Assignment 5

Submittors: Moran Neptune 208474544 and Bar Perlman 305026882

Part 1: Theoretical Questions:

```
<u>2. C.</u>
?- female(Y), relatives(joseph, Y)
1. s = \{\}
              Eqs = [mn(ttt,M,mn,V,HO,mn,br(M)) = mn(ttt,br(HO),mn,br(M),b)
                    r(V),mn,br(M)
           2. s = \{ \}
              Eqs = [ttt=ttt, M=br(HO), mn=mn, V=br(M), HO=br(V), mn=mn,
                     br(M)=br(M)
           3. s = \{\} \cdot \{M = br(HO)\} = \{M = br(HO)\}
              Eqs = [V=br(br(HO)), HO=br(V)]
           4. s = \{M = br(HO)\} \cdot \{V = br(br(HO))\}
              Eqs = [HO=br(br(br(HO)))]
הינה מעגלית. HO=br(br(br(HO))) הינה מעגלית. אפשרית, הקשירות
\underline{B}. Unify[\mathbf{m}(M,N), \mathbf{n}(M,N)]
                              -מכיוון שלא מדובר באף בדיקה מהאלגוריתם (משום שלא מדובר ב
משתנה וקבוע, או בשני המשוואות יש אטומיים שווים, ומשום שהסמל של שניהם שונה
(m < --> n)
אזי נחזיר שגיאה.
C. Unify[kmr(pt(pt),V,pt,g,kmr(TF),pt), kmr(pt(V),V,pt,g,kmr(TF),HO)]
           1. s = \{\}
              Eqs = [kmr(pt(pt), V, pt, g, kmr(TF), pt) = kmr(pt(V), V, pt, g, kmr(TF), HO)]
           2. s = \{ \}
              Eqs = [pt(pt)=pt(V), V=V, pt=pt, kmr(TF)=kmr(TF), pt=HO]
           3. s = \{\} \cdot \{pt(pt) = pt(V)\} = \{pt(pt) = pt(V)\} = \{pt = V\}
              Eqs = [V=HO]
           4. s = \{pt=V\} \cdot \{V=HO\} = \{pt=HO\}
              Eqs = []
           5. s = \{pt=HO\}
\underline{D}. Unify[\mathbf{mn}([T]), \mathbf{mn}(T)]
           1. s = \{ \}
              Eqs = [mn([T]) = mn(T)]
           2. s = \{ \}
              Eqs = [[T]=T]
           לכן, ההצבה לא אפשרית, הקשירות [T] הינה מעגלית.
```

 \underline{E} . Unify[g([br | [KK | K]]), g([[br | KK] | K])]

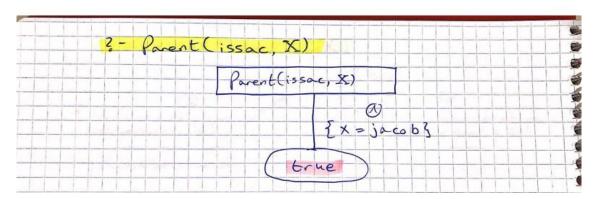
1.
$$s = \{\}$$

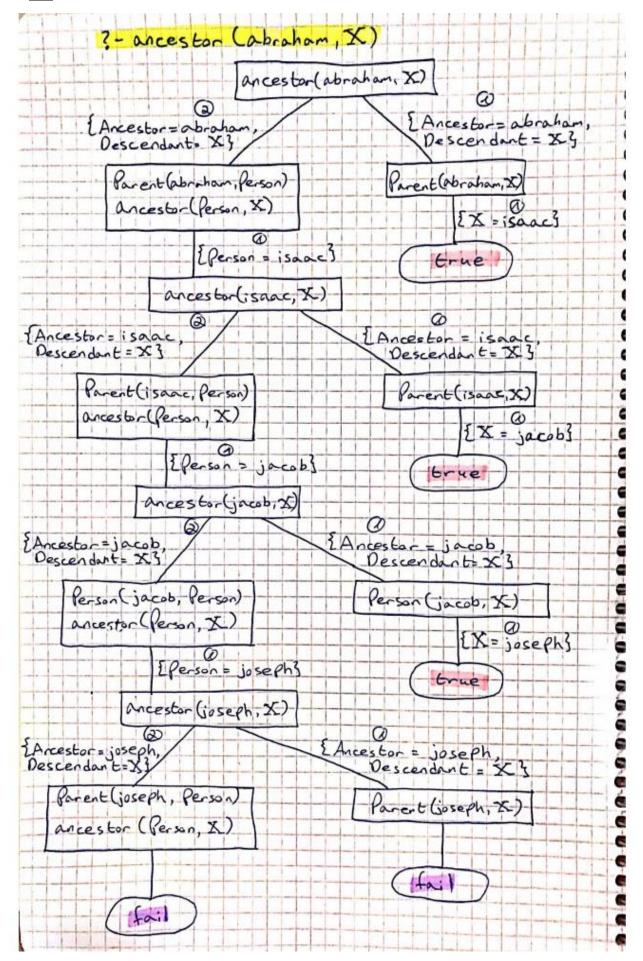
Eqs = [g([br | [KK | K]]) = g([[br | KK] | K])]
2. $s = \{\}$
Eqs = [[br | [KK | K]] = [[br | KK] | K]]
3. $s = \{\} \cdot \{br = [br | KK]\}$
Eqs = []KK|k]=K]

. הינה מעגלית br=[br|KK] הינה אפשרית, אפשרית, לכן, ההצבה לא

<u>.4</u>

:A.1





```
\underline{B}: 1. X = Jacob
```

X = Isaac.2

X = Jacob

X = Joseph

<u>C:</u> These trees are finite, because they are not infinite - there is not an infinite route, no route is from the version of: p(X) := p(Y), q(X, Y)

<u>D:</u> These trees are success proof trees. Because, there is at least one success route in .each tree

<u>A5..</u>

CompoundC --> regularFormC

'regularFormC --> '(' StringC SC* ')

'lambdaC --> '(' 'lambda' '(' StringC* ')' '(' SC* ')' '('SC*')' ')

NumberC --> < Number>

BooleanC --> #t | #f

StringC --> <String>