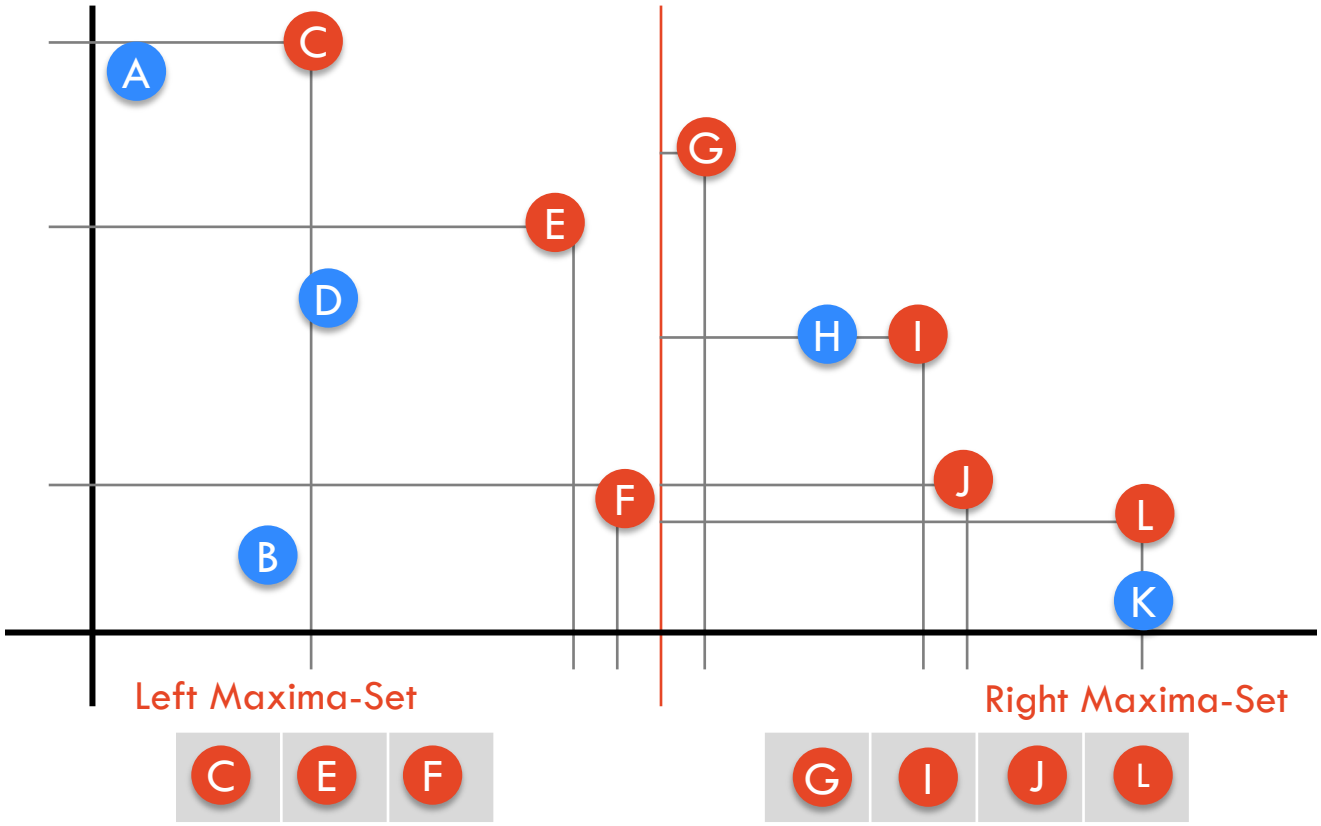
Maxima-Set

Recur recursively find the Maxima-Set of each half.



Maxima-Set

Recur recursively find the Maxima-Set of each half.



Maxima-Set

Conquer

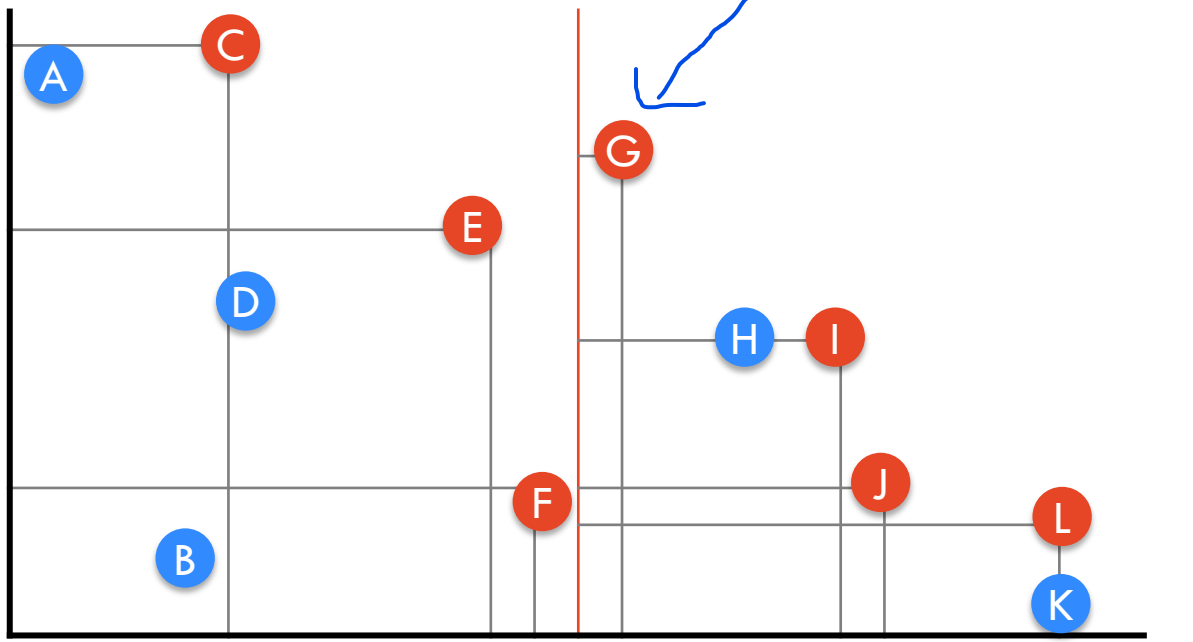
1. Find the highest point p in the Right MS

因为 left 所有x must 小于 p

$p =$ G

Observations:

1. Every point in MS of the whole is in Left MS or Right MS
2. Every point in Right MS is in MS of the whole
3. Every point in Left MS is either in MS of the whole or is dominated by p



C E F

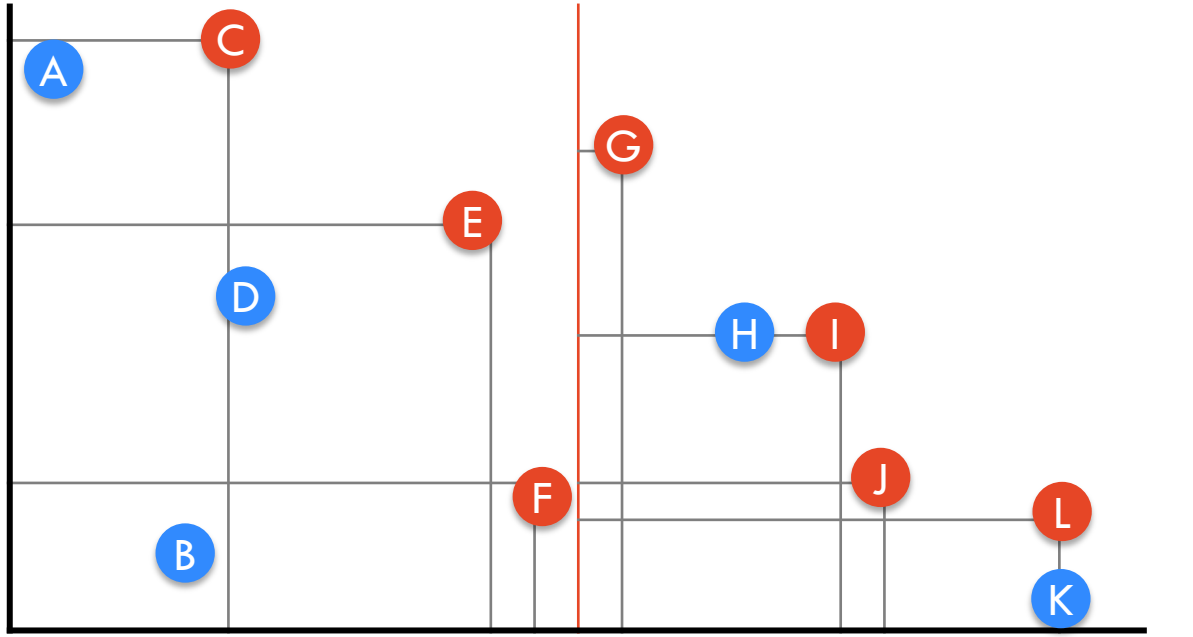
G I J L

Maxima-Set

Conquer

1. Find the highest point p in the Right MS
2. Compare every point q in the Left MS to this point.
   If $q.y > p.y$, add q to the Merged MS
3. Add every point in the Right MS to the Merged MS

$p =$ G



Merged Maxima-Set

C G I J L

Left Maxima-Set

C E F

Right Maxima-Set

G I J L

Maxima-Set

Base case a single point.

The MS of a single point is the point itself.

Maxima-Set 

Maxima-Set: Analysis

Preprocessing Sort the points by increasing x coordinate and store them in an array. Note: we only do this once. Break ties in x by sorting by increasing y coordinate. — $O(n \log n)$

Divide sorted array into two halves. — $O(n)$

Recur recursively find the MS of each half. — $2T(n/2)$

Conquer compute the MS of the union of Left and Right MS — $O(n)$

1. Find the highest point p in the Right MS
2. Compare every point q in the Left MS to this point.
   If q.y > p.y, add q to the Merged MS
3. Add every point in the Right MS to the Merged MS

$$T(n) = 2T(n/2) + O(n) = O(n \log n)$$

Overall Running Time: pre-processing + $T(n) = O(n \log n)$

Maxima-Set: Correctness

Preprocessing Sort the points by increasing x coordinate and store them in an array. Note: we only do this once. Break ties in x by sorting by increasing y coordinate.

Divide sorted array into two halves.

Recur recursively find the MS of each half.

Conquer compute MS of union of Left/Right MS

1. Find the highest point p in the Right MS
2. Compare every point q in the Left MS to this point. If $q.y > p.y$, add q to the Merged MS
3. Add every point in the Right MS to the Merged MS

Observations:

1. Every point in MS of the whole is in Left MS or Right MS
2. Every point in Right MS is in MS of the whole
3. Every point in Left MS is either in MS of the whole or is dominated by p

Integer multiplication

Given two n-digit integers x and y

Problem compute the product x y

While this seems like recreational mathematics, it does have real applications: Public key encryption is based on manipulating integers with thousands of bits.

Integer multiplication: Naïve approach

Given two n-digit integers x and y

Problem compute the product x y

Suppose we wanted to do it by hand. We assume that two digits can be multiplied or added in constant time

In primary school we all learn an algorithm for this problem that performs $\Theta(n^2)$ operations

Integer multiplication: Divide and conquer

Let $x = x_1\ 2^{n/2} + x_0$ and $y = y_1\ 2^{n/2} + y_0$

Then $x\ y = x_1\ y_1\ 2^n + x_1\ y_0\ 2^{n/2} + x_0\ y_1\ 2^{n/2} + x_0\ \ y_0$

We can compute the product of two n-digit numbers by making 4 recursive calls on $n/2$-digit numbers and then combining the solutions to the subproblems.

Integer multiplication: Divide and conquer

```
def multiply(x, y):
  // x and y are positive integers represented in binary
  if x = 0 or y = 0 then return 0
  if x = 1 then return y
  if y = 1 then return x

  // recursive case
  let x_1 and x_0 be such that x = x_1 2^(n/2) + x_0
  let y_1 and y_0 be such that y = y_1 2^(n/2) + y_0

  return multiply(x_1, y_1) 2^n +
         (multiply(x_1, y_0) + multiply(x_0, y_1)) 2^(n/2) +
         multiply(x_0, y_0)
```

Integer multiplication: Correctness

Let $x = x_1\ 2^{n/2} + x_0$ and $y = y_1\ 2^{n/2} + y_0$

Then $x\ y = x_1\ y_1\ 2^n + x_1\ y_0\ 2^{n/2} + x_0\ y_1\ 2^{n/2} + x_0\ \ y_0$

Straight forward application of induction to prove that `multiply(x, y)` = x y

Integer multiplication: Complexity analysis

Recall $x\ y = x_1\ y_1\ 2^n + (x_1\ y_0 + x_0\ y_1)\ 2^{n/2} + x_0\ \ y_0$

Divide step (produce halves) takes $O(n)$

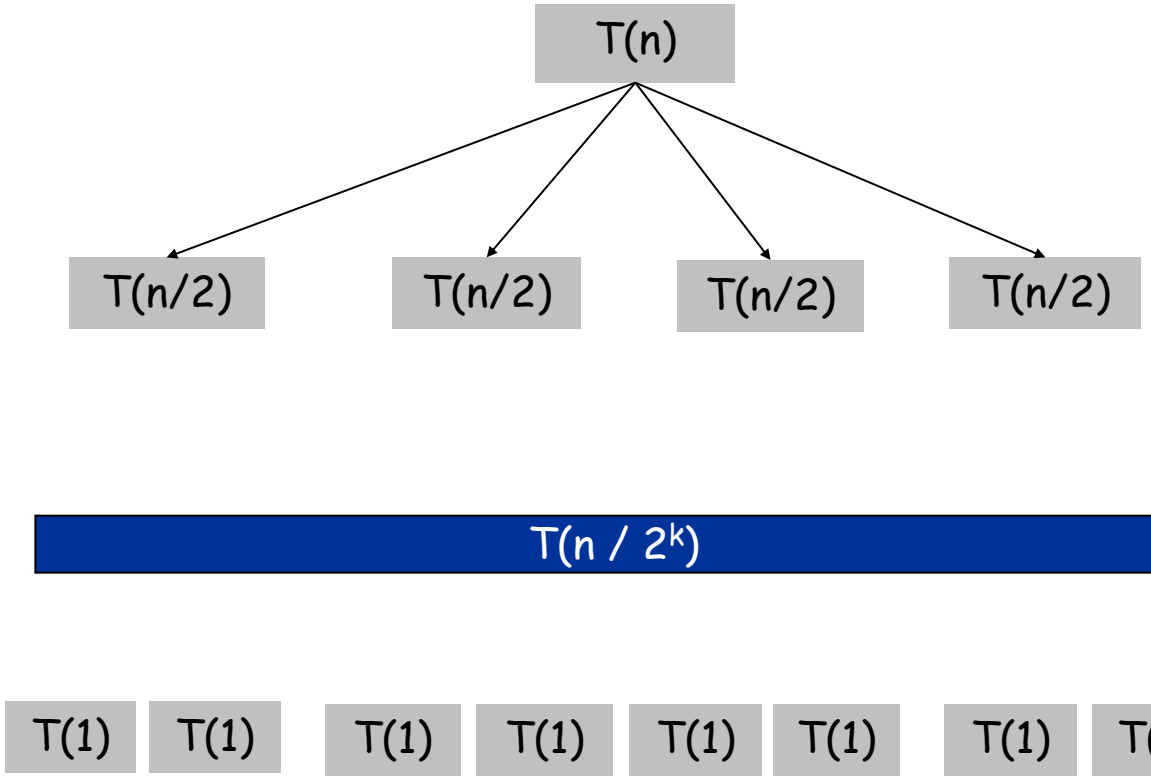
Recur step (solve subproblems) takes $4\ T(n/2)$

Conquer step (add up results) takes $O(n)$

$$T(n) = \begin{cases} 4\ T(n/2) + O(n) & \text{for } n > 1 \\ O(1) & \text{for } n = 1 \end{cases}$$

This solves to $T(n) = O(n^2)$. No better than naïve!!!

Proof by unrolling: T(n) = 4 T(n/2) + O(n)



1 (of size n)

4 (of size n/2)

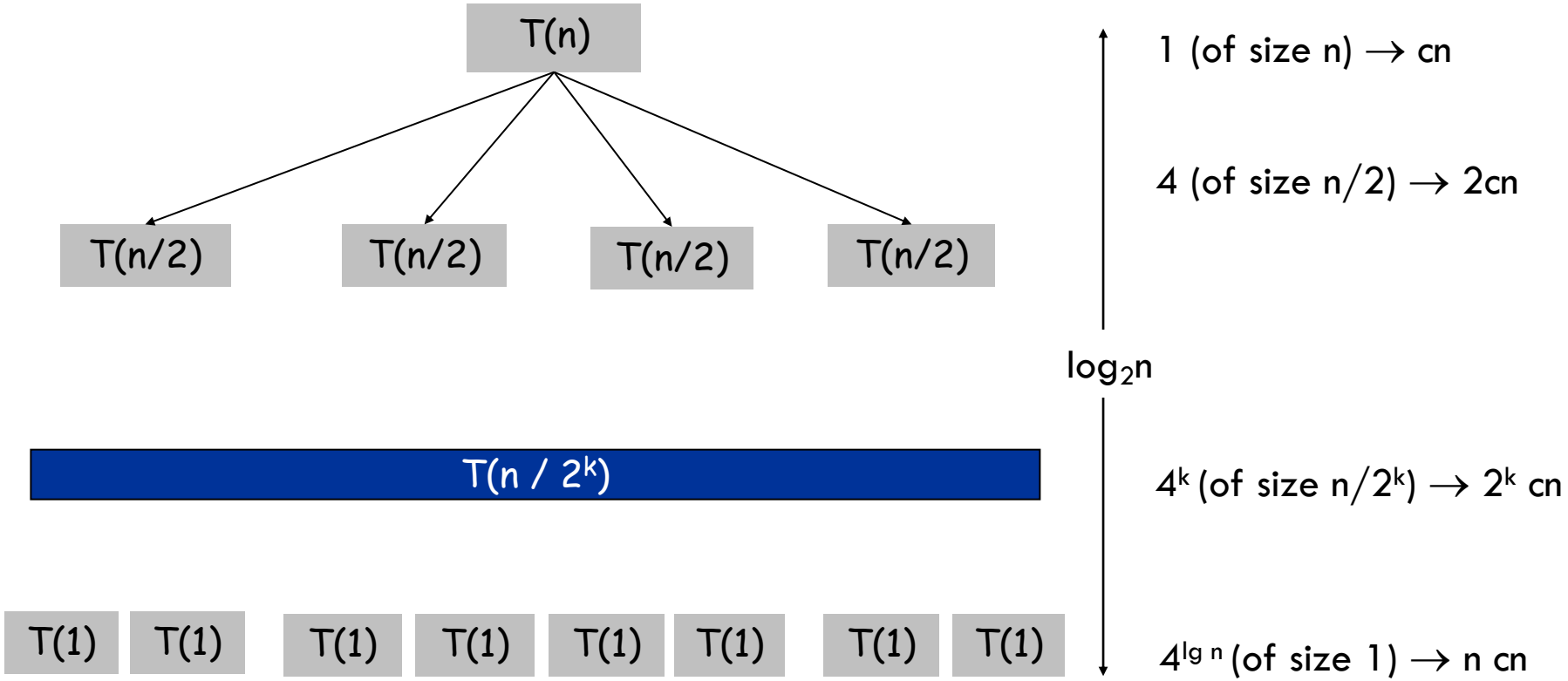
. . .

4^k (of size $n/2^k$)

. . .

$4^{\lg n}$ (of size 1)

Proof by unrolling: T(n) = 4 T(n/2) + O(n)



Integer multiplication: Divide and conquer v2.0

Let $x = x_1\ 2^{n/2} + x_0$ and $y = y_1\ 2^{n/2} + y_0$

$$x\ y = x_1\ y_1\ 2^n + (x_1\ y_0 + x_0\ y_1)\ 2^{n/2} + x_0\ y_0$$

$$(x_1 + x_0)\ (y_1 + y_0) = x_1\ y_1 + x_1\ y_0 + x_0\ y_1 + x_0\ y_0$$

We can compute the product of two n-digit numbers by making 3 recursive calls on $n/2$-digit numbers and then combining the solutions to the subproblems.

Integer multiplication: Divide and conquer

```
def multiply(x, y):
  // base case
  ⋮
  // recursive case
  let x_1 and x_0 be such that x = x_1 2^(n/2) + x_0
  let y_1 and y_0 be such that y = y_1 2^(n/2) + y_0

  first_term ← multiply(x_1, y_1)
  last_term ← multiply(x_0, y_0)
  other_term ← multiply(x_1 + x_0, y_1 + y_0)

  return first_term 2^n +
         (other_term - first_term - last_term) 2^(n/2) +
         last_term
```

Integer multiplication: Complexity analysis

Divide step (produce halves) takes $O(n)$

Recur step (solve subproblems) takes $3\ T(n/2)$

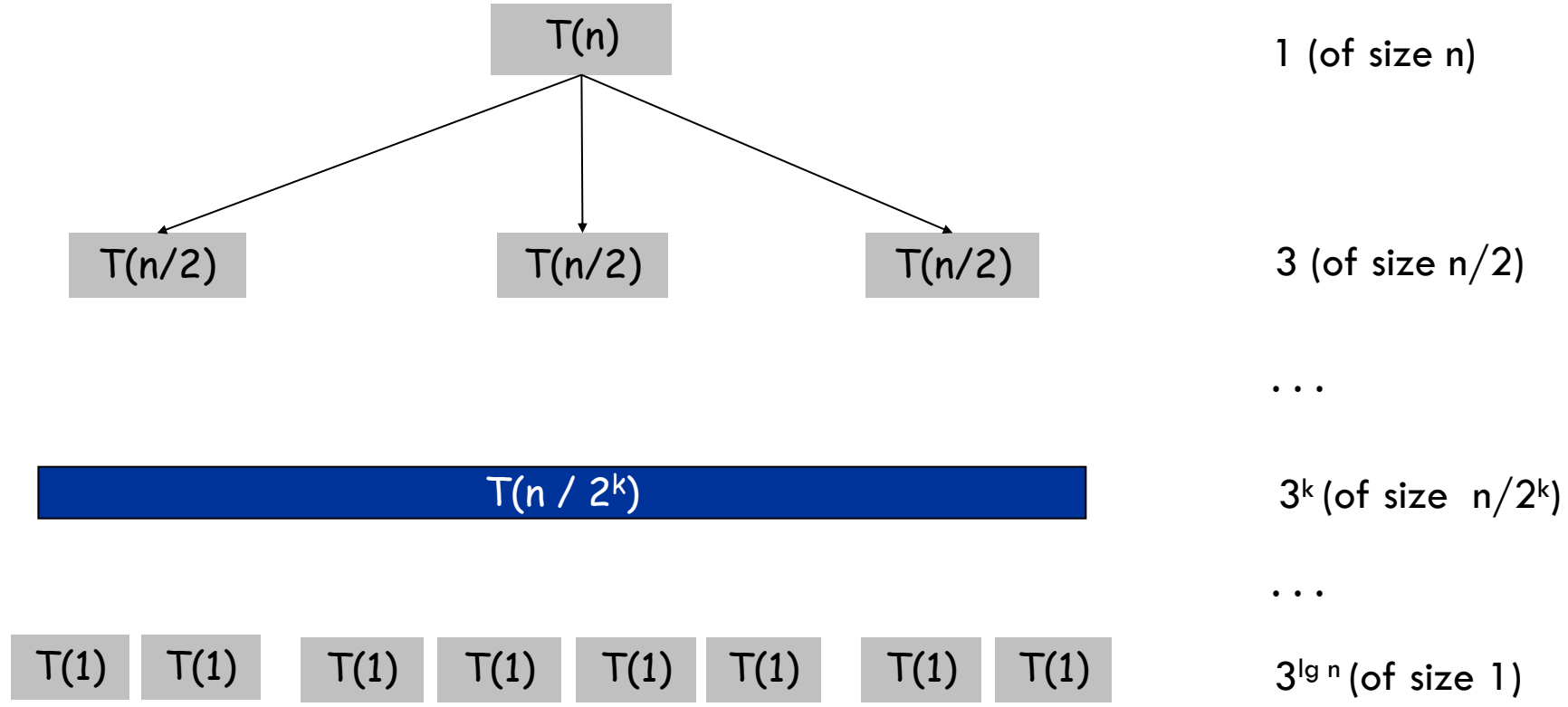
Conquer step (add up results) takes $O(n)$

$$T(n) = \begin{cases} 3\ T(n/2) + O(n) & \text{for } n > 1 \\ O(1) & \text{for } n = 1 \end{cases}$$

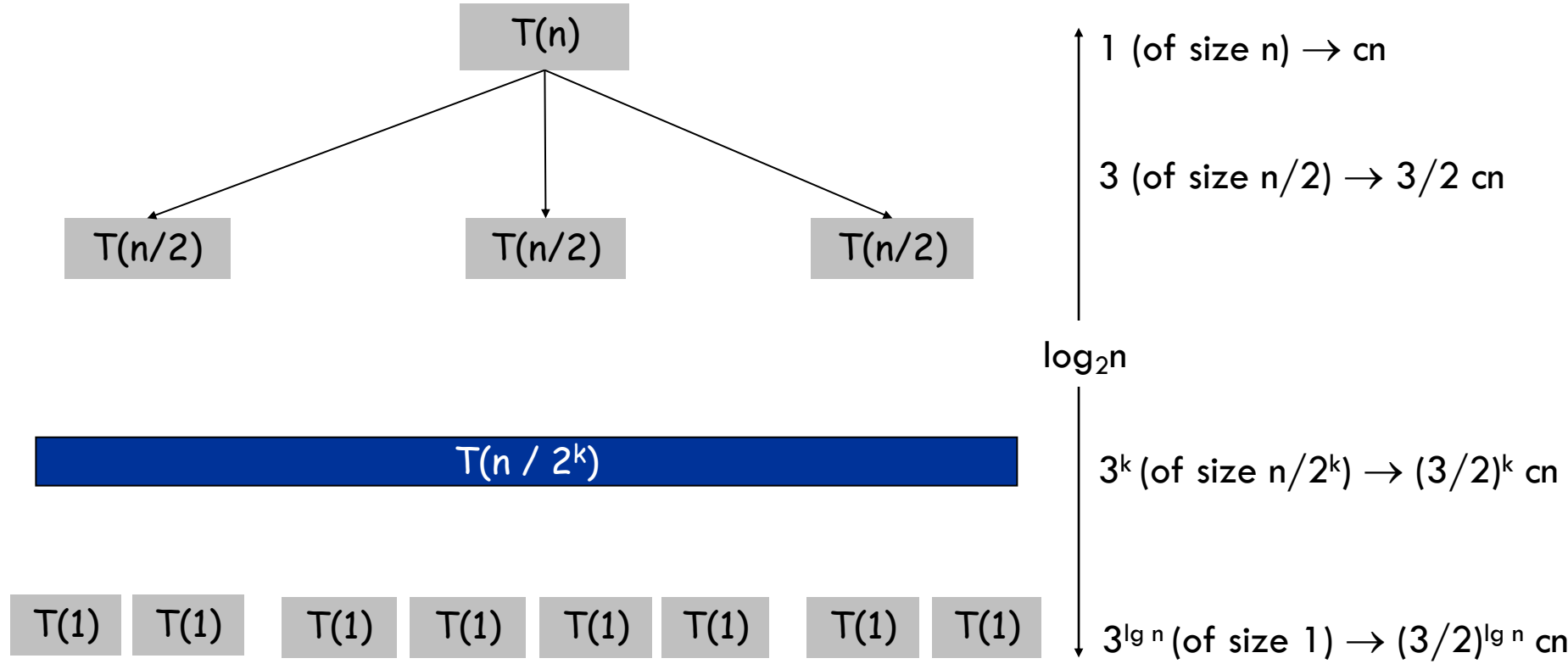
This solves to $T(n) = O(n^{\log_2 3})$, where $\log_2 3 \approx 1.6$

Better than naïve!!!

Proof by unrolling: T(n) = 3 T(n/2) + O(n)



Proof by unrolling: T(n) = 3 T(n/2) + O(n)



Geometric series facts

Let r be a positive real and k a positive integer then

$$1 + r + r^2 + \ldots + r^k = (r^{k+1} - 1)/(r-1)$$

Consequently if $r > 1$ then

$$1 + r + r^2 + \ldots + r^k < r^{k+1} / (r-1)$$

and if $r < 1$ then

$$1 + r + r^2 + \ldots + r^k < 1 / (1-r)$$

Logarithms facts

Base exchange rule:

$$\log_a x = (\log_b x)/(\log_b a)$$

Product rule:

$$\log_a (xy) = (\log_a x) + (\log_a y)$$

Power rule:

$$\log_a x^b = b \log_a x$$

Master Theorem

Let f(n) and T(n) be defined as follows:

$$T(n) = \begin{cases} a\,T(n/b) + f(n) & \text{for } n \geq d \\ c & \text{for } n < d \end{cases}$$

Depending on a, b and f(n) the recurrence solves to:

1. if $f(n) = O(n^{\log_b a - \varepsilon})$ for $\varepsilon > 0$ then $T(n) = \Theta(n^{\log_b a})$, Upper
2. if $f(n) = \Theta(n^{\log_b a} \log^k n)$ for $k \geq 0$ then $T(n) = \Theta(n^{\log_b a} \log^{k+1} n)$,
3. if $f(n) = \Omega(n^{\log_b a + \varepsilon})$ and $a\,f(n/b) \leq \delta\,f(n)$ for $\varepsilon > 0$ and $\delta < 1$ then $T(n) = \Theta(f(n))$, Lower

Note: If f(n) is given as big-O, you can only conclude T(n) as big-O (not Θ).

Note: You should be able to solve all recurrences in this class using unrolling, but if you are comfortable using the Master Theorem, go for it.

The Master Theorem

$$T(n) = \begin{cases} a\,T(n/b) + f(n) & \text{for } n \geq d \\ c & \text{for } n < d \end{cases}$$

Examples

1. $T(n) = 8T(n/2) + n^2$

   $a=8, b=2, f(n)=n^2, \log_b(a) \rightarrow \log_2(8) = 3$; so $f(n) = O(n^{\log_b a - \varepsilon})$ (case 1)

   $T(n) \in \Theta(n^3)$

2. $T(n) = 2T(n/2) + O(n)$

   $a=2, b=2, f(n)=n, \log_b(a) \rightarrow \log_2(2) = 1$; so $f(n) = \Theta(n^{\log_b a} \log^k n)$ (case 2 with $k=0$)

   $T(n) \in O(n \log(n))$

3. $T(n) = 2T(n/2) + O(n^2)$

   $a=2, b=2, f(n)=n^2, \log_b(a) \rightarrow \log_2(2) = 1$; so $f(n) = \Omega(n^{\log a + \varepsilon})$ (case 3)

   $T(n) \in O(n^2)$

Selection

Given an unsorted array A holding n numbers and an integer k, find the kth smallest number in A

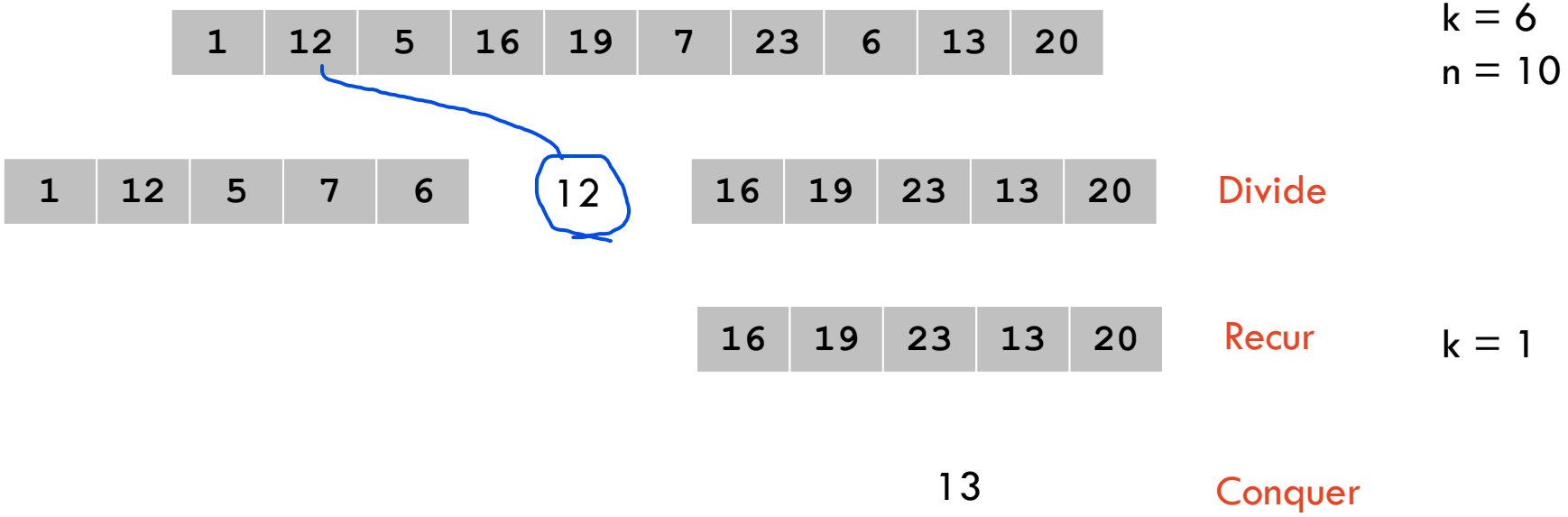
Trivial solution: Sort the elements and return kth element

Can we do better than O(n log n)?

Yes, with divide and conquer!

First attempt

1. Divide find the median ($\lfloor n/2 \rfloor$th element for simplicity) and split array on the halves, $\leq$ and $>$ than the median
2. Recur if k $\leq \lfloor n/2 \rfloor$ find kth element on smaller half
   if k $> \lfloor n/2 \rfloor$ find (k-$\lfloor n/2 \rfloor$)th element on larger half
3. Conquer return value of the recursive call



Selection time complexity

Divide step (find median and split) takes at least O(n)

Recur step (solve left or right subproblem) takes T(n/2)

Conquer step (return recursive result) takes O(1)

If we could compute the median in O(n) time then:

$$T(n) = \begin{cases} T(n/2) + O(n) & \text{for } n > 1 \\ O(1) & \text{for } n = 1 \end{cases}$$

This solves to T(n) = O(n) but only if we can solve the median problem, which is in fact a special case of selection with k=$\lfloor n/2 \rfloor$

Second attempt: Approximating the median

We don't need the exact median. Suppose we could find in $O(n)$ time an element x in A such that

Array

$$|A| / 3 \leq \text{rank}(A, x) \leq 2\,|A| / 3$$

任意在这个区间的都行

"median"

Then we get the recurrence

$$T(n) = \begin{cases} T(2n/3) + O(n) & \text{for } n > 1 \\ O(1) & \text{for } n = 1 \end{cases}$$

Which again solves to $T(n) = O(n)$

To approximate the median we can use a recursive call!

Median of 3 medians

Consider the following procedure

- Partition A into $|A| / 3$ groups of 3
- For each group find the median (brute force)
- Let x be the median of the medians (computed recursively)

We claim that x has the desired property

$$|A| / 3 \leq \text{rank}(A, x) \leq 2|A| / 3$$

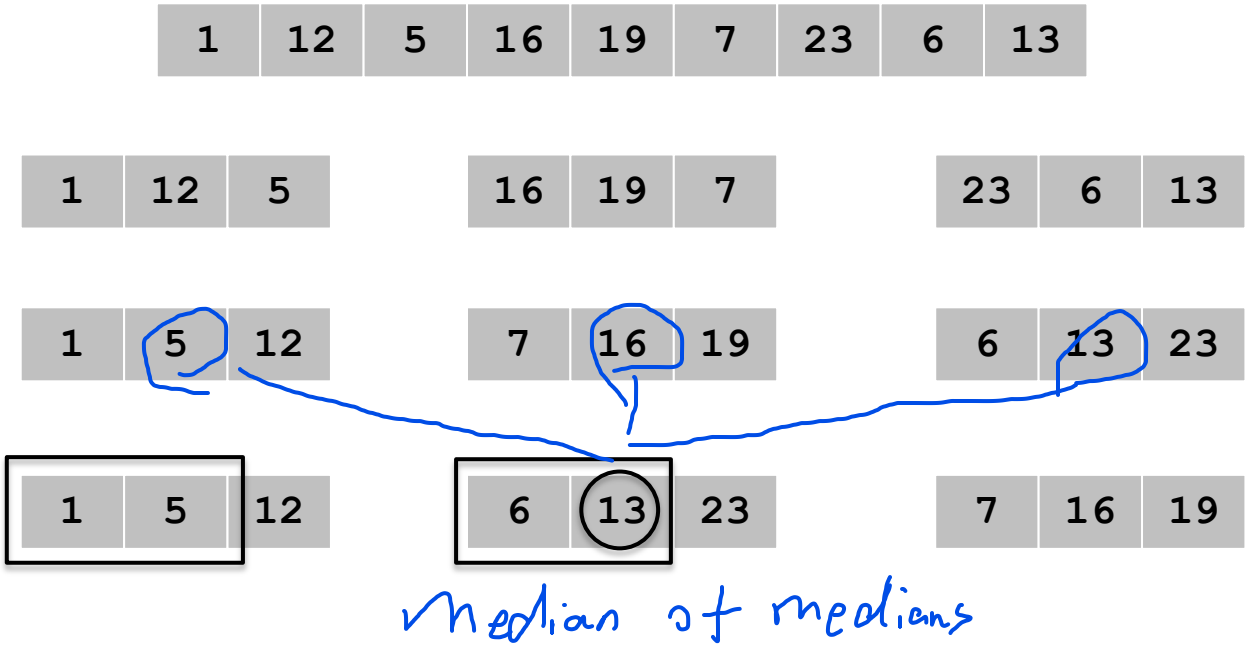
Half of the groups have a median that is smaller/larger than x, and each group has two elements smaller/larger than x, thus

$$\#\text{ elements smaller than } x > 2\,(|A| / 6) = |A| / 3$$
$$\#\text{ elements greater than } x > 2\,(|A| / 6) = |A| / 3$$

Median of 3 medians

Let x be the median of the medians, then

$$|A| / 3 \leq \text{rank}(A, x) \leq 2|A| / 3$$



elements smaller than x > 2 (|A| / 6) = |A| / 3

elements greater than x > 2 (|A| / 6) = |A| / 3

Median of 3 median time complexity

We don't need the exact median. With a recursive call on $n/3$ elements, we can find x in A such that

$$|A|/3 < \text{rank}(A, x) < 2|A|/3$$

Then we get the recurrence

$$T(n) = T(2\,n\,/\,3) + T(n\,/\,3) + O(n)$$

Which solves to $T(n) = O(n \log n)$

No better than sorting!

Median of 5 medians

We don't need the exact median. With a recursive call on $n/5$ elements, we can find x in A such that

$$3|A|/10 < \text{rank}(A, x) < 7|A|/10$$

Then we get the recurrence

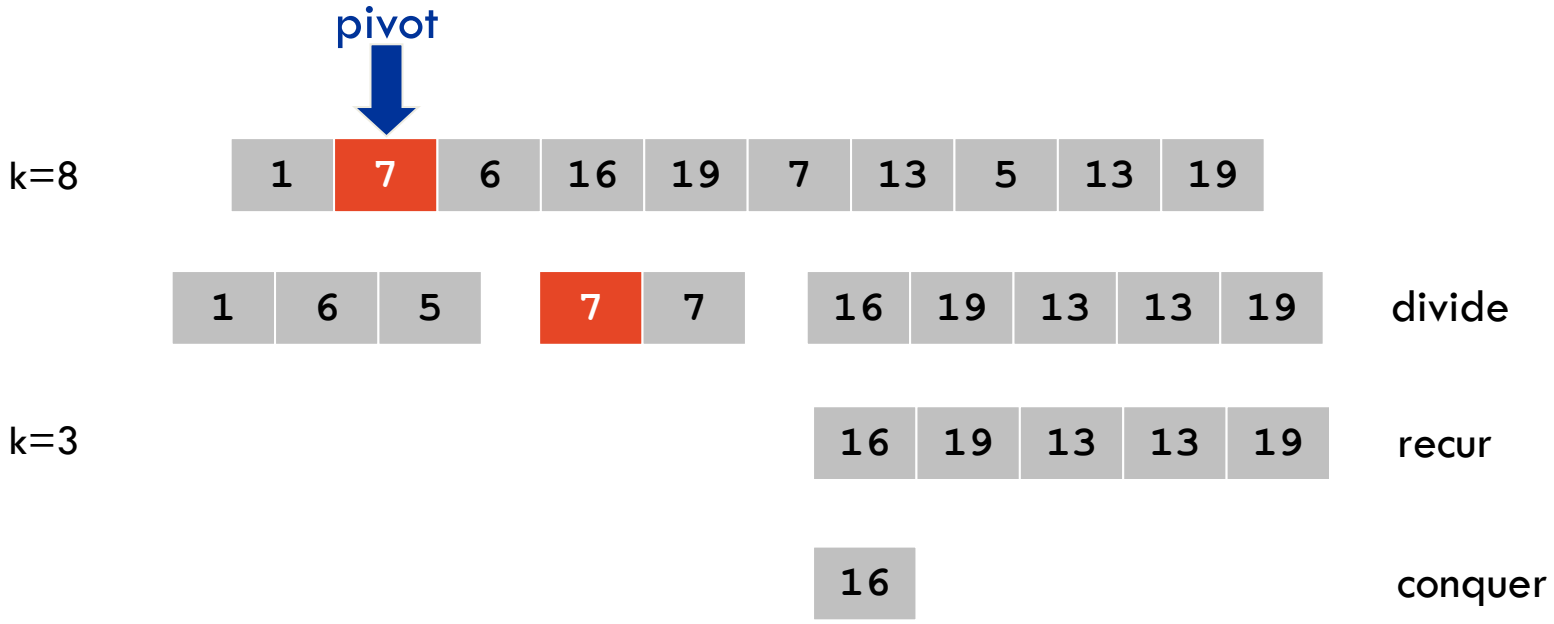
$$T(n) = T(7\,n\,/\,10) + T(n\,/\,5) + O(n)$$

Which solves to $T(n) = O(n)$

Asymptotically faster than sorting!

Quick selection

1. Divide Choose a random element from the list as the pivot
   Partition the elements into 3 lists:
   (i) less than, (ii) equal to and (iii) greater than the pivot
2. Recur Recursively select right element from correct list
3. Conquer Return solution to recursive problem



Quick selection complexity analysis

Divide step (pick pivot and split) takes O(n)

Recur step (solve left and right subproblem) takes T(n')

Conquer step (return solution) takes O(1)

Now we can set up the recurrence for T(n):

$$E[T(n)] = \begin{cases} E[T(n')] + O(n) & \text{for } n > 1 \\ O(1) & \text{for } n = 1 \end{cases}$$

This solves to $E[T(n)] = O(n)$

(details available on the textbook but not examinable)