

# Lecture 6 - Graph Theory and Coloring

A **graph**  $G$  is a pair of sets  $(V, E)$  where

$V$  is a nonempty set of items called **vertices** or **nodes**

$E$  is a set of 2-item subsets of  $V$  called **edges**.



$x-y$  is another notation for an edge.

Two nodes  $x$  and  $y$  are **adjacent** if they are connected by edge.

An edge  $e = \{x, y\}$  is **incident** to  $x$  and  $y$ .

The number of edges incident to a node is the **degree** of the node. Denoted by  $\deg(x)$ .

A graph is **simple** if it has no loops or multiple edges:

- **Loop** is an edge only connects up one node.
- **Multiple edge** is we have got two edges connects the same endpoints.

## Graph Coloring Problem

### Definition

Given a graph  $G$  and  $K$  colors, assign a color to each node so adjacent nodes get different colors.

### Chromatic Number

The minimum value of  $K$  for which such a coloring exists is the **Chromatic Number** of  $G$  denoted by  $\chi(G)$

## Basic Graph Coloring Alg

for  $G = (V, E)$

1. order the nodes
2. order the colors
3. For  $i=1,2,\dots,n$  assign lowest legal color.

A *bipartite graph* can be colored with two colors such that no two adjacent vertices share the same color. This means we can divide the graph's vertices into two distinct sets where:

- All edges connect vertices from one set to vertices in the other set.
- No edges exist between vertices within the same set.

