# GTU Department of Computer Engineering CSE 312 – Spring 2024 HW 1 Report

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# HW 1 Part A

# syscalls.cpp

- getPid(): Gets the process ID
- waitpid(): Waits for a child process with the given ID
- sys exit(): Exits the current process
- fork(): Creates a child process (without returning the child's PID)
- fork(int \*pid): Creates a child process and returns the child's PID in the provided pointer
- **exec(void entrypoint())**: Executes and replaces curren task.
- addTask(void entrypoint()): Adds a new task to the system
- **printf**: Prints the desired char sequence to screen. Implementation is in **kernel.cpp**.

```
int myos::getPid()
    int pId = 1;
    asm("int $0x80" : "=c"(pId) : "a"(SYSCALLS::GETPID));
void myos::waitpid(common::uint8 t wPid)
    asm("int $0x80" : : "a"(SYSCALLS::WAITPID), "b"(wPid));
void myos::sys_exit()
    asm("int $0x80" : : "a"(SYSCALLS::EXIT));
// Creates a child process (without returning the child's PID)
void myos::fork()
    asm("int $0x80" : : "a"(SYSCALLS::FORK));
// Creates a child process and returns the child's PID in the provided pointer
void myos::fork(int *pid)
    asm("int $0x80" : "=c"(*pid) : "a"(SYSCALLS::FORK));
int myos::exec(void entrypoint())
    int result;
    asm("int $0x80" : "=c"(result) : "a"(SYSCALLS::EXEC), "b"((uint32_t)entrypoint));
    return result;
int myos::addTask(void entrypoint())
    int result;
    asm("int $0x80" : "=c"(result) : "a"(SYSCALLS::ADDTASK), "b"((uint32_t)entrypoint));
    return result:
```

```
enum SYSCALLS
{
    EXIT,GETPID,WAITPID,FORK,EXEC,PRINTF,ADDTASK
};
```

# Syscalls HandleInterrupt

Acts as a bridge between user-space system call requests and the kernel's system call implementation a.k.a interrupts.

```
uint32_t SyscallHandler::HandleInterrupt(uint32_t esp)
   CPUState *cpu = (CPUState *)esp;
   switch (cpu->eax)
   case SYSCALLS::EXEC:
       esp = InterruptHandler::sys_exec(cpu->ebx);
       break;
    case SYSCALLS::FORK:
       cpu->ecx = InterruptHandler::sys_fork(cpu);
       return InterruptHandler::HandleInterrupt(esp);
       break;
       printf((char *)cpu->ebx);
       break;
   case SYSCALLS::EXIT:
       if (InterruptHandler::sys_exit())
           return InterruptHandler::HandleInterrupt(esp);
       break;
   case SYSCALLS::WAITPID:
       if (InterruptHandler::sys_waitpid(esp))
           return InterruptHandler::HandleInterrupt(esp);
    case SYSCALLS::GETPID:
       cpu->ecx = InterruptHandler::sys_getpid();
   case SYSCALLS::ADDTASK:
       cpu->ecx = InterruptHandler::sys addTask(cpu->ebx);
   default:
       break:
   return esp;
```

## interrupts.cpp(Changes)

- sys\_getpid(): Gets the process ID of the current task by calling task managers GetPId.
- sys\_exec(common::uint32\_t entrypoint): Executes a new task by calling task managers ExecTask.
- sys\_addTask(common::uint32\_t entrypoint) : Adds a new task by calling task managers
- sys\_fork(CPUState \*cpustate): Forks the current task by calling task managers ForkTask.
- sys exit(): Exits the current task by calling task managers ExitCurrentTask.
- **sys\_waitpid(common::uint32\_t pid):** Wait for a child task to exit by calling task managers **WaitTask.**

```
common::uint32_t InterruptHandler::sys_getpid()
{
    return interruptManager->taskManager->GetPId();
}

common::uint32_t InterruptHandler::sys_exec(common::uint32_t entrypoint)
{
    return interruptManager->taskManager->ExecTask((void (*)())entrypoint);
}

common::uint32_t InterruptHandler::sys_addTask(common::uint32_t entrypoint)
{
    return interruptManager->taskManager->AddTask((void (*)())entrypoint);
}

common::uint32_t InterruptHandler::sys_fork(CPUState *cpustate)
{
    return interruptManager->taskManager->ForkTask(cpustate);
}

bool InterruptHandler::sys_exit()
{
    return interruptManager->taskManager->ExitCurrentTask();
}

bool InterruptHandler::sys_waitpid(common::uint32_t pid)
{
    return interruptManager->taskManager->WaitTask(pid);
}
```

• Interrupt handlers HandleInterrupt function calls task manager scheduler function **Schedule**.

```
uint32_t InterruptHandler::HandleInterrupt(uint32_t esp)
{
    return (uint32_t)interruptManager->taskManager->Schedule((CPUState *)esp);
}
```

## multitasking.cpp

- ForkTask(CPUState \*cpustate): ForkTask creates a new process.
  - It checks for available slots and copies the parent task's state.
  - Increases pId by using static pIdCounter variable.
  - It copies current tasks stack to new created one.
  - On success, it returns the new task's process ID (pId).

```
common::uint32_t TaskManager::ForkTask(CPUState *cpustate)
{
    if (taskCount >= 256)
    {
        printf("There are 256 processes are active. No more can be created.\n");
        return 0;
    }
} tasks[taskCount].taskState = READY;
tasks[taskCount].pId = (tasks:[currentTask].pId;
tasks[taskCount].pId = (task::pIdCounter)++;

// Copy the parent task's stack to the new task's stack
for (int i = 0; i < sizeof(tasks[currentTask].stack]; i++)
{
        tasks[taskCount].stack[i] = tasks[currentTask].stack[i];
}

// Calculate the offset of the CPU state within the parent task's stack
common::uint32_t currentTaskOffset = (((common::uint32_t)cpustate - (common::uint32_t)tasks[currentTask].stack));

// Set the new task's CPU state pointer based on the offset and its own stack
tasks[taskCount].cpustate = (CPUState *)(((common::uint32_t)tasks[taskCount].stack) + currentTaskOffset);

// Initialize specific CPU register (ECX) to 0
tasks[taskCount].cpustate->ecx = 0;
taskCount++;

// Set the new task's time slice (timeRemaining)
tasks[taskCount].timeRemaining = timeQuantum;
return tasks[taskCount].timeRemaining = timeQuantum;
return
```

- AddTask(Task \*task): Adds a new task with the provided task object.
  - It checks for available slots and copies the state, cpu registers, program counter and specific flags from the provided task.
  - Sets the time of the new task with timeQuantum.
  - On success, it returns true.

```
TaskManager::AddTask(Task *task)
if (taskCount >= 256)
tasks[taskCount].taskState = READY;
tasks[taskCount].pId = task->pId;
tasks[taskCount].cpustate = (CPUState *)(tasks[taskCount].stack + 4096 - sizeof(CPUState));
tasks[taskCount].cpustate->eax = task->cpustate->eax;
tasks[taskCount].cpustate->ebx = task->cpustate->ebx;
tasks[taskCount].cpustate->ecx = task->cpustate->ecx;
tasks[taskCount].cpustate->edx = task->cpustate->edx;
tasks[taskCount].cpustate->edi = task->cpustate->edi;
tasks[taskCount].cpustate->ebp = task->cpustate->ebp;
tasks[taskCount].cpustate->eip = task->cpustate->eip;
tasks[taskCount].cpustate->cs = task->cpustate->cs;
tasks[taskCount].cpustate->eflags = task->cpustate->eflags;
tasks[taskCount].timeRemaining = timeQuantum;
taskCount++;
```

- AddTask(void entrypoint()): Adds a new task with the provided task object.
  - It checks for available slots and copies the state, cpu registers, program counter and specific flags from the provided entry point function.
  - Sets the time of the new task with timeQuantum.
  - On success, it returns true.

```
common::uint32_t TaskManager::AddTask(void entrypoint())
{
    tasks[taskCount].taskState = READY;
    tasks[taskCount].pId = (Task::pIdCounter)++;
    // Set the new task's CPU state pointer to the top of its stack with an offset for CPUState size tasks[taskCount].cpustate = (CPUState *)(tasks[taskCount].stack + 4096 - sizeof(CPUState));

// Initialize specific CPU registers (EAX, EBX, ECX, EDX) to 0
    tasks[taskCount].cpustate->eax = 0;
    tasks[taskCount].cpustate->ebx = 0;
    tasks[taskCount].cpustate->ecx = 0;
    tasks[taskCount].cpustate->edx = 0;

// Initialize additional CPU registers (ESI, EDI, EBP) to 0
    tasks[taskCount].cpustate->esi = 0;
    tasks[taskCount].cpustate->ebp = 0;

// Set the program counter (EIP) to the entry point function address
    tasks[taskCount].cpustate->eip = (uint32_t)entrypoint;
    // Set the code segment selector from the global descriptor table (gdt)
    tasks[taskCount].cpustate->cs = gdt->CodeSegmentSelector();
    // Set specific status flags (EFLAGS)
    tasks[taskCount].cpustate->eflags = 0x202;
    taskCount++;
    tasks[taskCount].timeRemaining = timeQuantum;
    return tasks[taskCount].timeRemaining = timeQuantum
```

- WaitTask(common::uint32\_t esp): Suspends the current task until another task with the specified process ID (pid) finishes.
  - It prevents a task from waiting for itself and checks for invalid targets or finished tasks.
  - Set the current task's CPU state pointer
  - Set the current task's state to WAITING.
  - Set the process ID to wait.
  - On success, it returns true.

```
bool TaskManager::WaitTask(common::uint32_t esp)
{
    CPUState *cpustate = (CPUState *)esp;
    // Extract the process ID to wait for from a CPU register (EBX)
    common::uint32_t pid = cpustate->ebx;
    if (tasks[currentTask].pId == pid || pid == 0) // prevention self waiting
        return false;
    int index = getIndex(pid);
    // Check if the target task exists and isn't already finished
    if (index == -1)
        return false;
    if (taskCount <= index || tasks[index].taskState == FINISHED)
        return false;
    tasks[currentTask].cpustate = cpustate;
    tasks[currentTask].waitPid = pid;
    tasks[currentTask].taskState = WAITING;
    return true;
}</pre>
```

- ExecTask(void entrypoint()): Replaces the current task's state with the provided entry point function(Task).
  - Set the current task's CPU state pointer to the top of its stack with an offset for CPUState size.
  - Initialize specific and additional CPU registers, program counter, code segment selector and flags.
  - Returns a pointer to the CPU state for context switching.

```
common::uint32_t TaskManager::ExecTask(void entrypoint())
{
    tasks[currentTask].taskState - READY;
    // Set the current task's CPU state pointer to the top of its stack with an offset for CPUState size tasks[currentTask].cpustate = (CPUState *)(tasks[currentTask].stack + 4096 - sizeof(CPUState));

// Initialize specific CPU registers (EAX, EBX, ECX, EDX) to 0
    tasks[currentTask].cpustate->eax = 0;
    tasks[currentTask].cpustate->ebx = 0;
    tasks[currentTask].cpustate->edx = 0;
    tasks[currentTask].cpustate->edx = 0;

// Initialize additional CPU registers (ESI, EDI, EBP) to 0
    tasks[currentTask].cpustate->edi = 0;
    tasks[currentTask].cpustate->ebp = 0;

// Set the program counter (EIP) to the entry point function address
    tasks[currentTask].cpustate->eip = (uint32_t)entrypoint;

// Set the code segment selector from the global descriptor table (gdt)
    tasks[currentTask].cpustate->cs = gdt->CodeSegmentSelector();

// Set specific status flags (EFLAGS)
    tasks[currentTask].cpustate->eflags = 0x202;
    return (uint32_t)tasks[currentTask].cpustate;
}
```

- ExitCurrentTask(): Marks the current task as finished.
  - Prints process table when task is ended.

```
bool TaskManager::ExitCurrentTask()
{
   tasks[currentTask].taskState = FINISHED;
   PrintProcessTable();
   return true;
}
```

• **GetPID():** Returns the process ID of the current task.

```
common::uint32_t TaskManager::GetPID()
{
    return tasks[currentTask].pId;
}
```

• getIndex(common::uint32 t pid): Finds the index of a task based on its process ID (pId).

```
int TaskManager::getIndex(common::uint32_t pid)
{
    int index = -1;
    for (int i = 0; i < taskCount; i++)
    {
        if (tasks[i].pId == pid)
        {
            index = i;
            break;
        }
    }
    return index;
}</pre>
```

• PrintProcessTable: Prints process table with their pid, parent pid and their states

```
void TaskManager::PrintProcessTable()
   printf("\n^^^^^^^^^^^^^^^^^^^^^\n");
   printf("PID\tPPID\tSTATE\n");
   for (int i = 1; i < taskCount; i++)</pre>
       printNum(tasks[i].pId);
       printf("\t ");
       printNum(tasks[i].pPId);
       printf("\t ");
       if (tasks[i].taskState == TaskState::READY)
           if (i == currentTask)
               printf("RUNNING");
               printf("READY");
       else if (tasks[i].taskState == TaskState::WAITING)
           printf("WAITING");
       else if (tasks[i].taskState == TaskState::FINISHED)
            printf("FINISHED");
       printf("\n");
```

- **Scheduler**: Round-robin scheduling algorithm for tasks. It selects the next task to run based on the following logic:
  - Check for tasks: It first verifies if there are any tasks available for scheduling (checks task count). If not, it simply returns the current CPU state.
  - **Save current task state:** If tasks exist, it saves the CPU state of the currently running task (if any).
  - **Initialize next task search:** It starts searching for the next ready task by setting an index to the task after the current one (wrapping around if needed).
  - Loop for ready task: It iterates through tasks until a task in the "READY" state is found.
    - Waiting tasks: During this loop, it checks for tasks waiting on another task to finish (state: WAITING). If a waiting task is encountered, it searches for the task it's waiting on (using waitPid).
      - Waited-on task finished: If the waited-on task is finished, the
        waiting task's state is updated to "READY" and it is removed
        from the waiting list.
    - **Wrap around:** The search for a ready task is circular, meaning it wraps back to the beginning of the task list if no ready tasks are found after iterating through all tasks.
    - o **Break on loop back:** The loop breaks if it reaches the current task index again, indicating no ready tasks were found.
  - Handle current task time: If the current task was in a "READY" state, it reduces its remaining time quantum.
    - **Time quantum expired:** If the time quantum reaches zero, the current task's time is reset, and the scheduler moves to the next task.
  - Move to next task (default): If the current task wasn't ready or its time quantum didn't expire, the scheduler simply moves to the next task.
  - Final loop for ready task (optional): In case the initial loop for a ready task didn't find one due to waiting tasks becoming ready later, it does another loop with the same logic as step 4.
  - **Update current task and return:** Finally, it updates the current task index to the one found ready and returns the CPU state of the newly selected task for context switching.

- kernel.cpp
  - printNum: Converts the provided integer to char sequence and prints.

```
void printNum(int num)
{
    char numberStr[10];
    itoa(num, numberStr, 10);
    printf(numberStr);
}
```

• printf: Prints the provided char sequence to the screen

• **printArray**: Prints the array with n integer elements to the screen.

```
void printArray(int arr[], int n)
{
    printf("Array : {");
    for (int i = 0; i < n; i++)
    {
        printNum(arr[i]);
        if (i + 1 != n)
        {
            printf(",");
        }
        printf("} ");
}</pre>
```

#### forExample

- The code utilizes a loop to call fork six times, each time storing the returned value in a corresponding pid variable.
- The return value of fork determines the execution path:

#### • Parent Process (return value > 0):

- o The parent process continues execution after the **fork** call.
- o It uses waitpid to wait for the child process (identified by pid) to finish before proceeding.

#### • Child Process (return value == 0):

- The child process executes the code within the if block.
- It prints a message indicating its purpose using printf.
- It calls custom functions like TaskCollatz1, TaskCollatz2, TaskCollatz3, and long\_running\_program. These functions likely perform specific tasks.
- Finally, the child process exits using sys\_exit.

```
roid forkExample()
                                                fork(&pid4);
  int pid1 = 0;
                                                if (pid4 == 0)
  int pid2 = 0;
                                                    printf("LongRunning_10000 : ");
                                                    printNum(long_running_program(14000));
  int pid6 = 0;
                                                    printf("\n");
                                                    sys_exit();
  int parentPid = getPid();
  printf("Task Pid:
  printNum(parentPid);
                                                waitpid(pid4);
  printf("\n");
                                                fork(&pid5);
  fork(&pid1);
  if (pid1 == 0)
                                                if (pid5 == 0)
      printf("Collatz_1 : ");
                                                    printf("LongRunning_20000 : ");
      TaskCollatz1();
                                                    printNum(long_running_program(18000));
      sys_exit();
                                                    printf("\n");
  waitpid(pid1);
                                                    sys_exit();
  fork(&pid2);
  if (pid2 == 0)
                                                waitpid(pid5);
      printf("Collatz_2 : ");
                                                fork(&pid6);
      TaskCollatz2();
                                                if (pid6 == 0)
      sys_exit();
  waitpid(pid2);
                                                    printf("LongRunning_15000 : ");
                                                    printNum(long_running_program(16000));
  fork(&pid3);
                                                    printf("\n");
  if (pid3 == 0)
                                                    sys_exit();
      printf("Collatz_3 : ");
      TaskCollatz3();
                                                waitpid(pid6);
      sys_exit();
                                                sys_exit();
   waitpid(pid3);
```

- kernelMain (Changes)
  - Task task1(&gdt, EmptyTask);: Creates a Task object named task1 using the GDT and assigns it the EmptyTask function. This is a igniter task. It must be added to run next tasks.
  - Task task2(&gdt, forkExample);: Creates a Task object named task2 using the GDT and assigns it the forkExample function.(This function explained above)
  - taskManager.AddTask(&task1);: Adds task1 to the task manager.
  - taskManager.AddTask(&task2);: Adds task2 to the task manager.

```
TaskManager taskManager;

Task task1(&gdt, EmptyTask);
Task task2(&gdt, forkExample);

taskManager.AddTask(&task1);
taskManager.AddTask(&task2);
```

# **System Enhancements**

- System Call Implementations
  - New System Call Functions:
    - sys\_getpid(), sys\_exec(), sys\_addTask(), sys\_fork(), sys\_exit(), sys\_waitpid().
  - Handling System Calls:
    - The HandleInterrupt function is used to interrupt calls.
- Multitasking Enhancements
  - Task Management:
    - Added pIdCounter for process IDs, taskState for managing task states, and waitPid for tracking wait conditions.
    - ForkTask() duplicates the current task; AddTask() initializes new tasks.
  - Round Robin Scheduling:
    - Implemented in Schedule(), cycling through tasks based on state and timeQuantum.
  - PrintProcessTable() method prints task states.
- Kernel Improvements
  - Initialization and Task Management:
    - TaskManager manages tasks; forkExample initializes tasks and processes.
  - Interrupt Management

# **Implementation Strategies**

- System Call Implementation:
  - Added system call handling methods and integrated them within the
  - InterruptManager.
- Multitasking Enhancements:
  - Improved task management and scheduling.
- Kernel Improvements:
  - Added support for multiple tasks and system calls.
  - Implemented priority-based task switching

# **General Flow**

- 1. kernel.cpp uses system calls from syscalls.h
- 2. syscalls.cpp calls interrupts.cpp functions
- 3. interrupts.cpp uses multitasking.cpp functions for system calls and scheduler.