

## THE ALGEBRA OF REGULAR EXPRESSIONS

### *Reminder of Basic Definitions and Some Basic Proofs*

(1) For languages  $L, M \subseteq \Sigma^*$ ;  $L+M$ ,  $L.M$  and  $L^*$  are interpreted as follows :

$L+M = L \cup M$ ;  $L.M = \{w \mid w = u.v; u \in L; v \in M\}$ ;  $L^* = \bigcup_{i=0,+\infty} L^i$  where  $L^i := L.L \dots .L$  ( $i$  times)

(2)  $(L+M)^* = (L^*.M^*)^*$

#### ***Proof of (2):***

Let  $u \in (L+M)^*$  then by definition  $u = u_1.u_2 \dots .u_k$  for some integer  $k \geq 0$  where for each  $j$   $u_j \in L+M$ . But  $L \subseteq L^* \subseteq L^*.e \subseteq L^*.M^*$  and  $M \subseteq M^* \subseteq e.M^* \subseteq L^*.M^*$  and thus  $u_j \in L^*.M^* + L^*.M^* = L^*.M^*$  and therefore  $(L+M)^* \subseteq (L^*.M^*)^*$

Conversely let  $u \in (L^*.M^*)^*$  then by definition  $u = u_1.u_2 \dots .u_k$  where  $u_j \in L^*.M^*$  hence

$u_j = v_j^1.v_j^2 \dots .v_j^{l(j)}.w_j^1.w_j^2 \dots .w_j^{p(j)}$  where  $v_j^m \in L$  and  $w_j^m \in M$ . Hence

$u = z_1.z_2 \dots .z_q$  where  $q = \sum_{j=1,k} l(j)+p(j)$  and each  $z_j \in L+M$ . This proves that  $(L^*.M^*)^* \subseteq (L+M)^*$  which proves that  $(L+M)^* = (L^*.M^*)^*$

(3)  $(L+M)^* = (L^*+M^*)^*$

#### ***Proof of (3):***

Since  $L \subseteq L^*$  and  $M \subseteq M^*$  it follows that  $(L+M)^* \subseteq (L^*+M^*)^*$ .

Conversely let  $u \in (L^*+M^*)^*$  then  $u = (v_1+w_1). \dots .(v_k+w_k)$  where for each  $j$

$v_j \in L^*$  and  $w_j \in M^*$ . We show that  $u \in (L^*.M^*)^*$  by using induction on  $k$ . For  $k=1$

$v_1 \in L^* \subseteq L^*.e \subseteq L^*.M^* \subseteq (L^*.M^*)^*$  similarly  $w_1 \in M^* \subseteq e.M^* \subseteq L^*.M^* \subseteq (L^*.M^*)^*$  hence  $v_1+w_1 \in (L^*.M^*)^*$ . Now assume statement holds for  $k-1$ , hence

$z := (v_1+w_1). \dots .(v_{k-1}+w_{k-1}) \in (L^*.M^*)^*$ . But using the above reasoning for  $v_k+w_k$  it follows that

$v_k+w_k \in (L^*.M^*)^*$  and therefore  $u = z.(v_k+w_k) \in (L^*.M^*)^*$ .  $(L^*.M^*)^* = (L^*.M^*)^*$  using the

obvious identity  $K^*.K^* = K^*$  for any language  $K$ . This proves that  $(L^*+M^*)^* \subseteq (L^*.M^*)^*$ , but by (2)

$(L+M)^* = (L^*.M^*)^*$  hence  $(L^*+M^*)^* \subseteq (L+M)^*$  and (3) is proved.

(4)  $(L.M)^* \subseteq (L^*M^*)^*$  and  $(L.M)^* = (L^*M^*)^*$  iff  $e \in L$  and  $e \in M$

#### ***Proof of (4):***

First statement is obvious using  $L \subseteq L^*$  and  $M \subseteq M^*$ .

To prove the second one assume  $e \in L$  and  $e \in M$

and let  $u \in (L^* \cdot M^*)^*$  then  $u = v_1 \cdot w_1 \dots \cdot v_k \cdot w_k$  where  $v_j \in L^*$  and  $w_j \in M^*$  therefore

$v_j = y_j^1 \cdot \dots \cdot y_j^{l(j)}$  and  $w_j = z_j^1 \cdot \dots \cdot z_j^{p(j)}$  with  $y_j^m \in L$  and  $z_j^m \in M$ . Hence

$u = q_1 \cdot \dots \cdot q_r$  where  $r = \sum_{j=1,k} (l(j) + p(j))$  where each  $q_i \in L$  or  $q_i \in M$ . Using the assumption we can write  $u = q'_1 \cdot \dots \cdot q'_{r'}$  by adding an empty string in between the  $q_j$  strings, if necessary, so that we have for each  $j=1, \dots, r'$ ,  $q'_j \in L$  and  $q'_{j+1} \in M$ . This proves that  $u \in (L.M)^*$ . To prove the converse result we present counter-examples that violate the assumption  $e \in L$  and  $e \in M$ .

Suppose  $e \notin L$  choose  $L = 0.0^*$  and  $M = 1^*$  then  $1 \in (L^*.M^*)^*$  whereas  $1 \notin (L.M)^*$ ; alternatively if  $e \notin M$  choose  $L = 0^*$  and  $M = 1.1^*$  then  $0 \in (L^*.M^*)^*$  whereas  $0 \notin (L.M)^*$ .

## **Homework #2**

**(1)** Using either the results or the techniques used above try to simplify the following expressions and prove your simplification.

(i)  $(0+1)^*.1.(0+1) + (0+1)^*.1.(0+1)$

(ii)  $((0^*.1^*)+1)^*(0+1)^*$

(iii)  $(L+M^*)^*$

(iv)  $(L.M^*)^*$

**(2)** Convert the regular expression  $((0.0^*.1)^* + 0.1)^*$  into an  $\epsilon$ -NFA

**(3) Problems from the textbook**

3.1.1 (b) and (c)

3.1.4 (b) and (c)

3.2.1 (c) and (d)

3.2.3