

# Database Management Systems (CSN-351)

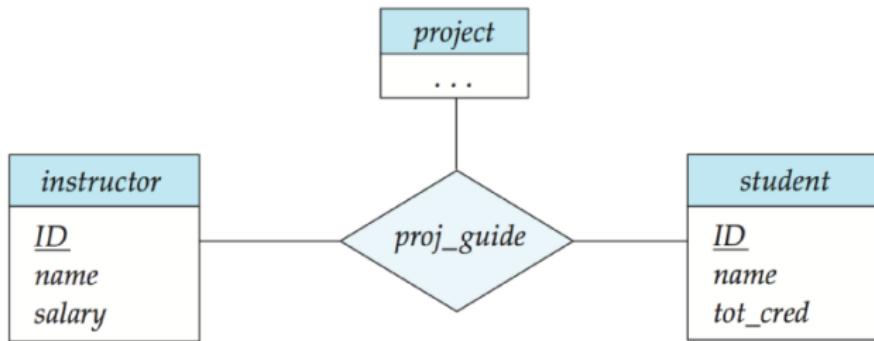
## Entity Relationship Model (contd.)

**BTech 3rd Year (CS) + Minor + Audit**

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# E-R Diagram with a Ternary Relationship



# Cardinality Constraints on Ternary Relationship

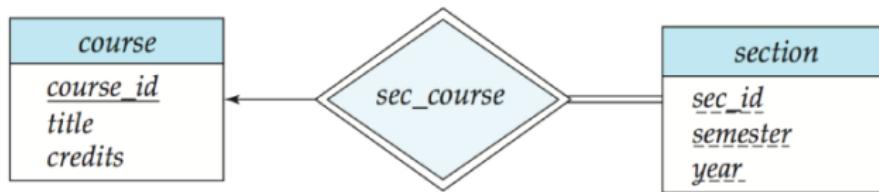
- At most one arrow is allowed out of a ternary (or greater degree) relationship to indicate a cardinality constraint.
- If there are more than one arrow, there are multiple ways of defining the meaning.

# Weak Entity Sets

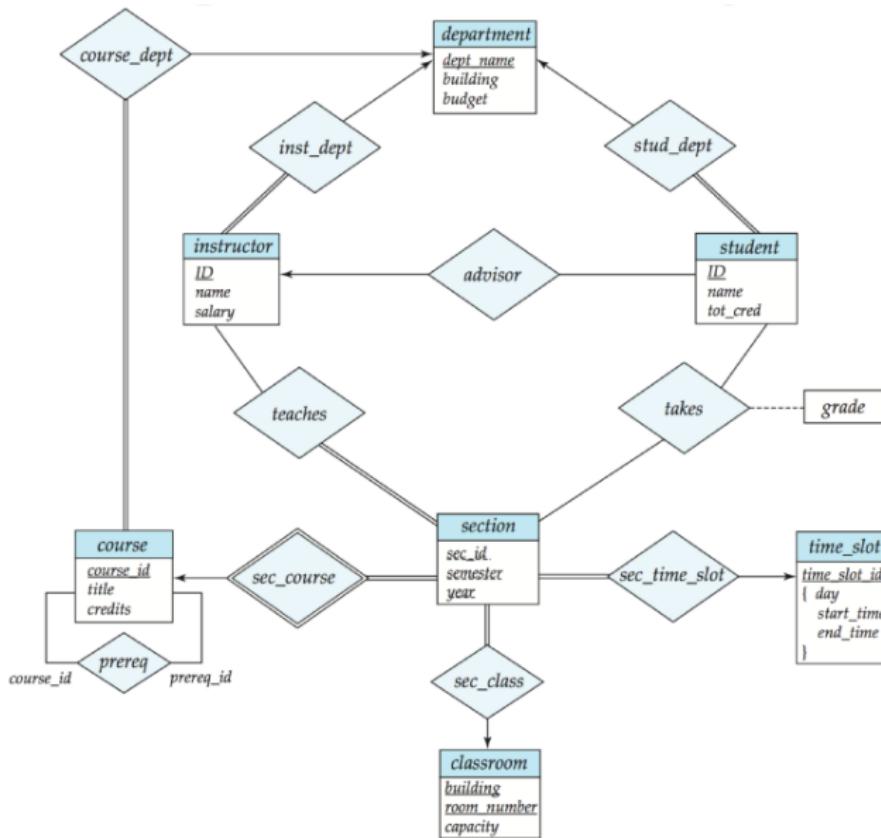
- An entity set that does not have a *primary key* is referred to as a weak entity set.
- The existence of a weak entity set depends on the existence of an *identifying entity set*.
  - It must relate to the identifying entity set via a *total, one-to-many relationship* set from the identifying to the weak entity set.
  - Identifying relationship depicted using a *double diamond*.

# Weak Entity Sets (contd.)

- The *discriminator* (or *partial key*) of a weak entity set is the set of attributes that distinguishes among all the entities of a weak entity set that depend on one particular strong entity.
- The *primary key of a weak entity set* is formed by the primary key of the strong entity set on which the weak entity set is existence dependent, plus the weak entity set's discriminator.



# E-R Diagram for a University

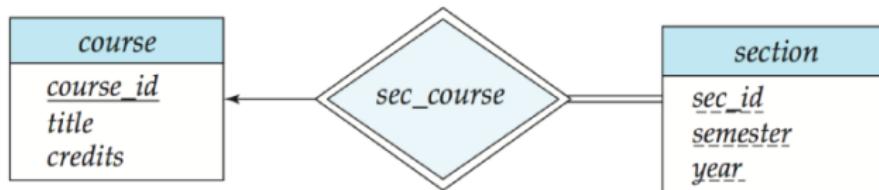


# Reduction to Relation Schemas

- Entity sets and relationship sets can be expressed uniformly as *relation schemas* that represent the contents of the database.
- A database which conforms to an E-R diagram can be represented by *a collection of schemas*.
- For each entity set and relationship set there is a *unique schema* that is assigned the name of the corresponding entity set or relationship set.
- Each schema has a number of *columns* (generally corresponding to attributes), which have unique names.

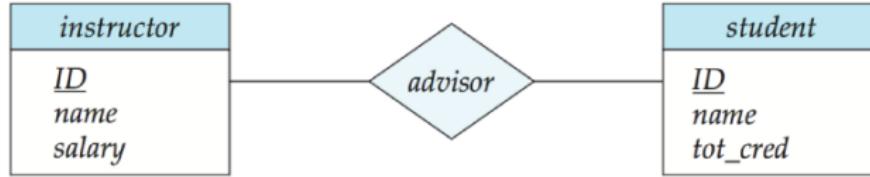
# Representing Entity Sets With Simple Attributes

- A strong entity set reduces to a schema with the same attributes  
*student*(ID, name, tot\_cred)
- A weak entity set becomes a table that includes a column for the primary key of the identifying strong entity set  
*section*(course\_id, sec\_id, sem, year)



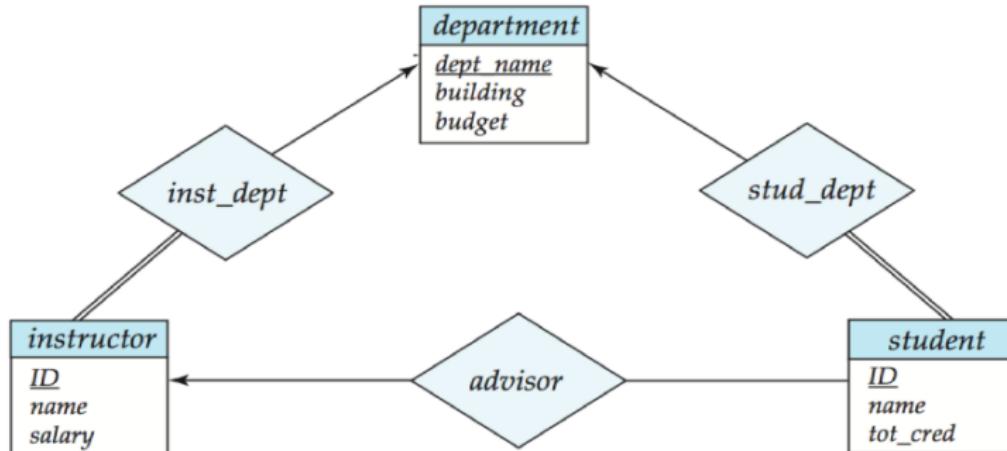
# Representing Relationship Sets

- A many-to-many relationship set is represented as a schema with attributes for the primary keys of the two participating entity sets, and any descriptive attributes of the relationship set.
- Example: schema for relationship set advisor  
 $advisor(s\_id, i\_id)$



# Redundancy of Schemas

- Many-to-one and one-to-many relationship sets that are total on the many-side can be represented by adding an extra attribute to the *many* side, containing the primary key of the *one* side
- Example: Instead of creating a schema for relationship set *inst\_dept*, add an attribute *dept\_name* to the schema arising from entity set *instructor*.



## Redundancy of Schemas (contd.)

- For one-to-one relationship sets, either side can be chosen to act as the *many* side
- If participation is partial on the *many* side, replacing a schema by an extra attribute in the schema corresponding to the *many* side could result in null values
- The schema corresponding to a relationship set linking a weak entity set to its identifying strong entity set is redundant.
  - Example: The *section* schema already contains the attributes that would appear in the *sec\_course* schema

# Composite Attributes

<i>instructor</i>
<u>ID</u>
<i>name</i>
<i>first_name</i>
<i>middle_initial</i>
<i>last_name</i>
<i>address</i>
<i>street</i>
<i>street_number</i>
<i>street_name</i>
<i>apt_number</i>
<i>city</i>
<i>state</i>
<i>zip</i>
{ <i>phone_number</i> }
<i>date_of_birth</i>
<i>age ()</i>

- Composite attributes are flattened out by creating a separate attribute for each component attribute
- Example: given entity set *instructor* with composite attribute *name* with component attributes *first\_name* and *last\_name* the schema corresponding to the entity set has two attributes *name\_first\_name* and *name\_last\_name*
- Ignoring multivalued attributes, extended *instructor* schema is  

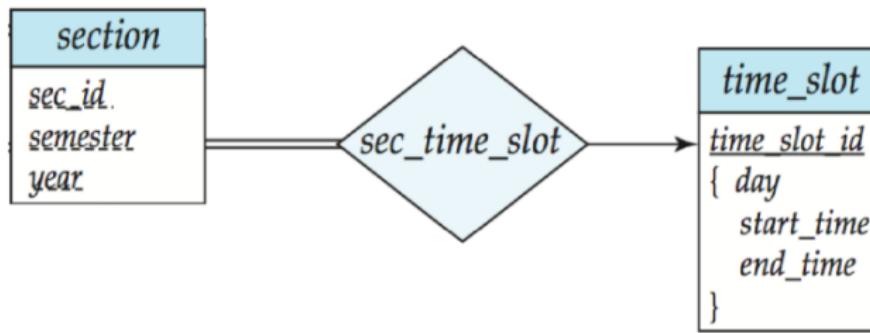
$$\textit{instructor}(\underline{\textit{ID}}, \textit{first\_name}, \textit{middle\_initial}, \textit{last\_name}, \textit{street\_number}, \textit{street\_name}, \textit{apt\_number}, \textit{city}, \textit{state}, \textit{zip\_code}, \textit{date\_of\_birth})$$

# Multivalued Attributes

- A multivalued attribute M of an entity E is represented by a separate schema EM.
- Schema EM has attributes corresponding to the primary key of E and an attribute corresponding to multivalued attribute M.
- Example: Multivalued attribute *phone\_number* of instructor is represented by a schema: *inst\_phone*(ID, phone\_number)
- Each value of the multivalued attribute maps to a separate tuple of the relation on schema EM.
  - For example, an instructor entity with primary key 22222 and phone numbers 456-7890 and 123-4567 maps to two tuples: (22222, 456-7890) and (22222, 123-4567)

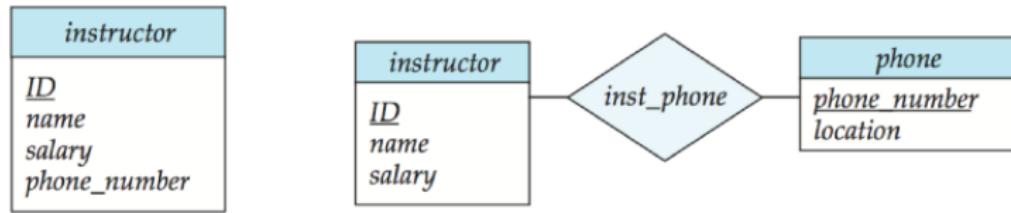
# Multivalued Attributes (contd.)

- Special case: entity *time\_slot* has only one attribute other than the primary-key attribute, and that attribute is multivalued.
- Optimization: Don't create the relation corresponding to the entity, just create the one corresponding to the multivalued attribute
  - time\_slot*(time\_slot\_id, day, start\_time, end\_time)
- Limitation: *time\_slot* attribute of section (from *sec\_time\_slot*) cannot be a foreign key due to this optimization



# Design Issues

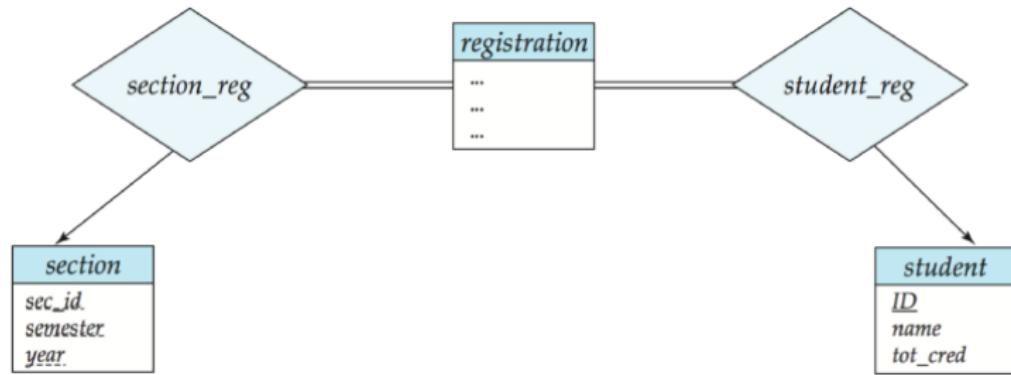
- Use of entity sets vs. attributes



- Use of phone as an entity allows extra information about phone numbers (plus multiple phone numbers)

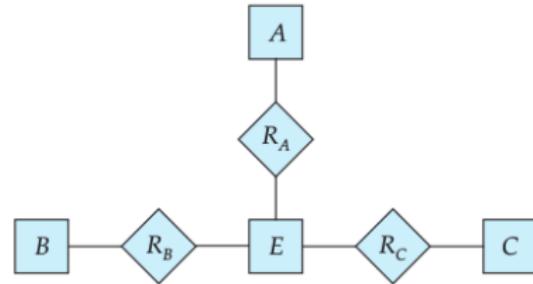
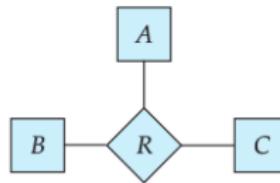
# Design Issues

- Use of entity sets vs. relationship sets



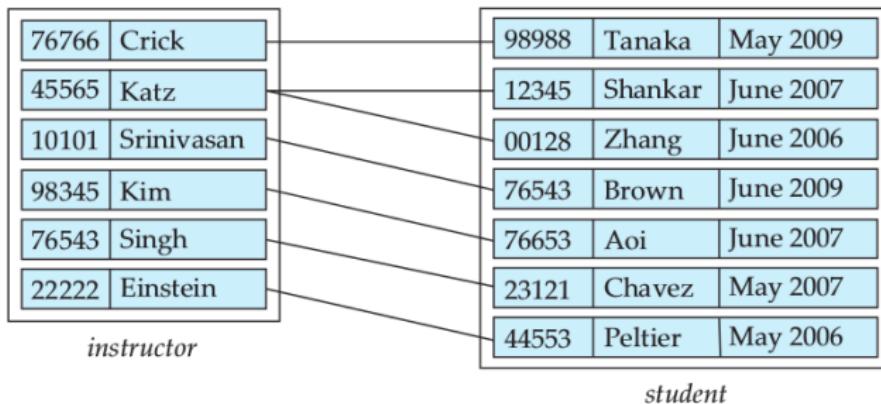
# Design Issues

- Binary versus  $n$ -ary relationship sets



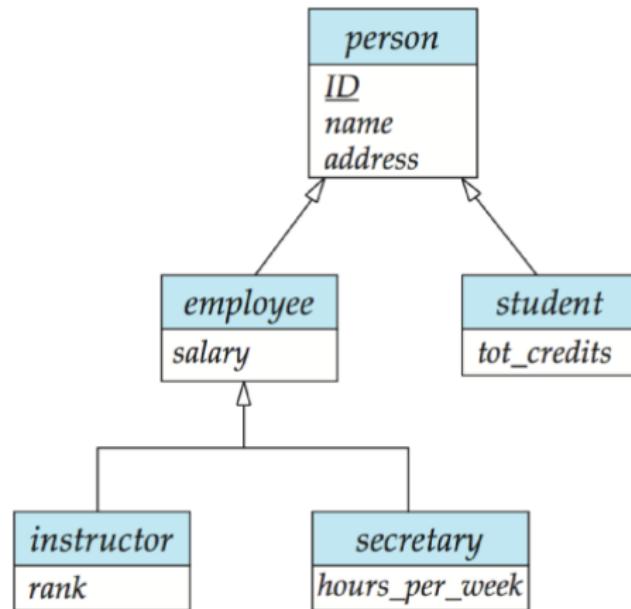
# Design Issues

- Placement of relationship attributes



# Extended E-R Features: Specialization

- Example of Specialization



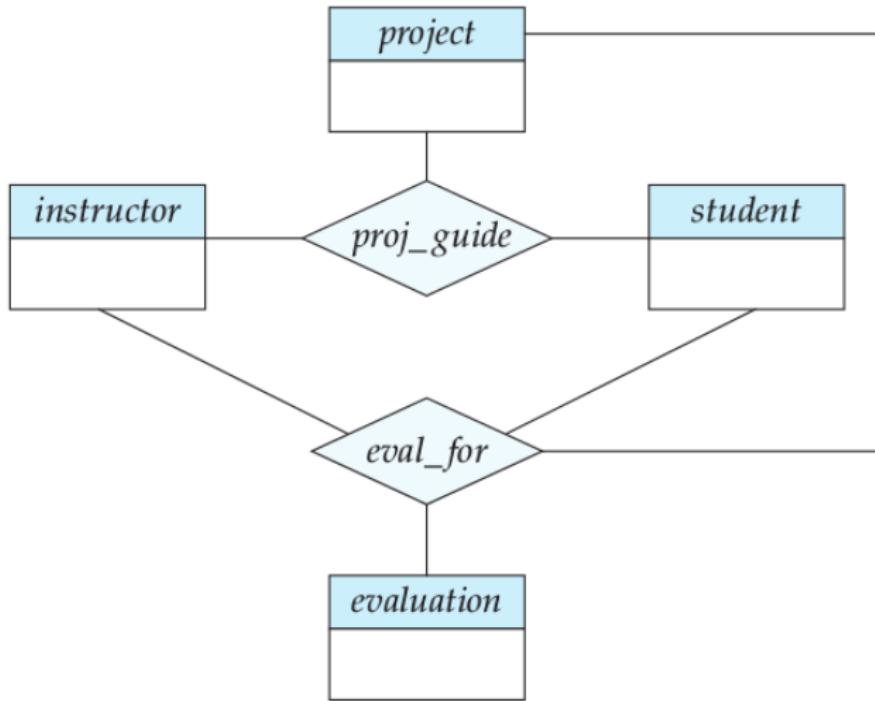
# Specialization (contd.)

- Top-down design process; we designate subgroupings within an entity set that are distinctive from other entities in the set.
- These subgroupings become lower-level entity sets that have attributes or participate in relationships that do not apply to the higher-level entity set.
- Depicted by a triangle component labeled ISA (e.g., instructor “is a” person).
- *Attribute inheritance* — a lower-level entity set inherits all the attributes and relationship participation of the higher-level entity set to which it is linked.

# Extended E-R Features: Generalization

- A bottom-up design process — combine a number of entity sets that share the same features into a higher-level entity set.
- Specialization and generalization are simple inversions of each other; they are represented in an E-R diagram in the same way.
- The terms specialization and generalization are used interchangeably.

# Extended E-R Features: Aggregation



# Aggregation (contd.)

