## Lab 2

## Functional dependencies and Normal forms

Student name: Ngô Văn Phương

Student ID: HE140784

Class: SE1431

## **EXERCISES**

1. Consider the relation scheme with attributes S (store), D (department), I (item), and M (manager), with functional dependencies  $SI \rightarrow D$  and  $SD \rightarrow M$ .

$$R[S,D,I,M]$$
  
 $F[SI \to D,SD \to M]$ 

a) Find all keys for SDIM.

$$L:S,I$$
 $M:D$ 
 $key:[S,I,D]$ 

b) Show that SDIM is in second normal form but not third normal form.

R is in 2NF because D and M is fully dependent on PK But R is not 3NF because M is not a prime attribute of R

2. Consider the relation scheme with attributes CITY, ST, and ZIP, which we here abbreviate C, S, and Z. We observed the dependencies CS → Z and Z → C. The decomposition of the relation scheme CSZ into SZ and CZ has a lossless join. Does this decomposition reserve dependencies?

$$R[C,S,Z]$$
 $F[CS \rightarrow Z,Z \rightarrow C]$ 
 $L:S$ 
 $M:C,Z$ 
 $key:[S,C] or[S,Z]$ 
 $R1[SZ]$  PK: $[SZ]$ 
 $R2[CZ]$  PK: $[Z]$ 

R1 are  $PK \Rightarrow$  the decomposition preserve dependencies

3. Let 
$$F = \{AB \rightarrow C, A \rightarrow D, BD \rightarrow C\}$$
.

a) Find a minimal cover for F.

$$A \rightarrow D, BD \rightarrow C \Rightarrow AB \rightarrow C$$
 (pseudo transitivity)  
 $G[A \rightarrow D, BD \rightarrow C]$ 

b) Give a 3NF, dependency-preserving decomposition of ABCD into only two schemes (with respect to the set of functional dependencies F).

$$R1[A,D]$$
  $F1[A \rightarrow D]$   
 $R2[B,C,D]$   $F2[BD \rightarrow C]$ 

c) What are the projected dependencies for each of your schemes?

$$F1 \mid A \rightarrow D \mid F2 \mid BD \rightarrow C \mid$$

d) Does your answer to (a) have a lossless join? If not, how could you modify the database scheme to have a lossless join and still preserve dependencies?

(a) doesn't have lossless join because it have no key [A, B, D]

$$R1[A,D]$$
  $F1[A \rightarrow D]$   
 $R2[B,C,D]$   $F2[BD \rightarrow C]$   
 $R3[A,B,D]$ 

4. Let 
$$F = \{AB \rightarrow C, A \rightarrow B\}$$
.

a) Find a minimal cover for F.

$$F[AB \to C, A \to B]$$

$$A \to B, AB \to C \Rightarrow A \to C \text{ (pseudo transitivity)}$$

$$G[A \to B, A \to C]$$

b) When (a) was given on an exam at a large western university, more than half the class answered  $G = \{A \rightarrow B, B \rightarrow C\}$ . Show that answer is wrong by giving a relation

that satisfies F but violates G.

$$G[A \rightarrow B, A \rightarrow C]$$

$$A \rightarrow C \Rightarrow AB \rightarrow C$$

G keep all dependencies of F but F have nothing to find  $B \rightarrow C$ 

- 5. Suppose we are given relation scheme ABCD with functional dependencies (A  $\rightarrow$  B, B  $\rightarrow$  C, A  $\rightarrow$  D, D  $\rightarrow$  C}. Let p be the decomposition (AB,AC,BD).
- a) Find the projected dependencies for each of the relation schemes of p.

$$p[AB, AC, BD]$$
 $R1[A, B]$ 
 $F1[A \rightarrow B]$ 
 $R2[A, C]$ 
 $F2[A \rightarrow C]$ 
 $R3[B, D]$ 
 $F3[$ 

b) Does p preserve the given dependencies?

p does not preserve the given dependencies because  $F1 \cup F2 \cup F3 \neq F$ 

6. Consider the relation scheme ABCD with dependencies

$$F = \{A \rightarrow B, B \rightarrow C, D \rightarrow B\}$$

We wish to find a lossless-join decomposition into BCNF.

a) Suppose we choose, as our first step, to decompose ABCD into ACD and BD. What are the projected dependencies in these two schemes?

$$R1[B,D]$$
  $F1[D \rightarrow B]$   
 $R2[A,C,D]$   $F2[AD \rightarrow C]$ 

b) Are these schemes in BNCF? If not, what further decomposition is necessary?

These schemes in BNCF because D is the super - key of R1 and AD are the super - key of R2