

# Indexes as Access Paths

- A single-level index is an auxiliary file that makes it more efficient to search for a record in the data file.
- The index is usually specified on one field of the file (although it could be specified on several fields)
- One form of an index is a file of entries <**field value, pointer to record**>, which is ordered by field value
- The index is called an access path on the field.

# Indexes as Access Paths (contd.)

- The index file usually occupies considerably less disk blocks than the data file because its entries are much smaller
- A binary search on the index yields a pointer to the file record
- Indexes can also be characterized as dense or sparse
  - A **dense index** has an index entry for every search key value (and hence every record) in the data file.
  - A **sparse (or nondense) index**, on the other hand, has index entries for only some of the search values

# Indexes as Access Paths (contd.)

- Example: Given the following data file EMPLOYEE(NAME, SSN, ADDRESS, JOB, SAL, ... )
- Suppose that:
  - record size  $R=150$  bytes                      block size  $B=512$  bytes     $r=30000$  records
- Then, we get:
  - blocking factor  $Bfr = B \div R = 512 \div 150 = 3$  records/block
  - number of file blocks  $b = (r/Bfr) = (30000/3) = 10000$  blocks
- For an index on the SSN field, assume the field size  $V_{SSN}=9$  bytes, assume the record pointer size  $P_R=7$  bytes. Then:
  - index entry size  $R_I = (V_{SSN} + P_R) = (9+7) = 16$  bytes
  - index blocking factor  $Bfr_I = B \div R_I = 512 \div 16 = 32$  entries/block
  - number of index blocks  $b = (r/Bfr_I) = (30000/32) = 938$  blocks
  - binary search needs  $\log_2 b = \log_2 938 = 10$  block accesses
  - This is compared to an average linear search cost of:
    - $(b/2) = 30000/2 = 15000$  block accesses
  - If the file records are ordered, the binary search cost would be:
    - $\log_2 b = \log_2 30000 = 15$  block accesses

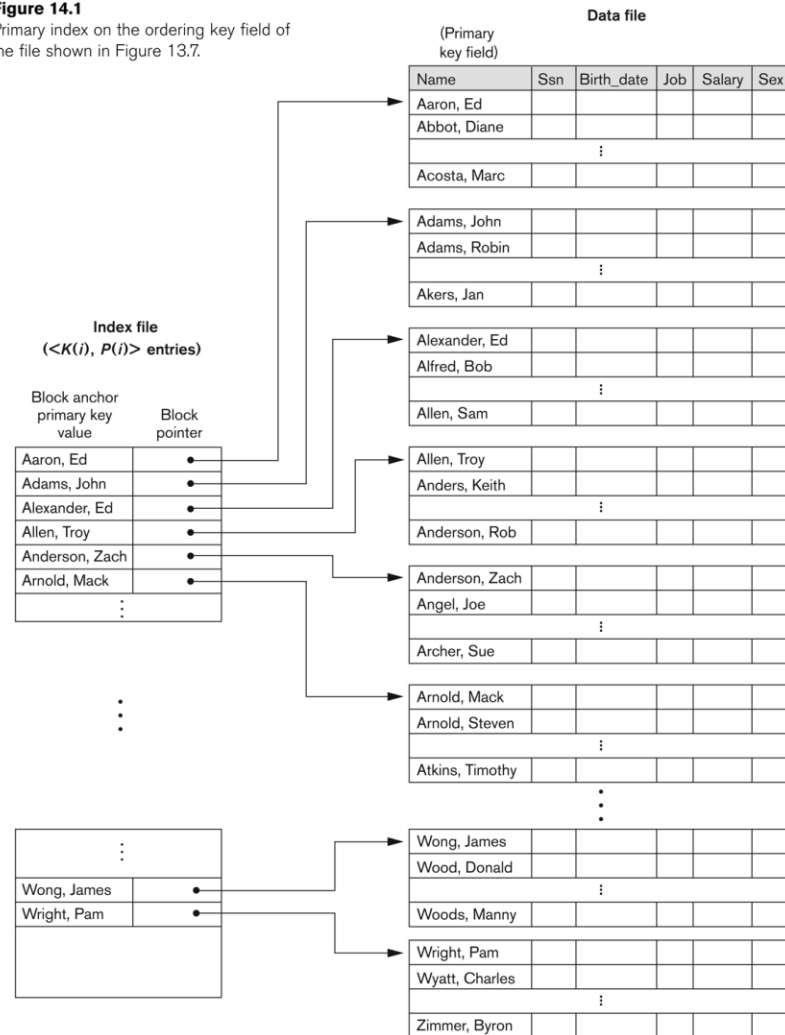
# Types of Single-Level Indexes

- Primary Index
  - Defined on an ordered data file
  - The data file is ordered on a **key field**
  - Includes one index entry *for each block* in the data file; the index entry has the key field value for the *first record* in the block, which is called the *block anchor*
  - A similar scheme can use the *last record* in a block.
  - A primary index is a nondense (sparse) index, since it includes an entry for each disk block of the data file and the keys of its anchor record rather than for every search value.

# Primary index on the ordering key field

**Figure 14.1**

Primary index on the ordering key field of the file shown in Figure 13.7.

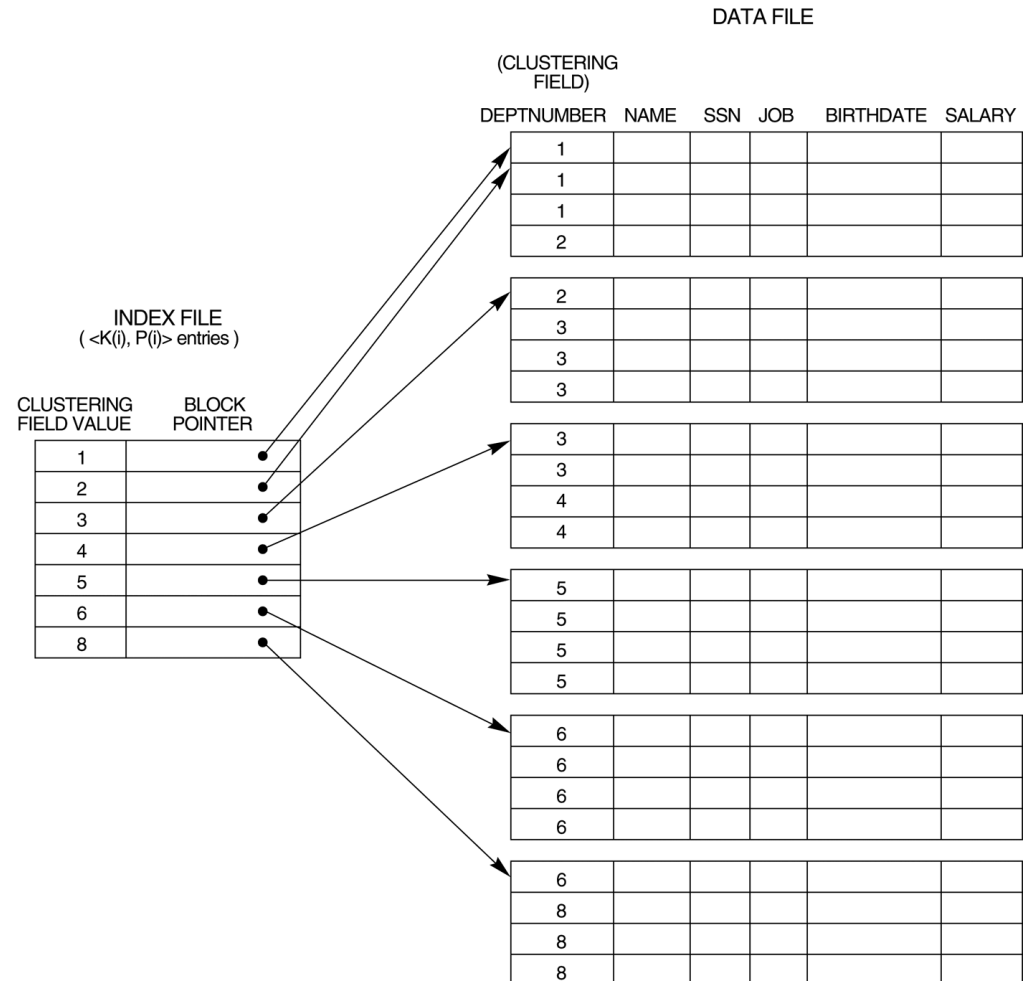


# Types of Single-Level Indexes

- Clustering Index
  - Defined on an ordered data file
  - The data file is ordered on a *non-key field* unlike primary index, which requires that the ordering field of the data file have a distinct value for each record.
  - Includes one index entry *for each distinct value* of the field; the index entry points to the first data block that contains records with that field value.
  - It is another example of *nondense* index where Insertion and Deletion is relatively straightforward with a clustering index.

# A Clustering Index Example

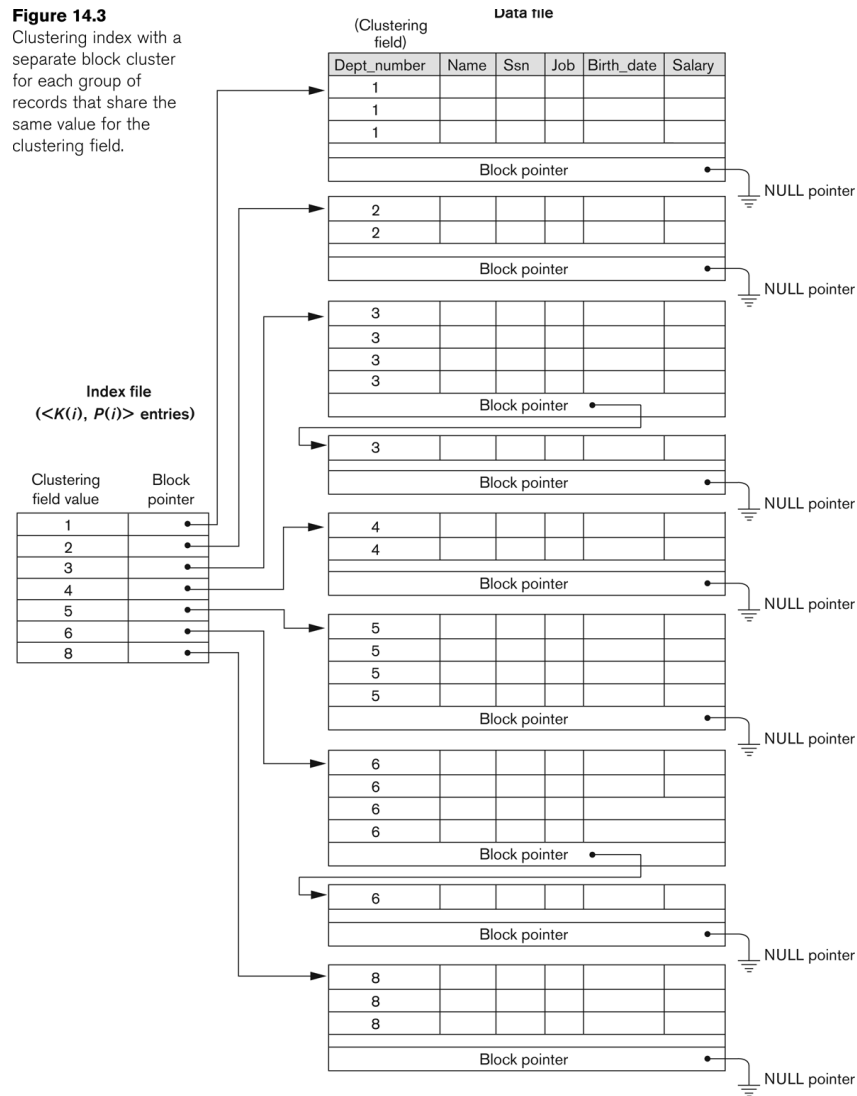
- **FIGURE 14.2**  
A clustering index on the DEPTNUMBER ordering non-key field of an EMPLOYEE file.



# Another Clustering Index Example

**Figure 14.3**

Clustering index with a separate block cluster for each group of records that share the same value for the clustering field.





# Types of Single-Level Indexes

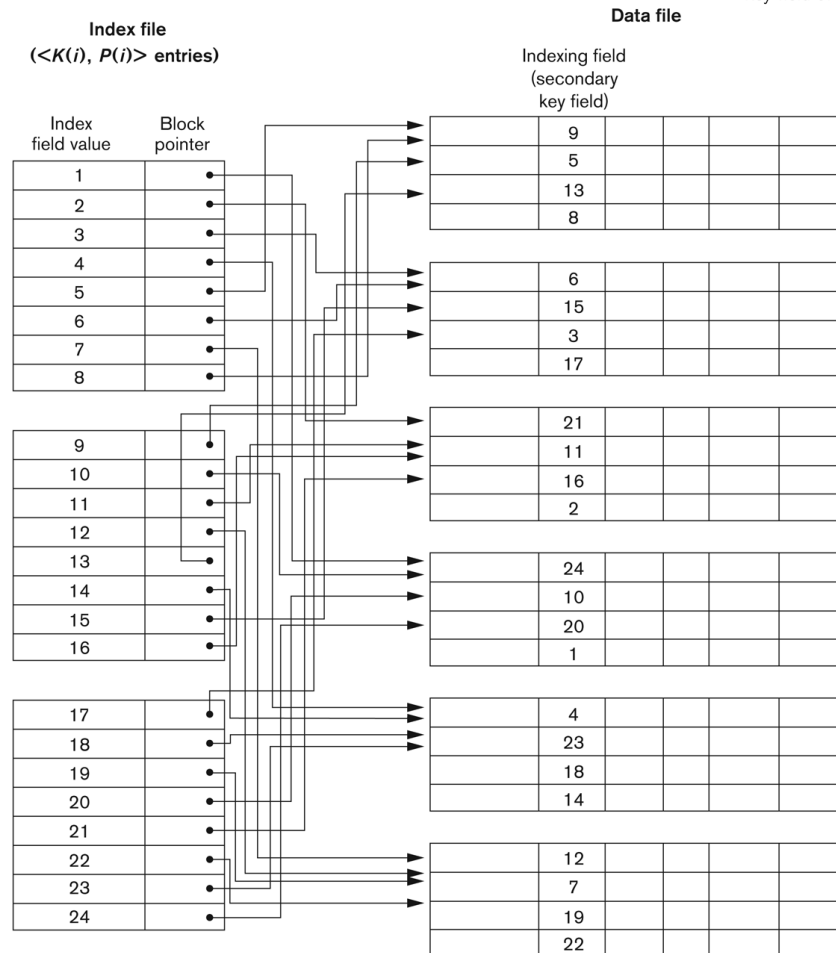
- Secondary Index

- A secondary index provides a secondary means of accessing a file for which some primary access already exists.
- The secondary index may be on a field which is a candidate key and has a unique value in every record, or a non-key with duplicate values.
- The index is an ordered file with two fields.
  - The first field is of the same data type as some **non-ordering field** of the data file that is an indexing field.
  - The second field is either a **block** pointer or a record pointer.
  - There can be *many* secondary indexes (and hence, indexing fields) for the same file.
- Includes one entry *for each record* in the data file; hence, it is a *dense index*

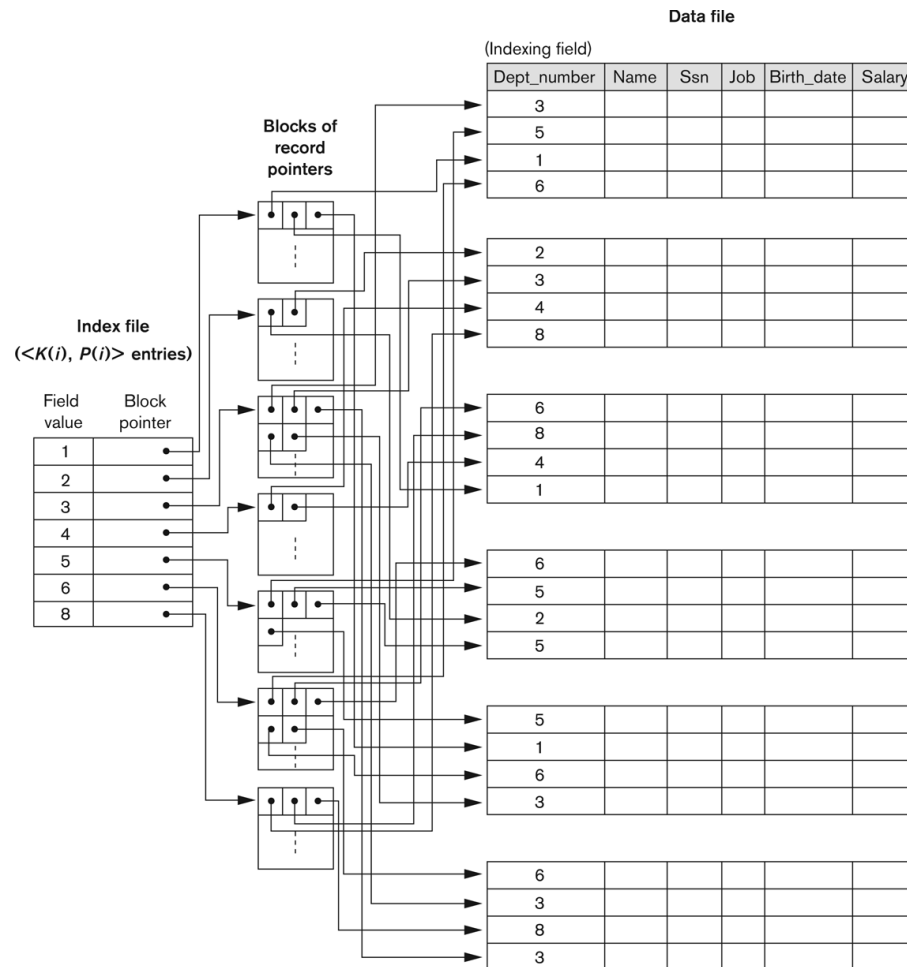
# Example of a Dense Secondary Index

**Figure 14.4**

A dense secondary index (with block pointers) on a nonordering key field of a file.



# An Example of a Secondary Index



**Figure 14.5**

A secondary index (with record pointers) on a nonkey field implemented using one level of indirection so that index entries are of fixed length and have unique field values.

# Properties of Index Types

**TABLE 14.2 PROPERTIES OF INDEX TYPES**

TYPE OF INDEX	NUMBER OF (FIRST-LEVEL) INDEX ENTRIES	DENSE OR NONDENSE	BLOCK ANCHORING ON THE DATA FILE
Primary	Number of blocks in data file	Nondense	Yes
Clustering	Number of distinct index field values	Nondense	Yes/no <sup>a</sup>
Secondary (key)	Number of records in data file	Dense	No
Secondary (nonkey)	Number of records <sup>b</sup> or Number of distinct index field values <sup>c</sup>	Dense or Nondense	No

<sup>a</sup>Yes if every distinct value of the ordering field starts a new block; no otherwise.

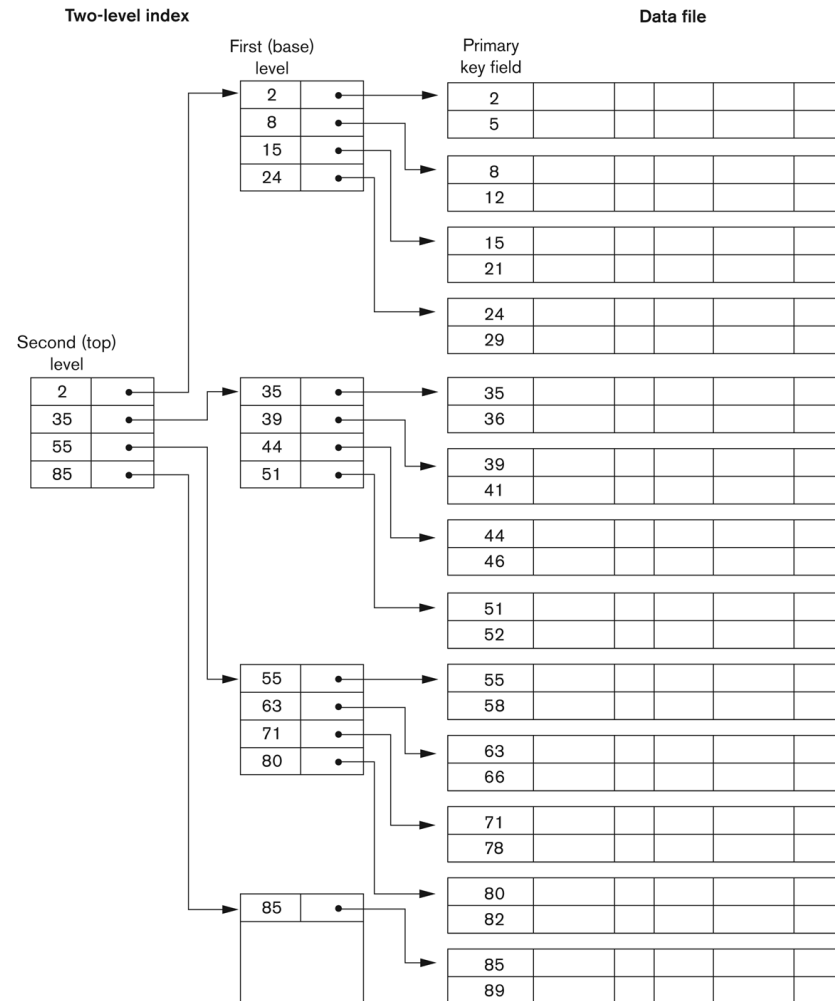
<sup>b</sup>For option 1.

<sup>c</sup>For options 2 and 3.

# Multi-Level Indexes

- Because a single-level index is an ordered file, we can create a primary index *to the index itself*;
  - In this case, the original index file is called the *first-level index* and the index to the index is called the *second-level index*.
- We can repeat the process, creating a third, fourth, ..., top level until all entries of the *top level* fit in one disk block
- A multi-level index can be created for any type of first-level index (primary, secondary, clustering) as long as the first-level index consists of *more than one* disk block

# A Two-level Primary Index



**Figure 14.6**  
A two-level primary index resembling ISAM (Index Sequential Access Method) organization.