

R-B Trees

Three Properties of a Red-Black Tree

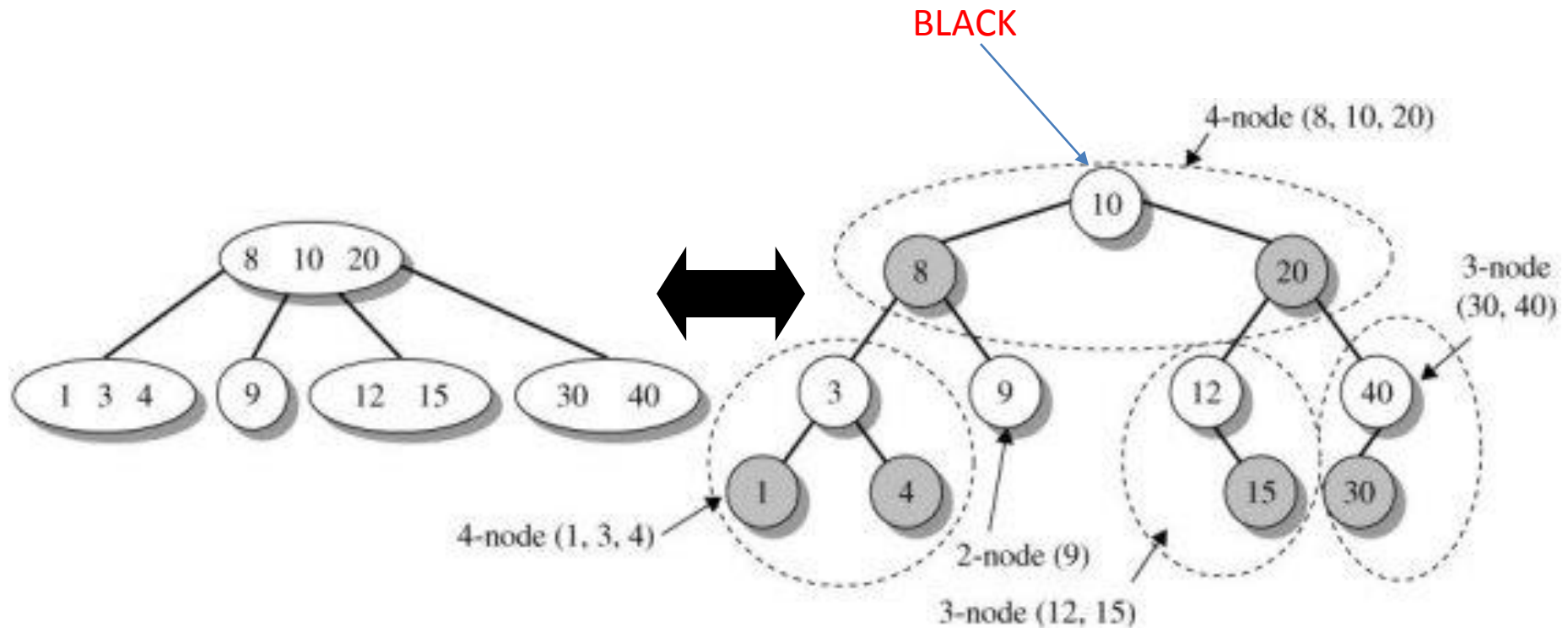
that must always be true for the tree to be red-black

- 1. The root must always be BLACK
(white in our pictures)
- 2. A RED parent never has a RED child
 - in other words: there are never two successive RED nodes in a path

continued

- 3. Every path from a node to an null leaf (node) contains the same number of BLACK nodes
 - called the *black height*
- We can use black height to measure the balance of a red-black tree.

Check the Example Properties



The black-height of the red-black tree is 2. Each path from the root contains exactly two BLACK nodes.

Three things
to do.

Inserting a Node

1. Search down the tree to find the insertion point
2. Once the insertion point is found, add the new item as a RED *leaf node*
 - this may create two successive RED nodes
 - again use rotation and recoloring to reorder/rebalance the tree
3. Keep the root as a BLACK node.

Inserting a New node

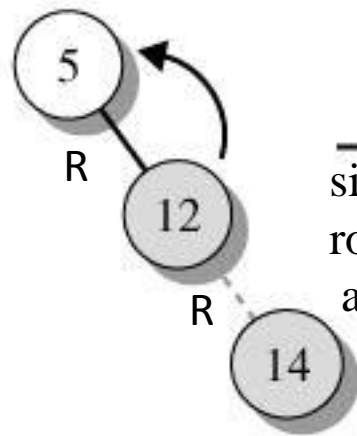
- Always add a new item to a tree as a RED leaf node
 - this may create two successive RED nodes, which breaks property 2
 - Assumption: Null is treated as Black node

RB-INSERT(T, x)

```
1  TREE-INSERT( $T, x$ )
2   $color[x] \leftarrow \text{RED}$ 
3  while  $x \neq \text{root}[T]$  and  $color[p[x]] = \text{RED}$ 
4      do if  $p[x] = \text{left}[p[p[x]]]$ 
5          then  $y \leftarrow \text{right}[p[p[x]]]$ 
6              if  $color[y] = \text{RED}$ 
7                  then  $color[p[x]] \leftarrow \text{BLACK}$  ▷ Case 1
8                       $color[y] \leftarrow \text{BLACK}$  ▷ Case 1
9                       $color[p[p[x]]] \leftarrow \text{RED}$  ▷ Case 1
10                      $x \leftarrow p[p[x]]$  ▷ Case 1
11                 else if  $x = \text{right}[p[x]]$ 
12                     then  $x \leftarrow p[x]$  ▷ Case 2
13                     LEFT-ROTATE( $T, x$ ) ▷ Case 2
14                      $color[p[x]] \leftarrow \text{BLACK}$  ▷ Case 3
15                      $color[p[p[x]]] \leftarrow \text{RED}$  ▷ Case 3
16                     RIGHT-ROTATE( $T, p[p[x]]$ ) ▷ Case 3
17                 else (same as then clause
                        with “right” and “left” exchanged)
18   $color[\text{root}[T]] \leftarrow \text{BLACK}$ 
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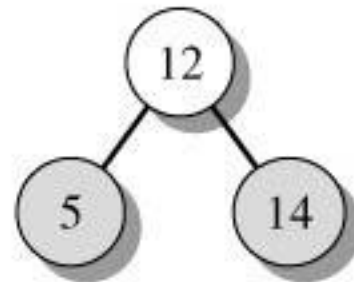
Example: Insert 14

Red-black tree after
inserting 14



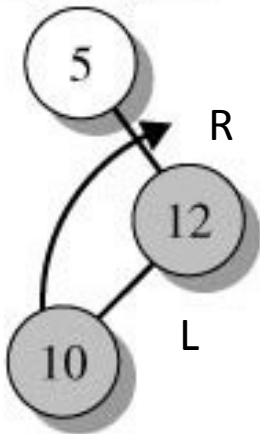
single left
rotation of 12
and color flip

Red-black tree after
single left rotation



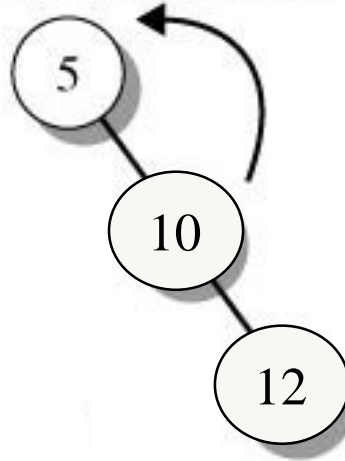
Insert 10 instead of 14

Red-black tree after inserting 10



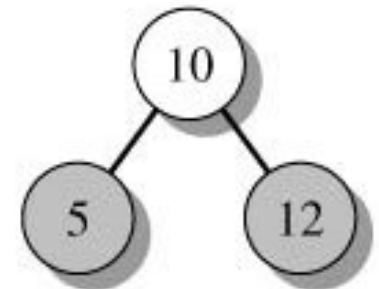
single right rotation of 10

Red-black tree after a single right rotation about 10



single left rotation and color flip

Red-black tree after a single left rotation about 10 and recoloring



RL = double left rotation of node 10
(right then left)

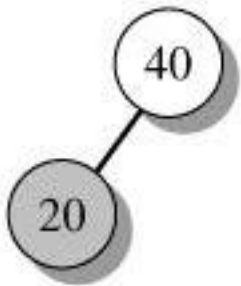
Building a Red-Black Tree



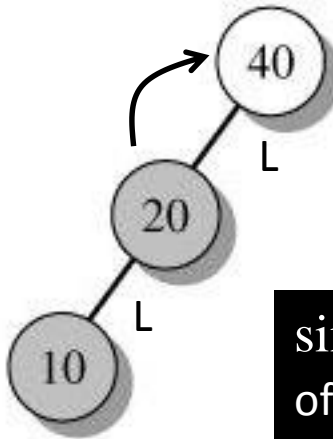
Insert 40
as a RED node



As the root node
make it BLACK

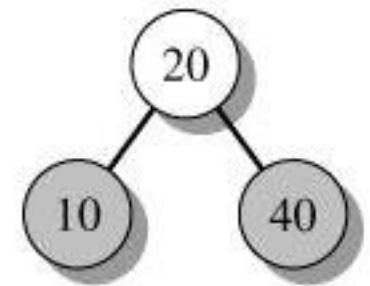


Insert 20

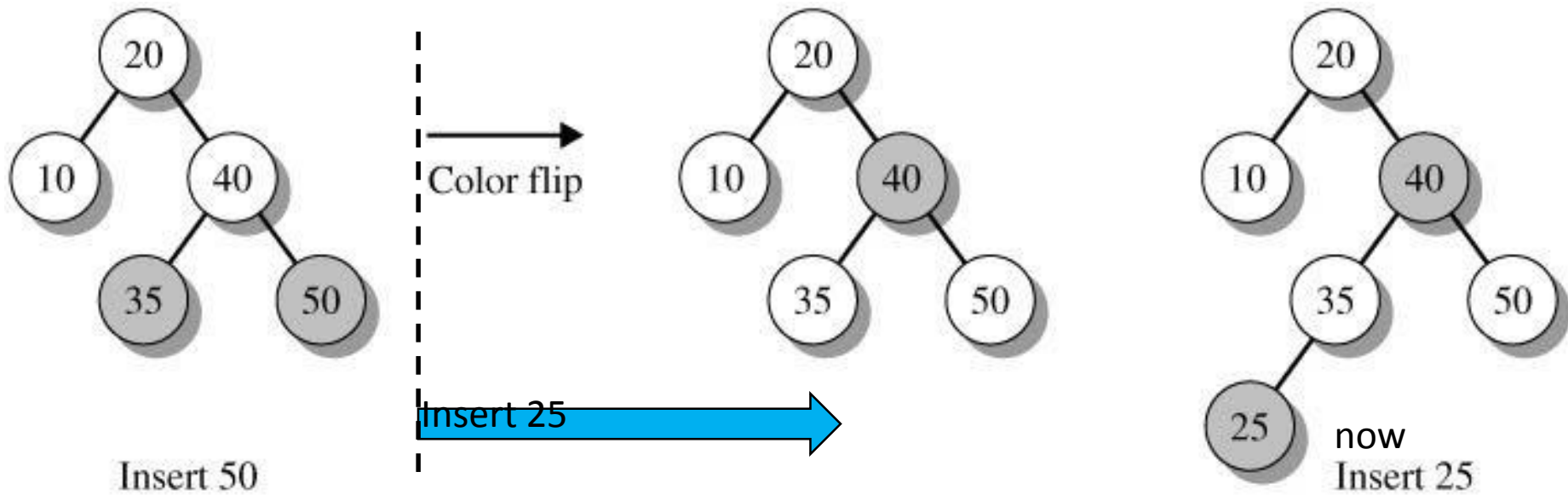
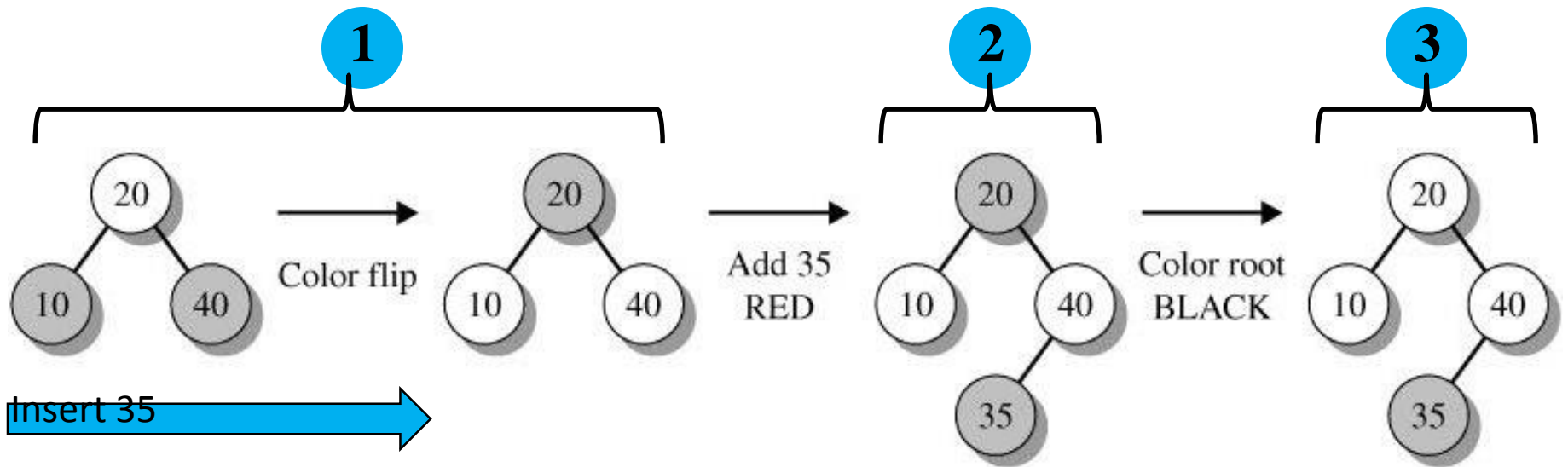


Insert 10

single right rotation
of 20 and color flip

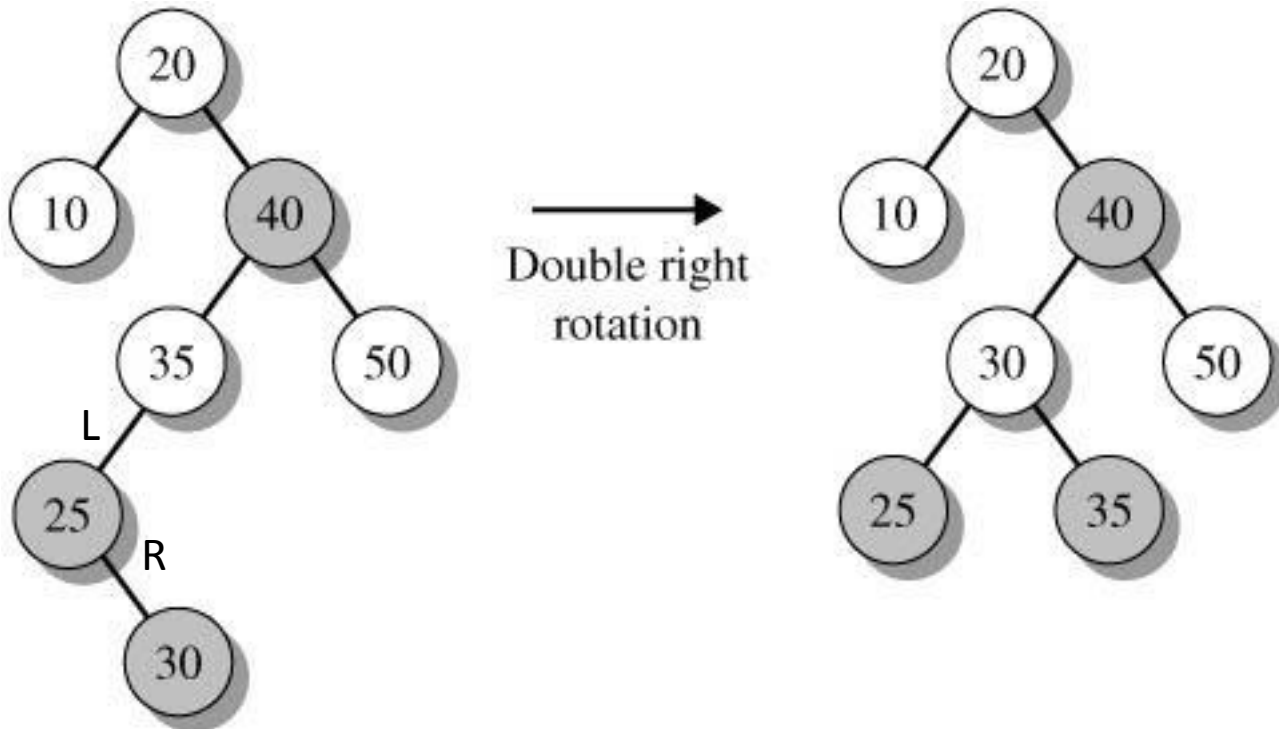


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Insert 30



LR = double right rotation of node 30
(left then right)