

SABANCI UNIVERSITY
Faculty of Engineering and Natural Sciences
CS 302 Automata Theory
Fall 2016

NOTES ON THE ANATOMY OF A *PDA*

Let $P = (Q, \Sigma, \Gamma, \delta, s, Z_0, F)$ be a given *PDA* and let

$$(q, w, \gamma) \vdash^n (p, e, e) \dots\dots\dots(1)$$

be a computation of P . In what follows all computations follow the same path as in (1), that is, the same transition sequence used in (1) is used.

(Fact 1) For any $u \in \Sigma^*, \beta \in \Gamma^*$ we have $(q, wu, \gamma\beta) \vdash^n (p, u, \beta)$

Proof of Fact 1

Given

$$(q, w, \gamma) \vdash^n (p, e, e) \dots\dots\dots(1)$$

we prove the result by induction on n

n=1 Case

If $n=1$ then we must have $w = a$ where $a \in \Sigma$ or $a = e$; $\gamma = X \in \Gamma$ and $(p, e) \in \delta(q, a, X)$

which implies that $(q, au, X\beta) \vdash^n (p, u, \beta)$ by definition of the δ function.

Now assume that $(q, w, \gamma) \vdash^n (p, e, e)$ implies that $(q, wu, \gamma\beta) \vdash^n (p, u, \beta)$

(induction hypothesis) and prove that :

$$(q, w, \gamma) \vdash^{n+1} (p, e, e) \text{ implies that } (q, wu, \gamma\beta) \vdash^{n+1} (p, u, \beta)$$

But letting $(q, w, \gamma) := (q, aw', Y\gamma') \vdash (q', w', \xi\gamma')$ where we assume that $(q', \xi) \in \delta(q, a, Y)$ we have

$(q, w, \gamma) := (q, aw', Y\gamma') \vdash (q', w', \xi\gamma') \vdash^n (p, e, e)$ and by induction hypothesis

$(q', w'u, \xi\gamma'\beta) \vdash^n (p, u, \beta)$ and for the case $n=1$ applied to the first step of the computation above

$$(q, wu, \gamma\beta) \vdash (q', w'u, \xi\gamma'\beta) \text{ we conclude that } (q, wu, \gamma\beta) \vdash^{n+1} (p, u, \beta)$$

(Fact 2)

Conversely if

$$(q, wu, \gamma\beta) \vdash^n (p, u, \beta) \dots\dots\dots(2)$$

Then again using the same path in (2) we conclude that

$$(q, w, \gamma) \vdash^n (p, e, e) \dots\dots\dots(3) \text{ provided that :}$$

$$(q, wu, \gamma\beta) \dashv\vdash^k (q_k, v_k u, \gamma_k \beta) \text{ with } |\gamma_k| > 0 \text{ for all } 0 \leq k < n \dots\dots\dots(4)$$

(Condition (4) is referred to as the *no top-of-the-stack exposure of β* throughout the computation path)

Proof of Fact 2

To prove the converse fact we again use induction on n . The case for $n=1$ is trivial and we omit the proof. Hence given

$$(q, wu, \gamma\beta) \dashv\vdash^{n+1} (p, u, \beta) \dots\dots\dots(2)$$

using the same mechanism in the proof of **Fact 1** and the implication $|\gamma| > 0$ of the additional assumption (*no top-of-the-stack exposure of β*) of **Fact 2** we have

$$(q, wu, \gamma\beta) := (q, aw'u, Y\gamma'\beta) \dashv\vdash (q', w'u, \xi\gamma'\beta) \dashv\vdash^n (p, u, \beta)$$

and hence by the induction hypothesis

$$(q', w', \xi\gamma') \dashv\vdash^n (p, e, e)$$

Also using the $n=1$ case

$$(q, w, \gamma) := (q, aw', Y\gamma') \dashv\vdash (q', w', \xi\gamma') \dashv\vdash^n (p, e, e) \text{ which proves the result.}$$

(Fact 3)

$$\text{Let } (q, w, X_1 X_2 \dots X_m) \dashv\vdash^n (p, e, e) \dots\dots\dots(5)$$

where $X_j \in \Gamma$ for $j = 1, 2, \dots, m$

then along the computational path defined by (5) above, for each $j = 1, 2, \dots, m$ there exists a *unique* integer $n(j)$, a segmentation $u_j v_j$ of w (i.e. $w = u_j v_j$) and a state q_j such that

$$(q, w, X_1 X_2 \dots X_m) \dashv\vdash^{n(j)} (q_j, v_j, X_j X_{j+1} \dots X_m) \dots\dots\dots(6)$$

and

$$(q, w, X_1 X_2 \dots X_m) \dashv\vdash^k (q', w' v_j, \gamma X_j X_{j+1} \dots X_m) \dots\dots\dots(7)$$

with $|\gamma| > 0$ for all $0 \leq k < n(j)$. Note $n(1) = 0$ and $q_1 = q$.

Proof of Fact 3

Given

$$(q, w, X_1 X_2 \dots X_m) \dashv\vdash^n (p, e, e) \dots\dots\dots(3)$$

we use induction on n for the proof.

For $n=1$ case, (3) implies that $w = a$ where $a \in \Sigma$ or $a = e$; $m = 1$ and $(p, e) \in \delta(q, a, X_1)$

The conclusion then trivially holds where $n(1)=0$; $q_1 = q$ and $w = a = e$ ($u_1 = e, v_1 = a$) since $(q, a, X_1) \dashv\vdash^0 (q, a, X_1)$. Now assume the conclusion holds for n and prove it for $n+1$. Thus we assume that

$$(q, w, X_1 X_2 \dots X_m) \dashv\vdash^{n+1} (p, e, e) \dots\dots\dots(4)$$

and expand the first transition in (4) as

$(q, w, X_1 X_2 \dots X_m) = (q, aw', X_1 X_2 \dots X_m) \vdash^{---} (q', w', Y_1 Y_2 \dots Y_p X_2 X_3 \dots X_m)$ where in general $(q', Y_1 Y_2 \dots Y_p) \in \delta(q, a, X_1)$. Now we apply the induction hypothesis to the case $(q', w', Y_1 Y_2 \dots Y_p X_2 X_3 \dots X_m) \vdash^{--- n} (p, e, e) \dots \dots \dots (5)$

with n transitions. If we let the corresponding unique integers for $Y_1, Y_2, \dots, Y_p, X_2, X_3, \dots, X_m$ as $n'(1)$ to $n'(p+m-1)$ then the unique integer corresponding to X_j (from state q) is $n'(p+j-1)+1$ which implies that $n(j) = n'(p+j-1)+1$ for $j=2, \dots, m$ and similarly if q'_k are the states and u'_k and v'_k are the segmentations of w' for each $k=1, \dots, p+m-1$ we conclude that, using also the initial step before the last n steps, $q_j = q'_{p+j-1}$ and $w = a u'_{p+j-1} v'_{p+j-1}$ and hence $u_j = a u'_{p+j-1}$; $v_j = v'_{p+j-1}$ for $j = 2, \dots, m$. The remaining details follow from the induction hypothesis.

(Fact 4)

Suppose that (5) above holds then there is a unique segmentation

$w = w_1 w_2 \dots w_m$ such that for $j = 1, 2, \dots, m$

$(q_j, w_j, X_j) \vdash^{--- k(j)} (q_{j+1}, e, e) \dots \dots \dots (8)$

with $k(j) := n(j+1) - n(j)$ and $q_{m+1} := p$ and $n(m+1) := n$ and w_j are related to u_j and v_j

(see Fact 3 above) by :

$u_j = w_1 \dots w_j$ and $v_j = w_{j+1} \dots w_m$

Proof of Fact 4

We use Fact (3) to prove Fact (4). This time we use induction on m , that is, the number of active stack symbols. For $m=1$ the result trivially follows from Fact (3). Assume the result holds for m and prove it for $m+1$. Now using Fact (3)

$(q, w, X_1 X_2 \dots X_{m+1}) \vdash^{--- n(2)} (q_2, v_1, X_2 X_3 \dots X_{m+1}) \dots \dots \dots (6)$

where $w = u_1 v_1$. Now we can use Fact (2) with $u = v_1$ and $\beta = X_2 X_3 \dots X_m$ to conclude that

$(q, u_1, X_1) \vdash^{--- n(2)} (q_2, e, e)$ where $k(1) = n(2) - n(1) = n(2) - 0 = n(2)$ as required for $j=1$.

Now using the induction hypothesis we conclude that the original premiss

$(q_2, v_1, X_2 \dots X_{m+1}) \vdash^{--- n-n(2)} (p, e, e)$

yields the required results for $j=2, \dots, m+1$ where $k(j) = n(j+1) - n(j)$.

The input segmentation for v_1 is $v_1 = w_2 \dots w_{m+1}$ which follows from the induction hypothesis and the fact that $w_1 = u_1$ completes the entire segmentation. Remaining minor details are left to the reader.

(Fact 5)

If for $j = 1, 2, \dots, m$

$$(q_j, w_j, X_j) \vdash^{k(j)} (q_{j+1}, e, e) \dots \dots \dots (9)$$

then

$$(q, w, X_1 X_2 \dots X_m) \vdash^n (p, e, e)$$

where $w := w_1 w_2 \dots w_m$, $p := q_{m+1}$ and $q := q_1$.

Proof of Fact 5

We are given that

$$(q_j, w_j, X_j) \vdash^{k(j)} (q_{j+1}, e, e) \dots \dots \dots (7)$$

for $j = 1, 2, \dots, m$ and asked to prove that

$$(q, w, X_1 X_2 \dots X_m) \vdash^n (p, e, e)$$

where $n := k(1) + k(2) + \dots + k(m)$, $q_1 := q$, $q_{m+1} := p$ and $w := w_1 w_2 \dots w_m$

We prove this by using induction on m . For $m=1$ the result is trivial since the conclusion and the premiss are identical. Now assume the statement is true for m and assuming (7) holds for $j=1, 2, \dots, m+1$ we shall prove that

$$(q, w, X_1 X_2 \dots X_{m+1}) \vdash^n (p, e, e)$$

Using **Fact (2)** and (7) for $j=1$

$$(q, w_1 u, X_1 \beta) \vdash^{k(1)} (q_2, u, \beta) = (q_2, w_2 w_3 \dots w_{m+1}, X_2 X_3 \dots X_{m+1}) \dots \dots \dots (8)$$

where we chose $u = w_2 w_3 \dots w_{m+1}$ and $\beta = X_2 X_3 \dots X_{m+1}$ in applying **Fact (2)**

Now we can invoke the induction hypothesis since the depth of the stack is m and write

$$(q_2, w_2 w_3 \dots w_{m+1}, X_2 X_3 \dots X_{m+1}) \vdash^{n'} (p, e, e) \dots \dots \dots (9)$$

where $n' = k(2) + k(3) + \dots + k(m+1)$. Combining (8) and (9) we arrive at

$$(q, w_1 w_2 \dots w_{m+1}, X_1 X_2 \dots X_{m+1}) \vdash^{k(1)+n'} (p, e, e)$$

which proves **Fact (5)**.

The facts below show that the constructed equivalent CFG to a PDA is indeed equivalent !

For the given $PDA P = (Q, \Sigma, \Gamma, \delta, s, Z_0, F)$ define a CFG , $G = (V, \Sigma, R, S)$ where

$$V = \{S\} \cup \{[q X p] \mid p, q \in Q, X \in \Gamma\} \dots\dots\dots(10)$$

In addition to the initial productions namely

$$S \rightarrow [s Z_0 p]$$

where s is the initial state and p takes values over Q and we have the following productions in R :

whenever $(q_1, X_1 X_2 \dots X_m) \in \delta(q, a, X)$ there are productions

$$[q X p] \rightarrow a [q_1 X_1 q_2] \dots [q_m X_m p] \dots\dots\dots(11)$$

for all possible $q_2, \dots, q_m, p \in Q$ and if $(q_1, e) \in \delta(a, q, X)$

$[q X q_1] \rightarrow a$ is a production.

(Fact 6)

If $(q, w, X) \vdash_P^* (p, e, e)$ then $[q X p] \Rightarrow_G^* w$

Proof of Fact 6

Again we use induction on n . For $n=1$

$(q, w, X) \vdash (p, e, e)$ implies $w=a$ or $w=e$ and $(p, e) \in \delta(q, a, X)$ hence

$[q X p] \rightarrow a$ which proves the result.

Now given $(q, w, X) \vdash^{n+1} (p, e, e)$ in order to use induction on n we first evaluate the first step of the computation as below :

$(q, aw', X) \vdash (q_1, w', X_1 X_2 \dots X_m)$ where we assumed that

$(q_1, X_1 X_2 \dots X_m) \in \delta(q, a, X)$. Now by definition of the productions of G we have

$$[q X p] \rightarrow a [q_1 X_1 q_2] \dots [q_m X_m p] \dots\dots\dots(12)$$

are productions for all possible $q_2, q_3, \dots, q_m, p \in Q$. Since

$$(q_1, w', X_1 X_2 \dots X_m) \vdash^n (p, e, e)$$

which by **Fact 4** implies for $j = 1, \dots, m$ there are states q'_j for $j = 1, \dots, m+1$ where

$$q'_{m+1} := p, q'_1 = q_1 \text{ such that}$$

$$(q'_j, w'_j, X_j) \vdash^{k(j)} (q'_{j+1}, e, e) \text{ where each } k(j) < n+1 \text{ and by the induction hypothesis}$$

and choosing $q_j := q'_j$ for $j=1, 2, \dots, m$ and $p := q'_{m+1}$ in (12) $[q'_j X_j q'_{j+1}] \rightarrow w'_j$ and the result

follows by recalling that $w' = w'_1 w'_2 \dots w'_m$; $w = a w'$ and a *leftmost derivation* of (12) yields w .

(Fact 7)

If $[q X p] \Rightarrow_G^* w$ then $(q, w, X) \vdash_{P^*} (p, e, e)$

Proof of Fact 7

The proof is similar to that of **Fact 6**. This time we use induction on the steps of a leftmost derivation that yields w . For $n=1$ we have $[q X p] \Rightarrow w$ which implies that $w=a$ or $w=e$ and $(p, e) \in \delta(q, a, X)$ that further implies $(q, a, X) \vdash (p, e)$

Now we prove the case for $[q X p] \Rightarrow^{n+1} w$. We perform a first step of the derivation which in general yields

$$[q X p] \Rightarrow a [q_1 X_1 q_2] [q_2 X_2 q_3] \dots [q_m X_m p]$$

where q_2, \dots, q_m, p take values in Q and q_1 takes a value dictated by a transition

$$(q_1, X_1 X_2 \dots X_m) \in \delta(q, a, X)$$

Hence we have

$$(q, aw', X) \vdash (q_1, w', X_1 X_2 \dots X_m)$$

Now since the steps of derivation for each

$$[q_j X_j q_{j+1}] \Rightarrow w'_j, \quad j=1, \dots, m-1 \text{ and } [q_m X_m p] \Rightarrow w'_m$$

is necessarily less than $n+1$, the induction hypothesis can be applied which leads to

$$(q_j, w'_j, X_j) \vdash (q_{j+1}, e, e) \text{ for } j=1, \dots, m-1 \text{ and}$$

$$(q_m, w'_m, X_m) \vdash (p, e, e)$$

which by **Fact 4** implies that

$$(q_1, w'_1 w'_2 \dots w'_m, X_1 X_2 \dots X_m) \vdash (p, e, e) \text{ and therefore}$$

$$[q X p] \Rightarrow aw'_1 w'_2 \dots w'_m = w \text{ implies that}$$

$$(q, w, X) = (q, aw', X) \vdash (q_1, w', X_1 X_2 \dots X_m) \vdash (p, e, e)$$

which proves **Fact 7**

(Fact 8)

$$L_P = L_G$$

Proof of Fact 8

To prove that $L_P = L_G$ we first use **Fact 6** with $q=s$ (initial state) and $X = Z_0$ (initial stack) and use the *empty stack acceptance criterion* which proves that $L_P \subseteq L_G$. Conversely we use **Fact 7** after the first derivation step which is $S \Rightarrow [s Z_0 p]$ for an appropriate p which yields a string w , that is, $S \Rightarrow [s Z_0 p] \Rightarrow^* w$, and apply **Fact 7** which proves that $L_G \subseteq L_P$.

This concludes the proofs for all the **Facts**.

An Illustrative Example

Consider the single state *PDA*

$$P = (\{q\}, \{a,b\}, \{a,b,Z_0\}, \delta, q, Z_0, \{q\})$$

that accepts by empty stack the language

$L = \{\omega \in (a+b)^* \mid \#as = \#bs\}$ with the following δ - transitions :

$$(q, a, Z_0) \rightarrow (q, aZ_0) \dots\dots(1)$$

$$(q, b, Z_0) \rightarrow (q, bZ_0) \dots\dots(2)$$

$$(q, a, a) \rightarrow (q, aa) \dots\dots(3)$$

$$(q, b, b) \rightarrow (q, bb) \dots\dots(4)$$

$$(q, a, b) \rightarrow (q, e) \dots\dots(5)$$

$$(q, b, a) \rightarrow (q, e) \dots\dots(6)$$

$$(q, e, Z_0) \rightarrow (q, e) \dots\dots(7)$$

Using the definition justified by *Fact 1* to *Fact 8* above $G = (V, T, P, S)$ is

$$T = \{a, b\}$$

$$V = \{S, [q Z_0 q], [qaq], [qbq]\}$$

Letting $Z := [q Z_0 q]$, $A := [qaq]$, $B := [qbq]$

$$V = \{S, Z, A, B\}$$

and P becomes the set :

$$S \rightarrow Z ; Z \rightarrow a A Z \mid b B Z \mid e \text{ (using (1), (2) and (7))} ; A \rightarrow a A A \mid b \text{ (using (3) and (5))} ; B \rightarrow b B B \mid a \text{ (using (4) and (6))}$$