### Mathematics Summer

# Graph Theory

# from Diestel.

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#### 1. Basics

#### 1.1. Graphs

**Definition 1.1.1.** A graph a pair G = (V, E) such that  $E \subseteq [V]^2$ . For clarity, we assume that  $V \cap E = \emptyset$ . The elements of V are the vertices of the graph G and the elements of E are its edges.

**Definition 1.1.2**. The *order* of a graph, written |G|, is the number of vertices of G. The number of edges of G is denotes by |G|. Graphs are *finite*, *infinite*, *countable* and so on according to their order.

*Example*: The *empty graph* is  $(\emptyset, \emptyset)$ , also denotes as  $\emptyset$  simply.

Example: A graph of order 0 or 1 is also known as a trivial graph.

**Definition 1.1.3.** A vertex v is *incident* with an edge e if  $v \in e$ ; then e is an edge at v. The two vertices incident with an edge are its *endvertices* or *ends*, and an edge *joins* its ends.

**Definition 1.1.4**. An edge  $\{x,y\}$  is usually written as xy (or yx). If  $x \in X$  and  $y \in Y$ , then xy is an X-Y edge. The set of all X-Y edges in a set E is denoted by E(X,Y).

*Remark*: Instead of  $E(\{x\}, Y)$  and  $E(X, \{y\})$ , we write E(x, Y) and E(X, y). The set of all the edges in E at a vertex v is denoted by E(v).

**Definition 1.1.5**. Two vertices x, y of G are adjacent or neighbors if xy is an edge of G. Two edges  $e \neq f$  are adjacent if they have an end in common.

*Example*: If all the vertices of G are pairwise adjacent, then G is *complete*. A complete graph on n vertices is denoted  $K^n$ .

**Definition 1.1.6**. A set of vertices or edges is *independent* if no two of its elements are adjacent. Independent sets of vertices are also called *stable*.

**Definition 1.1.7** . Let G=(V,E) and G'=(V',E'). A map  $\varphi:V\to V'$  is a homomorphism from G to G' if it preserves the adjacency of vertices, that is, if  $\{\varphi(x),\varphi(y)\}\in E'$  whenever  $\{x,y\}\in E$ .

**Lemma 1.1.1:** For every vertex x' in the image of  $\varphi: G \to G'$ , its inverse image  $\varphi^{-1}(x')$  is an independent set of vertices in G.

**Definition 1.1.1.** If  $\varphi$  is bijective and its inverse  $\varphi^{-1}$  is also a homomorphism (i.e.  $xy \in E \iff \varphi(x)\varphi(y) \in E'$  for all  $x, y \in V$ ), we call  $\varphi$  an *isomorphism*. We also say G and G' are isomorphic as denoted by  $G \simeq G'$  (or even simpler, G = G', when we only care about the *isomorphism type* of a given graph)

**Definition 1.1.9** . An isomorphism from G to itself is an *automorphism* of G.

**Definition 1.1.10**. A class of graphs that is closed under isomorphism is called a *graph property*.

Example: Containing a triangle is a graph property.

**Definition 1.1.11**. A map taking graphs as arguments is called a *graph invariant* if it assigns equal values to isomorphic graphs.

*Example*: The number of vertices and the number of edges are graph invariants. The greatest number of pairwise adjacent vertices is also another one.

**Definition 1.1.12** . We define  $G \cup G' = (V \cup V', E \cup E')$  and  $G \cap G' = (V \cap V', E \cap E')$ . If  $G \cap G' = \emptyset$  then G and G' are *disjoint*.

**Definition 1.1.13**. If  $V' \subseteq V$  and  $E' \subseteq E$ , then G' is subgraph of G (and G a supergraph of G'), written  $G' \subseteq G$ . If  $G' \subseteq G$  but  $G' \neq G$  then G' is a proper subgraph of G.

#### Remark:

**Definition 1.1.14** . If  $G' \subseteq G$  and G' contains all the edges  $xy \in E$  with  $x, y \in V'$ , then G' is an induced subgraph of G; we say that V' induces or spans G' in G.

Remark: If  $U \subseteq V$  is any set of vertices, then G[U] denotes the graph on U whose edges are precisely the edges of G with both ends in U.

### 2. Matchings, Covering, Packing

- 3. Connectivity
- 4. Flows
- 5. Hamilton Cycles