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| *Cheatsheet*  *version*  *20210604*  **RCTTypeCLSAG (Type 5)** conceptual (not serialized) components | | | | | | | |
| **Input** (an instance of this section **for each actual** input UTXO) | | | | | **Output UTXO “ t ”** | **TX Fee** | **Outputs Unlock Time** |
|  | decoy UTXOs  **offsets** | **actual** UTXO  **offset** | decoy UTXOs  **offsets** |  | **X** stealth address | **f**  plain value, so to make miners able to evaluate it  expressed, as “**a**” in commitments, in **atomic units** (1/1012 XMR) | absolute time, expressed as:   * **Lock disabled** (0) * **Block height** (< 500.000.000) * **Epoch time** (≥ 500.000.000)   protocol enforces a default relative lock of **10 blocks**, too |
| level j=1  CLSAG pubkeys:  **X i** | **X …** | **X i = 𝜋** | **X …** | CLSAG  **signature**  (**c 1** , **r i**)  whole TX is  signed except  signatures’ fields | **C ≜ b G + a J** commitment  with:  **a** : moneroj in the UTXO (**plain** value)  **b** ≜ H( “commitment\_mask” , H(r Vi , t) )  = H( “commitment\_mask” , H(v0 R , t) )  usual payer/payee “exchange  trick” via transaction and view keys |
| level j=2  CLSAG pubkeys:  **C i - C pseudo** | (C i = 𝜋 – C pseudo) = (b - b’) G  being able to sign with this pub/priv EC keys pair hence proves that C i = 𝜋 and C pseudo commit to the same “a”  **C …** | **C i = 𝜋 = b G + a J** | **C …** | **Bulletproof** | **Extra** |
| **C pseudo ≜ b’G + a J** | | | output commitments’ “**a**” range proofs to avoid:  inputs outputs  5 **J** + 6 **J** = 21 **J** + (**𝑙** -10) **J**  = (21-10) **J** + **𝑙 J**  *to be continued…* | * **R** transaction pubkey(s) * encrypted **payID** (if any) * **. . .**   NOTE: loosely structured field, signed but not part of consensus (⟹ wallets duties) |
| key images: |  | effective **X\***  artifact **C\***  ( ⟹ CLSAG’s  pseudo **W\*** ) |  | | **amount**  **encrypted** value:  inflation by cyclic group ↻ *overflow*  ≜ **a** ⊕ H( “amount” , H(r Vi , t) )  = **a** ⊕ H( “amount” , H(v0 R , t) ) |

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| **Pedersen Commitments 101** | **CLSAG recap & now** | **Secret value flow via Monero’s UTXOs** |
| Commitments are a way to bind to a value, without revealing it (maybe postponing the disclosure). Pedersen ones in EC form:  **C ≜ b G + a J**  where:  **a** : committed value  **b** : random blinding factor, introducing entropy to make  hard to get “a”by means of a rainbow table  **G** : common generator point  **J** : an EC point for which “j” in: J ≜ j G is unknown (DLP)  PROPERTIES  **Theoretically hiding** : many (a,b)couples can give the  same commitment value C  **Computationally binding** : pretending to have committed  a fake different value is equivalent to solving DLP:  C(a2 , b2) = C(a , b) ⟹ j = (b – b2)/(a2 – a) ⟹ DLP solved  **Homomorphism** : C(a1 , b1) + C(a2 , b2) = C(a1+a2 , b1+b2) | * CLSAG is a sort of multisignature schema like MLSAG, given that each “squeezed” layer features its own actual key and a set of decoys * Differently from MLSAG, effective key image protecting against double-spending is available only for layer 1, so each input UTXO needs a separate CLSAG (but MLSAGs were used this way as well, due to anonymity concerns) * In RingCT two-levels-CLSAGs are used, where layer 2 signature is built to be a proof of equivalence between committed values of the actual UTXO and of the pseudo output commitment Cpseudo : the equivalence subsists *iff* a new specific public/private elliptic keys pair exists, so signing the CLSAG’s 2nd layer with that pair proves the pair exists and so the equivalence | Monero UTXOs encode their value in two fields: the commitment **C** and the **amount**:   * **amount** is the way payer and payee secretly share the value: the payer calculates it from the plain value **a** by the given definition; the payee can get back and verify the plain value **a** by just XORing the **amount** with the same hash function used in its definition (two equal XORs elide themselves). To have an hash computable by both payer and payee, transaction and view keys are used in the same Diffie-Hellman-like way they are in stealth addresses generation * commitments **C** are the way by which every network node can verify that the sum of transaction input UTXOs values are equal to the sum of transaction output UTXOs values plus fee, all without knowing the actual values (apart from fee’s, which is plain by design)   In RingCT schema:   * the payee verifies that output **amount** and **C** express the same value **a** (binding the payer/payee exchanged value to the one validated by the network), thanks to the -already documented- way by which **amount** and **b** are defined * for each input, payer defines a pseudo output commitment Cpseudo bound to same value “a” of the actual UTXO (without revealing its position in the ring, being decoupled by CLSAG signature) * to validate the value flow, network nodes hence only have to also check balance & bulletproof:   ⟹ Σ inputs a J = Σ outputs a J + f J  Σ inputs a = Σ outputs a + f  Σ inputs Cpseudo = Σ outputs C + f J  Bullet  proof  check  Σ inputs b’ = Σ outputs b ( imposed by the payer ) |

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| **Credits**  **& useful sources** | * Many [Monero Stack Exchange](https://monero.stackexchange.com/) posts (e.g. “[Complete extra field structure ...](https://monero.stackexchange.com/questions/11888/complete-extra-field-structure-standard-interpretation)”) * [MoneroBlocks APIs](https://localmonero.co/blocks/api) (e.g. [this call](https://localmonero.co/blocks/api/get_transaction_data/b0b3c299bd14f846bed560ac457b6fb8b84184c72629aad81910fc33526d4949) to inspect a TX – JSON parsing via [jq](https://stedolan.github.io/jq/) suggested) * Monero CLI [Source Code](https://github.com/monero-project/monero) (e.g. [/src/ringct/rctTypes.h](https://github.com/monero-project/monero/blob/master/src/ringct/rctTypes.h) with new signature type 5), hopefully explored via a code-inspector software (e.g. [Sourcetrail](https://www.sourcetrail.com/)) * This cheatsheet relies on previous ones dealing with [Monero Addresses](https://github.com/baro77/MoneroAddressesCS) and [Ring Signatures](https://github.com/baro77/RingsCS) * [Zero to Monero: 2nd Edition](https://www.getmonero.org/library/Zero-to-Monero-2-0-0.pdf) chapters 5, 6 and appendices A, B (even if still about MLSAG-signed  TXs, CLSAG case is straightforward once you master differences between the two Rings flavours) * [From Zero (Knowledge) to Bulletproofs](https://github.com/AdamISZ/from0k2bp) (first 6 pages dealing with commitments) |