Queues

- Linear list.
- One end is called front.
- Other end is called rear.
- Additions are done at the rear only.
- Removals are made from the front only.













rear











front

rear













front

rear



The Interface Queue

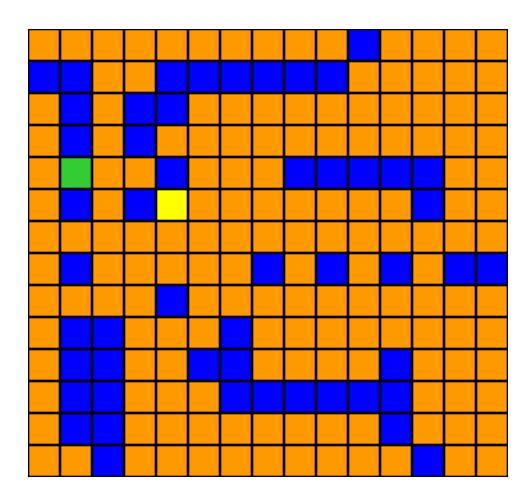
```
public interface Queue
public boolean isEmpty();
public Object getFrontEelement();
public Object getRearEelement();
public void put(Object theObject);
public Object remove();
```

Revisit Of Stack Applications

- Applications in which the stack cannot be replaced with a queue.
 - Parentheses matching.
 - Towers of Hanoi.
 - Method invocation and return.
 - Try-catch-throw implementation.

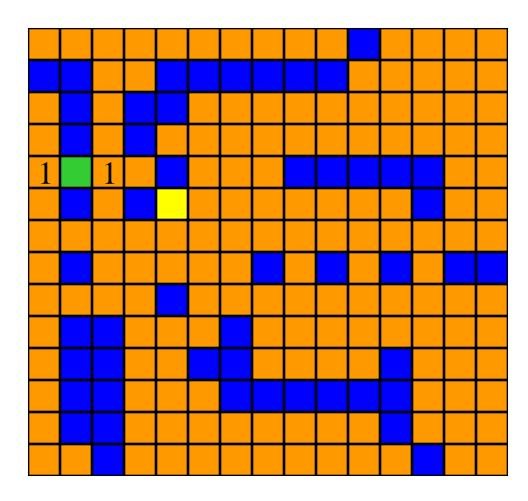
- Application in which the stack may be replaced with a queue.
 - Rat in a maze.
 - Results in finding shortest path to exit.

- start pin
- end pin



Label all reachable squares 1 unit from start.

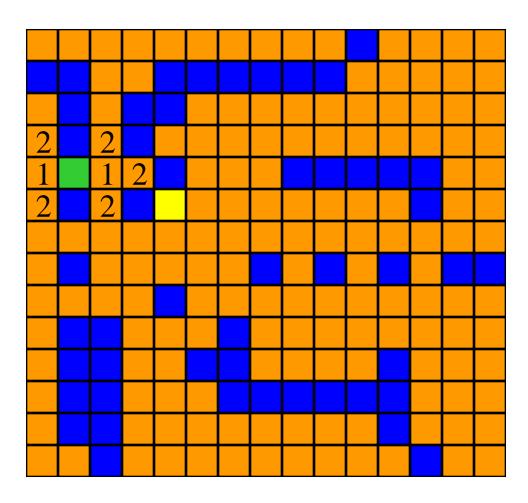
- start pin
- end pin



Label all reachable unlabeled squares 2 units from start.

start pin

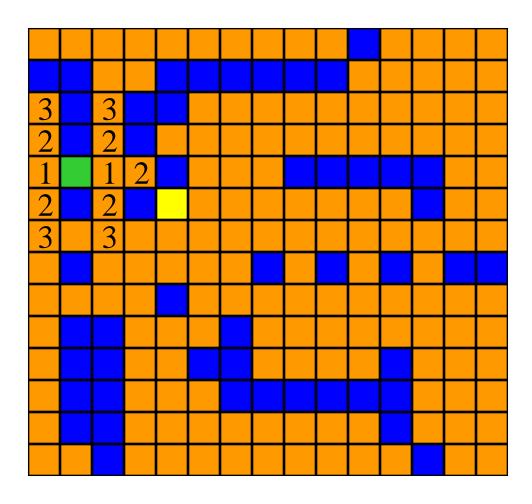
end pin



Label all reachable unlabeled squares 3 units from start.

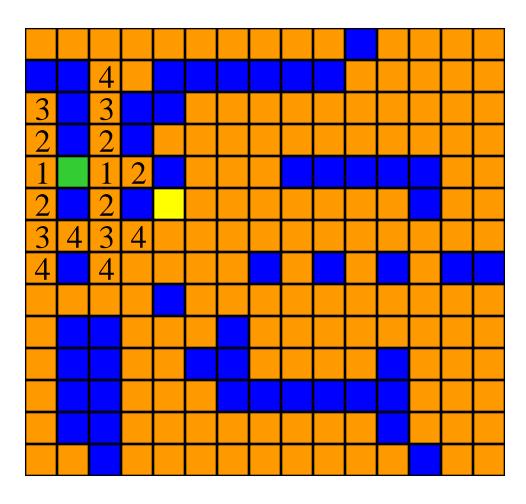
start pin

end pin



Label all reachable unlabeled squares 4 units from start.

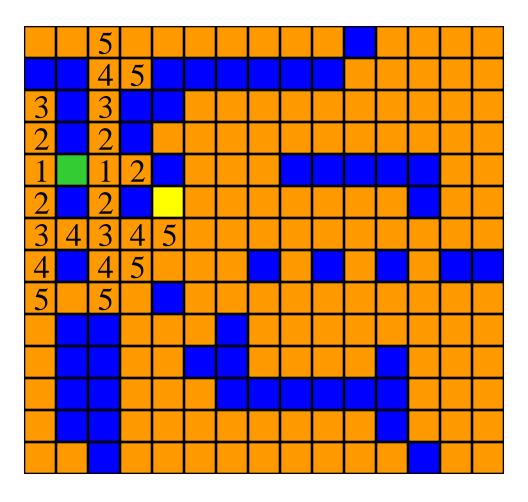
- start pin
- end pin



Label all reachable unlabeled squares 5 units from start.

start pin

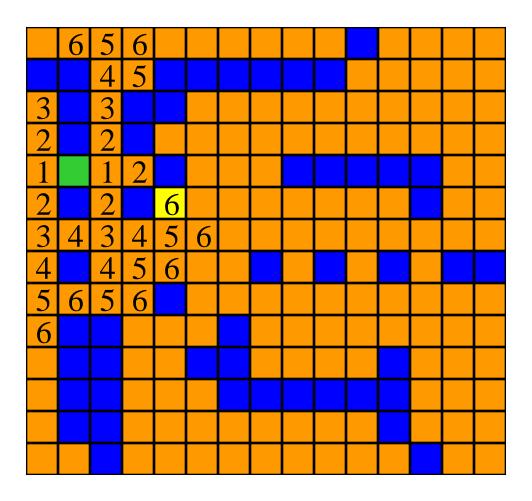
end pin



Label all reachable unlabeled squares 6 units from start.

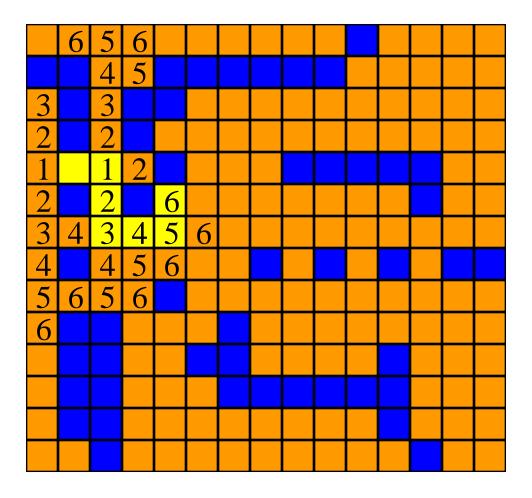
start pin

end pin



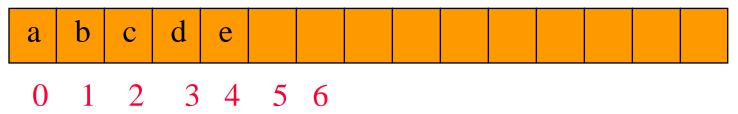
End pin reached. Traceback.

- start pin
- end pin



End pin reached. Traceback.

Derive From ArrayLinearList



- > when front is left end of list and rear is right end
 - Queue.isEmpty() => super.isEmpty()
 - -O(1) time
 - getFrontElement() => get(0)
 - -O(1) time
 - getRearElement() => get(size() 1)
 - -O(1) time
 - put(theObject) => add(size(), theObject)
 - -O(1) time
 - $remove() \Rightarrow remove(0)$
 - O(size) time

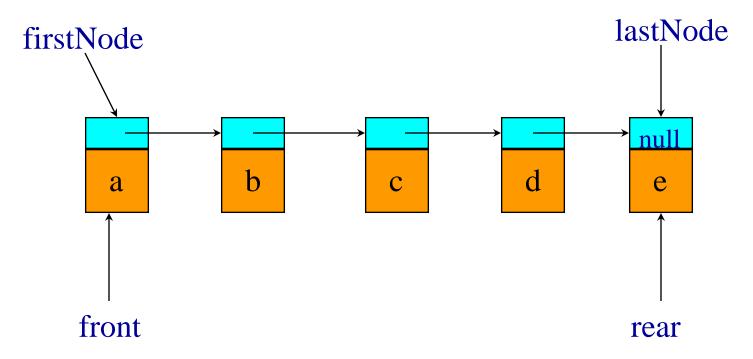
Derive From ArrayLinearList



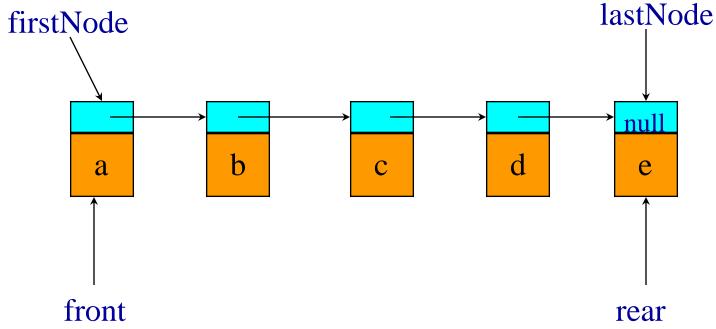
- when rear is left end of list and front is right end
 - Queue.isEmpty() => super.isEmpty()
 - -O(1) time
 - getFrontElement() => get(size() 1)
 - -O(1) time
 - getRearElement() => get(0)
 - -O(1) time
 - put(theObject) => add(0, theObject)
 - O(size) time
 - remove() => remove(size() 1)
 - -O(1) time

Derive From ArrayLinearList

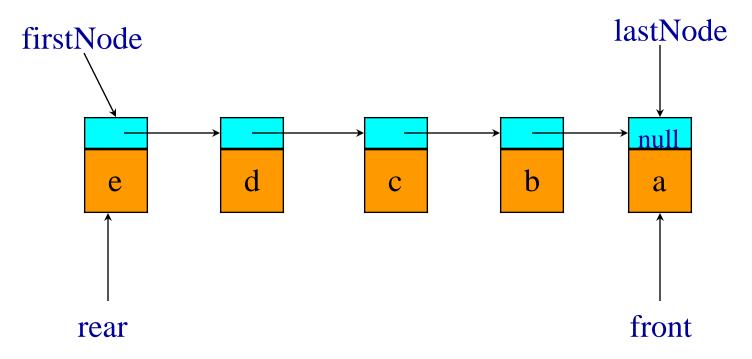
• to perform each opertion in O(1) time (excluding array doubling), we need a customized array representation.



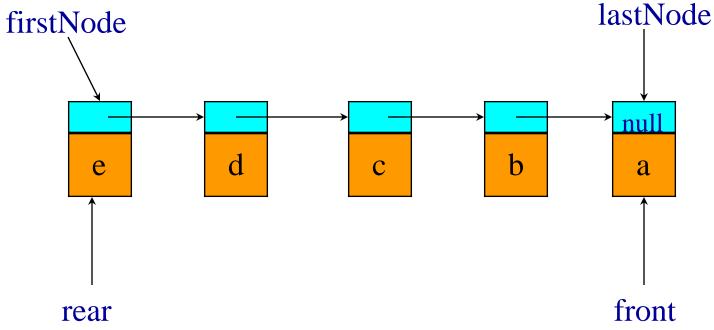
- > when front is left end of list and rear is right end
 - Queue.isEmpty() => super.isEmpty()
 - -O(1) time
 - getFrontElement() => get(0)
 - -O(1) time



- getRearElement() => getLast() ... new method
 - -O(1) time
- put(theObject) => append(theObject)
 - -O(1) time
- remove() \Rightarrow remove(0)
 - -O(1) time



- > when front is right end of list and rear is left end
 - Queue.isEmpty() => super.isEmpty()
 - -O(1) time
 - getFrontElement() => getLast()
 - -O(1) time

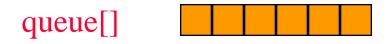


- getRearElement() => get(0)
 - -O(1) time
- put(theObject) => add(0, theObject)
 - -O(1) time
- remove() => remove(size-1)
 - O(size) time

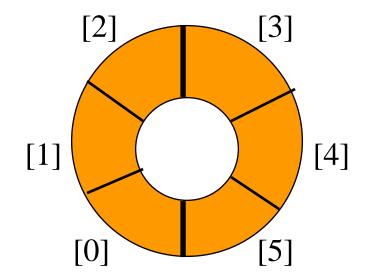
Custom Linked Code

 Develop a linked class for Queue from scratch to get better performance than obtainable by deriving from Linked List.

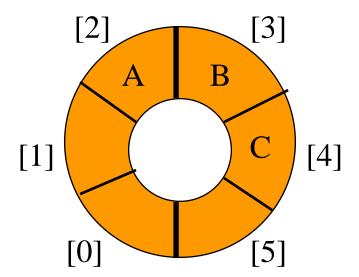
• Use a 1D array queue.



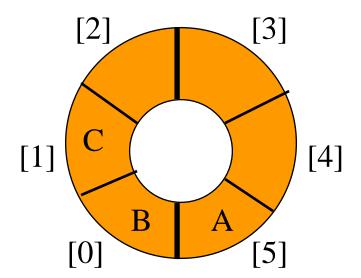
Circular view of array.



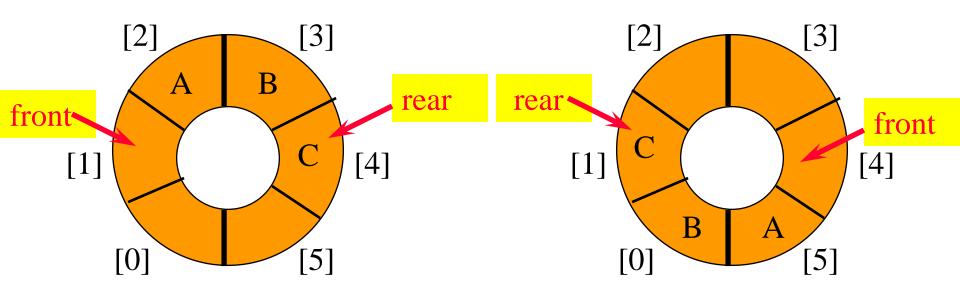
• Possible configuration with 3 elements.



• Another possible configuration with 3 elements.

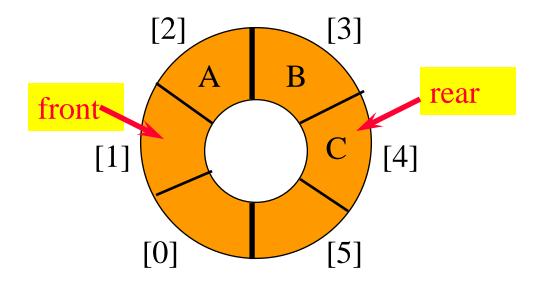


- Use integer variables front and rear.
 - front is one position counterclockwise from first element
 - rear gives position of last element



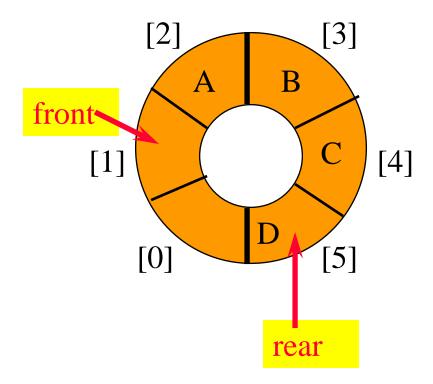
Add An Element

• Move rear one clockwise.



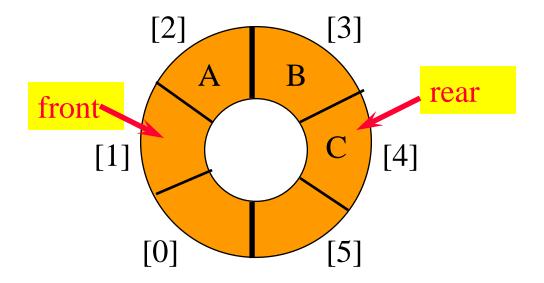
Add An Element

- Move rear one clockwise.
- Then put into queue[rear].



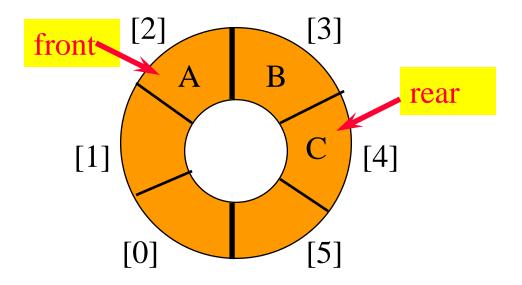
Remove An Element

Move front one clockwise.



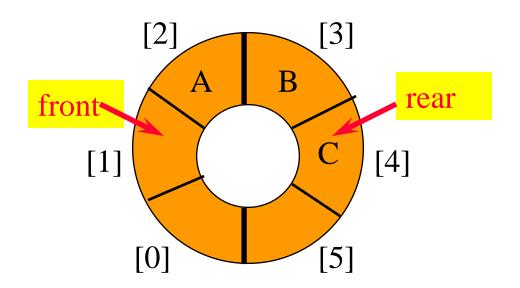
Remove An Element

- Move front one clockwise.
- Then extract from queue[front].

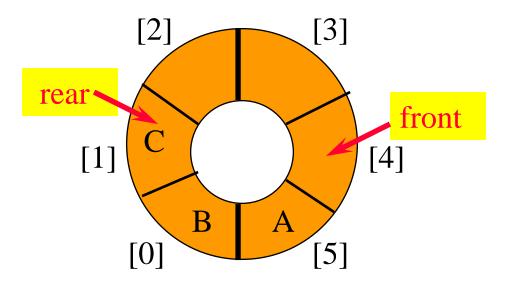


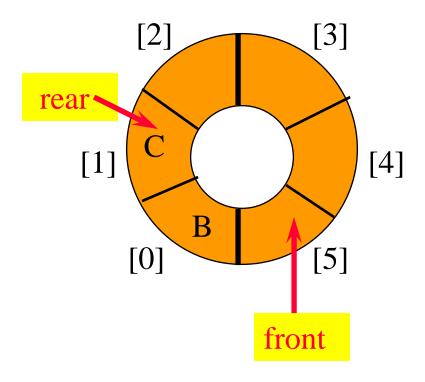
Moving rear Clockwise

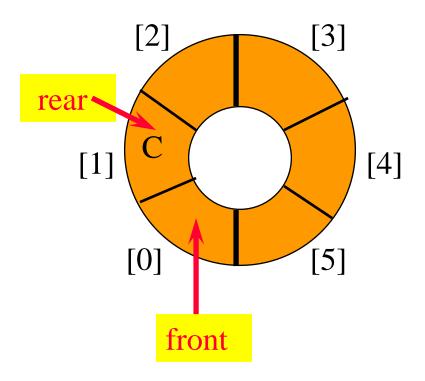
rear++;if (rear = = queue.length) rear = 0;

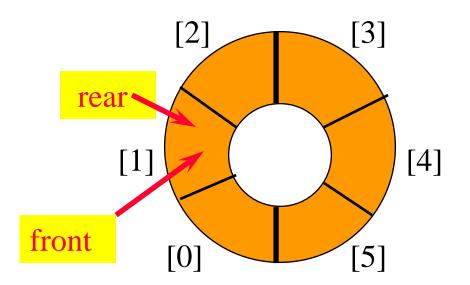


• rear = (rear + 1) % queue.length;

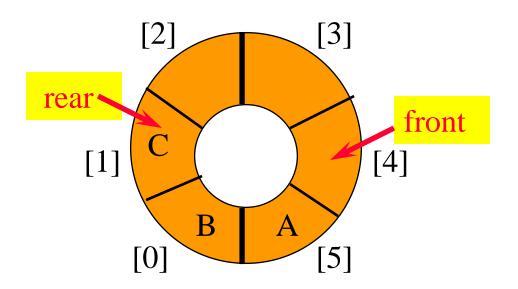


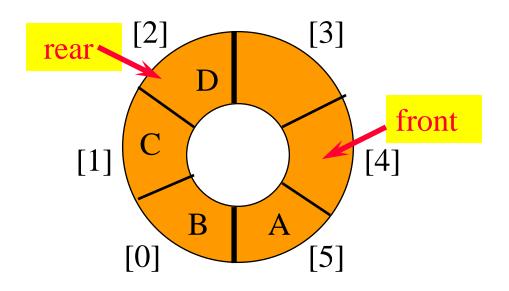


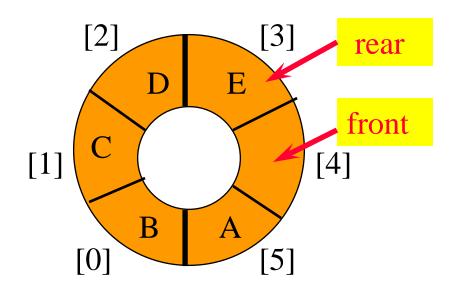


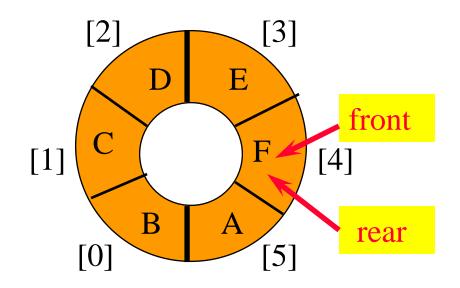


- When a series of removes causes the queue to become empty, front = rear.
- When a queue is constructed, it is empty.
- So initialize front = rear = 0.









- When a series of adds causes the queue to become full, front = rear.
- So we cannot distinguish between a full queue and an empty queue!

Ouch!!!!!

- Remedies.
 - Don't let the queue get full.
 - When the addition of an element will cause the queue to be full, increase array size.
 - This is what the text does.
 - Define a boolean variable lastOperationIsPut.
 - Following each put set this variable to true.
 - Following each remove set to false.
 - Queue is empty iff (front == rear) && !lastOperationIsPut
 - Queue is full iff (front == rear) && lastOperationIsPut

Ouch!!!!!

- Remedies (continued).
 - Define an integer variable size.
 - Following each put do size++.
 - Following each remove do size---.
 - Queue is empty iff (size == 0)
 - Queue is full iff (size == queue.length)
 - Performance is slightly better when first strategy is used.