Homework #5 due November 22, 2016, Tuesday

(1) Let *M* be the *PDA* defined by $Q = \{q, q_0, q_1, q_2\}$, $\Sigma = \{a,b\}$, $\Gamma = \{a\}$, $F := \{q, q_1\}$.

$$\delta(q_0, a, Z_0) = \{(q, Z_0)\}$$

$$\delta(q, a, Z_0) = \{(q, aZ_0)\}$$

$$\delta(q, a, a) = \{(q, aa)\}$$

$$\delta(q, b, a) = \{(q_1, e)\}$$

$$\delta(q_1, b, a) = \{(q_1, e)\}$$

$$\delta(q_1, b, Z_0) = \{(q_2, e)\}$$

- a) Describe the language accepted by M.
- b) Trace all computations of the strings aab, abb, aba in M.
- c) Show that *aaabb*, $aaab \in L(M)$.
- (2) Construct PDAs that accept each of the following languages.

a)
$$\{a^{i}b^{j} \mid 0 \leq i \leq j\}$$

b)
$$\{a^{i}c^{j}b^{i}\mid i,j\geq 0\}$$

d)
$$\{a^{i}b^{j}c^{k} \mid i+k=j\}$$

$$e) \quad \{a^{\dot{i}}b^{\dot{j}} \mid 0 \le i \le j \le 2i\}$$

$$f(a^{i+j}b^ic^j | i,j > 0)$$

(3) $L = \{w \in \{a, b\}^* \mid at \ least \ one \ prefix \ of \ w \ contains \ strictly \ more \ b$'s than a's.}.

For example, baa, abb, abbbaa are in L, but aab, aabbab are not in L.

- a) Construct a PDA that accepts L by final state.
- **b)** Construct a PDA that accepts \boldsymbol{L} by empty stack.
- (4) From the main text Exercises 6.1.1 (b),(c), 6.2.6, 6.3.2