

Assignment 1

Types

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1.

$$M ::= \lambda x.M \mid MM \mid \text{true} \mid \text{false} \mid \text{if } M \text{ then } M \text{ else } M \mid \langle M, M \rangle \mid \text{spread}(M, M) \mid \text{left}(M) \mid \text{right}(M)$$

$$T ::= \mathbb{B} \mid T \rightarrow T \mid T \times T \mid T + T$$

- (a) Extend call-by-value small-step operational semantics of the Curry-style Simply Typed λ -Calculus

Values, $V ::= \lambda x : T.M \mid \text{true} \mid \text{false} \mid \langle M, M \rangle \mid \text{left}(M) \mid \text{right}(M)$ Evaluation Contexts

$$C ::= \bullet \mid CM \mid VC \mid \text{if } C \text{ then } M \text{ else } M \mid \text{spread}(C, M) \mid \text{case}(C, M, M)$$

Rules:

$$\begin{array}{c} \frac{}{(\lambda x : T.M)V \rightarrow_v M[x \setminus V]} \beta \\[10pt] \frac{}{\text{if true then } M \text{ else } N \rightarrow_v M} \text{ITET} \\[10pt] \frac{}{\text{if false then } M \text{ else } N \rightarrow_v N} \text{ITEF} \\[10pt] \frac{M \rightarrow_v N}{C[M] \rightarrow_v C[N]} \text{CTX}_C \\[10pt] \frac{M \rightarrow_v \langle P, Q \rangle}{\text{spread}(M, N) \rightarrow_v NPQ} \text{SP} \\[10pt] \frac{M \rightarrow_v \text{left}(N)}{\text{case}(M, P, Q) \rightarrow_v NP} \text{CL} \end{array}$$

$$\frac{M \rightarrow_v \text{right}(N)}{\text{case}(M, P, Q) \rightarrow_v NQ} \text{CR}$$

(b) Prove: **Note: unless stated otherwise Π_n stands for the RHS of the bottom of the previous proof tree**

$$\text{spread}((\lambda x. \langle x, x \rangle) \text{true}, \lambda y. \lambda z. \text{if } y \text{ then left}(z) \text{ else right}(z)) \rightarrow_v^* \text{left}(\text{true})$$

$$\frac{\frac{\langle \text{true}, \text{true} \rangle}{\lambda x. \langle x, x \rangle \text{true}} \beta}{\Pi_1} \text{spread}(\bullet, M) \quad (1)$$

Where Π_1 is defined as:

$$\frac{\text{spread}((\lambda x. \langle x, x \rangle) \text{true}, \lambda y. \lambda z. \text{if } y \text{ then left}(z) \text{ else right}(z)) \rightarrow_v \text{spread}(\langle \text{true}, \text{true} \rangle, \text{if } y \text{ then left}(z) \text{ else right}(z))}{\Pi_1 \rightarrow_v (\lambda y. \lambda z. \text{if } y \text{ then left}(z) \text{ else right}(z)) \text{true true}} \text{SP} \quad (2)$$

$$\frac{\Pi_1 \rightarrow_v \lambda z. \text{if true then left}(z) \text{ else right}(z)}{\Pi_2 \rightarrow_v \lambda z. \text{if true then left}(z) \text{ else right}(z)} \beta \quad (3)$$

$$\frac{\Pi_2 \rightarrow_v \text{if true then left}(\text{true}) \text{ else right}(\text{true})}{\Pi_3 \rightarrow_v \text{if true then left}(\text{true}) \text{ else right}(\text{true})} \beta \quad (4)$$

$$\frac{\Pi_3 \rightarrow_v \text{left}(\text{true})}{\Pi_4 \rightarrow_v \text{left}(\text{true})} \text{ITET} \quad (5)$$

(c) Typing Rules:

$$\frac{}{\Gamma, x : T \vdash x : T} \text{VAR}$$

$$\frac{\Gamma, x : T \vdash M : U}{\Gamma \vdash \lambda x : T. M : T \rightarrow U} \text{ABS}$$

$$\frac{\Gamma \vdash M : T \rightarrow U \quad \Gamma \vdash N : T}{\Gamma \vdash MN : U} \text{APP}$$

$$\frac{}{\Gamma \vdash \text{true} : \mathbb{B}} \text{T}$$

$$\frac{}{\Gamma \vdash \text{false} : \mathbb{B}} \text{F}$$

$$\frac{\Gamma \vdash M : \mathbb{B} \quad \Gamma \vdash N : T \quad \Gamma \vdash P : T}{\Gamma \vdash \text{if } M \text{ then } N \text{ else } P : T} \text{ITE}$$

$$\frac{\Gamma \vdash M : T \quad \Gamma \vdash M : U}{\Gamma \vdash \langle M, N \rangle : T \times U} \text{PAIR}$$

$$\frac{\Gamma \vdash M : T \times U \quad \Gamma \vdash N : T \rightarrow U \rightarrow Z}{\Gamma \vdash \text{spread}(M, N) : Z} \text{ SP}$$

$$\frac{\Gamma \vdash M : T}{\Gamma \vdash \text{left}(M) : T + U} \text{ L}$$

$$\frac{\Gamma \vdash M : U}{\Gamma \vdash \text{right}(M) : T + U} \text{ R}$$

$$\frac{\Gamma \vdash M : T \quad \Gamma \vdash P : T \rightarrow Z}{\Gamma \vdash \text{case}(\text{left}(M), P, Q) : Z} \text{ CASEL}$$

$$\frac{\Gamma \vdash M : U \quad \Gamma \vdash Q : U \rightarrow Z}{\Gamma \vdash \text{case}(\text{right}(M), P, Q) : Z} \text{ CASER}$$

(d) Prove:

`spread(($\lambda x. \langle x, x \rangle$)true, $\lambda y. \lambda z. \text{if } y \text{ then left}(z) \text{ else right}(z)$)`

has type $\mathbb{B} + \mathbb{B}$

$$\frac{\overline{\text{true} : \mathbb{B}} \text{ T}}{\text{left}(\text{true}) : \mathbb{B} + U} \text{ L}$$

Where U is of **any** type, \therefore we can say $U \mapsto \mathbb{B}$ giving us $\text{left}(\text{true}) : \mathbb{B} + \mathbb{B}$