

CS 302
QUIZ 3

27 February, 2020

ANSWERS

(a) (5 pts) See the relevant slide.

(b) (5 pts) The language $L = \{w \in (0+1)^* \mid w = 0^k 1^m, k \geq m+1; k, m \geq 0\}$

is **not** a regular language. Assume it is regular and choose $w = 0^{n+1} 1^n \in L$, where n is as given by the pumping lemma (PL). By the PL, $w = x.y.z$; $|x.y| \leq n$; $|y| > 0$ and $x.y^j.z \in L$ for $j=0,1,\dots$; in particular for $j=0$, $x.z \in L$. But by construction

$x.y = 0^p$, $p \leq n$, $y = 0^q$, $q > 0$ and $z = 0^{n+1-p} 1^n$; therefore

$x.z = 0^{p-q} 0^{n+1-p} 1^n = 0^{n+1-q} 1^n \notin L$ since $q > 0$ and so $k = n+1-q \leq n = m$

This violation of the PL contradicts the assumption that L is a regular language.