# Sets

Lesson 3

#### Named Sets

We'll use the following names for subsets of real numbers:

 $\mathbf{R}$  = the set of real numbers,

 $\mathbb{R}$  = the set of real numbers for handwritten text,

$$\mathbf{R}_{>0} = \{ x \in \mathbf{R} \mid x > 0 \},\,$$

 $\mathbf{R}_{\neq 0} = \{x \in \mathbf{R} \mid x \neq 0\}, \text{ (and similarly for other subscripts)}$ 

 $\mathbf{Z} = \mathsf{the} \; \mathsf{set} \; \mathsf{of} \; \mathsf{integers},$ 

 $\mathbb{Z} =$  the set of integers for handwritten text,

 $\mathbf{Q}$  = the set of rational numbers,

 $\mathbb{Q} =$  the set of rational numbers for handwritten text,

 $\emptyset = A$  set with no members, that is the empty set

# Set Operators

#### **Definition**

Let A and B be sets. Define the set *union*, intersection, and difference

$$A \cap B = \{x \mid (x \in A) \land (x \in B)\},\$$
  

$$A \cup B = \{x \mid (x \in A) \lor (x \in B)\},\$$
  

$$A \setminus B = \{x \mid (x \in A) \land (x \notin B)\},\$$

respectively.

3

# Set (an) example

### Example

We have

```
\{6, 107\} \cap \{28, 107\} = \{107\},\
\{6, 107\} \cup \{28, 107\} = \{6, 28, 107\},\
\{6, 107\} \setminus \{28, 107\} = \{6\},\
\{28, 107\} \setminus \{6, 107\} = \{28\}.
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- lacktriangle The last two examples show that in general  $A \setminus B \neq B \setminus A$ .
- ▼ The set difference is so much like real number subtraction, sometimes the symbol "-" is used instead of \.

# Set predicates

#### **Definition**

Let A and B be sets. Define

$$A \subset B \equiv (\forall x \in A)(x \in B),$$
  
 $A = B \equiv (A \subset B) \land (B \subset A).$ 

Specializing  $A\subset B$  to  $A=\varnothing$  gives

$$[\varnothing\subset B]\equiv (\forall x\in\varnothing)(x\in B)\equiv {\sf true}.$$

We've shown that:

## Proposition

Thus for all sets A and for any empty set  $\varnothing$ , we have  $\varnothing \subset A$ .

5

## Set equality

To show that sets A and B are equal, we almost always prove that  $A\subset B$  and  $B\subset A$ . If a proposition has the form

### Proposition

If  $H_1, H_2, \ldots$ , and  $H_n$ , then A = B.

where  $H_1, H_2, \dots H_n$  is the hypothesis, a template for proving the theorem is

#### Proof

Suppose  $x \in A$ . We'll show that  $x \in B$ . Since  $x \in A, H_1, H_2, \ldots$  and  $H_n$ , we have  $\ldots$ ; therefore  $x \in B$ .

Suppose  $x \in B$ . We'll show that  $x \in A$ . Since  $x \in B, H_1, H_2, \ldots$  and  $H_n$ , we have  $\ldots$ ; therefore  $x \in A$ .

① Notice how in the first case we append  $x \in A$  to the hypothesis; and in the second case, we append  $x \in B$ .

## Establish notation

#### **Proposition**

The set union is associative.

#### Proof

Let A, B, and C be sets. We'll show that  $A \cup (B \cup C) = (A \cup B) \cup C$ . Our proof uses the fact that the disjunction is associative; we have

- The statement of the proposition doesn't introduce notation, so the proof must do so.
- Alternatively, we can show that  $A \cup (B \cup C) \subset (A \cup B) \cup C$  and  $(A \cup B) \cup C) \subset A \cup (B \cup C)$ .

7

# Alternative proofs

#### **Proof**

Let A,B, and C be sets. We'll show that  $A\cup (B\cup C)=(A\cup B)\cup C$ . We have

$$\begin{aligned} x \in A \cup (B \cup C) &\implies (x \in A) \vee (x \in B \cup C), \\ &\implies (x \in A) \vee ((x \in B) \vee (x \in C)), \\ &\implies ((x \in A) \vee (x \in B)) \vee (x \in C), \\ &\implies x \in (A \cup B) \cup C. \end{aligned}$$

Similarly, we can show that  $x \in (A \cup B) \cup C \implies x \in A \cup (B \cup C)$ .

# The uniqueness of emptiness

### Proposition

There is at most one empty set.

#### Proof

Let O and O' be empty sets. Since O is empty, we have  $O \subset O'$ . Similarly since O' is empty, we have  $O' \subset O$ . We have shown that  $O \subset O'$  and  $O' \subset O$ ; therefore O = O'.

- With impunity, we can now refer to **the** empty set.
- ② A clumsy way to proof this is by contradiction. The proof assumes that there are empty sets O and O', but  $O \neq O'$ .

### Conflation

**Question:** True or false:  $\emptyset = \{\emptyset\}$ .

**Answer:** It's false. The set  $\{\varnothing\}$  is a set that has (exactly) one member, namely its member is the empty set. But the empty set has no members, so  $\varnothing \neq \{\varnothing\}$ 

We can write this as

# Counterexample ( $\varnothing = \{\varnothing\}$ )

We have  $\emptyset \in \{\emptyset\}$ , but  $\emptyset \notin \emptyset$ ; therefore  $\emptyset \neq \{\emptyset\}$ .

A counterexample is a proof of the negation of some statement. Usually to be considered a counter example, the proof examines one particular case.

## A unique template

If a proposition has the form

### Proposition

If  $H_1, H_2, \ldots$ , and  $H_n$ , there is at most one object X.

A template for its proof is

#### Proof

Let X and X' be such objects. Since  $H_1, H_2, \ldots$ , and  $H_n$ , we have  $\ldots$ ; therefore X = X'.

• When X and X' are real numbers, we might prove X=X' by showing that both  $X \leq X'$  and  $X' \leq X$  are true. Together, these inequalities prove that X=X'.

### Generalized unions

Let I be a set. And suppose that every member of I is a set. Since every member of I is a set, we can find the union of all of its members. We define

$$\underset{x \in I}{\cup} x = \{ a \mid (\exists x \in I) (a \in x) \}.$$

In this context, we say that the set I is an index set.

## Example

Define  $I = \{\{1,2\},\{107\}\}$ . Then I is a set and each member of I is a set. We have

### Finite unions

### Proposition

Let  $A_1,A_2,\ldots,A_n$  be sets. Define an index set I by  $I=\{A_1,A_2,\ldots,A_n\}$ . Then

$$\bigcup_{x \in I} x = A_1 \cup A_2 \cup \dots \cup A_n.$$

- **1** The set union is associative and commutative, so the meaning of  $A_1 \cup A_2 \cup \cdots \cup A_n$  unambiguous.
- 2 An index set neededn't be finite.

### Nonfinite unions

#### Example

The index set needn't be finite—here is an example. For  $x \in \mathbf{R}$ , define  $I = \{(-\infty, x) \mid x \in \mathbf{R}\}$ . Our index set is a set of open intervals. We claim that

$$\underset{x \in I}{\cup} x = \mathbf{R}.$$

#### **Proof**

Suppose  $a\in \underset{x\in I}{\cup} x$ . We'll show that  $a\in \mathbf{R}$ . Since  $a\in \underset{x\in I}{\cup} x$ , there is  $z\in I$  such that  $a\in z$ . But  $z\subset \mathbf{R}$ , so  $a\in \mathbf{R}$ ; we've shown that  $\ \cup x\subset \mathbf{R}$ .

Suppose  $a \in \mathbf{R}$ . We'll show that  $a \in \bigcup_{x \in I} x$ . We have  $a \in (-\infty, a+1)$ .

Further  $(-\infty, a+1) \in I$ ; therefore  $a \in \bigcup_{x \in I} x$ .

- **1** Notice that  $a \notin (-\infty, a)$ . But it is true that  $a \in (-\infty, a+1)$ .
- 2 It's also true that  $a \in (-\infty, a + 107 \pi^2)$ .

### Generalized intersections

Let I be a set. And suppose that every member of I is a set. Since every member of I is a set, we can find the union of all of its members. We define

$$\underset{x \in I}{\cap} x = \{ a \mid (\forall x \in I) (a \in x) \}.$$

## Example

Define  $I = \{\{1,2\},\{107\}\}$ . Then I is a set and each member of I is a set. We have

### Finite intersections

### Proposition

Let  $A_1,A_2,\ldots,A_n$  be sets. Define an index set I by  $I=\{A_1,A_2,\ldots,A_n\}$ . Then

$$\bigcap_{x \in I} x = A_1 \cap A_2 \cap \dots \cap A_n.$$

- **①** The set intersection is associative and commutative, so the meaning of  $A_1 \cap A_2 \cap \cdots \cap A_n$  unambiguous.
- 2 An index set needn't be finite.

### Nonfinite intersections

### Example

For  $x\in {\bf R}$ , define  $I=\{(-\infty,x)\mid x\in {\bf R}\}$ . Our index set is a set of open intervals. We claim that

$$\bigcap_{x \in I} x = \varnothing.$$

#### Proof

We'll prove this using contradiction. Suppose  $\bigcap\limits_{x\in I}x$  has at least one member; say  $a\in\bigcap\limits_{x\in I}x$ . We have

$$(\forall x \in \mathbf{R})(a \in (-\infty, x)).$$

In particular, we have  $a\in (-\infty,a)$ . But  $a\in (-\infty,a)$  is false; therefore  $\underset{x\in I}{\cap} x$  cannot have a member, so  $\underset{x\in I}{\cap} x$  is the empty set.

#### Alternative notation

Sometimes we take the index set to be a subset of  ${\bf R}$  and we denote the sets members by subscripts. Say  $I\subset {\bf R}$  and  $A_x$  is a set for each  $x\in I$ . This notation is particlarly popular when  $I={\bf Z}_{>0}$ . For example

$$\bigcap_{k \in \mathbf{Z}_{>0}} A_k = \{ a \mid (\forall n \in \mathbf{Z}_{>0}) (a \in A_n) \}$$

And

$$\bigcup_{k \in \mathbf{Z}_{>0}} A_k = \{ a \mid (\exists n \in \mathbf{Z}_{>0}) (a \in A_n) \}$$

When the index set is uncountable, maybe it's just me, but definitions such as

$$A_x = (-\infty, x)$$
 for all  $x \in \mathbf{R}$ 

are semi-bazaar looking. For such cases, I think it's more clear to define the index set to be a set of sets:

$$I = \{(-\infty, x) \mid x \in \mathbf{R}\}.$$

#### **Functions**

To define a function F with domain A and formula blob, we can write

$$F=x\in A\mapsto \mathsf{blob}.$$

In the rare cases that it's important to give the function a codomain, we can write

$$F = x \in A \mapsto \mathsf{blob} \in B$$
,

where  $\operatorname{codomain}(F) = B$ . Generically for a function F with domain A and codomain B, we say that F is a function from A to B.

#### Example

The notation

$$F = x \in [-1, 1] \mapsto 2x + 1$$

is our compact way of writing: Define F(x) = 2x + 1, for  $-1 \le x \le 1$ .

# Function signature

The notation  $F: A \rightarrow B$  means

- F is a function.
- $\bigcirc$  dom(F) = A.
- $\odot$  codomain(F) = B.

We'll say that  $A \to B$  is the *signature* of a function. The signature of a function doesn't tell us its formula. It does tell us the domain of a function and it indicates what the outputs of the function can be.

## Range

#### **Definition**

For any function, we define

$$\operatorname{range}(F) = \left\{ F(x) | x \in \operatorname{dom}(F) \right\}.$$

Thus range(F) is the set of all outputs.

#### **Fact**

Let F be a function. Then

$$[y \in \text{range}(F)] \equiv (\exists x \in \text{dom}(F)) (y = F(x)).$$

## Example

Define  $F = x \in [-1,1] \mapsto 2x + 1$ . Then  $\frac{3}{2} \in \operatorname{range}(F)$  because  $\frac{1}{4} \in \operatorname{dom}(F)$  and  $F(\frac{1}{4}) = \frac{3}{2}$ .

#### **Ontoness**

The codomain of a function tells us something about its outputs, but remember that the range and the codomain of a function need not be the same. For all functions F, we have

range  $F \subset \operatorname{codomain}(F)$ .

#### **Definition**

A function is *onto* if its range and codomain are equal.

## Example

**Question**: Is the sine function onto? **Answer** It is if its codomain is [-1,1]. But if its codomain is  $\mathbf{R}$ , then no it's not onto. There is no standard value for the codomain of the trigonometric functions, so the asking "Is the sine function onto?" is rubbish.

## Equality

#### **Definition**

Functions F and G are equal dom(F) = dom(G) and for all  $x \in dom(F)$ , we have F(x) = G(x). Equivalently

$$(F=G)\equiv (\mathrm{dom}(F)=\mathrm{dom}(G))\wedge (\forall x\in \mathrm{dom}(F))(F(x)=G(x)).$$

• The definition of function equality does not involve the codomain of the function. Thus two functions can be equal, but have unequal codomains.

### Example

The functions  $F=x\in[-1,1]\mapsto x\in[-1,1]$  and  $G=x\in[-1,1]\mapsto x\in\mathbf{R}$  are equal, but F is onto and G is not onto.

Thus ontoness isn't a property of a function.

# Apply a function to a set

#### **Definition**

Let  $F: A \to B$ . For any subset A' of A define

$$F(A') = \{F(x) | x \in A'\}.$$

Equivalently, we have

$$y \in F(A') \equiv (\exists x \in A')(y = F(x)).$$

## Proposition

For all functions F, we have  $F(\operatorname{dom} F) = \operatorname{range}(F)$ . Further  $F(\emptyset) = \emptyset$ .

# Inverse image

#### **Definition**

Let  $F: A \to B$ . For any subset B' of B define

$$F^{-1}(B') = \{ x \in A | F(x) \in B \}.$$

Equivalently, we have

$$x \in F^{-1}(B') \equiv F(x) \in B.$$