

6. Natural Deduction (continued)



Language & Logic

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Module syllabus

- Syntax of formal & natural languages
 - grammars, parsing
- Propositional logic
 - truth tables, semantics, proofs via natural deduction
- Predicate calculus
 - proofs via natural deduction
- Program correctness
 - structural induction

Recap

- Natural deduction for propositional logic
 - introduction and elimination rules for each connective
- New rules
 - \forall -introduction
 - \rightarrow -introduction, i.e. conditional proof
 - \forall -elimination, i.e. case analysis
 - \neg -introduction, i.e. proof by contradiction
- New proof techniques/concepts
 - sub-proofs, hypotheses (assumptions) & discharging

Today

- Proof strategies
 - i.e. which inference rules to apply and when
- Primitive vs. derived inference rules
 - i.e. which rules do we assume and which can we prove?
- Theorems & theorem introduction
- Soundness and completeness

Basic proof strategy

- Choose the right inference rules
 - use **introduction** rules to add connectives to the conclusion
 - use **elimination** rules to work with premises
 - only apply them to the **main connective**
- Plan the proof
 - work **backwards** from the conclusion
 - and **forwards** from the premises
 - look carefully at the structure of the propositions
 - e.g. proving $Q \rightarrow (P \rightarrow R)$ vs. $Q \wedge (P \wedge R)$
- And remember each sub-proof changes the context...

Proof strategy: Rules of thumb

- Some **rules of thumb** (so-called “Golden Rules”) for proofs:
 1. If the main connective in the conclusion is an **implication**, use **implication introduction**
 - i.e., assume its antecedent and try to prove the consequent
 2. If the main connective in any of the premises is a **disjunction**, try to use **disjunction elimination**
 - i.e., assume each disjunct separately and try to prove the same conclusion with both
 3. If all else fails, try **negation introduction**
 - i.e., assume the opposite of what you want to prove and deduce a contradiction

Example

- $\neg(P \wedge (\neg Q \vee R)) : P \rightarrow Q$

1.	P	Hypothesis	{1}
2.	$\neg Q$	Hypothesis	{2}
3.	$\neg Q \vee R$	\vee -introduction ₂	{2}
4.	$P \wedge (\neg Q \vee R)$	\wedge -introduction _{1,3}	{1,2}
5.	$\neg(P \wedge (\neg Q \vee R))$	Premise	{5}
6.	\perp	\wedge -introduction _{4,5}	{1,2,5}
7.	$\neg\neg Q$	\neg -introduction _{2,6}	{1,5}
8.	Q	$\neg\neg$ -elimination ₇	{1,5}
9.	$P \rightarrow Q$	\rightarrow -introduction _{1,8}	{5}

Inference rules for propositional logic

$$\frac{A \quad B}{A \wedge B} \wedge\text{-introduction}$$

$$\frac{A \wedge B}{A} \wedge\text{-elimination}$$

$$\frac{A}{A \vee B} \vee\text{-introduction}$$

$$\frac{A \vee B \quad A \vdash C \quad B \vdash C}{C} \vee\text{-elimination}$$

$$\frac{A \vdash B}{A \rightarrow B} \rightarrow\text{-introduction}$$

$$\frac{A \rightarrow B \quad A}{B} \rightarrow\text{-elimination}$$

$$\frac{A \vdash \perp}{\neg A} \neg\text{-introduction}$$

$$\frac{\neg\neg A}{A} \neg\neg\text{-elimination}$$

(symmetric rules omitted for clarity)

Inference rules

- Which inference rules do we need to prove **any** valid argument? And do we need **all** of them?
- **Primitive vs derived** inference rules
 - **primitive** inference rules allow us to prove things that we would not be able to without them
 - **derived** rules do not allow us to prove anything new
 - and can be proved correct using primitive ones

- For example:

$$\frac{A \rightarrow B \quad \neg B}{\neg A} \quad \text{“modus tollens”}$$

Some derived inference rules

$$\frac{A \rightarrow B \quad \neg B}{\neg A} \quad \text{“modus tollens”}$$

$$\frac{A}{\neg\neg A} \quad \neg\neg\text{-introduction}$$

$$\frac{A \vee B \quad \neg A}{B} \quad \text{“disjunctive syllogism”}$$

$$\frac{\perp}{A} \quad \text{“}\perp\text{ proves anything”}$$

Proof (disjunctive syllogism)

- $A \vee B, \neg A : B$

1.	$A \vee B$	Premise	{1}
2.	$\neg A$	Premise	{2}
3.	A	Hypothesis	{3}
4.	$\neg B$	Hypothesis	{4}
5.	\perp	\wedge -introduction _{2,3}	{2,3}
6.	$\neg\neg B$	\neg -introduction _{4,5}	{2}
7.	B	$\neg\neg$ -elimination ₆	{2}
8.	B	Hypothesis	{8}
9.	B	\vee -elimination _{1,3,7,8,8}	{1,2}

Hypothesis
not used
at all

Theorems

- A **theorem** is a formula (a proposition) which can be proved
 - i.e., a valid argument with no premises
- Examples of theorems
 - $\vdash P \rightarrow P$
 - $\vdash P \vee \neg P$ (“Law of Excluded Middle”)
- We can rewrite valid arguments as theorems
 - $\vdash ((A \rightarrow B) \wedge \neg B) \rightarrow \neg A$ (“modus tollens”)
 - $\vdash ((A \vee B) \wedge \neg A) \rightarrow B$ (“disjunctive syllogism”)
- Closely related: a **tautology** is a formula that is always true
 - i.e., is true for any possible truth valuation (of atomic propositions)

Theorem introduction

- In general (exercises, exam questions, etc.)
 - use only the inference rules you are given/allowed
 - so what use are derived inference rules?
- Theorem introduction
 - insert a known theorem directly into a proof
 - may yield shorter, clearer or more elegant proofs
 - and provide a separate proof of the theorem (e.g. in exam)
- Using theorem introduction
 - apply a **uniform substitution** (like when applying inference rules)
 - e.g., if we know: $\vdash ((A \rightarrow B) \wedge \neg B) \rightarrow \neg A$ (modus tollens)
 - we can insert, e.g.: $((\neg(P \wedge Q) \rightarrow R) \wedge \neg R) \rightarrow \neg(P \wedge Q)$

Example

- Prove: $P \vee \neg Q, P \rightarrow R, S \rightarrow Q \vdash R \vee \neg S$

1.	$P \vee \neg Q$	Premise	{1}
2.	P	Hypothesis	{2}
3.	$P \rightarrow R$	Premise	{3}
4.	R	\rightarrow -elimination _{3,2}	{2,3}
5.	$R \vee \neg S$	\vee -introduction ₄	{2,3}
6.	$\neg Q$	Hypothesis	{6}
7.	$S \rightarrow Q$	Premise	{7}
8.	$((S \rightarrow Q) \wedge \neg Q) \rightarrow \neg S$	Theorem (modus tollens)	{6,7}
9.	$(S \rightarrow Q) \wedge \neg Q$	\wedge -introduction _{7,6}	{6,7}
10.	$\neg S$	\rightarrow -elimination _{9,8}	{6,7}
11.	$R \vee \neg S$	\vee -introduction ₁₀	{6,7}
12.	$R \vee \neg S$	\vee -elimination _{1,2,5,6,11}	{1,3,7}

Using:

$\vdash ((A \rightarrow B) \wedge \neg B) \rightarrow \neg A$



Proofs

- Theorems
 - useful to know, even just to confirm validity or to guide the proof
- Equivalences
 - similarly, **equivalences** can be very helpful, e.g.:
 - $\neg(A \wedge B) \equiv \neg A \vee \neg B$
 - $\neg(A \vee B) \equiv \neg A \wedge \neg B$
 - $A \rightarrow B \equiv \neg A \vee B$
- Natural deduction for propositional logic (using the set of inference rules) is **sound** and **complete**
 - **soundness** means, if we can prove an argument, then it is valid
 - **completeness** means, if it is valid, then we can prove it

Semantic and syntactic validity

- We've seen two methods to show validity of an argument
 - **truth table** construction & analysis (semantic validity)
 - **natural deduction** proofs (syntactic validity)
- **Semantic validity**
 - an argument is valid iff it cannot be the case that all the premises are true and the conclusion false at the same time
- **Syntactic validity**
 - an argument is valid iff the conclusion can be derived from the premises by means of stipulated rules of inference
- **Notation**
 - semantic validity: $A \rightarrow B, \neg B \models \neg A$
 - syntactic validity: $A \rightarrow B, \neg B \vdash \neg A$

Soundness & completeness

- Soundness

- any syntactically valid argument is semantically valid
- i.e. $A_1, \dots, A_n \vdash B$ implies $A_1, \dots, A_n \models B$
- (alternatively: all theorems are tautologies)

- Completeness

- $A_1, \dots, A_n \models B$ implies $A_1, \dots, A_n \vdash B$
- i.e., we can construct a proof of any valid argument in propositional logic using natural deduction
- (alternatively: all tautologies are theorems)

Summary

- Proof strategies: rules of thumb
 - if there is a \rightarrow in the conclusion, try \rightarrow -introduction
 - if there is a premise of the form $A \vee B$, try \vee -elimination
 - otherwise, try \neg -introduction
- Further concepts of natural deduction/propositional logic
 - primitive vs. derived inference rules
 - theorems (and tautologies)
 - soundness and completeness
- Practice / revision
 - exercise classes on Tue/Thur: practice on some harder proofs
 - see also examples in, e.g., Tomassi book