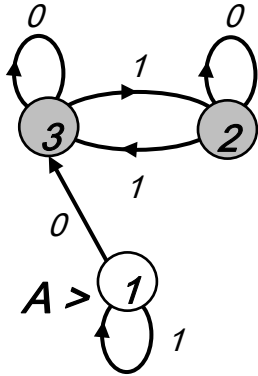


SABANCI UNIVERSITY
Faculty of Engineering and Natural Sciences
CS 302 Automata Theory

ANSWERS TO THE FINAL EXAM

Question 1 (20 points)



$A = (Q_A, \Sigma, \delta_A, q_{0A}, F_A)$ is given by the figure above where $q_{0A} = 1$ and $F_A = \{2, 3\}$.

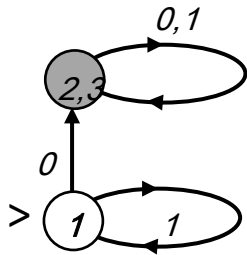
(a)

q	σ	q'
1	0	3
1	1	1
2	0	2
2	1	3
3	0	3
3	1	2

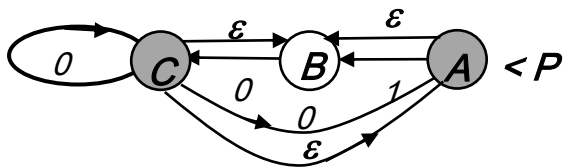
(b) By simple inspection : $L_A = 1^*.0.(0+1)^*$; $L_A^c = 1^*$

(c) Using Table Filling Algorithm below we observe that states 2 and 3 are equivalent hence minimal state DFA is as shown below.

	1	2	3
1		x	x
2			-----
3			



Question 2 (20 points)



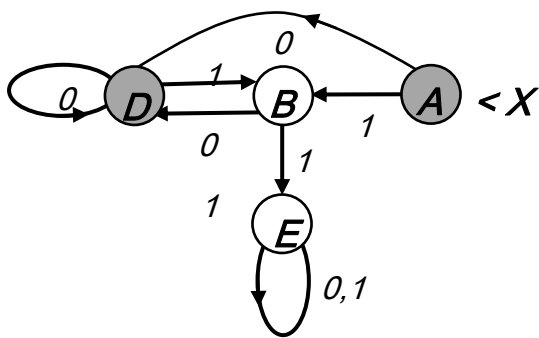
(a)

q	σ	q'
A	0	ABC
A	1	B
B	0	ABC
B	1	\emptyset
C	0	A,B,C
C	1	B

(b) DFA table below

q	σ	q'
A^*	0	ABC
A	1	B
B	0	ABC
B	1	\emptyset
$ABC(D^*)$	0	ABC
ABC	1	B
$\emptyset(E)$	0	\emptyset
\emptyset	1	\emptyset

DFA X is as below



Question 3 (20 points)

CFG is $G = (\{S, A, B\}, \{0, 1\}, R, S)$ where the production set R is given below :

$R : S \rightarrow AB ; A \rightarrow 0A1 / e ; B \rightarrow 1B0 / e$

(a)

null production elimination :

$S \rightarrow AB / A / B ; A \rightarrow ZERO A ONE / ZERO ONE ; B \rightarrow ONE B ZERO / ONE ZERO$

unit production elimination :

$S \rightarrow AB / ZERO A ONE / ZERO ONE / ONE B ZERO / ONE ZERO ;$

$A \rightarrow ZERO A ONE / ZERO ONE ; B \rightarrow ONE B ZERO / ONE ZERO$

and the CNF is as below :

$S \rightarrow AB / X ONE / ZERO ONE / Y ZERO / ONE ZERO ; A \rightarrow X ONE / ZERO ONE ;$

$B \rightarrow Y ZERO / ONE ZERO ; X \rightarrow ZERO A ; Y \rightarrow ONE B ; ZERO \rightarrow 0 ; ONE \rightarrow 1$

(b) Note that L_G is given by $L_G = \{0^n 1^{n+k} 0^k ; n, k \geq 0\}$ and the PDA P :

$P = (\{q, f, p, r, s\}, \{0, 1\}, \{0, 1, Z_0, Z_1\}, \delta, q, Z_0, \{f\})$

accepts L_G where δ is given below

$(q, e, Z_0) \rightarrow (f, Z_1)$

$(f, 0, Z_1) \rightarrow (q, 0Z_1)$

$(q, 0, 0) \rightarrow (q, 00)$

$(q, 1, 0) \rightarrow (p, e)$

$(p, 1, 0) \rightarrow (p, e)$

$(p, 1, Z_1) \rightarrow (r, 1Z_1)$

$(r, 1, 1) \rightarrow (r, 11)$

$(r, 0, 1) \rightarrow (s, e)$

$(s, 0, 1) \rightarrow (s, e)$

$(s, e, Z_1) \rightarrow (f, Z_1)$

OR using the standard procedure of converting a CGF to a PDA P by empty stack :

$P = (\{q_0\}, \{0, 1\}, \{0, 1, Z_0\}, \delta, q_0, Z_0)$

where δ is given below

$(q_0, e, Z_0) \rightarrow (q_0, SZ_0)$

$(q_0, e, S) \rightarrow (q_0, AB)$

$(q_0, e, A) \rightarrow (q_0, 0A1)$

$(q_0, e, A) \rightarrow (q_0, e)$

$(q_0, e, B) \rightarrow (q_0, 1B0)$

$(q_0, e, B) \rightarrow (q_0, e)$

$(q_0, 0, 0) \rightarrow (q_0, e)$

$(q_0, 1, 1) \rightarrow (q_0, e)$

$(q_0, e, Z_0) \rightarrow (q_0, e)$

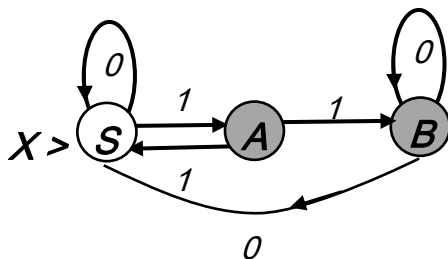
(c) For the first solution P is already a **DPDA** ; for the second it clearly is **not** a **DPDA**.

Question 4 (20 points)

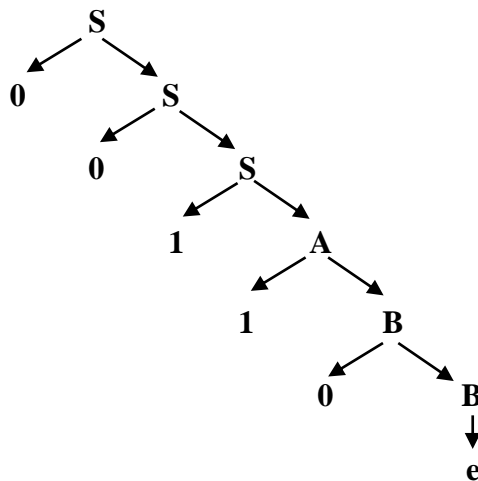
CFG is given as $G = (\{S, A, B\}, \{0, 1\}, R, S)$ where R is given below

$R : S \rightarrow 0S / 1A ; A \rightarrow 1B / 1S / e ; B \rightarrow 0B / 0S / e$

(a) CFG is a *right linear grammar* and is identified with an NFA as below where variables in G are identified with states in X and thus L_G is regular.



(b) There is a unique parse tree for the string 00110; however, there are other strings for which it is possible to construct two different parse trees (e.g. 11001). Hence, G is an *ambiguous grammar*.



Note: The answer “This is a non-ambiguous grammar, because there is a unique parse tree for the string 00110” will be accepted as well.

Question 5 (20 points)

(a) $(s, \underline{\#} 0^n 1^n) \dashv\dashv |_T (h, \underline{\#} 1^n 0^n)$

<i>TM</i>	<i>Condition</i>	<i>Next TMs</i>
$> A = R$	$\sigma = 0$	$1 . A$
	$\sigma = 1$	$0 . A$
	$\sigma = \#$	$L_{\#} . h$

(b) We use a 2 tape TM as below and initial condition is : $(s, \underline{\#} w, \underline{\#})$

<i>TM</i>	<i>Condition</i>	<i>Next TMs</i>
$> A = R^1 R^2$	$\sigma^1 = \#$	C
	$\sigma^1 = 0$	$0^2 . A$
	$\sigma^1 = 1$	$L^1 . B$
$B = R^1 . L^2$	$\sigma^1 = 1 \wedge \sigma^2 = 0$	B
	$\sigma^1 = \# \wedge \sigma^2 = \#$	h
	<i>else</i>	C
$C = \#^l$	-	C