#### THE ALGEBRA OF REGULAR EXPRESSIONS

Reminder of Basic Definitions and Some Basic Proofs

(1) For languages L,  $M \subseteq \Sigma^*$ ; L+M, L.M and  $L^*$  are interpreted as follows:

$$L+M=L\cup M$$
;  $L.M=\{w\mid w=u.v;\ u\in L;v\in M\}$ ;  $L^*=\bigcup_{i=0,+\infty}L^i$  where  $L^i:=L.L.\ldots L$  (i times)

(2) 
$$(L+M)^* = (L^*. M^*)^*$$

## **Proof of (2):**

Let  $u \in (L+M)^*$  then by definition  $u = u_1 . u_2 . ... . u_k$  for some integer  $k \ge 0$  where for each j  $u_j \in L+M$ . But  $L \subseteq L^* \subseteq L^*$ .  $e \subseteq L^*$ .  $M^*$  and  $M \subseteq M^* \subseteq e$ .  $M^* \subseteq L^*$ .  $M^*$  and thus  $u_j \in L^*$ .  $M^* + L^*$ .  $M^* = L^*$ .  $M^*$  and therefore  $(L+M)^* \subseteq (L^* . M^*)^*$  Conversely let  $u \in (L^*.M^*)^*$  then by definition  $u = u_1 . u_2 . ... . u_k$  where  $u_j \in L^*.M^*$  hence  $u_j = v_j^{-1} . v_j^{-2} ... . v_j^{-1(j)} . w_j^{-1} . w_j^{-2} ... . w_j^{-p(j)}$  where  $v_j^{-m} \in L$  and  $w_j^{-m} \in M$ . Hence  $u = z_1 . z_2 ... . z_q$  where  $q = \sum_{j=1,k} l(j) + p(j)$  and each  $z_j \in L+M$ . This proves that  $(L^*.M^*)^* \subseteq (L+M)^*$  which proves that  $(L+M)^* = (L^*.M^*)^*$ 

(3) 
$$(L+M)* = (L*+M*)*$$

### **Proof of (3):**

Since  $L \subset L^*$  and  $M \subset M^*$  it follows that  $(L+M)^* \subset (L^*+M^*)^*$ .

Conversely let  $u \in (L^*+M^*)^*$  then  $u = (v_I+w_I)$ . ... .  $(v_k + w_k)$  where for each j  $v_j \in L^*$  and  $w_j \in M^*$ . We show that  $u \in (L^*, M^*)^*$  by using induction on k. For k=1  $v_I \in L^* \subseteq L^*$ .  $e \subseteq L^*.M^* \subseteq (L^*, M^*)^*$  similarly  $w_I \in M^* \subseteq e$ .  $M^* \subseteq L^*.M^* \subseteq (L^*, M^*)^*$  hence  $v_I+w_I \subseteq (L^*, M^*)^*$ . Now assume statement holds for k-1, hence  $z := (v_I+w_I)$ . ...  $(v_{k-1} + w_{k-1}) \in (L^*, M^*)^*$ . But using the above reasoning for  $v_I+w_I$  it follows that  $v_k+w_k \in (L^*, M^*)^*$  and therefore u = z.  $(v_k+w_k) \in (L^*, M^*)^*$ .  $(L^*, M^*)^* = (L^*, M^*)^*$  using the obvious identity  $K^*$ .  $K^* = K^*$  for any language K. This proves that  $(L^*+M^*)^* \subseteq (L^*, M^*)^*$ , but by (2)  $(L+M)^* = (L^*, M^*)^*$  hence  $(L^*+M^*)^* \subseteq (L+M)^*$  and (3) is proved.

(4)  $(L.M)^* \subseteq (L^*M^*)^*$  and  $(L.M)^* = (L^*M^*)^*$  iff  $e \in L$  and  $e \in M$ 

## **Proof of (4):**

First statement is obvious using  $L \subseteq L^*$  and  $M \subseteq M^*$ .

To prove the second one assume  $e \in L$  and  $e \in M$ 

and let  $u \in (L^*, M^*)$  then  $u = v_1, w_1, \dots, v_k, w_k$  where  $v_i \in L^*$  and  $w_i \in M^*$  therefore

 $vj = v_i^1 \dots v_i^{l(j)}$  and  $wj = z_i^1 \dots z_i^{p(j)}$  with  $v_i^m \in L$  and  $z_i^m \in M$ . Hence

 $u = q_1 \dots q_r$  where  $r = \sum_{j=1,k} (l(j) + p(j))$  where each  $q_i \in L$  or  $q_i \in M$ . Using the assumption we can write  $u = q'_1 \dots q'_r$  by adding an empty string in between the  $q_j$  strings ,if necessary, so that we have for each  $j=1, \dots, r'$ ,  $q'_j \in L$  and  $q'_{j+1} \in M$ . This proves that  $u \in (L.M)^*$  To prove the converse result we present counter-examples that violate the assumption  $e \in L$  and  $e \in M$ .

Suppose  $e \not\in L$  choose  $L = 0.0^*$  and  $M = 1^*$  then  $1 \in (L^*.M^*)^*$  whereas  $1 \not\in (L.M)^*$ ; alternatively if  $e \not\in M$  choose  $L = 0^*$  and  $M = 1.1^*$  then  $0 \in (L^*.M^*)^*$  whereas  $0 \not\in (L.M)^*$ .

# Homework #2 due March 25, Thursday 2021, before recitation

- (1) Using either the results or the techniques used above try to simplify the following expressions and prove your simplification.
- (i) (0+1)\*.1.(0+1) +(0+1) \*.1.(0+1)
- (ii) **(((0\*.1\*)+1)\*(0+1)\*)\***
- (iii) (L+M\*)\*
- (iv) **(L.M\*)\***
- (2) Convert the regular expression  $((0.0^*.(1.1)) + 0.1)^*$  into an  $\varepsilon$ -NFA
- (3) Problems from the textbook
- 3.1.1 (b) and (c)
- 3.1.4 (b) and (c)
- 3.2.1 (c) and (d)
- 3.2.3