

# **Topics in Applied Optimization**

Optimization for ML and Data Sciences

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Dr. Pawan Kumar

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IIIT, Hyderabad

## Convex Optimization Problem: Local Optima = Global Optima

**Fact:** For a convex optimization problem, any local optima is a global optima

**Ans:** Proof on chalkboard!

# Convex Optimization Problem: Optimality Criteria

$$f(y) \geq f(x^*) + \gamma$$

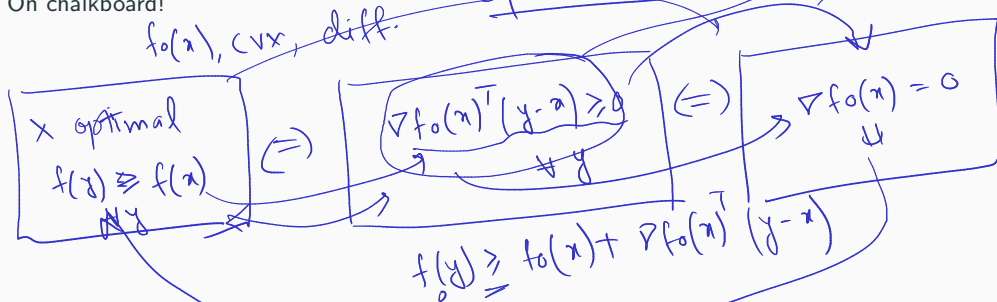
leaf

**Fact:** If  $f_0$  in a convex optimization problem is **differentiable**, then the point  $x$  is **optimal** if

$$\nabla f_0(x)^T (y - x) \geq 0 \quad \text{for all } y \in X$$

**Proof:** On chalkboard!

Equivalent Cases:



# Convex Optimization Problem: Optimality for Unconstrained problems

Fact: For an unconstrained problem, ( $m = p = 0$ ), the optimality condition

$$\nabla f_0(x)^T (y - x) \geq 0 \quad \text{for all } y \in X$$

$$f(y) \leq f(x^*)$$

reduces to the well known necessary and sufficient condition

$$\nabla f_0(x) = 0$$

Proof on chalkboard!

Suppose  $x$  is optimal  $\Rightarrow x \in \text{dom } f$  for all feasible  $y$  we have  $\nabla f_0(x^*)^T (y - x) \geq 0$

$$\nabla f_0(x)^T (y - x) \geq 0$$

Let us take  $y = x + \alpha d$

$$\begin{aligned} f(y) &\geq f(x) + \nabla f(x)^T (y - x) \\ \Rightarrow f(y) &\geq f(x) + \nabla f(x)^T (\alpha d) \\ &= f(x) + \alpha \nabla f(x)^T d \\ &= f(x) + \alpha \cdot 0 = f(x) \end{aligned}$$

$\Rightarrow x^*$  is optimal

## Scratch Space

So we prove that if  $\exists x^*$  st  $\boxed{\nabla f(x^*) = 0} \Rightarrow x^*$  is optimal

Again  $\Rightarrow \nabla f_0(x)^T (y-x) \geq 0$

(Claim: Assume that  $\nabla f_0(x)^T (y-x) \geq 0 \quad \forall y$ , then

$$\Rightarrow f_0(y) \geq f_0(x) + \underbrace{\nabla f_0(x)^T (y-x)}_{\geq 0} \leftarrow (\text{Convexity})$$

$$\Rightarrow \boxed{f_0(y) \geq f_0(x)} \Rightarrow x \text{ is } \underline{\text{optimal}}$$

Conversely. Suppose  $x$  is optimal and on contrary assume  
that  $\nabla f_0(x)^T (y-x) < 0 \rightarrow \textcircled{*}$

## Scratch Space

Consider that point:  $z(t) = ty + (1-t)x$ ,  $t \in [0,1]$

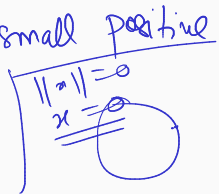
$z(t)$  is on the line joining  $x$  &  $y$  &  $x, y \in$  feasible set which is conv.  $\oplus$

$\Rightarrow z(t) \in$  feasible set.

$$\left. \frac{d}{dt} f_0(z(t)) \right|_{t=0} = \nabla f_0(x)^T z'(t) \Big|_{t=0} = \nabla f_0(x)^T (y-x) < 0$$

$\Rightarrow f_0(z(t))$  decreases as  $t$  increases.  $\Rightarrow \exists$  small positive  $t$  for which  $f_0(z(t)) < f_0(z(0)) = f_0(x)$

A contradiction to the fact that  $x$  was optimal



Take  $y = \underbrace{x - t \nabla f_0(x)}_{\substack{\nabla f_0(x)=0 \\ \text{if } x \text{ is optimal}}}$ . For small  $t$ , positive  $y$  is feasible

$$\nabla f_0(x)^T (y-x) = \nabla f_0(x)^T (-t \nabla f_0(x)) = -t \|\nabla f_0(x)\|_2^2 \geq 0$$

# Unconstrained Quadratic Optimization

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Consider the problem of minimizing

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where  $P \in S_+^n$ .

## Unconstrained Quadratic Optimization

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where  $P \in S_+^n$ . The **necessary and sufficient** condition for  $x$  to be minimizer of  $f_0$  is

$$\nabla f_0(x) = \underbrace{P x + q}_{=0} = 0. \quad (\text{Why?})$$

$\checkmark f_0$  is diff. (poly in  $x$ )  
① Is  $f_0$  convex?

$$\nabla f_0(x) = P x + q$$

$$\nabla^2 f_0(x) = P \succeq 0$$

$\boxed{P \succeq 0} \Rightarrow f_0$  is conv.

## Unconstrained Quadratic Optimization

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**Several** cases may occur:

1. If  $\underline{q \neq \mathcal{R}(P)}$ , then there is **no** solution.  $\underline{f_0}$  is unbounded below

$Pa = -q \leftarrow$  Can never be satisfied

$$\begin{bmatrix} p_1 & p_2 & \dots & p_n \\ \vdots & \vdots & & \vdots \end{bmatrix} \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix} = x_1 p_1 + x_2 p_2 + \dots + x_n p_n = \underline{-q}$$

$\exists q \in \mathcal{R}(P)$  if  $\underline{p_i}$ 's

# Unconstrained Quadratic Optimization

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1. If  $q \notin \mathcal{R}(P)$ , then there is **no** solution.  $f_0$  is unbounded below

2. If  $P \succ 0$ , then there is a **unique** minimizer,  $x^* = -P^{-1}q$

$\Rightarrow P \in S_{++}^n \leftarrow \text{strictly p.d.} \Rightarrow \text{non-singularity.}$

# Unconstrained Quadratic Optimization

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**Several** cases may occur:

1. If  $q \notin \mathcal{R}(P)$ , then there is **no** solution.  $f_0$  is unbounded below
2. If  $P > 0$ , then there is a **unique** minimizer,  $x^* = -P^{-1}q$
3. If  $P$  is **singular**, but  $q \in \mathcal{R}(P)$ .

# Unconstrained Quadratic Optimization

Consider the problem of minimizing

$$f_0(x) = (1/2)x^T Px + q^T x + c$$

where  $P \in S_+^n$ . The **necessary and sufficient** condition for  $x$  to be minimizer of  $f_0$  is

$$\nabla f_0(x) = Px + q = 0. \quad (\text{Why?})$$

Several cases may occur:

1. If  $q \notin \mathcal{R}(P)$ , then there is **no** solution.  $f_0$  is unbounded below
2. If  $P > 0$ , then there is a **unique** minimizer,  $x^* = -P^{-1}q$
3. If  $P$  is **singular**, but  $q \in \mathcal{R}(P)$ .

- set of **optimal** points is  $X_{\text{opt}} = -P^+q + \mathcal{N}(P)$ , where  $P^+$  denotes the pseudoinverse of  $P$

Handwritten notes and calculations:

$$Q = \begin{bmatrix} x_1^2 - 2x_1x_2 + x_2^2 \end{bmatrix} = (x_1 - x_2)^2 \geq 0$$

$$Q = 0 \Rightarrow x_1 = x_2$$

$$\begin{bmatrix} x_1 & x_2 \end{bmatrix} \begin{bmatrix} 1 & -1 \\ -1 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$$

$$\begin{pmatrix} P \\ \hline \end{pmatrix}$$

$$\lambda_1 P \geq 0 \Rightarrow P \in S_+^2$$

Handwritten note:

$$x \neq 0$$

no sol<sup>n</sup>  
unique  
infinitely

## Problems with Equality Constraints



## Problems with Equality Constraints

Consider the following optimization problem with **equality constraints only**

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & Ax = b\end{array}$$

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Consider the following optimization problem with **equality constraints only**

minimize  $f_0(x)$

subject to  $Ax = b$

- Feasible set is **affine**, recall that  $x$  is **feasible** if it satisfies

$$\nabla f_0(x)^T (y - x) \geq 0, \quad \text{for all } y \text{ such that } Ay = b$$

*along*

$$f(y) \geq f(x) + \dots$$

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- Every **feasible**  $y$  has the form:  $y = \underbrace{x}_{\downarrow} + \underbrace{v}_{\downarrow}$ , (**Why?**) for some  $v \in \underline{\mathcal{N}(A)}$

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- Every feasible  $y$  has the form:  $y = x + v$ , (Why?) for some  $v \in \mathcal{N}(A)$
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## Problems with Equality Constraints

$g$

Consider the following optimization problem with **equality constraints only**

$$\begin{aligned} &\text{minimize} && f_0(x) \\ &\text{subject to} && Ax = b \end{aligned}$$

$g(v) = c^T v > 0$   
+  $v \in S$

- Feasible set is **affine**, recall that  $x$  is **feasible** if it satisfies

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- Optimality condition** is

$$\nabla f_0(x)^T v \geq 0 \quad \text{for all } v \in \mathcal{N}(A)$$

- $\nabla f_0(x)^T v$  is **linear** in  $v$ , and **nonnegative** on a subspace, hence

$g(v) = \nabla f_0(x)^T v$

Linear in  $v$

$g(v) \geq 0 \rightarrow v \in S$

$\Rightarrow g(-v) = -g(v) \geq 0 \rightarrow \text{circle}$

$\Rightarrow g(v) \leq 0 \rightarrow \text{circle}$

$\Rightarrow g(v) = 0$

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  - $\nabla f_0(x)^T v = 0$  for all  $v \in \mathcal{N}(A)$ . (**Why?**)

## Problems with Equality Constraints

Consider the following optimization problem with **equality constraints only**

$$\begin{aligned} &\text{minimize } f_0(x) \\ &\text{subject to } Ax = b \end{aligned}$$

$$x \in S = \{x \mid Ax = b\}$$

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$$\nabla f_0(x)^T v = 0 \text{ for all } v \in \mathcal{N}(A). \text{ (Why?)}$$

$$\nabla f_0(x) \perp \mathcal{N}(A) \implies \nabla f_0(x) \in \mathcal{R}(A^T), \text{ (Why?) i.e., } \exists v \text{ such that}$$

$$\begin{aligned} &\nabla f_0(x)^T v = 0 \quad \forall v \in \mathcal{N}(A) \\ &\Rightarrow \nabla f_0(x) \perp \mathcal{N}(A) \\ &\mathcal{N}(A) \oplus \mathcal{R}(A^T) \\ &x = v + v^T \end{aligned}$$



# Problems with Equality Constraints

$$\frac{1}{2} p^T p + \frac{1}{2} x^T x + \frac{1}{2} y^T y$$

$$\begin{bmatrix} p \\ x \\ y \end{bmatrix} = \begin{bmatrix} c \\ b \end{bmatrix}$$

Consider the following optimization problem with **equality constraints only**

$$\begin{aligned} & \text{minimize} && f_0(x) \\ & \text{subject to} && Ax = b \end{aligned}$$

$$\nabla f_0(x) = p + A^T \lambda$$

$$p + A^T \lambda = -c$$

$$p + A^T \lambda = b$$

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  - $\nabla f_0(x)^T v = 0$  for all  $v \in \mathcal{N}(A)$ . (**Why?**)
  - $\nabla f_0(x) \perp \mathcal{N}(A) \implies \nabla f_0(x) \in \mathcal{R}(A^T)$ , (**Why?**) i.e.,  $\exists v$  such that
 
$$\nabla f_0(x) + A^T v = 0$$
- The above together with  $Ax = b$ . **Lagrange condition**.

$$\nabla f_0(x) \in \mathcal{R}(A^T)$$

$$\implies A^T v = \nabla f_0(x)$$

$$\nabla f_0(x) + A^T v = 0$$

$$\begin{bmatrix} \nabla f_0(x) + A^T v = 0 \\ Ax = b \end{bmatrix}$$