# BOĞAZİÇİ UNIVERSITY

# CMPE 321 DATABASE SYSTEMS

# STORAGE MANAGER IMPLEMENTATION PROJECT 2

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## 1 Introduction

In the first project of the course, we are asked to design a storage manager which will be implemented later in the following projects. In this project, the storage manager implementation has been carried out. A storage manager is a system which is responsible from the definition and the manipulation of the data and the data storage inside the memory. The project report consists of the assumptions and constraints for the storage manager, the properties and the relations of the storage structures, and the DDL & DML operations that are supported by the storage manager.

## 2 Assumptions & Constraints

## 2.1 Assumptions

#### 2.1.1 Base Assumptions

This subsection contains the base assumptions that are already mentioned in the project description.

- Page Size = 1024 Bytes
- File Size = 20488 Bytes
- The content of the File Header:
  - # of Total Pages
  - # of Active Pages
- The content of the Page Header:
  - Boolean: isActive -> States if the page is active/deleted.
  - # of Total Records
  - # of Active Records
- The content of the Record Header:
  - Boolean: isActive -> States if the record is active/deleted.
  - # of Fields
  - Primary Key (Record ID)
- Maximum number of fields a type can have = 10
- Maximum length of a type name = 10 characters
- Maximum length of a field name = 10 characters
- User always enters valid input.
- All fields shall be integers. However, type and field names shall be alphanumeric.
- A disk manager already exists that is able to fetch the necessary pages when addressed.

#### 2.1.2 Extra Assumptions

This subsection contains the additional assumptions that are determined by the assignee of the project.

• The contents of the fields are no larger than 4 Bytes.

### 2.1.3 Added Assumptions

This subsection contains the assumptions added during the second project.

• The system catalogue is a single page which has a dynamical size.

## 2.2 Constraints

#### 2.2.1 Base Constraints

This subsection contains the base constraints that are already mentioned in the project description.

- The data must be organized in the pages and pages must contain records.
- Storing all pages in the same file is not allowed. Plus, a file must contain multiple pages.
- Although a file contains multiple pages, it must read page by page when it is deleted. Loading the whole file to RAM is not allowed.

#### 2.2.2 Extra Constraints

This subsection contains the additional constraints that are determined by the assignee of the project.

- Two records with the same primary key values cannot be created.
- A file contains the records of only one type.
- Every record of a type has a unique primary key (record id).
- Every file has a unique file name.

• If the user creates a new type, the file name will be in the form of "typeName1.txt". In case of the lack of free space in the file holding the type, the new file that will contain new records of the same type will be named in the form of "typeName2.txt" and this procedure will apply for the next file creations of the same type.

#### 2.2.3 Added Constraints

This subsection contains the constraints added during the second project.

• The sizes of the data storage units are fixed (system catalogue excluded).

## 3 Storage Structures

The system has is composed of two primary storage structures that are the system catalogue which contains the metadata about the actual data, and the data storage structures that contains the actual data.

## 3.1 System Catalogue

System catalogue is a file that stores metadata, data about the actual data. The abstraction of the database has been carried out by the system catalogue. The user must first acces into this file before making an operation in the database. The system catalogue consists of only one page which holds the records consisting of the abstraction of the data within the files.

• System Catalogue Header

− # of Files: 4 Bytes

4 Bytes

# of Files

• System Catalogue Content

The system catalogue consists of a single page which holds the records corresponding to the abstraction of each prospected files to be created in case of necessity. The content of each record is as follows:

- Boolean: isActive : 1 Byte

-> States if the file is active/deleted.

File Name: 16 BytesType Names: 10 Bytes

- # of Fields of The Records of The Corresponding File: 4 Bytes

31 Bytes Per Record (+4 Bytes due to the file processing device used)

isActive	File Name	Type Name	# of Fields
boolA	File A	TypeA	$\#\mathrm{A}$
boolB	FileB	TypeB	#B
boolC	FileC	TypeC	#C

Table 1: The Structure of The System Catalogue

## 3.2 Data Storage Units

- 1. Files: 20488 Bytes
  - File Header: 8 Bytes
    - # of Total Pages: 4 Bytes # of Active Pages: 4 Bytes

# of Total Pages | # of Active Pages

• File Content: A file contains maximum 20 pages.

20x1024 = 20480 Bytes per File

	Page #1
	Page $\#2$
ſ	Page #20

- 2. Pages: 1024 Bytes
  - Page Header: 9 Bytes
    - Boolean isActive: 1 Byte
    - # of Total Records: 4 Bytes
    - − # of Active Records: 4 Bytes

is Active  $\mid \#$  of Total Records  $\mid \#$  of Active Records

• Page Content: 1015 Bytes

Record #1
Record #2
...
Record #N

- 3. Records: 9 + (# of Fields (Primary Key Excluded) x 4) Bytes
  - $\bullet\,$  Record Header: 9 Bytes
    - Boolean isActive: 1 Byte

- # of Fields: 4 Bytes
- Primary Key (Record ID): 4 Bytes

• Record Content: A record contains maximum 10 fields. (Primary Key field is considered as the part of the record header)

( of Fields) x 4 Bytes per Record

Field#1 | Field#2 | ... | Field#10

## 4 Operations

## 4.1 DDL Operations

DDL stands for the abbreviation of Data Definition Language which consists of the operations that are used to make creations, deletions, modifications, etc. on the database schema.

### 4.1.1 Create a Type

## Algorithm 1 Create a Type

```
declare newType //initialize the fields of the newType newType.typeName \leftarrow input newType.primaryKey \leftarrow input newType.numOfFields \leftarrow input for i=1 to newType.numOfFields do | newType.field#i \leftarrow input end systemCatalogue.append(newType) systemCatalogue.numOfFiles++
```

## 4.1.2 Delete a Type

```
Algorithm 2 Delete a Type

typeName ← input
listOfIndices ← getList(typeName)

if listOfIndices.size() == 0 then

| System.out.println("Type Not Found!")
| break

end

else

| for i = listOfIndices.size() to 1 do
| index ← listOfIndices[i-1]
| systemCatalogue.types.remove(index)
| systemCatalogue.header.numOfFiles—
| systemCatalogue.files.delete(typeName#i)
| end
| end
```

#### 4.1.3 List All Types

### Algorithm 3 List All Types

```
if systemCatalogue.header.numOfFiles == 0 then

System.out.println("No Types Exist!")

break
end
else

numOfFiles \leftarrow systemCatalogue.header.numOfFiles

for i = 1 to numOfFiles do

curLine \leftarrow "systemCatalogue.txt".nextLine

print curLine[1]

end
end
```

## 4.2 DML Operations

DML stands for the abbreviation of Data Manipulation Language which consists of the operations that are used to make creations, deletions, modifications, etc. on the data itself.

#### 4.2.1 Create a Record

```
Algorithm 4 Create a Record
declare newRecord
newRecord.content.typeName \leftarrow input
newRecord.content.numOfFields \( -\system Catalogue.get \) NumOfFields \( (type Name) \)
for i = 1 to newRecord.numOfFields do
| newRecord.field#i \leftarrow input
end
if systemCatalogue.header.numOfFiles == 0 then
   newFile ← new File("newRecord.content.typeName#i.txt")
   newPage \leftarrow new Page(), newPage.records.add(newRecord)
   newPage.numOfRecords++
   newFile.pages.add(newPage), newFile.numOfPages++
   systemCatalogue.files.add(newFile),systemCatalogue.numOfFiles++
end
else
   curFile \leftarrow systemCatalogue.header.firstFile
   while curFile.nextFile != NULL do
       if curFile.getType().equals(newRecord.typeName)&&curFile.numOfFreeBytes
       != \theta then
          foreach curPage in curFile do
             foreach curRecord in curPage do
                 if curRecord == NULL then
                    curRecord \leftarrow newRecord, curPage.numOfRecords++
                    return
                 end
             end
          end
       end
   end
   if systemCatalogue.numOfFiles < 20 then
       newFile \leftarrow new File("newRecord.content.typeName#i.txt")
       newPage \leftarrow new Page(), newPage.records.add(newRecord)
       newPage.numOfRecords++
      newFile.pages.add(newPage), newFile.numOfPages++
      systemCatalogue.files.add(newFile),systemCatalogue.numOfFiles++
   end
   else
      System.out.println("Not Enough Space!")
   end
\quad \text{end} \quad
```

#### 4.2.2 Delete a Record

## Algorithm 5 Delete a Record $primaryKey \leftarrow input$ $typeName \leftarrow input$ if systemCatalogue.header.numOfFiles == 0 then System.out.println("No Records Exist!"), break end else $curFile \leftarrow systemCatalogue.header.firstFile$ while true do while !curFile.typeName.equals(typeName) do if curFile.nextFile != NULL then $curFile \leftarrow curFile.nextFile$ end else return end end foreach curPage in curFile do foreach curRecord in curPage do if curRecord.primaryKey == primaryKey then curPage.numOfRecords curRecord = NULLif curPage.numOfRecords == 0 then curFile.pages.remove(curPage) curFile.numOfPages-if curFile.numOfPages == 0 then systemCatalogue.files.remove(curFile) systemCatalogue.header.numOfFiles end end return end $\mathbf{end}$ $curFile \leftarrow curFile.nextFile$ end end

#### 4.2.3 Search for a Record

```
Algorithm 6 Search a Record (by Primary Key)
primaryKey \leftarrow input
typeName \leftarrow systemCatalogue.getTypeName(primaryKey)
curFile \leftarrow systemCatalogue.header.firstFile
while true do
   \mathbf{while} \ ! curFile.typeName.equals(typeName) \ \mathbf{do}
       {f if} \ curFile.nextFile != NULL \ {f then}
           curFile \leftarrow curFile.nextFile
       end
       else
       return
       end
   end
   foreach curPage in curFile do
       foreach curRecord in curPage do
           if curRecord.primaryKey == primaryKey then
            return curRecord
           end
       end
   end
   curFile \leftarrow curFile.nextFile
end
```

#### 4.2.4 Update a Record

```
Algorithm 7 Update a Record (by Primary Key)
primaryKey \leftarrow input
isFoun \leftarrow false if systemCatalogue.header.numOfFiles == 0 then
   System.out.println("No records exist!") return
end
else
    while isFound = = false do
       while !curFile.typeName.equals(typeName) do
           if curFile.nextFile != NULL then
            \vdash curFile \leftarrow curFile.nextFile
           end
           else
            ⊢ return
           end
       end
       foreach curPage in curFile do
           foreach curRecord in curPage do
               if curRecord.primaryKey == primaryKey then
                   updatedRecord \leftarrow curRecord
                   isFound \leftarrow true
               end
           \mathbf{end}
       \quad \text{end} \quad
       if isFound == false then
        | curFile ← curFile.nextFile continue
       end
   \quad \text{end} \quad
    numOfFields \leftarrow updatedRecord.numOfFields
   for i=1 to numOfFields do
       updatedRecord.field\#i \leftarrow input
    end
\quad \text{end} \quad
```

## 4.2.5 List All Records of A Type

## Algorithm 8 List All Records of A Type

```
\overline{\text{typeName} \leftarrow \text{input}}
if systemCatalogue.header.numOfFiles == 0 then
   System.out.println("No Records Exist!")
end
else
    curFile \leftarrow systemCatalogue.header.firstFile
    while curFile.nextFile != NULL do
       {f if}\ curFile.typeName.equals(typeName)\ {f then}
           foreach curPage in curFile do
               foreach curRecord in curPage do
                  printRecord(curRecord)
               end
           end
       end
   end
end
```

### 5 Conclusion & Assessment

- A well designed storage manager allows a reliable database storage, efficient operations, and secure access into the database. In this project, the design of the storage manager has been converted into the implementation.
- The content of the storage structures had been designed in order to enable a better view before implementation. By the time the implementation has been proceeded, some unnecessary fields of the headers or uses have been detected and removed out of the project.
- For further information about the headers, some fields might have been added such as remaining bytes, timestamp fields, etc. ,but in order to provide a less complex implementation the fields inside the headers have been simplified as much as possible.
- The files containing a single type of record allow the administrators to easily handle the storage of the data by providing a simple and efficient abstraction of data.

PS: Due to the restrictions caused by the use of algorithm block in LaTeX, a blank page has been automatically generated when the page size has not been enough for the algorithm block.

## 5.1 Changes

- 1. File header has been updated.
- 2. File size has been updated.
- 3. Page header has been updated.
- 4. Page size has been updated into a fixed size of 1024 Bytes.
- 5. Record header has been updated.
- 6. Added assumptions constraints have been denoted in the corresponding sections
- 7. System catalogue header has been updated. System catalogue has been designed to be a single page of dynamical size.

 $8. \ \,$  Create type operation pseudo-code has been updated.