

**Fall 2017**

**SABANCI UNIVERSITY**  
**Faculty of Engineering and Natural Sciences**  
**CS 302 Automata Theory**

***Final Examination***

***Closed (Book+Notes+All Electronic Devices)***

***Duration : 150 minutes***

<b><i>Q1</i></b>	
<b><i>Q2</i></b>	
<b><i>Q3</i></b>	
<b><i>Q4</i></b>	
<b><i>Total</i></b>	

**Question 1** (25 points)

- (a) (8 pts) State the full definition of a *nondeterministic finite automaton (NFA)*,  $A = (Q, \Sigma, \delta, s, F)$  including the *extended transition function*  $\delta^E$  and the condition for a string  $u \in \Sigma^*$  to be accepted by  $A$  (in your definition the domain and range of the transition functions  $\delta$  and  $\delta^E$  should be explicitly and clearly specified).
- (b) (5pts) Construct an *NFA*  $A$  with at most 3 states that accepts the *regular expression*  $E = (1+0)^*.1.(1+0)$  interpreted as a language.
- (c) (6 pts) Construct a *DFA*  $B$  that accepts the language  $E$  in part (b).
- (d) (6 pts) Construct a *minimal state DFA*  $C$  that accepts the language  $E$  in part (b).

**Question 2** (25 points)

- (a) (9 pts) Specify, clearly and with all the details, a *context-free grammar* (CFG),  $G = (V, T, R, S)$  that generates the set of *all regular expressions* over some finite set  $\Sigma$ .
- (b) (8 pts) Write down a regular expression  $E$  over the set  $\{0,1\}$  corresponding to the language  $L$  where for each string  $u$  in  $L$ ,  $010$  OR  $101$  is *substring* of  $u$ .
- (c) (8 pts) Sketch an *NFA* that accepts the language  $L$  in part (b) above .

**Question 3** (25 points)

(a) (7 pts) State clearly the *definition* of a *PDA*  $P$  accepting a language  $L$  by *final state* ( $L(P)$ ) ; and by *empty stack* ( $N(P)$ ) .

(b) (10 pts) Compute a *PDA*  $P$  that accepts the language

$$L = \{\omega \in \{0+1\}^* \mid \#1s = \#0s \text{ in } \omega\}$$

(c) (8 pts) Is your *PDA* a deterministic *PDA* (*DPDA*) ? If not try to modify it so that it is a *DPDA*.

**Question 4** (25 points)

Consider the language  $L = (\omega \in \{a,b,c\}^* \mid \omega = a^{k+1}b^k c^{k-1} ; k > 0)$

(a) (15 pts) Choosing a appropriate  $\omega \in L$  ( if you do not show explicitly your choice of  $\omega$  you get NO credit ! ) and using the *pumping lemma* show that  $L$  is NOT a context-free language.

(b) (10 pts) Compute a 2-tape NDTM TM  $M$  that decides  $L = (\omega \in \Sigma \mid \omega = uvu , u,v \in \Sigma^*, u \neq \epsilon)$