

Fall 2020

18.1.2021

SABANCI UNIVERSITY
Faculty of Engineering and Natural Sciences
CS 302 Automata Theory

Remote Final Examination

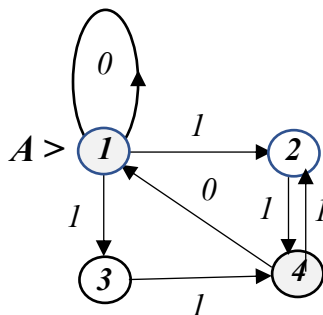
Closed (Book+Notes+All Electronic Devices)

Duration: 180 minutes

Question 1 (25 points)

(a) (5 pts) State the full definition of a **Nondeterministic Finite Automaton (NFA)** ; in particular the **domain** and **range** of the transition function δ and the extended transition function δ^E and the condition for an NFA to accept a string $u \in \Sigma^*$ in terms of δ^E and the set of final states F .

(b) (3 pts) State whether the NFA A below is a DFA with your reasons. **1** and **4** are the final states.



(c) (12 pts) Compute a **minimal state DFA C** that is equivalent to the NFA A in part (b).

(d) (5 pts) Try to express the language accepted by A above in simple natural language.

Question 2 (25 points)

Consider the CFG, $G = (V, T, R, S)$ where

$V = \{A, B, S\}$, $T = \{a, b\}$ and R is given by the following productions:

$S \rightarrow AB$

$A \rightarrow aaAb \mid e$

$B \rightarrow bbBa \mid a$

(a) (7 pts) Write down the CFL L_G generated by G above.

(b) (8 pts) Can the strings: (i) $u_1 = a^4b^2a$ and (ii) $u_2 = a^4b^4a^2$ be generated by G ? If so draw the corresponding **parse tree** for these strings.

(c) (10 pts) Construct a PDA $P = (Q, \Sigma, \Gamma, \delta, q_0, Z_0, F)$ that accepts the language L_G .

Question 3 (25 points)

(a) (15 pts) For the languages :

$L_1 = \{a^n b^{n+k} c^n; n, k \geq 0\}$ and $L_2 = \{a^n b^k c^n; n, k \geq 0\}$

prove whether each is a CFL or not by using either the **pumping lemma** or a CFG in **Chomsky Normal Form** that generates the language in question.

(b) (10 pts) For the right linear CFG $G = (\{A, B, S\}, \{0, 1\}, R, S)$ where R is given by :

$S \rightarrow 0A \mid 1B$; $A \rightarrow 0B \mid 1A \mid e$; $B \rightarrow 1S \mid 1$

draw an NFA X that accepts the language L_G generated by G and then convert your X to a **minimal state DFA** Y .

Question 4 (25 points)

(a) (10 pts) State the definition of a **Deterministic Turing Machine (DTM)** to: (i) **decide** a language ; (ii) **semidecide** a language ; and (iii) **compute** a function $f: \Sigma_0^* \rightarrow \Sigma_0^*$.

(b) (7 pts) Construct a single tape DTM M either in graphical or tabular form that decides the language $L = \{0^n 1^n; n \geq 0\}$. Assume initial configuration as $(s, \# w)$

(c) (8 pts) Repeat part (b) when $L = \{0^n 1^{2n}; n \geq 0\}$.